

MySQL - RDBMS

Agenda

- Module Overview
- DBMS vs RDBMS
- SQL
- Getting started
- Create Table
- Insert records
- Select records

Module Overview

- Syllabus
 - RDBMS (MySQL)
 - NoSQL Introduction (Mongo)
- Evaluation
 - Theory: 40 marks MCQ (CCEE)
 - Lab: 40 marks - SQL queries & PSM
 - 15-Oct / 16-Oct
 - Internals: 20 marks - Will be updated by CoCo.
- Module plan
 - C-DAC Source Book -- 72 hrs (Theory 36 hrs + Lab 36 hrs).
 - Sunbeam Database Technologies
 - 10 days (2 weeks) - 4-Oct to 16-Oct.
 - Theory: 40 hrs
 - 9.00 am to 1.30 pm.
 - Break: 11.00 am to 11.30 am.
 - Q & A: 8.45 am, 1.30 pm, 3.30 pm.
 - Lab: 36 hrs.
 - 4.00 pm to 7.00 pm.

DBMS vs RDBMS

SQL

MySQL

Installation

- Server = mysqld.exe
 - C:\Program Files\MySQL\MySQL Server 8.0\bin
- Client = mysql.exe
 - C:\Program Files\MySQL\MySQL Server 8.0\bin
- MySQL data directory/folder

- C:\ProgramData\MySQL\MySQL Server 8.0
- Default user of MySQL = "root" (administrator).
 - Password is given during installation = "manager".
- Set PATH.
 - Windows explorer.
 - This PC (right click) --> Properties --> Advanced System Settings --> Advanced --> Environment Variables
 - User Variables --> PATH --> EDIT
 - Click New --> Add MySQL PATH at the end "C:\Program Files\MySQL\MySQL Server 8.0\bin".
- Open "Command Prompt"
 - cmd> mysql --version

Getting started

- step 1. Login with "root" user.
 - cmd> mysql -u root -p
 - Password: manager
- step 2. Create a new user.
 - mysql> CREATE USER sunbeam@localhost IDENTIFIED BY 'sunbeam';
- step 3. Create a new database/schema.
 - mysql> CREATE DATABASE classwork;
 - mysql> SHOW DATABASES;
- step 4. Give all permissions to the new user on the new database.
 - mysql> GRANT ALL PRIVILEGES ON classwork.* TO sunbeam@localhost;
 - mysql> FLUSH PRIVILEGES;
- step 5. mysql> EXIT
- step 1. Login with new user and password on MySQL CLI.
 - cmd> mysql -u sunbeam -p
 - Password: sunbeam
- step 2. Execute queries -- DDL, DML, DQL, ...
 - mysql> SHOW DATABASES;
 - mysql> SELECT USER(), DATABASE();
 - mysql> USE classwork;
 - mysql> SELECT USER(), DATABASE();
 - mysql> SHOW TABLES;
- MySQL screen clear.
 - mysql> ! cls

- CREATE TABLE
 - CREATE TABLE tablename (col1 DATATYPE, col2 DATATYPE, col3 DATATYPE, ...);
- INSERT
 - INSERT INTO tablename VALUES (v1, v2, v3, ...);
 - INSERT INTO tablename VALUES (v1, v2, v3, ...), (v1, v2, v3, ...), (v1, v2, v3, ...), ..;
- SELECT
 - SELECT * FROM tablename;
 - ■ means All columns.

```
SHOW TABLES;
```

```
CREATE TABLE stud(id INT, name CHAR(20), marks DOUBLE);
```

```
SHOW TABLES;
```

```
DESCRIBE stud;
```

```
INSERT INTO stud VALUES (1, 'Soham', 98.20);
```

```
INSERT INTO stud VALUES (2, 'Sakshi', 97.40);
```

```
INSERT INTO stud VALUES (3, 'Prisha', 99.30), (4, 'Madhura', 96.29), (5, 'Om', 97.45);
```

```
SELECT * FROM stud;
```

```
INSERT INTO stud(name, marks, id) VALUES ('Pratham', 95.39, 6);
```

```
INSERT INTO stud(id, name) VALUES (7, 'Vedant');
```

```
SELECT * FROM stud;
```

```
CREATE TABLE students(id INT, name CHAR(20), marks DOUBLE);
```

```
SHOW TABLES;
```

```
SELECT * FROM students;
```

```
INSERT INTO students SELECT * FROM stud;
```

```
SELECT * FROM students;
```

```
DROP TABLE students;
```

```
SHOW TABLES;
```

SUNBEAM INFOTECH

MySQL - RDBMS

Agenda

- Logical vs Physical Layout
- MySQL data types
- CHAR vs VARCHAR vs TEXT
- SQL scripts
- DQL - SELECT
 - Projection
 - Computed Columns
 - DISTINCT
 - LIMIT
 - ORDER BY
 - WHERE
 - Relational Operators
 - NULL Operators
 - Logical Operators

