

07/26

*** know diff. btwn lea vs mov**

Intel Load Effective Address: \neq mov instruction

Putting the address into the register.

• Leal Src, DST

* Src has to be addr.

• leaq Src, DST

* DST has to be register!

• leal (%eax), %ebx = movl %eax, %ebx

↳ %ebx gets contents of %eax register

But this
Does not
go to memory

• movl (%eax), %ebx

↳ %ebx get 4 byte long from memory at addr in %eax

• leal 4(%eax), %ebx

↳ %ebx gets 4 + contents of %eax

• leal (%eax, %edx), %ebx

↳ %ebx gets sum of %eax, %edx contents

* Registers = numbers

↓

memory = addr.

* think of it as being similar
to pointers

```
int sum (int A, int B)
return A+B;
```

```
leal (%edi, %esi), %eax
ret.
```

```
* int fcn (int i) {
```

```
    int *p;
```

```
    p = &i;
```

```
    *p = 3;
```

```
    return 0;
```

get addr. of local var i

```
leaq -4(%rbp), %rbx
```

```
movl $3, (%rbx)
```

```
*p = 3
```

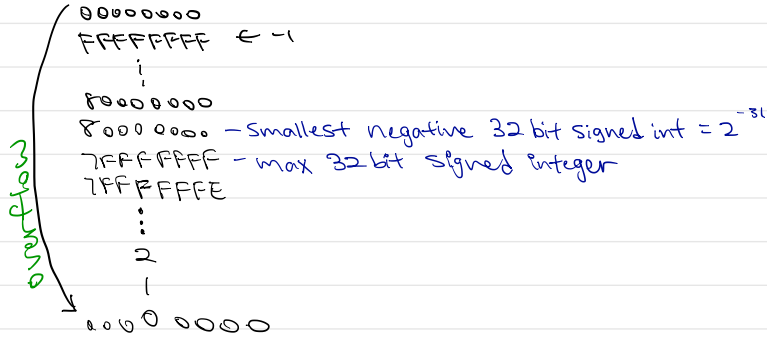
pointer
dereference

Lea not access memory



Overflow:

- occurs when you try to represent a number that does not fit
- * LSB gets stored and MSB gets lost



- Register size of processor determines range of values that can be represented.
 - Unsigned examples:

· 8 bit: $0 \dots 2^8 - 1$	· 32 bit: $0 \dots 2^{32} - 1$
· 16 bit: $0 \dots 2^{16} - 1$	· 64 bit: $0 \dots 2^{64} - 1$
- An arithmetic result produces a result larger than N bits.
- Overflow reduces result to modulo N, retaining only the least significant N bits, causing a wrap around.

addl src, dst dst += src
 (src > 0 && dst > 0 && sum < 0)
 || (src < 0 && dst < 0 && sum > 0)

overflow!

$$\begin{array}{r} +int \\ +int \\ \hline -int \end{array}$$

overflow!

$$\begin{array}{r} -int \\ +int \\ \hline -int \end{array}$$

ex.

$$\begin{array}{r} 7FFFFFFF \\ + \quad 1 \\ \hline 80000000 \end{array}$$

OVERFLOW

Subl Src, dst DST -= SRC

(src < 0 && dst > 0 && difference < 0)
|| (src > 0 && dst < 0 && difference > 0)

$$\begin{array}{r} + \text{int} \\ - - \text{int} \\ \hline - \text{int} \end{array} \quad \text{overflow!}$$

$$\begin{array}{r} - \text{int} \\ - + \text{int} \\ \hline + \text{int} \end{array} \quad \text{overflow!}$$

mult Src, dst DST *= SRC

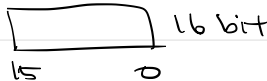
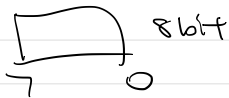
SRC - DST

(src > 0 && dst > 0 && product < 0)
|| (src < 0 && dst < 0 && product < 0)
|| (src < 0 && dst > 0 && product > 0)
|| (src > 0 && dst < 0 && product > 0)

$$\begin{array}{r} + \text{int} \\ \times + \text{int} \\ \hline - \text{int} \end{array} \quad \text{of!} \quad \left| \quad \begin{array}{r} - \text{int} \\ \times - \text{int} \\ \hline - \text{int} \end{array} \quad \text{of!} \quad \left| \quad \begin{array}{r} - \text{int} \\ \times + \text{int} \\ \hline + \text{int} \end{array} \quad \text{of!} \quad \left| \quad \begin{array}{r} + \text{int} \\ \times - \text{int} \\ \hline + \text{int} \end{array} \quad \text{of!} \right.$$

Bit Ordering:

A convention independent of byte ordering
(Endianness) ★





Bitwise operations: (And, or, Xor)

Bitwise And (operator: &)

int a, b, c;

c = A & B;

two opnds



$$\begin{array}{r} 1100 \\ \& 1010 \\ \hline 1000 \end{array}$$

Bitwise or (operator: |)

int a, b, c;

c = A | B;

two opnds



$$\begin{array}{r} 1100 \\ | 1010 \\ \hline 1110 \end{array}$$

Bitwise Xor (operator: ^)

int a, b, c;

c = A ^ B;

two opnds



$$\begin{array}{r} 1100 \\ ^ 1010 \\ \hline 0110 \end{array}$$

Bitwise Not (operator: ~)

• Flip the bits

int a, b;

a = ~b;



$$\begin{array}{r} \sim 01 \\ \hline 10 \end{array}$$



Logical operations: (And, or)

Logical And (operator: &&)

• result is 0 or 1

if (fC) && gU)

if left is 0, then right is not reached

Logical Or (operator: ||)

• result is 0 or 1

if (fC) || gU)

if left is 1, then right is not reached