

Linked Lists

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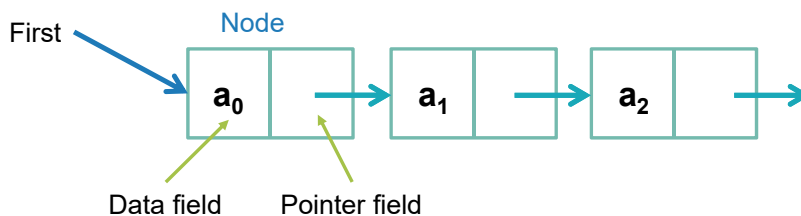
Array

- Store an ordered list
- Using sequential mapping
 - Element(node) a_i is stored in the location L_i of the array
 - Next node is at the location L_i+1
- Pros:
 - Suitable for random access
 - Efficient to insert/delete from the end
 - Adequate for special data structures, Stack and Queue
- Con:
 - Difficult to insert/delete nodes at arbitrary location

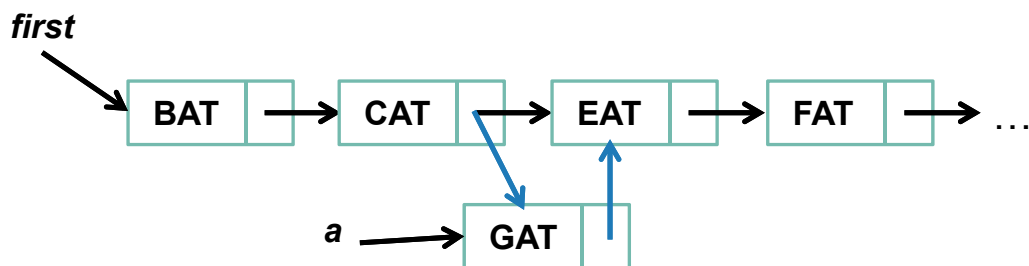
Why?

Linked Lists

- Nodes are no longer continue in the memory
- Each node stores the address or location of the next one
- Singly Linked List (SLL)
 - Each node has exactly one pointer field

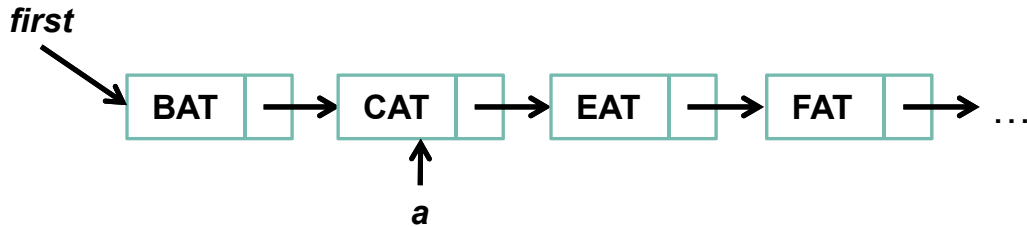


SLL Operation: Insert



- Steps to insert a "GAT" in between "CAT" and "EAT" nodes
 - Create a new node "a" and set data field to "GAT"
 - Set the link field of "a" to "EAT" node
 - Set the link field of "CAT" node to "a"

SLL Operation: Delete

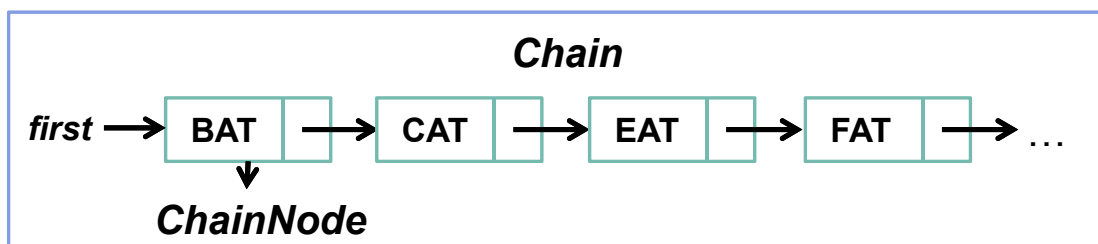


- Steps to delete a "EAT" node from the list
 - Locate the node "a" precedes the "EAT" node
 - Set the link field of "a" to node next to "EAT" node
 - Delete the "EAT" node

You do not need to move or shift any nodes!

