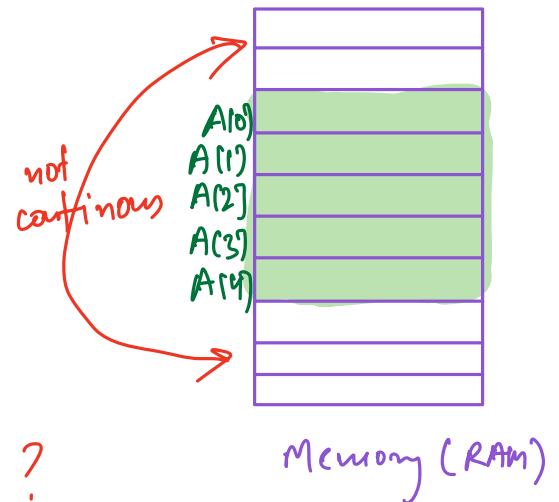


Linked list Basics

Array : $O(1)$ random access
 $A[i] \rightarrow O(1)$ TC

$\text{int } A[5] \rightarrow$ how it will be stored?



can you store $B[4]$ in the memory?

\rightarrow Can't create since we don't have
contiguous locations. [Memory fragmentation]

	At Start	At end	At K^{th} position (random)
Insertion	$O(N)$	$O(1)$	$O(N)$
Deletion	$O(N)$	$O(1)$	$O(N)$

3 \rightarrow *

1	2				
---	---	--	--	--	--

Insert at start

\Rightarrow

3	1	2			
---	---	---	--	--	--

shift existing elements
to right & then insert

Linked List

class Node {

int data;

Node next;

Node(int val) {

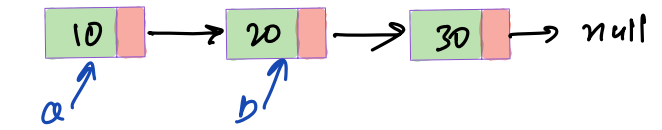
self.data = val

self.next = null

}

}

head ← variable name



Node a = new Node(10)

Node b = new Node(20)

a.next = b

b.next = new Node(30)

Create a list of 4 values → A(4)

Node head = new Node(A[0])

head.next = new Node(A[1])

head.next.next = new Node(A[2])

head.next.next.next = new Node(A[3])

Print

head.data

head.next.data

head.next.next.data

⋮
⋮
⋮

Iterate over linked list

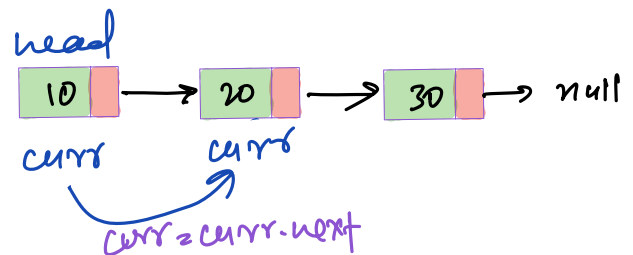
Node curr = head

while (curr != null) {

print(curr.data)

curr = curr.next

}



⇒ 10 20 30

Create LL using loop

```
Node head = new Node(A[0])
```

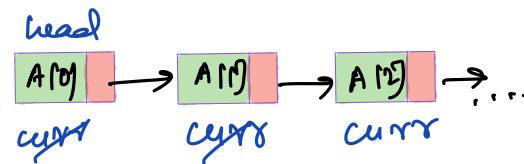
```
curr = head
```

```
for (i=1; i<n; ++i) {
```

```
    curr.next = new Node(A[i])
```

```
    curr = curr.next
```

```
}
```



Insert At Start

```
Node insertAtStart(head, val) {
```

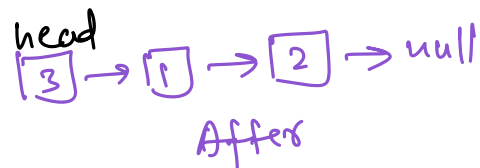
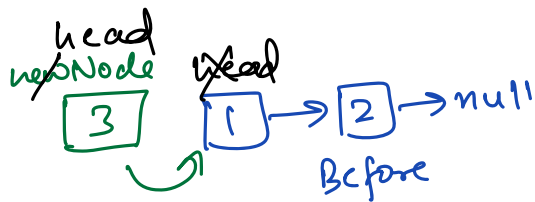
```
    Node newNode = new Node(val)
```

```
    newNode.next = head
```

```
    head = newNode
```

```
    return head
```

```
}
```