DBMS

- DBMS -- Database Management (CRUD)
 - Traditional databases - File based databases
 - Relational databases
 - NoSQL databases

Logical vs Physical Layout

MySQL data types

```
* CREATE TABLE stud(id INT, name CHAR(20), marks DOUBLE);
* id is by default signed.
* CREATE TABLE stud(id INT UNSIGNED, name CHAR(20), marks DECIMAL(5,2));
* id is now unsigned.
* DECIMAL(5,2) -- total 5 digits and 2 digits after decimal point.
```

```
CREATE TABLE temp(c1 CHAR(4), c2 VARCHAR(4), c3 TEXT(4));
```

```
DESCRIBE temp;
```

```
INSERT INTO temp VALUES('abcd', 'abcd', 'abcdef');
```

```
INSERT INTO temp VALUES('xy', 'xy', 'xy');
```

```
INSERT INTO temp VALUES('pqr', 'pqr', 'pqr');
```

```
INSERT INTO temp VALUES('pqrst', 'pqr', 'pqr');
-- error

INSERT INTO temp VALUES('pqr', 'pqrst', 'pqr');
-- error

SELECT * FROM temp;
```

CHAR vs VARCHAR vs TEXT

SQL scripts

```
USE classwork;

SELECT USER(), DATABASE();
-- sunbeam@localhost, classwork

SOURCE D:/classwork-db.sql

SHOW TABLES;

SELECT * FROM books;

SELECT * FROM dept;
```

