## Quiz 7

1. Consider the grammar  $G = (V = \{S, A, C, X, Y\}, \Sigma = \{a, b, c\}, R, S)$  where the set of rules R is as follows:

$$S \rightarrow AX|YC$$

$$A \rightarrow aA|\dot{\epsilon}$$

$$C \rightarrow cC | \epsilon$$

$$X \rightarrow bXc$$

$$\begin{array}{ccc} C & \to cC | \epsilon \\ X & \to bXc | \epsilon \\ Y & \to aYb | \epsilon \end{array}$$

Which of the following strings can be derived in **one** step from aaAbXc?

- (A) aaaAbbXcc
- (B) aaAbbXcc
- (C) aaAXbXc
- (D) Nothing because aaAbXc is not the start symbol of the grammar
- 2. Consider the grammar  $G = (V = \{S, A, C, X, Y\}, \Sigma = \{a, b, c\}, R, S)$  where the set of rules R is as follows:

$$S \rightarrow AX|YC$$

$$A \rightarrow aA|\epsilon$$

$$C \rightarrow cC | \epsilon$$

$$\begin{array}{ccc} C & \to cC \middle| \epsilon \\ X & \to bXc \middle| \epsilon \end{array}$$

$$Y \rightarrow aYb|\epsilon$$

Which of the following strings can be derived from S in zero or more steps?

- (A) aaba
- (B) aabbbc
- (C) aaAbXc
- (D) abbccc
- 3. Consider the grammar  $G = (V = \{S, A, C, X, Y\}, \Sigma = \{a, b, c\}, R, S)$  where the set of rules R is as follows:

$$\begin{array}{ccc} S & \to AX|YC \\ A & \to aA|\epsilon \end{array}$$

$$C \rightarrow cC | \epsilon$$

$$X \rightarrow bXc|\epsilon$$

$$Y \rightarrow aYb|\epsilon$$

The set of strings (over  $\Sigma$ ) derivable from A is

- (A)  $L(a^*)$
- (B) Strings with an even number of as
- (C)  $\emptyset$  because A is not the start symbol
- (D)  $\{a^n b^n \mid n \ge 0\}$

4. Consider the grammar  $G = (V = \{S, A, C, X, Y\}, \Sigma = \{a, b, c\}, R, S)$  where the set of rules R is as follows:

$$S \rightarrow AX|YC$$

$$A \rightarrow aA|\epsilon$$

$$C \rightarrow cC|\epsilon$$

$$X \rightarrow bXc|\epsilon$$

$$Y \rightarrow aYb|\epsilon$$

- $\mathbf{L}(G)$  is
- (A)  $\mathbf{L}(a^*b^*c^*)$
- (B)  $\{a^nb^nc^n \mid n \ge 0\}$
- (C)  $\{a^i b^j c^k \mid i = j \text{ or } j = k\}$
- (D)  $\{a^ib^jc^k \mid i=k\}$
- 5. Let  $G = (V, \Sigma, R, S)$  be a context-free grammar, where V is the set of variables,  $\Sigma$  is the set of terminals, R is the set of rules, and S is the start symbol. Which of the following is true about the language defined by G? Pick the most precise answer.
  - (A)  $\mathbf{L}(G) \subseteq V^*$
  - (B)  $\mathbf{L}(G) \subseteq \Sigma^*$
  - (C)  $\mathbf{L}(G) \subseteq (V \cup \Sigma)^*$
  - (D)  $\mathbf{L}(G) \subseteq V^* \cup \Sigma^*$
- 6. Consider the grammar  $G = (V = \{S, A, C, X, Y\}, \Sigma = \{a, b, c\}, R, S)$  where the set of rules R is as follows:

$$S \rightarrow AX|YC$$

$$A \rightarrow aA|\epsilon$$

$$C \rightarrow cC|\epsilon$$

$$X \rightarrow bXc|\epsilon$$

$$Y \rightarrow aYb|\epsilon$$

Which of the following statements is true about G?

- (A) G is ambiguous because there are at least two derivations from S producing abc.
- (B) G is ambiguous because there are at least two parse trees with root labelled S and yield abc.
- (C) G is not ambiguous because multiple derivations of abc from S does not imply ambiguity.
- (D) G may not be ambiguous because derivations and parse trees for a single string abc do not determine ambiguity.