

Today's Content

- * Length of longest substring with all distinct chars.
- * Anagrams Check.
- * Count of permutation of A in B.
- * Class of Anagrams.

Q1: Given string S. Find length of longest substring without repeating characters. $1 \leq |S| \leq 10^6$

Eg: s = a b c a b c b b - 3

s = a a b b - 2

s = a b c a b c d - 4

BF approach: Generate all possible substrings & check if it has all distinct chars. $\hookrightarrow \frac{N(N+1)}{2} * N \Rightarrow O(N^3)$

idea2: With every $s[i]$ as start index of substring, get max length window which contains all distinct chars.

	0	1	2	3	4	5	6	7
s =	a	b	a	c	e	a	f	f
	↑	↑	↑	↑	↑	↑	↑	↑
	2	4	3	4	3	2	1	1

ans = 0

Tc: $O(n^2)$

Sc: $O(n)$

```
for (i = 0; i < n; i++) {
```

```
    HashSet<char> hs // hs = set()
```

```
    for (j = i; j < n; j++) {
```

```
        if (hs.search(s[j]) == false) {
```

```
            hs.add(s[j])
```

```
        }
```

```
        else { break }
```

```
    }
```

```
    ans = max(ans, hs.size())
```

```
}
```

idea 3: Use 2 pointers approach

$$s = \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ a & b & a & c & e & a & f & g & f \end{matrix}$$

ans = 15

<u>l</u>	<u>r</u>	<u>str</u>	<u>s[r] in ts?</u>	<u>ts</u>
0	0	a	x	{ a b a c e a f g }
0	1	ab	x	
0	2	a ba		
1	2	b		
1	3	bac		
1	4	bace		
1	5	bace a		
2	5	ace a		
3	5	cea		
3	6	cea f		
3	7	cea fg		
3	8	cea fgf		
4	8	ea fgf		

∴ Variable Sliding Window

Pseudocode :

```
int longestSub(String S) {  
    HashSet <char> hs  
    ans = 0  
    l = 0, r = 0  
    while (r < S.length) {  
        if (hs.find(S[r]) == false) {  
            hs.add(S[r])  
            r++  
            ans = max(ans, hs.size())  
        }  
        else {  
            hs.delete(S[l]) // repeating char found!  
            l++  
        }  
    }  
    return ans  
}
```

TC: $O(n)$
SC: $O(n)$

idea 4: Binary Search

a) Target : Length of longest substring
b) Search space: $\frac{l}{1}$ to $\frac{r}{N}$

String with 10 char $\frac{l}{1}$ $\frac{r}{10}$ $\frac{m}{5}$

* is it possible to have substring of len = 5 with distinct chars?
↳ Sliding window! $\Rightarrow O(n)$

Pattern: 1 2 3 4 5 N
T T T T T F

TC: $O(N \log N)$, SC: $O(N)$

look for last
occurrence of True.

Qn: Given 2 strings A & B. Check if they are anagrams of each other.

Anagrams: 2 strings are said to be anagrams of each other if they are permutations of each other.

Eg:

tac	cat	✓
satya	yta	✗
lps	pls	✓
madam	dadam	✗

idea 1: Sort + Compare
TC: $O(n \log n + n)$

idea: Frequency of all chars in HM should be same

Step 1: Insert all chars of A in hm1 } $O(n)$

Step 2: Insert all chars of B in hm2 }

Step 3: Iterate over hashmap & compare if A & B have same chars or not. $\Rightarrow O(n)$

```
bool compare (HashMap <char, int> hm1, hm2) {
```

```
    if (hm1.size() == hm2.size()) {
```

```
        for (char k : hm1.keySet()) {
```

```
            if (hm2.find(k) == false ||
```

```
                hm1[k] != hm2[k]) {
```

```
                return false
```

```
            }
```

```
        }
```

```
        return true
```

```
    }
```

```
    return false
```

```
}
```

JAVA

TC: $O(n)$

PYTHON

```
for k, v in hm1.items():  
    if k not in hm2 or hm2[k] != v:  
        return False
```

$N \quad M \quad (N \geq M)$

Qn: Given 2 strings A & B.
{ lowercase alphabets }

Count no. of permutations of B in A
↳ no. of substrings in A which are
anagrams of B.

Eg: A : a b c a c c b a b c a c
(12)

B : a b a c c
(5)

idea1: For all substrings of len 5 in A, check if they are anagrams
of B.