Conceptual Design

- Defining a "ChainNode" class
 - Data field
 - Link field
- Designing a "Chain" class
 - Support various operation on ChainNodes



ChainNode & Chain Classes

■ Composite class

```
class ChainNode
{
    friend class Chain;
public:
    // Constructor
    ChainNode(int
value=0, ChainNode*
next=NULL) {
        data = value;
        link = next;
private:
    int data;
    ChainNode *link;
};
```

```
class Chain
{
public:
    // Create a chain with two nodes
    void Create2();

    // Insert a node with data=50
    void Insert50(ChainNode *x);

    // Delete a node
    void Delete(ChainNode *x, ChainNode *y);
private:
    ChainNode *first;
};
```

ChainNode & Chain Classes

■ Nested class

```
class Chain
{
public:
    // Create a chain with two nodes
    void Create2();

    // Insert a node with data=50
    void Insert50(ChainNode *x);

    // Delete a node
    void Delete(ChainNode *x, ChainNode *y);
private:
    class ChainNode{
    public:
        int data;
        ChainNode *link;
    }
    ChainNode *first;
};
```

Review Pointer Manipulation

■ Declaration

- `NodeA *a1=NULL,`
`*a2=NULL;`

■ Allocate memory

- `a1 = new NodeA;`
- `a2 = new NodeA[10];`

■ Delete memory

- `delete a1; a1=NULL;`
- `delete [] a2; a2=NULL;`

■ Dereference

- `NodeA &a1Ref = (*a1);`

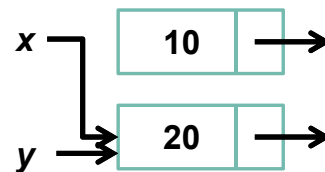
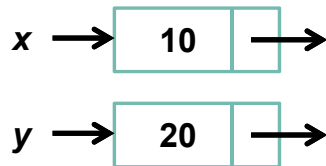
■ Access members

- `a1->memData;`
- `a1->memFunc();`
- `(*a1).memData;`
- `(*a1).memFunc();`

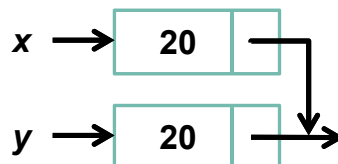
Pointer Assignment

■ ChainNode *x, *y;

- `x = y;`



- `*x = *y;`



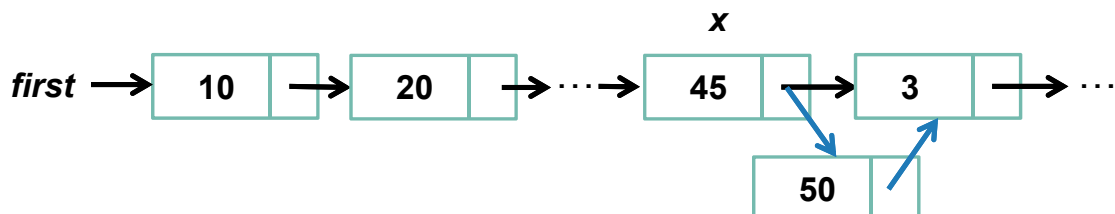
Chain Operations

```
void Chain::Create2()  
{  
    // Create and set the fields of 2nd node  
    ChainNode* second = new ChainNode(20,0);  
  
    // Create and set the fields of 1st node  
    first = new ChainNode(10,second);  
}
```



