Hashing-2

Content

→ Pair sum = K

→ Distinct elements in every window of len=K

Question 1

Given N array elements, check if there exists a pair (i,j) such that a[i] + a[j] = K and i != j

(K is given)

→ return true/false

Eg a[] = 8 9 1 -2 4 5 11 -6 7 5 (indices 0 1 2 3 4 5 6 7 8 9)

K = 11 ⇒ a[4] + a[8] = 4 + 7 = 11 True

K = 6 ⇒ a[2] + a[5] = 1 + 5 = 6 True

a[0] + a[3] = 8 - 2 = 6

K = 22 ⇒ a[6] + a[6] = 11 + 11 = 22 False

i == j ✗

Idea 1

Check all pairs sum = K

```
for (i=0; i<n; ++i) {
    a = a[i]      a+b = K      b = K-a = K-a[i]
    for (j=i+1, j<n; ++j) {
        if (a[i] + a[j] == K)  ⇒  if (a[j] == b)
            return true
    }
}
return false
```

TC : $O(N^2)$

SC : $O(1)$

Idea 2 : Use Hashset

hs : { 8 9 1 -2 4 5 11 -6 7 }

K = 11 ⇒ a+b = 11

a	b = K-a	check b is present in hashset or not ?
8	3	NO
9	2	NO
1	10	NO
-2	13	NO
4	7	YES { return true }

a+b = 5

a	b	present or not ?
8	-3	NO
9	-4	NO
1	4	YES { return true }

a + b = -4

a	b	present or not ?
8	-12	NO
9	-13	NO
1	-5	NO
-2	-2	YES { return true } ✗ wrong

Basically if a = b then occurence of a should be more than 1.

Note: frequency of elements are important to know.

Idea3 : Hashmap

arr = 8 9 -2 4 5 11 -6 7 5

hm = { <8,1> <9,1> <7,1> <-2,1> <4,1>
<5,2> <11,1> <-6,1> }

a + b = 10

a	b	present or NOT ?
8	2	NO
9	1	NO
-2	12	NO
4	6	NO
5	5	if(a==b & freq(a) > 1) { return true }

Code

```
bool pairSum (a[], K) {
    HashMap<int, int> hm
    Insert a[] → hm      // TODO → O(N)

    for (i=0; i<n; ++i) {
        a = a[i],  b = K - a[i]
        if (hm.search(b) == true) {
            if (a != b)
                return true
            if (hm[b] > 1)
                return true
        }
    }
    return false
}
```

T.C: O(N)
S.C: O(N)

Idea 4: Use HashSet again

at its index, hs will contain all elements from [0, i-1]

	0	1	2	3	4	5	6	7	8	9
a[] =	8	9	5	-2	11	5	7	-6	4	1

K = 22

a	b	ls	present or NOT ?
8	14	{}	NO
9	13	{8}	NO
5	17	{8,9}	NO
-2	24	{8,9,5}	NO
11	11	{8,9,5,-2}	NO → now code is working for a=b & freq(a)=1
		{8,9,5,-2,11}	

K = 10

a	b	ls	present or NOT
8	2	{}	NO
9	1	{8}	NO
5	5	{8,9}	NO
-2	12	{8,9,5}	NO
11	-1	{8,9,5,-2}	NO
5	5	{8,9,5,-2,11}	YES {return true}

Code

```
bool pairSum ( a[] , K ) {
    Hashset<int> hs
    for (i=0; i<n; ++i) {
        a = a[i] ,  b = K - a[i]
        if ( hs.search(b) == true)
            return true
        hs.insert(a)
    }
    return false
}
```

TC: O(N)

SC: O(N)

Question 2

Given N elements, calculate no. of distinct elements in every subarray of size K.

	0	1	2	3	4	5	6	7	8	9
eg a[10] =	2	4	3	8	3	9	4	9	4	10

K = 4

subarrays	distinct elements
[0, 3]	4
[1, 4]	3

[2,5]
[3,6]
[4,7]
[5,8]
[6,5]

3
4
3
2
3

⇒ Ans

Idea!: for every subarray of len = K insert into hashset & find size.

