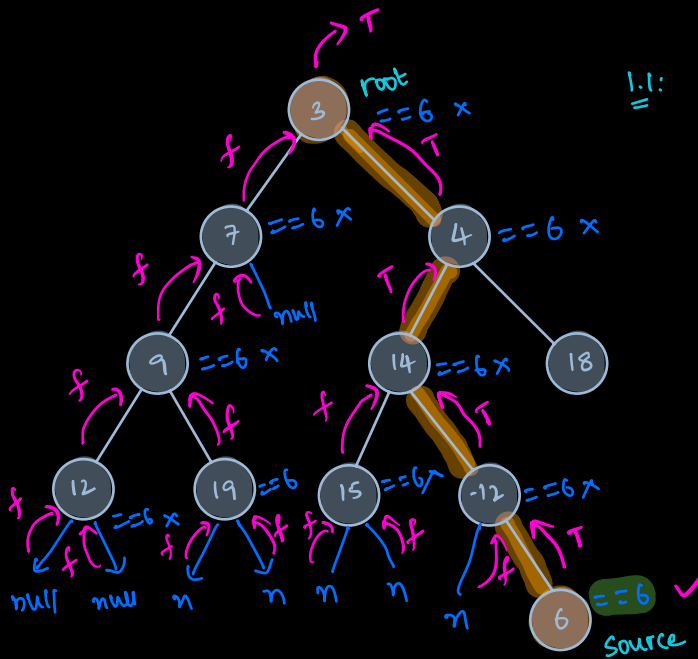


Q1. Given a binary tree with distinct values, find the path from root to source node.



1.1: Check if a given node is present in the tree or not.

```
bool check(Node root, int k)
{
    if(root == null) {return false;}
    if(root.data == k) {return true;}
    if(check(root.left, k) ||
        check(root.right, k))
        return true;
    return false;
}
```

Observation: If we store all the nodes that returned true, we will get path.

List<Node> path;

path : [6, 12, 14, 4, 3]

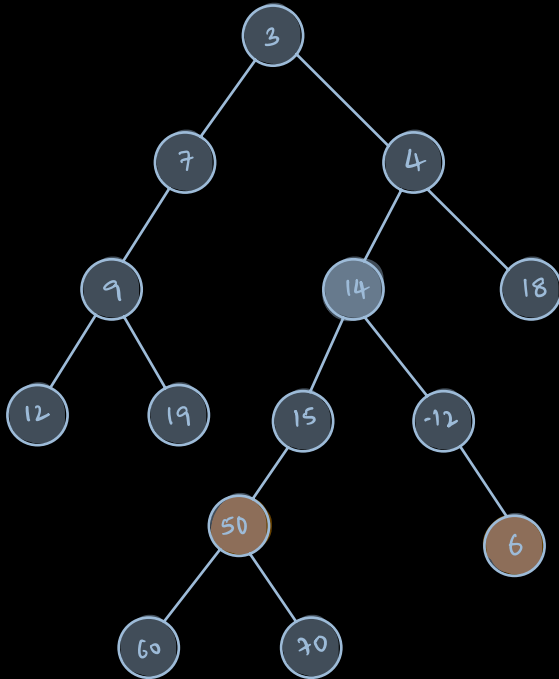
bool check(Node root, int k)

```
{
    if(root == null) {return false;}
    if(root.data == k) { path.add(root); return true; }
    if(check(root.left, k) ||
        check(root.right, k))
        path.add(root);
        return true;
    return false;
}
```

TC: $O(n)$
SC: $O(H)$

// reverse path if required.

Q2: LCA , least common ancestor.



$$LCA(14, 70) =$$

$$LCA(50, 6) = 14.$$

$$LCA(60, 70) = 50$$

$$LCA(19, 18) = 3.$$

→ path from root to 50 → [3, 4, 14, 15, 50]

path from root to 6 → [3, 4, 14, -12, 6]

[3, 4, 14, 15, 50]

[3, 4, 14, -12, 6]

ans.

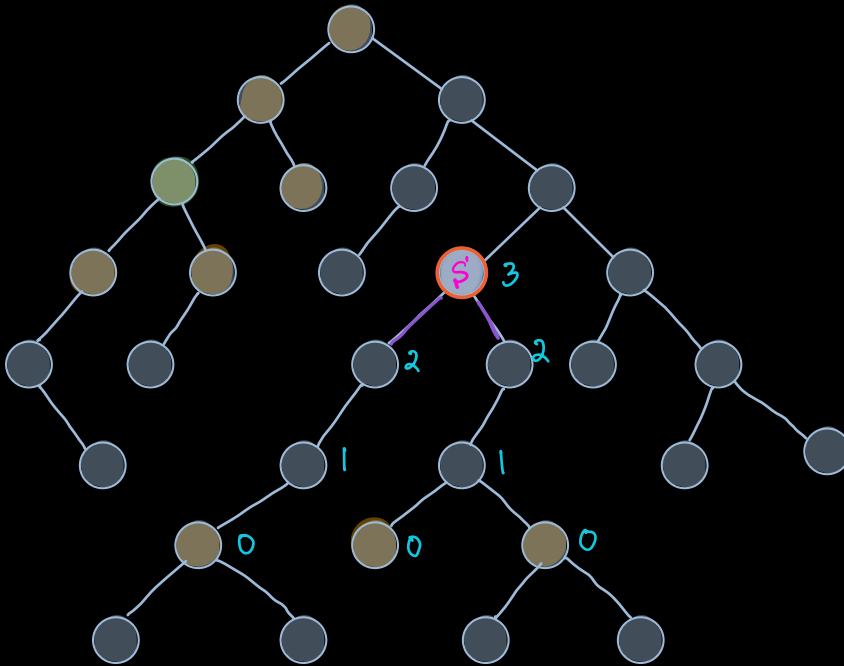
$$LCA(12, 18)$$

path from root to 12 → [3, 7, 9, 12]

path from root to 18 → [3, 4, 18]

38: Given a source node, find the no. nodes at distance 'k' from source node.