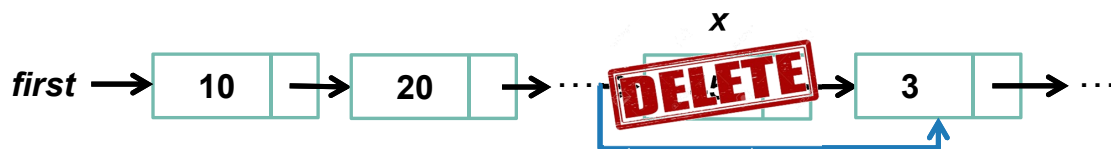
Chain Operations

```
void Chain::Insert50(ChainNode *x)  
{  
    if(first) // Insert after x  
        x->link = new ChainNode(50, x->link);  
    else // Insert into empty list  
        first = new ChainNode(50);  
}
```



Chain Operations

```
void Chain::Delete(ChainNode *x, ChainNode *y)
{ // x is the node to be deleted and y is the node
  // preceding x
  if(x==first) first = first->link;
  else y->link = x->link;
  delete x;
  x=NULL;
}
```



Template Chain Class

```
Template < class T > class Chain; // Forward declaration

template < class T >
class ChainNode {
friend class Chain <T>;
private:
    T data;
    ChainNode<T>* link;
};

template <class T>
class Chain {
public:
    // Constructor
    Chain(void) {first = last = NULL;}

    // Chain operations...

private:
    ChainNode<T> *first;
    ChainNode<T> *last;
};
```

Chain Operations

```
template < class T >
void Chain<T>::InsertBack(const T& e)
{
    if(first) { // Non-empty chain
        last->link = new ChainNode<T>(e);
        last = last->link;
    }
    else // Insert into an empty chain
        first = last = new ChainNode<T>(e);
}
```

```
template < class T >
void Chain<T>::Concatenate(Chain<T>& b)
{ // b is concatenated to the end of *this
    if ( first ) { last->link = b.first; last = b.last; }
    else { first = b.first; last = b.last; }
    b.first = b.last = 0;
}
```

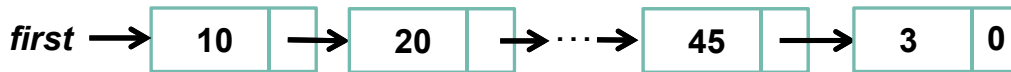
Chain Operations

- Reverse a chain, such that (a^1, a^2, \dots, a^n) turns into $(a^n, a^{n-1}, \dots, a^1)$



Chain Operations

- Reverse a chain, such that (a^1, a^2, \dots, a^n) turns into $(a^n, a^{n-1}, \dots, a^1)$

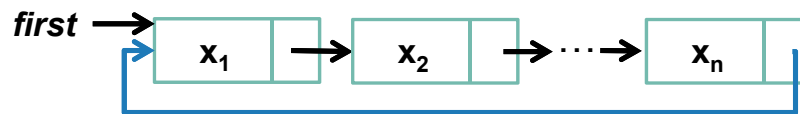


Chain Operations

```
template < class T >
void Chain<T>::Reverse(void)
{ // Turn a chain, (a1, ..., an) into (an, ..., a1)
  ChainNode<T> *current = first, *previous = NULL;
  while (current) {
    ChainNode<T> *r = previous;
    previous = current; // r is behind the previous
    current = current->link; // move current to next node
    previous->link = r; // link previous to previous node
  }
  first = previous;
}
```

Circular Lists

- A singly-linked circular list
- The link field of the last node points to the first node



- Check for the last node
 - if(current->link == **first**)
- You could visit a node from any position

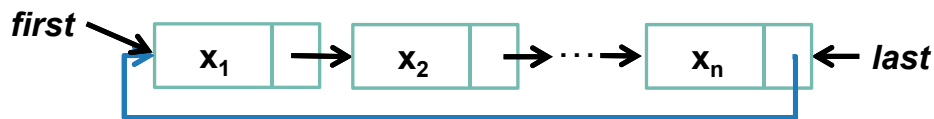
Circular Lists : Insert

- Suppose we want to insert a new node at the front of list
- Set link field of new node to **first** and set **first** to new node
- Go to the last node and set the link field to new node