TC: O(1)

Insert at end

```
Node insertAtEnd(head, val) {
```

```
    Node newNode = new Node(val)
```

```
    if (head == null)
        return newNode
```

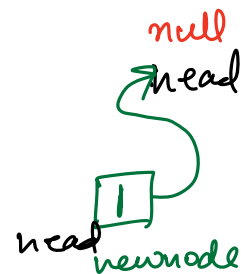
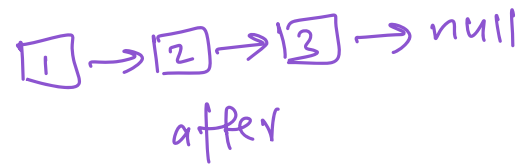
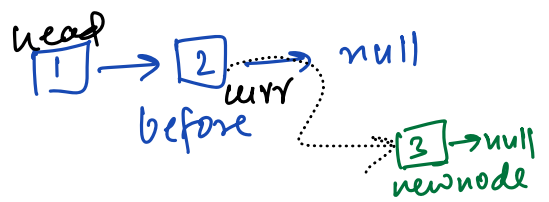
```
    Node curr = head
    while (curr.next != null) {
```

```
        curr = curr.next
    }
```

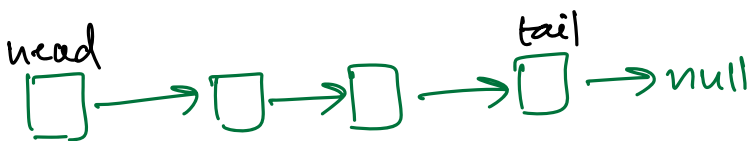
```
    curr.next = newNode
```

```
    return head
```

```
}
```

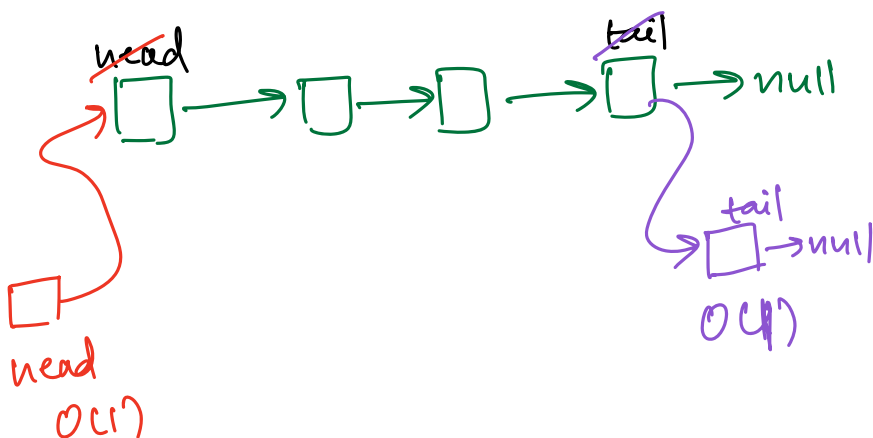


TC: $O(N)$

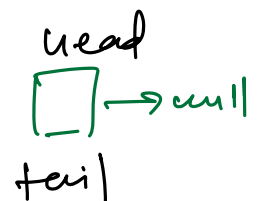


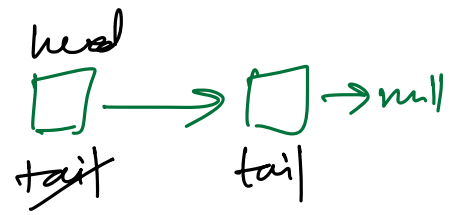
Using tail reference we can get $O(1)$

```
tail.next = newNode
tail = newNode
```

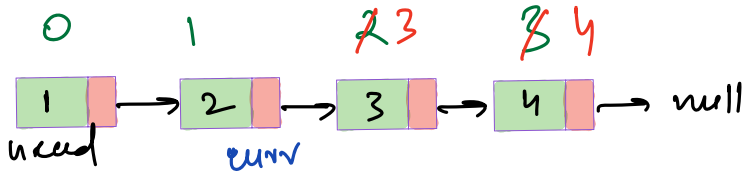


head = null
tail = null





Insert at K^{th} position (random)



