



Course code: ICT-4101

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### Submitted by

Name: MD. Amirul Hasan Shanto

ID: IT17027

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Dept. of ICT

MBSTU.

### Submitted To

Nazrul Islam

Assistant Professor

Dept. of ICT

MBSTU.

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1. a) What is Time Division Switching?  
b) Describe switching Network Design?  
c) Define TSI.
2. a) Define Time multiplexed space switch.  
b) Differentiate Router & switch.  
c) Describe simple network management protocol.
3. a) What is the necessity of QPGI message  
b) Describe both synchronous & asynchronous transmission.  
c) What is bridging?
4. a) What is link blocking?  
b) Estimate blocking for a pure distribution switch.  
c) How to solve the mixing stage?

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5.a) Describe the characterization of a Telephone Traffic.

b) How does the Erlang unit correspond to aces?

6.a) Justify Binomial Distribution Model.

b) Find probability of blocking through poission traffic model.

7.a) How does telecommunication network working?

b) Assume  $10^8$  number of devices covering total area of  $147570 \text{ km}^2$ . The radius of the wire is 2mm. Find the dept of wiring?

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8. a) How does a call get established?

b) What is pulse dialling?

c) Determine the cut-off frequency of 26 gauge cable loaded with 22mH inductance at 3000' spacing.  $r = 440 \Omega/\text{mile}$ ,  $\epsilon = 83 \text{nF}/\text{m}$

$$l = 1 \text{mH}, g = 0.2 \mu\text{s}.$$

Ans:  $\omega_c = \frac{1}{\sqrt{L(C + \frac{r}{2})}} = \frac{1}{\sqrt{1 \times 10^{-3} \times (83 \times 10^{-9} + \frac{440}{2 \times 1609.34} \times 10^6)}} = 1.57 \text{ rad/s}$

$\omega_c = 1.57 \text{ rad/s}$   
 $\omega_c = 1.57 \times 10^3 \text{ rad/s}$   
 $\omega_c = 1.57 \times 10^3 \times \frac{1}{2\pi} \text{ Hz}$   
 $\omega_c = 250 \text{ Hz}$

(d) L. OH resistance of the ground

Ans:  $R_L = \frac{V_o}{I_o} = \frac{10}{10^{-12}} = 10^{13} \Omega$

Answer to the Question NO: 1(a)

In digital time division switching system are divided by time slot but switching is still possible.

Switching is by a time-slot interchanger & is accomplished by rearranging the orders in which data is read out of the buffer. Incoming data enters a speech store while the outgoing channels indicate to the speech address memory which incoming timeslot it is assigned to.

During each time-slot, the outgoing circuit reads the speech store slot corresponding to the SAM.

Answer to the Question NO: 1(b)

For switching network design we focus on:

- i) Blocking yes versus non-blocking switches.
- ii) Number of cross-points.

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iii) Reliability

iv) overload

v) Growth

vi) Cost & technology

For trunk switch:

- i) One-to-one connection.
- ii) One specific must connect to one specific outlet.

Access switch:

- i) One-to-any connection
- ii) One specific inlet must connect to any free outlet.

Answer to the Question No: 1(e)

TSI is a time switch, switches one time slot channel on a single physical output.

Functionally equivalent to an  $n \times n$  space-divided switches where  $n$  is the number of time slot per frame

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### Answer to the Question No: 2(a)

A time multiplexed space switch is a space switch (with multiple physical inputs & outputs) that is potentially reconfigured entirely in every time slot at each frame.

(b)

Answer to the Question No: 2 (b)

Differences between routers & switches are follows:

Router	Switch
1. Appears in Network layers.	1. Appears in Data link layers.
2. PDU is Packet.	2. PDU is Frame.
3. Every port has its own Broadcast domain.	3. Switch has only one Broadcast domain.
4. Bandwidth sharing is dynamic.	4. No option for bandwidth sharing.
5. Have memory.	5. Switch has no memory.
6. Store IP address in the routing table.	6. Store MAC address in the lookup table.

Answer to the Question NO: 2(c)

Simple network management protocol is one of the UDP-based network protocols. It's monitoring network-attached devices for various administrative attentions. It has application layer - database schema, protocol & a group of data objects.