(All the nodes should be below source node).



$k = 3, ans = 3.$

idea: find source and
= apply level order
traversal.

idea 2:

distance between source and node = k .

distance between source child
-ren and node : $(k-1)$

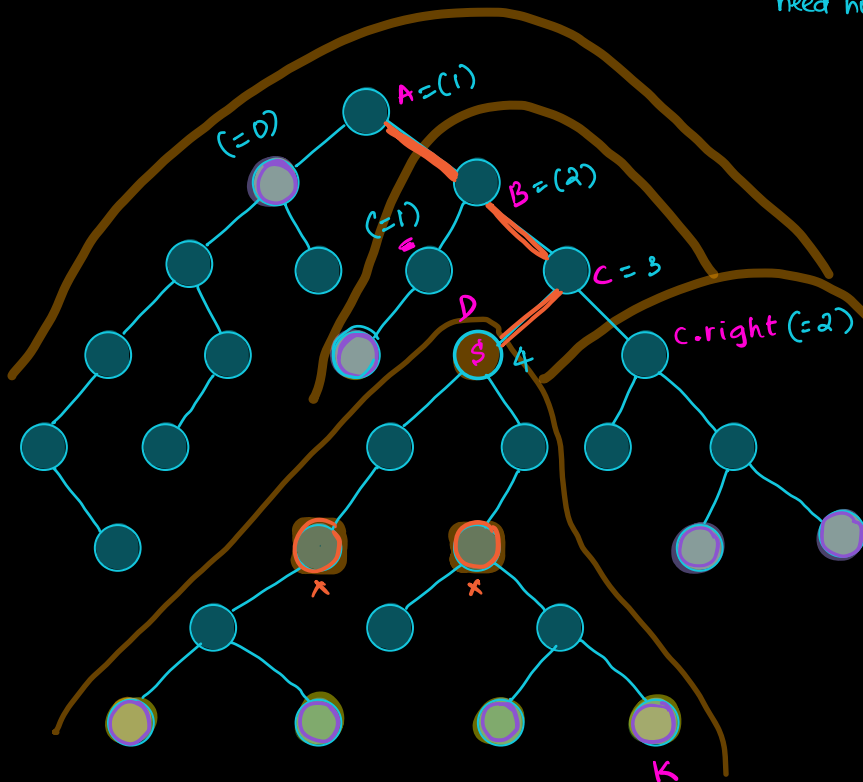
```

int countNodesBelow(Node s, int k)
{
    if (s == null || k < 0)
        return 0;
    if (k == 0)
        return 1;
    return countNodesBelow(s.left, k-1)
           + countNodesBelow(s.right, k-1);
}

```

49. Given a source node, find the no. of nodes at distance 'k' from source node.

↓ need not be just below.



$k=4$, ans=8.

path from source to root

0 1 2 3
node: [D C B A]
= $\rightarrow \rightarrow \rightarrow$

path[i].left = path[0]
C.left D.

path[2].left = path[1]
B.left C

i) Get the path from root to source \rightarrow path (array of nodes).

ii) Iterate on path array.

path[0]	path[1]	path[2]	path[3]
D	C	B	A
$k=4$, below(D, k)	$\swarrow \searrow$ D C.right path[1].left = path[0] C D. \Rightarrow below(C.right, k-2)	$\swarrow \searrow$ O C path[2].left = path[1] B C \Rightarrow below(B.left, k-3)	$\swarrow \searrow$ O B path[3].left = path[2] A B \Rightarrow below(A.left, k-4)

$$\begin{array}{l}
 1 \rightarrow k-2 \\
 2 \rightarrow k-3 \\
 3 \rightarrow k-4
 \end{array}
 \left. \vphantom{\begin{array}{l} 1 \rightarrow k-2 \\ 2 \rightarrow k-3 \\ 3 \rightarrow k-4 \end{array}} \right\} \Rightarrow i \rightarrow k-i-1$$

```

int countNodesAtKDistance(Node r, Node s, int k)
{
    List<Node> path; // To-Do.

    int c = countNodesBelow(path[0], k);

    int n = path.length();
    for (i = 1; i < n; i++)
    {
        // We're at the node path[i].
        // from path[i] we need to know whether to go left or right?
        if (k - i - 1 < 0) // happens when path length is more than k.
            break;
        if (path[i].left == path[i-1])
        {
            // I'm coming from left, look on right
            c = c + countNodesBelow(path[i].right, k - i - 1);
        }
        else
        {
            // I'm coming from right, look on left.
            c = c + countNodesBelow(path[i].left, k - i - 1);
        }
    }
    return c;
}

```

TC: $O(n)$ (\because We're visiting every node only once)

SC: $O(H)$

If $k=1$, how do we get $ans=3$.