DQL - SELECT

- Projection -- Display only given columns.

```
SELECT * FROM dept;
-- * all columns

SELECT deptno, dname FROM dept;

SELECT * FROM emp;

SELECT empno, ename, sal FROM emp;

SELECT sal, ename, deptno FROM emp;

CREATE TABLE newemp(id INT, name CHAR(30), sal DOUBLE);

INSERT INTO newemp SELECT * FROM emp;
-- error

INSERT INTO newemp SELECT empno, ename, sal FROM emp;

SELECT * FROM newemp;
```

```
INSERT INTO newemp(id, name, sal) SELECT empno, ename, sal FROM emp;

SELECT * FROM newemp;
```

```
-- display ename, sal, da=sal*0.5, gs=sal+sal*0.5
SELECT ename, sal, sal*0.5, sal+sal*0.5 FROM emp;
-- computed columns or pseudo columns.

SELECT ename AS name, sal, sal * 0.5 AS da, sal + sal * 0.5 AS gs FROM emp;
-- column alias

SELECT ename name, sal, sal * 0.5 da, sal + sal * 0.5 gs FROM emp;
-- writing AS is optional

SELECT ename name, sal, sal * 0.5 da, sal + sal * 0.5 gross sal FROM emp;
-- error

SELECT ename name, sal, sal * 0.5 da, sal + sal * 0.5 `gross sal` FROM emp;
```

```
-- display ename, deptno and dname (10=ACCOUNTING, 20=RESEARCH, 30=SALES).
SELECT ename, deptno FROM emp;

SELECT ename, deptno, CASE
WHEN deptno = 10 THEN 'ACCOUNTING'
WHEN deptno = 20 THEN 'RESEARCH'
WHEN deptno = 30 THEN 'SALES'
ELSE 'UNKNOWN'
END AS deptname
FROM emp;
```

```
SELECT DISTINCT job FROM emp;

SELECT DISTINCT deptno FROM emp;

SELECT deptno, job FROM emp;
-- 10 --> CLERK, MANAGER, PRESIDENT
-- 20 --> CLERK, ANALYST, MANAGER
-- 30 --> CLERK, MANAGER, SALESMAN

SELECT DISTINCT deptno, job FROM emp;
-- displays unique combination of deptno & job.
```

```
DESCRIBE emp;
```

```
SELECT * FROM emp;
-- show all rows

SELECT * FROM emp LIMIT 5;
-- show first 5 rows.

SELECT * FROM emp LIMIT 10;
-- show first 10 rows.

SELECT ename, sal FROM emp LIMIT 10;
-- show ename & sal of first 10 rows.

SELECT ename, sal FROM emp LIMIT 5;
-- show ename & sal of first 5 rows.

SELECT ename, sal FROM emp LIMIT 3, 5;
-- show ename & sal of 5 rows after first 3 rows.
```

```
SELECT * FROM emp;

SELECT * FROM emp ORDER BY sal DESC;

SELECT * FROM emp ORDER BY deptno ASC;

SELECT * FROM emp ORDER BY hire;
-- default sort = ASC.

SELECT * FROM emp ORDER BY deptno ASC, job ASC;
-- sort on multiple columns.

SELECT * FROM emp ORDER BY deptno DESC, job, ename;
-- sort on deptno (desc), job (asc), ename (asc)
```

```
-- display top 3 emps as per sal.
SELECT * FROM emp ORDER BY sal DESC LIMIT 3;

-- display emp whose name is last in alphabetically.
SELECT * FROM emp ORDER BY ename DESC LIMIT 1;

SELECT * FROM emp ORDER BY comm;

-- display emp with lowest sal.
SELECT * FROM emp ORDER BY sal LIMIT 1;

-- display emp with third lowest sal.
SELECT * FROM emp ORDER BY sal LIMIT 2,1;

-- display emp with second highest sal.
```



```
SELECT * FROM emp ORDER BY sal DESC;

SELECT * FROM emp ORDER BY sal DESC LIMIT 1, 1;

-- sort emp on da=sal*0.5
SELECT ename, sal, sal*0.5 da FROM emp;

SELECT ename, sal, sal*0.5 da FROM emp
ORDER BY sal*0.5;
-- order by expr

SELECT ename, sal, sal*0.5 da FROM emp
ORDER BY da;
-- order by column alias

SELECT ename, sal, sal*0.5 da FROM emp
ORDER BY 3;
-- order by column numer (in projection)
```

```
-- display emps of dept 30
SELECT * FROM emp WHERE deptno = 30;

-- display all emps with sal > 2000.0
SELECT * FROM emp WHERE sal > 2000.0;

-- display all ANALYST.
SELECT * FROM emp WHERE job = 'ANALYST';

-- display all emps not in dept 30.
SELECT * FROM emp WHERE deptno != 30;
SELECT * FROM emp WHERE deptno <> 30;
```

```
-- display emps in sal range from 1000 to 2000.
-- sal >= 1000 && sal <= 2000
SELECT * FROM emp WHERE sal >= 1000 AND sal <= 2000;

-- display ANALYSTs and MANAGERS.
-- job == 'ANALYST' || job == 'MANAGER';
SELECT * FROM emp WHERE job = 'ANALYST' OR job = 'MANAGER';

-- display emps who are not salesman.
SELECT * FROM emp WHERE job != 'SALESMAN';
SELECT * FROM emp WHERE job <> 'SALESMAN';

-- !(condition)
SELECT * FROM emp WHERE NOT job = 'SALESMAN';
```

```
-- display all emps hired in 1982.  
-- '1-1-1982' to '31-12-1982'  
SELECT * FROM emp WHERE hire >= '1982-01-01' AND hire <= '1982-12-31';
```

```
-- display all emps whose comm is null.  
SELECT * FROM emp WHERE comm = NULL;  
-- relational operators cannot be used with NULL.  
-- NULL is used with special operators.
```

```
SELECT * FROM emp WHERE comm IS NULL;
```

```
SELECT * FROM emp WHERE comm <=> NULL;
```

```
-- display all emps whose comm is not null.
```

```
SELECT * FROM emp WHERE comm IS NOT NULL;
```

```
SELECT * FROM emp WHERE NOT (comm IS NULL);
```

MySQL - RDBMS

Agenda

- DQL - SELECT
 - WHERE clause
 - IN operator
 - BETWEEN operator
 - LIKE operator
- DML - UPDATE
- DML - DELETE
- DDL - TRUNCATE
- DDL - DROP
- HELP
- DUAL table
- SQL Functions
 - String Functions

WHERE clause

- terminal> mysql -u sunbeam -psunbeam classwork
- IN operator
 - Similar to Logical OR for equality checking with multiple values (for same column).
 - NOT IN operator -- inverse of IN operator

```
SELECT USER(), DATABASE();
-- | sunbeam@localhost | classwork |

-- display all MANAGERS and emps from dept 10.
SELECT * FROM emp WHERE job = 'MANAGER' OR deptno = 10;

-- display all emps whose sal is less than 1000 or sal is more than 3500.
SELECT * FROM emp WHERE sal < 1000 OR sal > 3500;

-- display all analysts and managers.
SELECT * FROM emp WHERE job = 'ANALYST' OR job = 'MANAGER';

SELECT * FROM emp WHERE job IN ('ANALYST', 'MANAGER');

-- display emps whose names are JAMES, KING, MARTIN, FORD.
SELECT * FROM emp WHERE ename = 'JAMES' OR ename = 'KING' OR ename = 'MARTIN' OR
ename = 'FORD';

SELECT * FROM emp WHERE ename IN ('JAMES', 'KING', 'MARTIN', 'FORD');

-- display all emps whose are not salesman or manager.
SELECT * FROM emp WHERE job NOT IN ('SALESMAN', 'MANAGER');
```

```
SELECT * FROM emp WHERE NOT (job IN ('SALESMAN', 'MANAGER'));

SELECT * FROM emp WHERE NOT (job = 'SALESMAN' OR job = 'MANAGER');
```

