

In-Class Assignment

- Given,

$M = (Q, \Sigma, \delta, q_0, F)$ recognizing L

We define

$N = (Q', \Sigma, \delta', q_0', F')$ that recognizes
 $'/2L$

where,

s is a new start state

$$q_0' = s$$

$\Sigma' = Q \times Q \times Q \cup \{s\}$ where $s \notin Q$

$$Q' = \{ \langle q_i, q_j, q_k \rangle \mid q_i, q_j, q_k \in Q \}$$

$$\delta'(s, \epsilon) = \{ \langle q_0, q_i, q_j \rangle \mid q_i, q_j \in Q \}$$

$$\delta'(\langle q_i, q_j, q_k \rangle, a) = \{ \langle q_i, q_j, q_m \rangle \mid$$

$$\delta(q_i, a) = q_l, \exists b \in \Sigma.$$

$$\delta(q_k, b) = q_m \}$$

$$F' = \{ \langle q_i, q_j, q_k \rangle \mid q_i \in Q, q_j \in F \}$$

- In ~~$\langle q_i, q_j, q_k \rangle$~~ $\langle q_i, q_j, q_k \rangle$

q_i - left finger

q_j - initial guess for right finger

q_k - right finger

We move left finger based on input

- We move left finger from x

We read from x

Initial guess remains the same through out the process

Moving Right Finger:

Every time we read an input from x & change left finger, we will generate all possible reads for state on right

fingers & create multiple new right threads. We say that we accept x if by the end of execution of x , at least one thread has ~~one~~ right finger on one of final states.

- Initial Guess for Right Finger:
We will create ' n ' initial threads where ' n ' is number of states in DFA with right fingers on each of the states. If by the end of execution of x , at least one of ' n ' threads has right fingers on one of final states, then that input is accepted.