Circular Lists

- Computation complexity for finding the last one?
 - $O(N)$
- Computation complexity for finding the first one?
 - We could always access the first node via `[last->link]`
 - $O(1)$
- It is more convenient to store the last node of a circular list



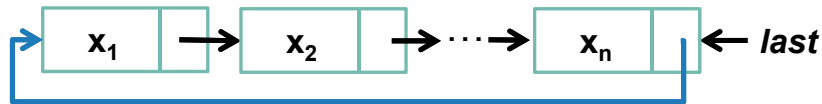
Circular Lists : Insert at Front

```
Template<class T>
void CircularList<T>::InsertFront(const T& e)
{
    ChainNode<T>* newNode = new ChainNode<T>(e);

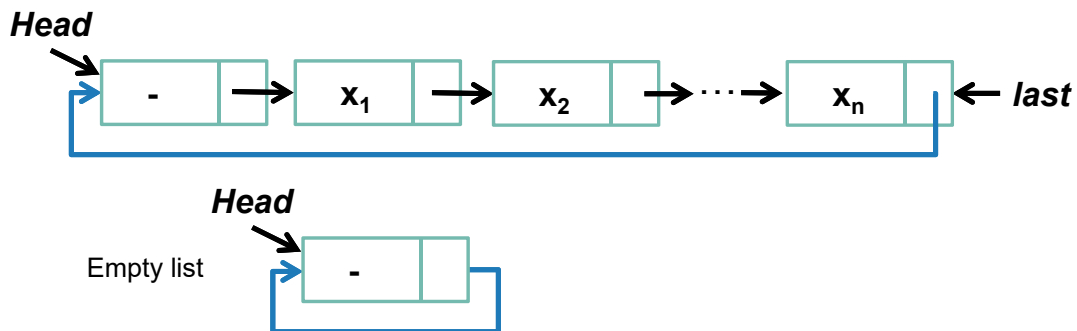
    if(last){ // nonempty list
        newNode->link = last->link;
        last->link = newNode;
    }
    else{ // empty list
        last = newNode;
        newNode->link = newNode;
    }
}
```

Circular Lists

- How to represent an “empty” list?



- Introducing a dummy node “Header”



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Sparse Matrix

$$a[6][6] = \begin{bmatrix} 15 & 0 & 0 & 22 & 0 & -15 \\ 0 & 11 & 3 & 0 & 0 & 0 \\ 0 & 0 & 0 & -6 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 91 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 28 & 0 & 0 & 0 \end{bmatrix}$$

A	row	col	value
smArray[0]	0	0	15
smArray[1]	0	3	22
smArray[2]	0	5	-15
smArray[3]	1	1	11
smArray[4]	1	2	3
smArray[5]	2	3	-6
smArray[6]	4	0	91
smArray[7]	5	2	28

- A matrix has many **zero** elements
- Devise a sequential array
 - store **non-zero** elements
 - **row-major** order
- Access specific column is difficult
- Using circular lists representation