- BETWEEN operator
 - range check (including both ends).
 - col BETWEEN start AND end
 - col >= start AND col <= end.

```
-- display manager of dept 20.
SELECT * FROM emp WHERE job = 'MANAGER' AND deptno = 20;

-- display all emps in dept 30 whose sal is less than 1500.
SELECT * FROM emp WHERE deptno = 30 AND sal < 1500;

-- display all emps in sal range 1500 to 3000.
SELECT * FROM emp WHERE sal >= 1500 AND sal <= 3000;

SELECT * FROM emp WHERE sal BETWEEN 1500 AND 3000;

-- display all emps hired in year 1982.
SELECT * FROM emp WHERE hire >= '1982-01-01' AND hire <= '1982-12-31';

SELECT * FROM emp WHERE hire BETWEEN '1982-01-01' AND '1982-12-31';
```

- Paperwork
 - ADAM
 - B
 - BLAKE
 - CLARK
 - J
 - JAMES
 - KING

```
INSERT INTO emp(ename) VALUES ('B'), ('J');

SELECT ename FROM emp ORDER BY ename;

-- display all emps whose names start with 'B' to 'J';
SELECT ename FROM emp WHERE ename BETWEEN 'B' AND 'J';
-- JAMES & JONES are not displayed because alphabetically they come after J.

SELECT ename FROM emp WHERE ename BETWEEN 'B' AND 'K';
-- will display all names starting from B to J + will also display "K" (if present).

SELECT ename FROM emp WHERE ename BETWEEN 'B' AND 'K' AND ename != 'K';
```

```
-- will display all names starting from B to J

-- display all emps whose sal is not in range 1000 to 2000.
SELECT * FROM emp WHERE sal NOT BETWEEN 1000 AND 2000;

-- display all emps whose sal is not in range 1500 to 3000.
SELECT * FROM emp WHERE sal NOT BETWEEN 1500 AND 3000;

-- display all emps between B to K and T to Z;
SELECT * FROM emp WHERE
ename BETWEEN 'B' AND 'K'
OR
ename BETWEEN 'T' AND 'Z';
```

- LIKE operator.
 - Used with strings for finding similar records.
 - Wildcard character to give pattern.
 - % : Any number of any char or Empty.
 - _ : Single occurrence of any char.
 - NOT LIKE : inverse of LIKE

```
-- find all emps whose name start with M.
SELECT * FROM emp WHERE ename LIKE 'M%';

-- find all emps whose name ends with H.
SELECT * FROM emp WHERE ename LIKE '%H';

-- find all emps whose name contain U.
SELECT * FROM emp WHERE ename LIKE '%U%';

-- find all emps whose name contains A twice.
SELECT * FROM emp WHERE ename LIKE '%A%A%';
-- ADAMS --> YES --> %=, %=D, %=MS
-- WARD --> NO
-- ANNA --> YES --> %=, %=NN, %=
-- RAAM --> YES --> %=R, %=, %=M

-- find all emps whose name is start with S to Z.
SELECT * FROM emp WHERE ename BETWEEN 'S' AND 'Z'
OR ename LIKE 'Z%';
```

```
-- display emps having 4 letter name.
SELECT * FROM emp WHERE ename LIKE '____';

-- display all emps whose name contains R on 3rd position.
SELECT * FROM emp WHERE ename LIKE '__R%';
```

```
-- display emps with 4 letter name and 3rd pos is R
SELECT * FROM emp WHERE ename LIKE '___R_';
```

- SELECT cols FROM tablename WHERE condition ORDER BY cols LIMIT n;

```
-- display emp with highest sal in range 1000 to 2000.
SELECT * FROM emp WHERE sal BETWEEN 1000 AND 2000
ORDER BY sal DESC LIMIT 1;

-- display CLERK with min sal.
SELECT * FROM emp
WHERE job = 'CLERK'
ORDER BY sal
LIMIT 1;

-- display fifth lowest sal from dept 20 and 30.
SELECT sal FROM emp WHERE deptno IN (20, 30);

SELECT sal FROM emp WHERE deptno IN (20, 30)
ORDER BY sal;

SELECT DISTINCT sal FROM emp
WHERE deptno IN (20, 30)
ORDER BY sal LIMIT 4, 1;
```