Answer to the Question NO: 3(a)

CPGI message is sent by the switch towards originator. When the switch receives the setup, the call is forwarded. After that the call is progressing. During the setup or active phase of call, a message

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is sent. This signifies the relayed to the originating has occurred.

Answer to the Question No: 3(b)

Synchronous transmission:

1. Synchronous transmission does not use start & stop bits, hence data transfer rate is quicker.
2. It uses clock signals that are built at each component.
3. Continual stream of data is sent between two nodes.
4. This method uses check digit instead of parity bits.

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### Asynchronous transmission:

1. It uses the start & stop bits.
2. It used when data need to be sent data intermittently.
3. It recognize the second packet at the information.
4. The starts & stop bits are supposed to be opposite polarity.

### Answer to the Question NO: 3 (i)

Bridging is one of the forwarded techniques to use in packet switched networks. Bridge makes no assumption about the address location of the network it's stored in MAC address table.

### Answer to the Question No: 4(a)

Because of the single link between each module & the modules in the next stage, there's a possibility of blocking.

4x4 → 4x4  
2 link block.

4x4 → 4x4

4x4 → 4x4  
2 link block.

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### Answer to the Question NO: 4(b)

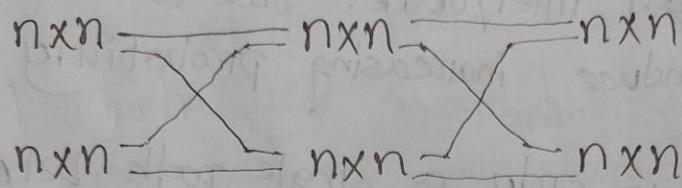
Estimate blocking: Distribution stages increase the overall inlet/outlet size at the switch but introduce increasing probability of blocking.

1. There is only a single path between any specific 1st stage inlet & any specific final stage outlet.
2. Mechanism of blockage is when an inter-stage link on required path is in use.
3. The greater the number of links in the path, the greater the probability that one of them is in use.
4. The more distribution stages we have, the greater the probability of blocking.

### Answer to the Question NO: 4(c)

The solution is to add a mixing stage that keeps the overall switch size the same.

but can reduce blocking by adding multiple paths through the switch.



### Answer to the Question NO: 5(a)

Characterization of telephone traffic:

Calling Rate: It is also called arrival rate, or attempts rate, etc.

1. Average number of calls initiated per unit time.
2. Each call arrival is independent of others call.
3. Call attempt arrivals are random in time.

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If receive  $\alpha$  calls from a terminal in time  $T$ :

$$\text{then } \gamma = \frac{\alpha}{T}$$

If receive  $\alpha$  calls from  $m$  terminals in time  $T$ :

 $\gamma$ :

$$\gamma_g = \frac{\alpha}{m \cdot T} \quad \downarrow \quad \gamma = \frac{\alpha}{m \cdot T}$$

Group calling rate      Perc terminal call rate.

Calling rate assumption:

Number of calls in time  $T$  is

Poisson distributed:

$$P(X) = \frac{e^{-\lambda} \cdot \lambda^X}{X!} \quad X = 0, 1, 2$$

$$\lambda = \gamma T$$

$$f(t) = \lambda \cdot e^{-\lambda t} \quad [0 \leq t \leq \infty]$$

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### Answer to the Question NO: 5(b)

From traffic intensity

$$A = Y \cdot h$$

For dimensionless Erlang unit

$$1 \text{ ccs/hour} = \frac{100 \text{ call. seconds}}{1 \text{ hr} \times 60 \text{ min} \times 60 \text{ sec}} = 0.027 E$$

$$\therefore 36 \text{ ccs/hour} = \frac{3600 \text{ call. seconds}}{1 \text{ hour} \times 60 \text{ min/hr} \times 60 \text{ sec/min}}$$

$$= 1 E$$

Percentage of time a terminal is busy

is equivalent to the traffic generated by

that terminal in Erlangs, or Average number of circuit in a group busy at any time.