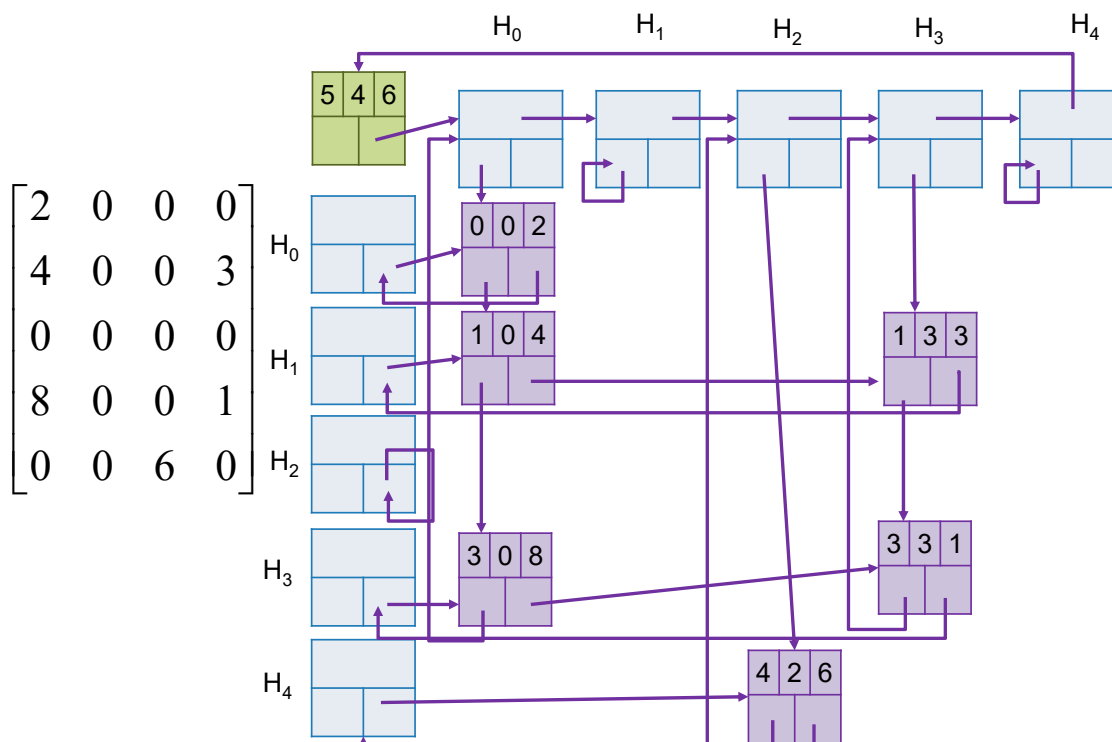
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Linked Structure

- Header node: for each row or column
 - **Down**: link to the 1st non-zero term in the column
 - **Right**: link to the 1st non-zero term in the row
 - **Next**: link to the next head node
 - The header node for row i is also the header node for column i
- Element node, each **non-zero** term that stores
 - Data of **row, col, and value**
 - A **down** field to link to the next non-zero term in the same **column**
 - A **right** field to link to the next non-zero term in the same **row**
- The header of header nodes (a circular list)
 - Store dimension of the matrix

Sparse Matrix in Linked Structure



Create a Sparse Matrix

- Given a $n \times m$ sparse matrix with r non-zero terms
 - the total number of required nodes are $\max\{n, m\} + r + 1$
- Input format
 - The 1st line gives the dimension of matrix and # of non-zero terms
 - Each subsequent input line is a triple of the form (i, j, a_{ij})
 - Triples are ordered by rows and within rows by columns

$$\begin{bmatrix} 2 & 0 & 0 & 0 \\ 4 & 0 & 0 & 3 \\ 0 & 0 & 0 & 0 \\ 8 & 0 & 0 & 1 \\ 0 & 0 & 6 & 0 \end{bmatrix}$$