DML - UPDATE

```
SHOW TABLES;

SELECT * FROM stud;

-- change marks to 95.00 for student with id 7.
UPDATE stud SET marks = 95.80 WHERE id = 7;

SELECT * FROM stud;

SELECT id,name,subject,price FROM books;

-- increase price of C Programming books by 50.
UPDATE books SET price = price + 50
WHERE subject = 'C Programming';

SELECT id,name,subject,price FROM books;

-- increase price of all books by 5%.
UPDATE books SET price = price + price * 0.05;

SELECT id,name,subject,price FROM books;
```

```
SELECT * FROM emp;

UPDATE emp SET empno=1, sal=1000, comm=NULL WHERE ename = 'B';

SELECT * FROM emp;
```

DML - DELETE

```
-- delete given rows from emp table -- ename = B / J
DELETE FROM emp WHERE ename IN ('B', 'J');

SELECT * FROM emp;

SELECT * FROM newemp;

-- delete all rows from emp table
DELETE FROM newemp;

SELECT * FROM newemp;

DESCRIBE newemp;

-- delete price of book with id 1001 -- edit/update
UPDATE books SET price = NULL WHERE id = 1001;

SELECT * FROM books;
```

DDL - TRUNCATE

- TRUNCATE is DDL query.
- TRUNCATE is to delete all rows - cannot use WHERE clause.
- Table structure is not deleted (similar to DELETE query).

```
-- delete all books
TRUNCATE TABLE books;

SELECT * FROM books;

DESCRIBE books;
```

DDL - DROP

- DROP command can be used for dropping/deleting whole table or database.
 - DROP TABLE tablename;
 - DROP DATABASE dbname;

```
SELECT * FROM dummy;

DESCRIBE dummy;

DROP TABLE dummy;

SHOW TABLES;

DESCRIBE dummy;
-- error
```

DELETE vs TRUNCATE vs DROP

- DELETE
 - Used to delete rows (not structure).
 - All rows or as per WHERE condition.
 - DML operation
 - DML ops can be rolled back (undo/discard) using transaction.
 - In most RDBMS, slower operation.
- TRUNCATE
 - Used to delete rows (not structure).
 - All rows.
 - DDL operation.
 - DDL ops cannot be rolled back.
 - In most RDBMS, faster operation.
- DROP
 - Used to delete rows as well as struct.
 - DDL operation.
 - DDL ops cannot be rolled back.
 - This is fastest operation.

HELP

```
HELP SELECT;

HELP INSERT;

HELP Functions;

HELP String Functions;

HELP UPPER;

SELECT UPPER('MySQL');

SELECT SUBSTRING('SUNBEAM', 4, 2);
```


DUAL Table

- In Oracle, DUAL table.
 - DUAL table is in memory single row single column virtual table to support SELECT query syntax.
 - SELECT cols FROM tablename;

```
SELECT 2 + 4 * 5 FROM DUAL;  
  
SELECT VERSION() FROM DUAL;  
  
SELECT NOW() FROM DUAL;  
  
SELECT LOWER('SunBeam') FROM DUAL;
```

- One cannot manipulate or drop DUAL table.

```
DESCRIBE DUAL;  
-- error  
  
DROP TABLE DUAL;  
-- error  
  
DELETE FROM DUAL;  
-- error
```

- ANSI SQL influenced by Oracle SQL.
- DUAL table is added in ANSI standard.
- In ANSI SQL, this table is optional.

```
SELECT 2 + 4 * 5;  
  
SELECT VERSION();  
  
SELECT NOW();  
  
SELECT LOWER('SunBeam');
```

SQL Functions

String Functions

```
SELECT LOWER('India');  
  
SELECT ename, LOWER(ename) FROM emp;
```

```
SELECT LEFT('Sunbeam', 2), RIGHT('Sunbeam', 2);

SELECT ename, LEFT(ename, 2), RIGHT(ename, 2) FROM emp;

-- display all emps whose name start with B to K.
SELECT ename, LEFT(ename,1) FROM emp
WHERE LEFT(ename,1) BETWEEN 'B' AND 'K';

SELECT SUBSTRING('SUNBEAM', 2, 3);
-- UNB (+ve pos -- pos from left)

SELECT SUBSTRING('SUNBEAM', -5, 3);
-- NBE (-ve pos -- pos from right)

SELECT SUBSTRING('SUNBEAM', 4, 0);
-- len=0, means no chars after given pos - empty

SELECT SUBSTRING('SUNBEAM', 4, -2);
-- len < 0, means no chars after given pos - empty

-- print chars 2-5 for all emp names.
SELECT ename, SUBSTRING(ename, 2, 4) FROM emp;

SELECT CONCAT('SUNBEAM', ' ', 'INFOTECH');

SELECT CONCAT('SUNBEAM', 2021);

SELECT CONCAT(ename, ' - ', job) FROM emp;

SELECT CONCAT(ename, ' is working in dept ', deptno, ' as ', job, '.') FROM emp;

SELECT LENGTH('   abcd ');

SELECT TRIM('   abcd ');

SELECT LENGTH(TRIM('   abcd '));
-- output of TRIM('   abcd ') is given as input to LENGTH().

SELECT LPAD('Sunbeam', 10, '*');
SELECT RPAD('Sunbeam', 10, '*');

SELECT RPAD(LPAD('Sunbeam', 10, '*'), 13, '*');
```

