CS 611: Theory of Computation

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Emptiness Problem

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Solution: Check if the start symbol S is generating.

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Solution: Check if the start symbol S is generating. How long does that take?

Determining generating symbols

Algorithm

Gen = $\{\}$
for every rule $A \to x$ where $x \in \Sigma^*$
    Gen = Gen $\cup$ $\{A\}$
repeat
    for every rule $A \to \gamma$
        if all variables in γ are generating then
            Gen = Gen $\cup$ $\{A\}$
until Gen does not change

Determining generating symbols

Algorithm

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Gen = {}
for every rule A → x where x ∈ Σ*
    Gen = Gen ∪ {A}
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- Both for-loops take $O(n)$ time where $n = |G|$.

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- Both for-loops take $O(n)$ time where $n = |G|$.
- Each iteration of repeat-until loop discovers a new variable. So number of iterations is $O(n)$. And total is $O(n^2)$.

Membership Problem

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Central question in parsing.

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- We will see an algorithm that runs in $O(n^3)$ time (the constant will depend on k).

First Ideas

Notation

Suppose $w = w_1 w_2 \cdots w_n$, where $w_i \in \Sigma$. Let $w_{i,j}$ denote the substring of w starting at position i of length j. Thus, $w_{i,j} = w_i w_{i+1} \cdots w_{i+j-1}$

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How do we determine if $A \stackrel{*}{\Rightarrow} w_{i,j}$ for every A, i, j?

Base Case

Substrings of length 1

Observation

For any A, i, $A \stackrel{*}{\Rightarrow} w_{i,1}$ iff $A \rightarrow w_{i,1}$ is a rule.

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Thus, for each A and i, one can determine if $A \overset{*}{\Rightarrow} w_{i,1}$.

Inductive Step

Longer substrings

Suppose for every variable X and every $w_{i,\ell}$ ($\ell < j$) we have determined if $X \overset{*}{\Rightarrow} w_{i,\ell}$

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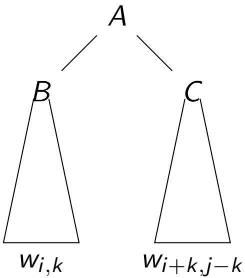

Suppose for every variable X and every $w_{i,\ell}$ ($\ell < j$) we have determined if $X \stackrel{*}{\Rightarrow} w_{i,\ell}$

- $A \stackrel{*}{\Rightarrow} w_{i,j}$ iff there are variables B and C and some $k < j$ such that $A \rightarrow BC$ is a rule, and $B \stackrel{*}{\Rightarrow} w_{i,k}$ and $C \stackrel{*}{\Rightarrow} w_{i+k,j-k}$

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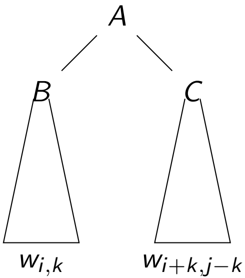

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- Since k and $j - k$ are both less than j, we can inductively determine if $A \stackrel{*}{\Rightarrow} w_{i,j}$.

Cocke-Younger-Kasami (CYK) Algorithm

Algorithm maintains $X_{i,j} = \{A \mid A \stackrel{*}{\Rightarrow} w_{i,j}\}$.

```
Initialize:  X_{i,1} = {A | A → w_{i,1}}
for j = 2 to n do
    for i = 1 to n − j + 1 do
        X_{i,j} = ∅
        for k = 1 to j − 1 do
            X_{i,j} = X_{i,j} ∪ {A | A → BC, B ∈ X_{i,k}, C ∈ X_{i+k,j−k}}
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`Initialize:` $X_{i,1} = \{A \mid A \rightarrow w_{i,1}\}$
for j `= 2 to` n **do**
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    $X_{i,j} = \emptyset$
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Correctness: After each iteration of the outermost loop, $X_{i,j}$ contains exactly the set of variables A that can derive $w_{i,j}$, for each i.Time $= O(n^3)$.

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Consider grammar
$S \rightarrow AB \mid BC,\ A \rightarrow BA \mid a,\ B \rightarrow CC \mid b,\ C \rightarrow AB \mid a$ Let $w = baaba$. The sets $X_{i,j} = \{A \mid A \overset{*}{\Rightarrow} w_{i,j}\}$:

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j/i	1	2	3	4	5
5					
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3					
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1	$\{B\}$	$\{A, C\}$	$\{A, C\}$	$\{B\}$	$\{A, C\}$
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All these problems are undecidable.