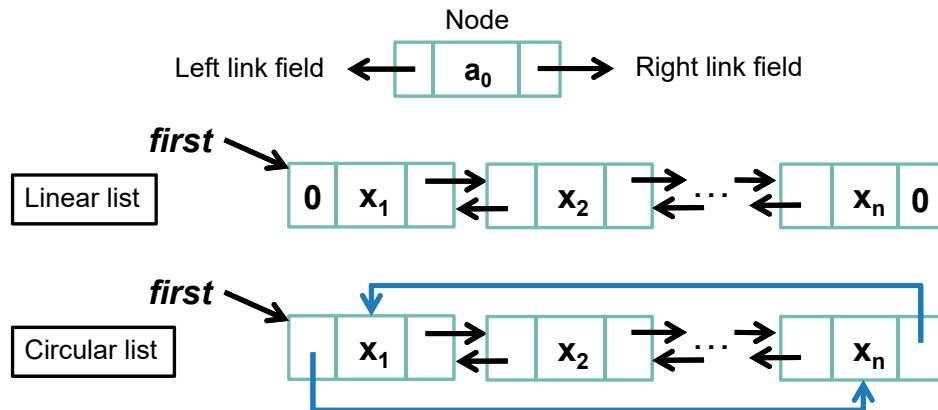
Input
5,4,6;
0,0,2;
1,0,4;
1,3,3;
3,0,8;
3,3,1;
4,2,6;

Create a Sparse Matrix

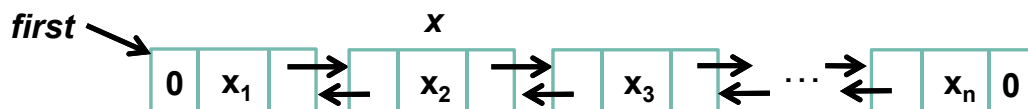
- Performance analysis
 - Set up header nodes, $O(\max\{n, m\})$
 - Set up non-zero nodes, $O(r)$
 - Close column lists, $O(\max\{n, m\})$
 - Link header nodes, $O(\max\{n, m\})$
- Total complexity: $O(\max\{n, m\} + r) = O(n + m + r)$

Double Linked Lists

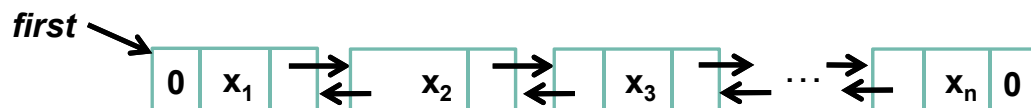
- Each node has **two** link fields
- Could move in **two directions** to visit nodes



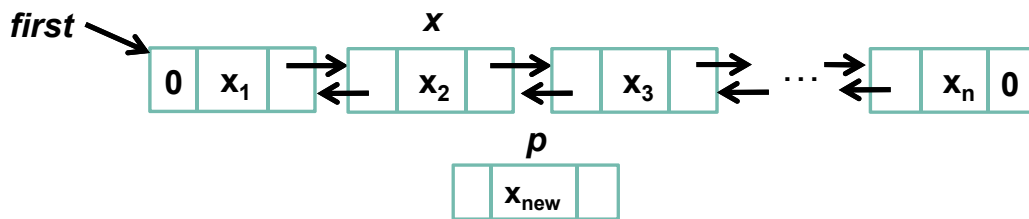
Double Linked Lists : Delete



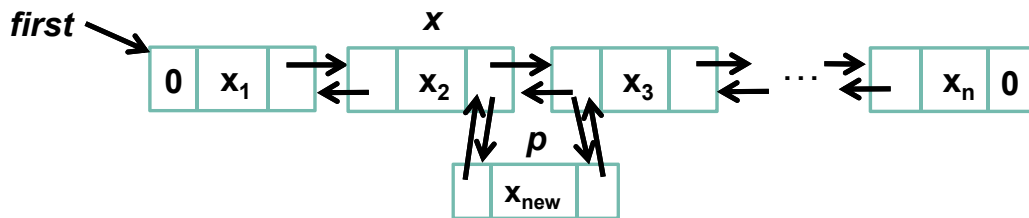
$x \rightarrow \text{left} \rightarrow \text{right} = x \rightarrow \text{right};$ $x \rightarrow \text{right} \rightarrow \text{left} = x \rightarrow \text{left};$



Double Linked Lists : Insert



$p \rightarrow \text{left} = x; p \rightarrow \text{right} = x \rightarrow \text{right}$



$x \rightarrow \text{right} \rightarrow \text{left} = p; \Rightarrow x \rightarrow \text{right} = p;$