MySQL - RDBMS

Agenda

- SQL Functions
 - Numeric Functions
 - Date and Time Functions
 - Control Flow Functions
 - Misc/Info Functions
 - List Functions
 - NULL Functions
 - Group Functions
- GROUP BY clause
- HAVING clause

SQL Functions

Numeric Functions

```
SELECT USER(), DATABASE();
-- | sunbeam@localhost | classwork |

HELP Numeric Functions;

SELECT POWER(2, 5);

SELECT SQRT(2);

SELECT RAND();

-- fetch rows in random order
SELECT empno, ename, sal FROM emp
ORDER BY RAND();

SELECT PI();

SELECT ROUND(3.141593, 2), ROUND(3.141593, 4);
-- 3.14, 3.1416

SELECT ROUND(314159.3, -2), ROUND(31415.93, -2);
-- 314200, 31400

SELECT ROUND(3.141593, -2);
-- 0

SELECT ROUND(7246851749, -5);

SELECT * FROM books;
```

```
SELECT id, name, ROUND(price,2) FROM books;
```

- CEIL -- nearest higher integer.
- FLOOR -- nearest lower integer.

```
SELECT CEIL(3.14), FLOOR(3.14);
-- 4, 3
```

```
SELECT CEIL(-3.14), FLOOR(-3.14);
-- -3, -4
```

Date and Time Functions

- DATE -- '1000-01-01' to '9999-12-31'
- TIME -- '-838:59:59' to '838:59:59'
- DATETIME -- '1000-01-01 00:00:00.000000' to '9999-12-31 23:59:59.999999'
- TIMESTAMP -- '1970-01-01 00:00:01' UTC to '2038-01-19 03:14:07' UTC
- YEAR -- '1901' to '2155'

```
HELP Date and Time Functions;
```

```
SELECT NOW(), SLEEP(5), SYSDATE();
```

```
SELECT DATE('2000-05-24 14:47:20'), TIME('2000-05-24 14:47:20');
```

```
SELECT DATE(NOW()), TIME(NOW());
```

```
SELECT DATE_ADD(NOW(), INTERVAL 4 DAY);
```

```
SELECT DATE_ADD(NOW(), INTERVAL 1 MONTH);
```

```
SELECT DATEDIFF(NOW(), '1983-09-28');
```

```
-- return number of days
```

```
SELECT TIMESTAMPDIFF(YEAR, '1983-09-28', NOW());
```

```
-- return number of years
```

```
SELECT ename, hire, TIMESTAMPDIFF(YEAR, hire, NOW()) exp_yrs, TIMESTAMPDIFF(MONTH, hire, NOW()) exp_mons FROM emp;
```

```
SELECT ename, hire, TIMESTAMPDIFF(YEAR, hire, NOW()) exp_yrs, TIMESTAMPDIFF(MONTH, hire, NOW()) % 12 exp_mons FROM emp;
```

```
SELECT DATE_FORMAT(NOW(), '%d-%b-%Y');
```

```
-- 07-Oct-2021
```

```
SELECT DAY(NOW()), MONTH(NOW()), YEAR(NOW()), HOUR(NOW()), MINUTE(NOW()),
```

```
SECOND(NOW()), WEEKDAY(NOW());

-- find all emps hired in 1982
SELECT * FROM emp WHERE YEAR(hire) = 1982;

-- MySQL standard date format: 'yyyy-mm-dd'
-- input date: 'dd/mm/yyyy' -- ??
SET @str = '28/09/2021';
SELECT STR_TO_DATE(@str, '%d/%m/%Y');
```

Control Functions

- IF(condition, expr_if_true, expr_if_false).

```
HELP IF FUNCTION;

-- display ename, sal and category
-- category = RICH if sal > 2500
-- category = POOR if sal <= 2500
SELECT ename, sal, IF(sal > 2500, 'RICH', 'POOR') AS category FROM emp;

-- print number is +ve or -ve or zero.
SET @num = 2;
-- MySQL user-defined variable -- session scope.
-- when MySQL CLI exit, variable will be destroyed

SELECT @num;

SELECT IF(@num > 0, '+ve', IF(@num < 0, '-ve', 'zero'));
```