Code

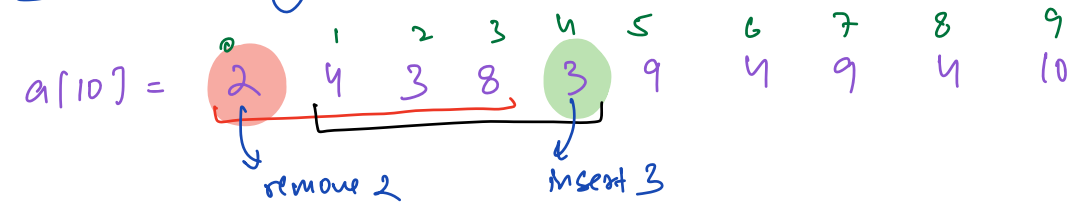
```
for (i=0; i<n-K; ++i) {          → (n-K+1) iterations
    HashSet<int) hs
    for (j=i; j<i+K; ++j) {      → K iterations
        hs.insert(a[j])
    }
    print(hs.size)
}
```

TC: $O((n-K+1)*K)$ ⇒ $O(N^2)$ SC: $O(K)$

1. $K=n$ → TC: $O((n-n+1)*n) = O(N)$
2. $K=1$ → TC: $O((n-1+1)*1) = O(N)$
3. $K=\frac{N}{2}$ → TC: $O\left(\left(n-\frac{n}{2}+1\right)*\frac{n}{2}\right) = O\left(\frac{N^2}{n}\right)$

$= O(N^2)$

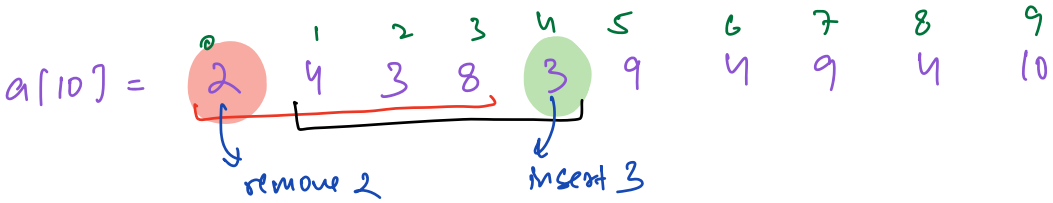
Idea 2 : Sliding Window using Hashset



		Hashset	hs.size
[0,3]	→	{2,4,3,8}	4
[1,4]	remove a[0], insert a[4]	{4,3,8}	3
[2,5]	remove a[1], insert a[5]	{3,8,9}	3
[3,6]	remove a[2], insert a[6]	{8,9,4}	3 ✗

In hashset we don't know the frequency of inserted elements. If we insert 2 times and remove 1 time, there should be 1 element in the hashset but we can't do it.

Idea 3 : Sliding window using Hashmap



s e		Hashmap	Size
[0,3]	⟶	{<2,1> <4,1> <3,1> <8,1>}	4
[1,4)	remove a[0], add a[4]	{<4,1> <3,2> <8,1>}	3
[2,5)	remove a[1], add a[5]	{<3,2> <8,1> <9,1>}	3
[3,6)	remove a[2], add a[6]	{<3,1> <8,1> <9,1> <4,1>}	4

Code

```
def distinctCount(a[], K) {
    n = a.length
    Hashmap<int, int> hm
    for (i=0; i<K; ++i) {
        if(hm.search(a[i]) == true)
            hm[a[i]]++
        else
            hm.insert({a[i], 1})
    }
    print(hm.size)  → for first subarray [0, K-1]

    s=1, e=K
```

```
while (e<n) {
  // subarray : [s, e]  →  remove a[s-1] , add a[e]

  hm[a[s-1]] --
  if(hm[a[s-1]] == 0)
      hm.remove(a[s-1])          → remove

  if( hm.search(a[e]) == true )
      hm[a[e]]++
  else
      hm.insert({a[e], 1})       → insert

  print(hm.size)

  s++, e++
}
}
```

TC : O(N)

SC : O(K)