Self-Study Topics

- Polynomial using linked lists
- Linked stacks and queues



Visit Elements in a Container

- Suppose we have a chain C of datatype Chain<int>.
 - Output all integers in C
 - Obtain the maximum, minimum or mean of all integers in C
 - Obtain the sum, product, or sum of squares of all integers in C
- All operations require to visit every element in the chain C

How to Visit a Container?

```
For each item in C
{
    currentItem = current item in C;
    //do something with currentItem;
}
```

- In an array representation

```
for (int i = 0; i < n; i++)
{
    int currentItem = a[i];
    // do something with currentItem;
}
```

How to Visit a Container?

```
For each item in C
{
    currentItem = current item in C;
    // do something with currentItem;
}
```

■ In a linked list representation

```
for (ChainNode<int> *ptr=first; ptr!=0; ptr=ptr->link)
{
    int currentItem = ptr->data;
    // do something with currentItem;
}
```

Visiting a Container using Iterator

- A powerful mechanism to visit a container with arbitrary data type
- Guarantee runtime range safety
- Applicable to all STL algorithms
- Suitable for team development
- Might scarify some amount of performance

```
// Possible implementation of STL copy algorithm
template < class Iterator >
void copy(Iterator start, Iterator end, Iterator to)
{ // copy from src[start, end) to dst[to, to+end-start)
    while (start != end)
    { *to = *start ; start++ ; to++; }
}
```

What is an Iterator ?

```
void main()
{
    int x [3] = {0,1,2};
    for (int* y = x; y != x+3; y++)
        cout << *y << endl;
}
```

- An **iterator** is a pointer to an element in a container
- Using dereferencing operator (*) to access an element
- Support pre- or post- increment operator (++)

```
void main()
{
    for (Iterator y = start; y != end; y++)
        cout << *y << endl;
}
```

C++ Iterators

- **Input** iterator
 - Read access, pre- and post- “++” operators.
- **Output** iterator
 - Write access, pre- and post- “++” operators.
- **Forward** iterator
 - pre- and post- “++” operators.
- **Bidirectional** iterator
 - pre- and post- “++” and “--” operators.
- **Random access** iterator
 - Permit pointer jumps by arbitrary amounts.
- All iterators supports “==”, “!=” and “*” operators

Forward Iterator for Chain

```
template <class T>
class Chain {
public:
    // Constructor
    Chain(void) {first = last = NULL;}

    // Chain operations...

    class ChainIterator{...};

    // Get the first element
    ChainIterator begin() {return ChainIterator(first);}

    // Get the end of the list
    ChainIterator end() {return ChainIterator(0);}
private:
    ChainNode<T> *first;
    ChainNode<T> *last;
};
```

Forward Iterator for Chain

■ General usage

```
void main()
{
    Chain<int> myChain;
    // do operations on myChain here...

    // print out every element in myChain
    Chain<int>::ChainIterator my_it;
    for (my_it = myChain.begin(); my_it != myChain.end(); ++my_it)
        cout << *my_it << endl;

    // Use STL algorithm to calculate the sum of myChain
    int sum = std::accumulate(myChain.begin(), myChain.end(), 0);
}
```

```

Class ChainIterator{ // A nested class within Chain
public:
    // Constructor
    ChainIterator(ChainNode<T>* startNode = 0)
        {current = startNode;}

    // Dereferencing operator
    T& operator*() const {return current->data;}
    T* operator->() const {return &current->data;}

    // Increment operator
    ChainIterator& operator++() // pre- "++"
    { current = current->link ; return *this; }

    ChainIterator operator++(int) // post- "++"
    {
        ChainIterator old = *this;
        current = current->link;
        return old;
    }

    // Equality operators
    bool operator!=(const ChainIterator right) const
    { return current != right.current; }

    bool operator==(const ChainIterator right) const
    { return current == right.current; }

private:
    ChainNode<T>* current;
};

```