List Functions

```
SELECT CONCAT('A', 12, 'B', 34.45);

SELECT GREATEST(23, 98, 53, 67);

SELECT LEAST(23, 98, 53, 67);

SELECT name, price, LEAST(price, 700) FROM books;

SELECT GREATEST('AEROPLANE', 'CAR');

SELECT LEAST('AEROPLANE', NULL, 'CAR');

SELECT CONCAT('A', 12, NULL, 'B', 34.45);
```

Misc Functions

```
SELECT VERSION();

SELECT SYSDATE();

SELECT USER(), DATABASE();
```

NULL Functions

- NULL is special value in RDBMS.
- It is irrespective of data type.
- NULL is not 0, 0.0, '\0', 'NULL'.
- NULL represent missing/absent/empty value.

```
SELECT COALESCE(NULL, NULL, 12, 'Nilesh');
-- return = first non-null value.

SELECT COALESCE(12.34, NULL, 'Nilesh');

-- display comm of emp and if no comm then display sal.
SELECT ename, comm, sal, COALESCE(comm, sal) FROM emp;

SELECT ename, comm, sal, IF(comm IS NULL, sal, comm) FROM emp;

SELECT ename, comm, sal, IFNULL(comm, sal) FROM emp;
-- if arg1 == NULL, then result = arg2, else arg1.

SELECT ename, sal, NULLIF(sal, 3000.0) FROM emp;
-- if arg1 == arg2, then result = NULL, else arg1.

SELECT IFNULL(NULL, 'Hello');
```

Group Functions

- Single Row Functions:
 - "n" Input Rows --> "n" Output Rows
 - Function execute once for each row.
- Group Functions/Multi Row Functions/Aggregate Functions
 - "n" Input Rows --> "1" Output Row
 - Aggregate value
 - COUNT(), SUM(), AVG(), MAX(), MIN()
 - STDEV(), COR(), ...

```
SELECT COUNT(sal), SUM(sal), AVG(sal), MAX(sal), MIN(sal) FROM emp;

SELECT COUNT(comm), SUM(comm), AVG(comm), MAX(comm), MIN(comm) FROM emp;
-- NULL values are ignored by Group Functions.
```

```
-- display max income and min income from emp.
-- income = sal + comm
SELECT ename, sal, comm, sal + IFNULL(comm,0) AS income FROM emp;

SELECT MAX(sal + IFNULL(comm,0)), MIN(sal + IFNULL(comm,0)) FROM emp;
```

- GREATEST vs MAX and LEAST vs MIN
 - GREATEST/LEAST -- single row function.
 - MAX/MIN -- group/aggregate function.
 - GREATEST/LEAST -- operate on multiple values from the same row.
 - MAX/MIN -- operate on multiple values in different rows (given column).
 - GREATEST/LEAST -- list function (multiple args).
 - MAX/MIN -- single arg (column name)
 - GREATEST/LEAST -- if any arg is NULL, result is NULL.
 - MAX/MIN -- if any row has NULL value in given column, that will be ignored.

LIMITATIONS OF GROUP FUNCTIONS

```
SET @@sql_mode='ONLY_FULL_GROUP_BY,STRICT_TRANS_TABLES,NO_ENGINE_SUBSTITUTION';

SELECT ename, MAX(sal) FROM emp;
-- error: cannot select any column with group fn.

SELECT LOWER(ename), MAX(sal) FROM emp;
-- error: cannot select single row fn with group fn.

SELECT * FROM emp WHERE sal = MAX(sal);
-- error: cannot use group fn in WHERE clause

SELECT SUM(MAX(sal)) FROM emp;
-- error: cannot nest group fn in each other.
```

- To set sql_mode permanently.
 - step 1: Run Notepad -- "Run as Administrator".
 - step 2: Open my.ini from C:\ProgramData\MySQL\MySQL Server 8.0.
 - step 3: Under [mysqld] change sql_mode.
 - sql-mode=ONLY_FULL_GROUP_BY,STRICT_TRANS_TABLES,NO_ENGINE_SUBSTITUTION
 - step 4: Restart MySQL server (or restart computer).

