

## Removing grammar ambiguity

1. Left recursion.
2. Left factoring.

Both problems prevents any LL parser from deciding deterministically which rule should be fired.

### Left Recursion

- We have to eliminate left recursion because top down parsing methods can not handle left recursive grammars.
- A grammar is left recursive if it has a nonterminal A such that there is a derivation  $A \rightarrow A\alpha$  for some string  $\alpha$
- Consider the left recursive grammar  
 $A \rightarrow A\alpha_1 \mid A\alpha_2 \mid \beta_1 \mid \beta_2$

Solution:

$$\begin{aligned} A &\rightarrow \beta_1 A' \mid \beta_2 A' \\ A' &\rightarrow \alpha_1 A' \mid \alpha_2 A' \mid \epsilon \end{aligned}$$

### Left Factoring

- In left factoring it is not clear which of two alternative productions to use to expand a nonterminal A.

$$\text{i.e. if } A \rightarrow \alpha\beta_1 \mid \alpha\beta_2$$

- We don't know whether to expand A to  $\alpha\beta_1$  or to  $\alpha\beta_2$
- To remove left factoring for this grammar replace all A productions containing  $\alpha$  as prefix by  $A \rightarrow \alpha A'$  then  $A' \rightarrow \beta_1 \mid \beta_2$

Example 1:

$$\begin{aligned} X &\rightarrow X + X \mid X * X \mid D \\ D &\rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \end{aligned}$$

Remove Left factoring and left recursion.

Solution:

First removing left factoring.

$$\begin{aligned} X &\rightarrow X B \mid D \\ B &\rightarrow + X \mid * X \\ D &\rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \end{aligned}$$

Still the grammar is ambiguous it has left recursion.

$$X \rightarrow D X'$$
$$X' \rightarrow B X' \mid \epsilon$$
$$B \rightarrow + X \mid * X$$
$$D \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$

## A Recursive-Descent Parser

- One parse method per non-terminal symbol
- A non-terminal symbol on the right-hand side of a rewrite rule leads to a call to the parse method for that non-terminal
- A terminal symbol on the right-hand side of a rewrite rule leads to “consuming” that token from the input token string
- $|$  in the CFG leads to “if-else” in the parser
- $()^*$  in the CFG leads to “while” in the parser

### Task:

Do a small parser program with the previous grammar (Example 1) after removing its ambiguity.

- 1- Remove left factoring and left recursion.
- 2- Tokenizer
- 3- Prepare node classes to build parse tree.
- 4- Apply recursive descent parser to generate parse tree.
- 5- Apply this program on this input.  
5 + 6 \* 10

### References:

- 1- Recursive descent parser <http://web.cse.ohio-state.edu/software/2231/web-sw2/extras/slides/27.Recursive-Descent-Parsing.pdf>
- 2- Recursive descent parser <https://www.youtube.com/watch?v=S7DZdn33eFY>
- 3- Left recursion [https://www.youtube.com/watch?v=K-mfuhx9B\\_0](https://www.youtube.com/watch?v=K-mfuhx9B_0)
- 4- Left factoring [https://en.wikipedia.org/wiki/LL\\_parser#Left\\_Factoring](https://en.wikipedia.org/wiki/LL_parser#Left_Factoring)