	A	B
0-4	a b c a c	a b a c c
1-5	b c a c c	a b a c c
2-6	c a c c b	a b a c c
3-7	a c c b a	a b a c c
4-8	c c b a b	a b a c c
5-9	c b a b c	a b a c c
6-10	b a b c a	a b a c c
7-11	a b c a c	a b a c c

Ans = 3

TC: $O(N * M)$
SC: $O(1)$

idea 2: Use sliding window

0 1 2 3 4 5 6 7 8 9 10 11
 A : a b c a c c b a b c a c
 (12)

B : a b a c c
 (5) HB
 $\{a:2, b:1, c:2\}$

	Add	Remove	HA:	HA == HB
0-4			$\{a:2, b:1, c:2\}$	✓
1-5	5 (c)	0 (a)	$\{a:1, b:1, c:3\}$	✗
2-6	6 (b)	1 (b)	$\{a:1, b:1, c:3\}$	✗
3-7	7 (a)	2 (c)	$\{a:2, b:1, c:2\}$	✓
4-8	8 (b)	3 (a)	$\{a:1, b:2, c:2\}$	✗
5-9	9 (c)	4 (c)	$\{a:1, b:2, c:2\}$	✗
6-10	10 (a)	5 (c)	$\{a:2, b:2, c:1\}$	✗
7-11	11 (c)	6 (b)	$\{a:2, b:1, c:2\}$	✓

Steps 1: a) Insert all chars of B in hm2 ans = 3

b) Insert FIRST B chars of A in hm1 & compare

c) $\text{for}(i = B.length; i < A.length; i++) \{$
 $c = A[i]$ // add

$r = A[i - B.length]$ // remove

$hm1[c]++$

$hm1[r]--$

$\text{if}(hm1[r] == 0) \{$ // freq = 0, remove it
 $hm1.delete[r]$
 $\}$

$\text{if}(\text{compare}(hm1, hm2)) \{ ans++ \}$

$\}$

return ans

TC: $O(N = \text{length of } A)$

SC: $O(26) \Rightarrow O(1)$

Qn: Given an array of strings $A[N]$. Return all groups of strings that are anagrams. Represent the group by 1-based indexing.

Eg:- $A = \begin{matrix} & 1 & 2 & 3 & 4 \\ \{ & \text{cat} & \text{dog} & \text{god} & \text{tac} \end{matrix}$
 $\text{ans} = \begin{matrix} \{ & [1, 4] \\ & [2, 3] \end{matrix}$

$A = \begin{matrix} & 1 & 2 & 3 \\ \{ & \text{rat} & \text{tar} & \text{art} \end{matrix}$
 $\text{ans} = \{ [1, 2, 3] \}$

$A = \begin{matrix} & 1 & 2 & 3 & 4 \\ \{ & \text{dog} & \text{cat} & \text{bat} & \text{rat} \end{matrix}$
 $\text{ans} = \begin{matrix} \{ & [1] \\ & [2] \\ & [3] \\ & [4] \end{matrix}$

idea: Sort the strings within the array & add them to $\text{HashMap} \langle \text{String}, \text{List} \langle \text{int} \rangle \rangle$

Iterate over the strings & compare with hashmap
 if key found: add to ans
 else: continue

$A = \begin{matrix} 0 & 1 & 2 & 3 \\ \{ & \text{cat} & \text{dog} & \text{god} & \text{tac} \end{matrix}$
 $\hookrightarrow \text{Sort: } \{ \text{act}, \text{dgo}, \text{dgo}, \text{act} \}$

~~$\{ \text{act}: [1, 4] \}$~~
 ~~$\text{dgo}: [2, 3] \}$~~

$\text{ans} = \{ [1, 4], [2, 3] \}$

Pseudocode :

anagrams (String A[N]) {

HashMap <String, List <int>> hm

for (i=0; i < N; i++) {

Sort (A[i])

$N * (M \log M)$

hm [A[i]].add(i+1) // 1 based indexing

}

List <List <int>> ans

for (i=0; i < N; i++) { — $O(N)$

if (hm.find(A[i]) {

ans.add(hm[A[i]])

hm.remove(A[i])

}

}

return ans

}

TC: $O(NM \log M)$

SC: $O(N * M + N) = O(N * M)$

N idx as val
being stored.