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### Answer to the Question NO: 6(a)

Let's assumptions:

- m source
- A Erlangs of offered traffic
- per source  $T_0 = A/m$

probability that a specific source is busy:  $P(B) = A/m$

For Binomial distribution to give the probability that a certain number ( $K$ ) of those  $m$  sources is busy:

$$P(K) = \binom{m}{K} \left(\frac{A}{m}\right)^K \left(1 - \frac{A}{m}\right)^{m-K}$$

$$= \left(\frac{m!}{K!(m-K)!}\right) \left(\frac{A}{m}\right)^K \left(1 - \frac{A}{m}\right)^{m-K}$$

For ~~N~~  $N$  serves: ( $N < m$ )

$$P(B) = P(k \geq N) = P(k=N) + P(k=N+1) + \dots + P(k=n)$$

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$$P(B) = \sum_{k=N}^m \binom{m}{k} \left(\frac{A}{m}\right)^k \left(1 - \frac{A}{m}\right)^{m-k}$$

$$= 1 - \sum_{k=0}^{N-1} \binom{m}{k} \left(\frac{A}{m}\right)^k \left(1 - \frac{A}{m}\right)^{m-k}$$

Answer to the Question No: 6(b)

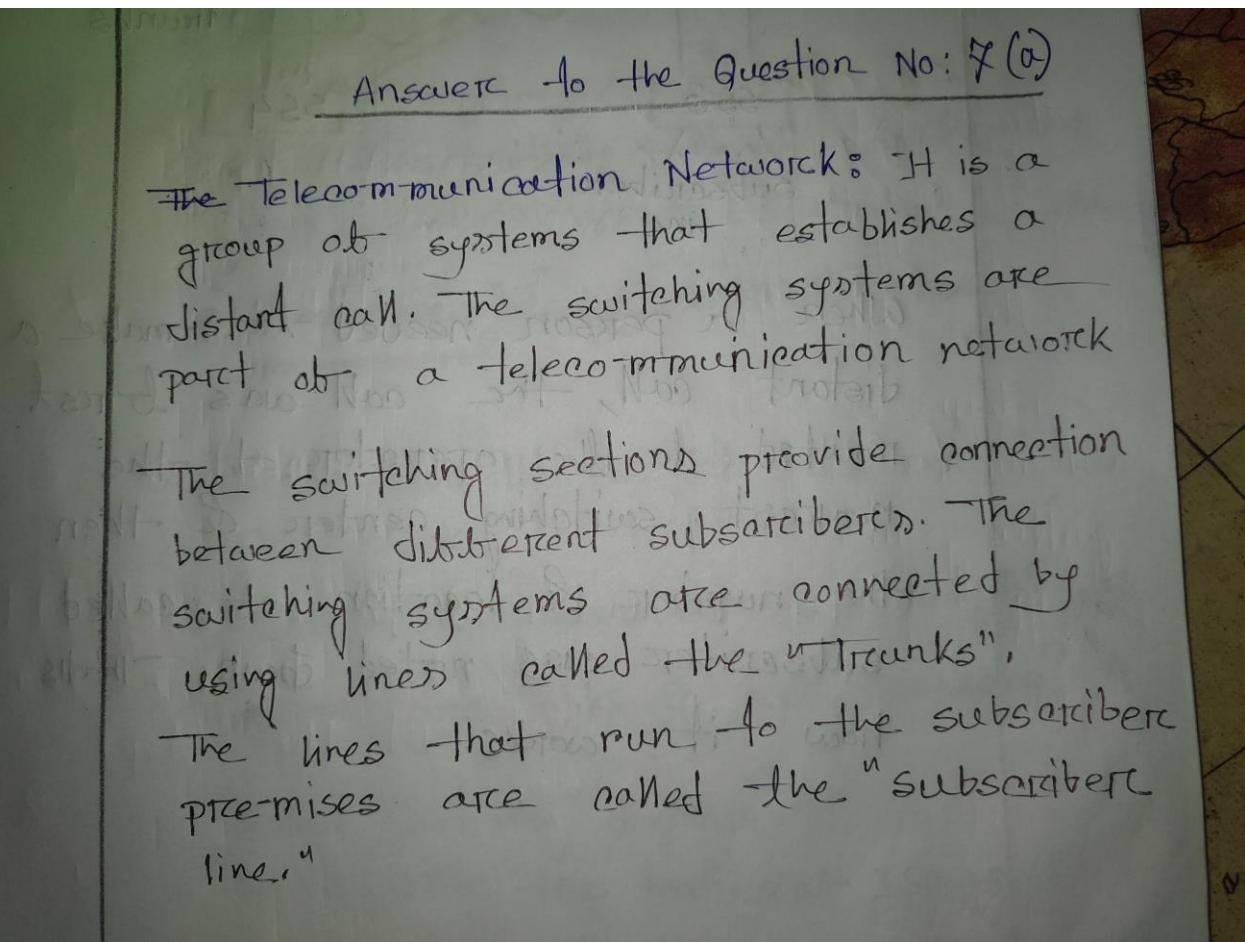
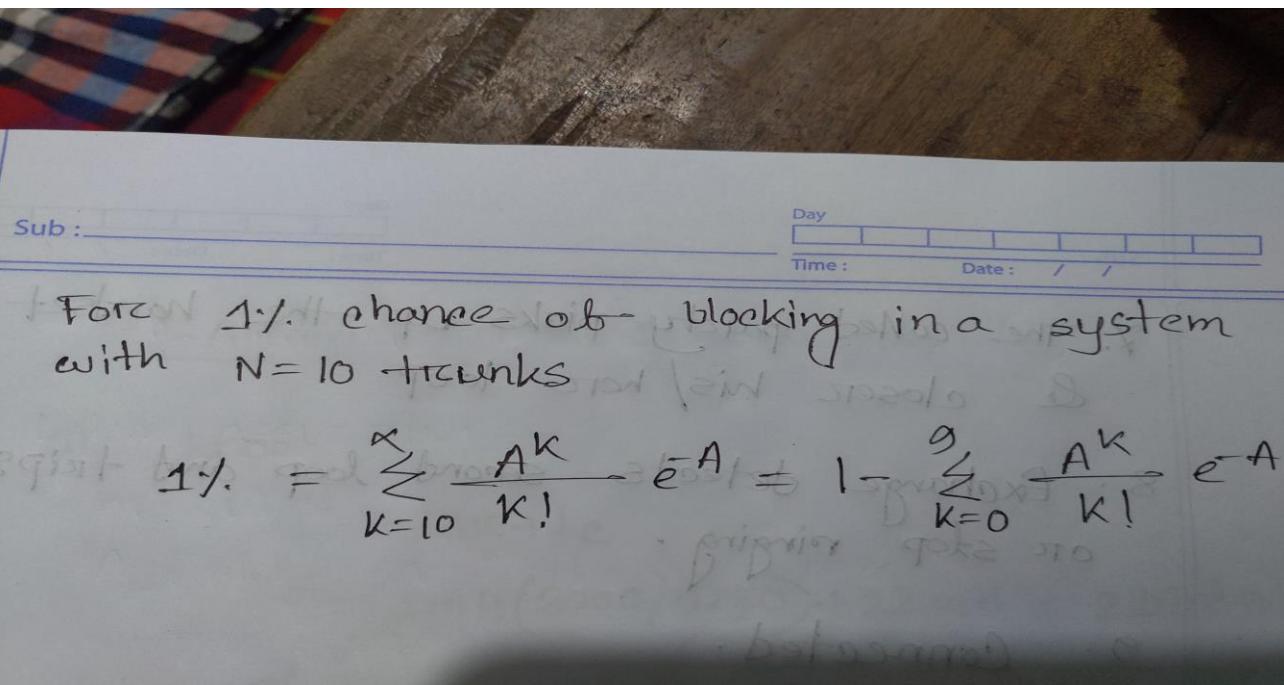
For probability of blocking:

$$P(B) = P(k \geq N) = P(N) + P(N+1) + \dots + P(\alpha)$$

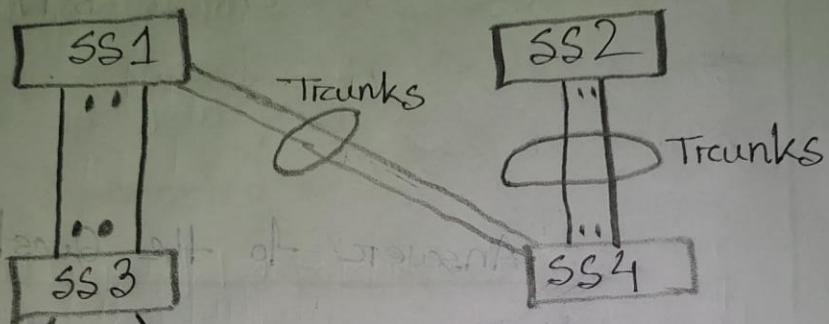
$$= \sum_{k=N}^{\alpha} \frac{e^{-A} A^k}{k!} = \sum_{k=N}^{\alpha} \frac{A^k}{k!} e^{-A}$$

$$= \left(1 - \sum_{k=0}^{N-1} \frac{A^k}{k!}\right) e^{-A}$$

$$1 - e^{-A} \cdot \left(1 + \frac{A}{1!} + \frac{A^2}{2!} + \dots + \frac{A^{N-1}}{(N-1)!}\right) = 1 - e^{-A} \cdot (1 + A + \frac{A^2}{2!} + \dots + \frac{A^{N-1}}{(N-1)!})$$



The following figure shows a telecommunication network.



When a person needed to make a distant call, the call was first routed to the operator at the nearest switching center & then the number & location of called subscriber was noted down. That's how it works.

### Answer to the Question No: 7(b)

$$\text{Hence } n = 10^8$$

$$\text{connection area} = 2 \times 10^3 \text{ km}^2$$

$$\& r = 2 \text{ m.m} = 2 \times 10^{-3} \text{ m} \quad \text{Total area} = 147570 \text{ km}^2$$

$$\text{Fully connected} = \frac{n(n-1)}{2}$$

$$\approx \frac{n^2}{2}$$

$$\approx \frac{n^2}{2} = \frac{(10^8)^2}{2} = \frac{10^{16}}{2}$$

$$= 5 \times 10^{15} \text{ pairs}$$

Avg distance pair ~~not cross section~~

$$= \pi r^2$$

$$= \pi (2 \times 10^3)^2$$

$$= 12.5 \times 10^{12} \text{ km}^2$$

$$\text{Therefore } \text{wiring volume} = 5 \times 10^{15} \times 12.5 \times 10^{12} \times 2 \times 10^3$$

$$= 125 \times 10^6 \text{ km}^3$$

$$\therefore \text{Dept of wiring} = \frac{\text{wiring volume}}{\text{Total area}}$$



Depth of wiring =  $\frac{125 \times 10^6}{147570500} = 847.06 \text{ km}$

Ans.

- Answer to the Question No: 8(a)
- Establishing a call:
1. Calling customer takes phone off hook.
  2. C.O. detects the loop & indicates readiness with dial tone.
  3. Calling customer hears dial tone & dials number etc.
  4. The network checks on the called party is alerted.
  5. Ring tone is returned to the caller.

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7. The called party picks up the handset & closes his/her loop.
8. Exchange detects second loop and trips or stop ringing.
9. Connected.

### Answer to the Question No: 8(b)

Pulse dialing is rapidly disconnected and reconnected in sequence with one pulse for digit value "1", two pulses for digit value "2". Each pulse lasts 0.1 seconds.

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Answer to the Question NO: 8(c)

We know

$$L' = L_{\text{air}} + L_C$$

$$= 1 \text{ mH} \left( \frac{3000}{5280} \right) + 22 \text{ mH} = 22.57 \text{ mH}$$

$$C'' = C_{\text{air}} = 83 \text{ nF} \left( \frac{3000}{5280} \right) = 42.2 \text{ nF}$$

$$\omega_C = \frac{2}{\sqrt{L' C''}} = 61.3 \times 10^3 \text{ rad/s}$$

$$f_C = \frac{\omega_C}{2\pi} = 9.76 \text{ kHz}$$

Ans.