insert at $K=2$

to reach curr, iterate $K-1$ times

Node insertAtK(node, val, K) {

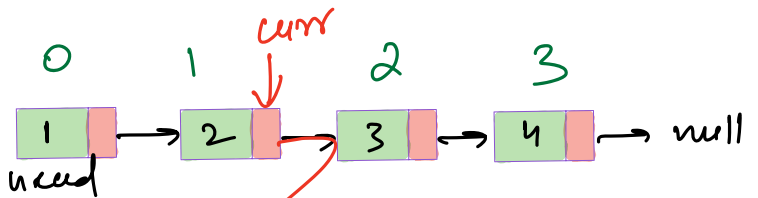
newNode = new Node(val)

curr = head

for(i=0; i<K-1; i++) {

curr = curr.next

}

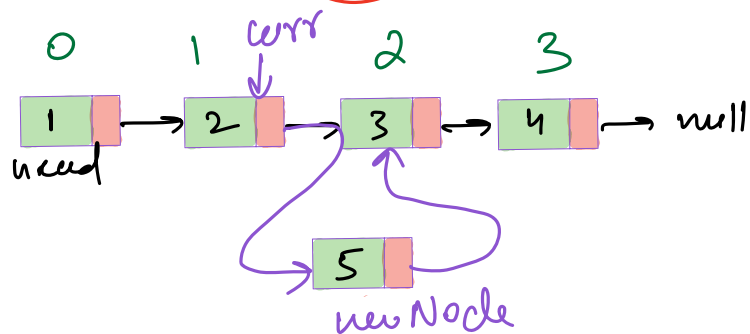


* { curr.next = newNode
newNode.next = curr.next

✓ { newNode.next = curr.next
curr.next = newNode

return head

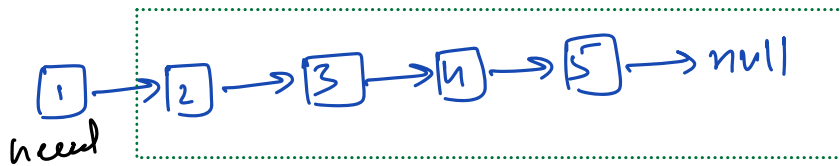
}



TC: $O(K) \Rightarrow O(N)$

Question

Print LL in reverse order.



Print (head)

main question



5 4 3 2 1

Print (head.next)

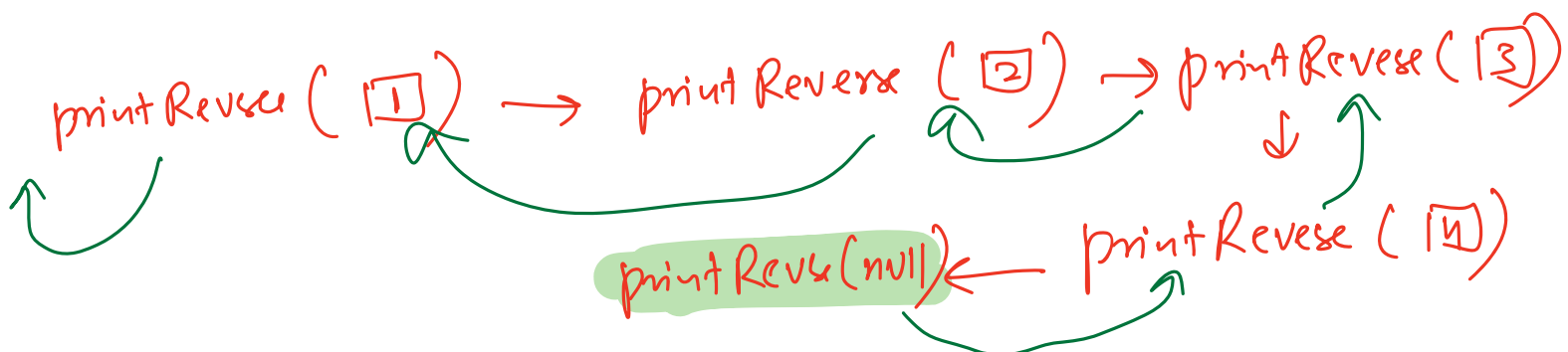


5 4 3 2

```
void printReverse (Node curr) {  
    if (curr == null) { return }  
    printReverse (curr.next)  
    print (curr.data)  
}
```

printReverse (head)

3



Print: 4 3 2 1

Recommendations

1. Write code on paper
2. Dry run
3. **Never use** "hit & trial"
4. Edge Cases →
 - LL is null? (head = null)
 - size of LL = 1?
 - size = 2/3?
 - problem specific