DQL - SELECT

GROUP BY clause

- By default GROUP functions work on all the rows.
- With GROUP BY we can use GROUP functions on group of rows.

```
SELECT deptno, COUNT(sal), SUM(sal), AVG(sal), MAX(sal), MIN(sal) FROM emp
GROUP BY deptno;
```

```
SELECT job, COUNT(sal), SUM(sal), AVG(sal), MAX(sal), MIN(sal) FROM emp
GROUP BY job;
```

```
SELECT deptno, COUNT(empno) FROM emp
GROUP BY deptno;
```

```
SELECT job, COUNT(empno) FROM emp
GROUP BY job;
```

```
SELECT empno, ename, deptno, job FROM emp
ORDER BY deptno, job;
```

```
SELECT deptno, job, COUNT(empno) FROM emp
GROUP BY deptno, job;
```

```
SELECT deptno, job, COUNT(empno) FROM emp
GROUP BY deptno, job
ORDER BY deptno, job;
```

```
-- deptno  job    count
-- 10      C       1
-- 10      M       1
-- 10      P       1
-- 20      A       2
-- 20      C       2
-- 20      M       1
-- 30      C       1
-- 30      M       1
-- 30      S       4
```

```
SELECT empno, COUNT(empno) FROM emp
GROUP BY empno;
```

```
-- deptwise total sal
SELECT deptno, SUM(sal) FROM emp GROUP BY deptno;
```

```
SELECT SUM(sal) FROM emp GROUP BY deptno;
-- it is not mandatory to project grouped column.
-- however output will be meaningless.
```

```
SELECT deptno, SUM(sal) FROM emp;
-- error: cannot select column with group fn.
```


HAVING clause

- Must be used with GROUP BY clause only.
- Mainly used to apply condition on aggregate values/results.
- HAVING clause vs WHERE clause
 - WHERE clause: evaluated for each row.
 - HAVING clause: evaluated for each group.
 - WHERE clause: can be used with column, single row fn, but not group fn.
 - HAVING clause: can be used with grouped column or group fn, but not on other columns.

```
-- display deptno in which total sal is more than 9000.
```

```
SELECT deptno, SUM(sal) FROM emp
GROUP BY deptno
HAVING SUM(sal) > 9000;
```

```
-- display jobs for which avg sal is more than 2500.
```

```
SELECT job, AVG(sal) FROM emp
GROUP BY job
HAVING AVG(sal) > 2500;
```

```
-- display max sal for each job for emps in deptno 10 and 20.
```

```
SELECT * FROM emp WHERE deptno IN (10,20);
```

```
SELECT job, MAX(sal) FROM emp
WHERE deptno IN (10,20)
GROUP BY job;
```

```
-- display max sal for each job for emps in deptno 10 and 20. display max sal only
if it is more than 2500.
```

```
SELECT job, MAX(sal) FROM emp
WHERE deptno IN (10,20)
GROUP BY job
HAVING MAX(sal) > 2500;
```

```
-- find avg sal for deptno 10 and 20.
```

```
SELECT deptno, AVG(sal) FROM emp
WHERE deptno IN (10, 20)
GROUP BY deptno;
```

```
-- more efficient
```

```
SELECT deptno, AVG(sal) FROM emp
GROUP BY deptno
HAVING deptno IN (10, 20);
```

```
-- less efficient
```

```
-- find the dept that spends max on emp sals.  
SELECT deptno, SUM(sal) FROM emp  
GROUP BY deptno;
```

```
SELECT deptno, SUM(sal) FROM emp  
GROUP BY deptno  
ORDER BY SUM(sal) DESC;
```

```
SELECT deptno, SUM(sal) FROM emp  
GROUP BY deptno  
ORDER BY SUM(sal) DESC  
LIMIT 1;
```

```
-- find the jobs which have lowest avg sal.  
SELECT job, AVG(sal) FROM emp  
GROUP BY job  
ORDER BY AVG(sal)  
LIMIT 1;
```

```
-- find the jobs which have lowest avg income.  
SELECT job, AVG(sal + IFNULL(comm,0.0)) FROM emp  
GROUP BY job  
ORDER BY AVG(sal + IFNULL(comm,0.0))  
LIMIT 1;
```

MySQL - RDBMS

Agenda

- Joins

Case sensitive string comparision

```
SELECT 'SunBeam' = 'SUNBEAM';

SELECT BINARY 'SunBeam' = BINARY 'SUNBEAM';

SELECT BINARY 'SUNBEAM' = BINARY 'SUNBEAM';

SELECT * FROM emp WHERE ename = 'King';
-- case insensitive search

SELECT * FROM emp WHERE BINARY ename = BINARY 'King';
-- case sensitive search -- do not match

SELECT * FROM emp WHERE BINARY ename = BINARY 'KING';
-- case sensitive search -- match
```

Joins

```
USE classwork;

SELECT USER(), DATABASE();
-- sunbeam@localhost, classwork

SOURCE D:/sep21/DAC/dbt/db/joins.sql

SHOW TABLES;

SELECT * FROM emps;

SELECT * FROM depts;

SELECT * FROM addr;

SELECT * FROM meeting;

SELECT * FROM emp_meeting;
```

Cross Joins

```
// for loop
for(int i=0; i<emp.length; i++) {
    Emp e = emp[i];
    for(int j=0; j<dept.length; j++) {
        Dept d = dept[j];
        System.out.println(e.ename + " -- " + d.dname);
    }
}
```

```
// for-each loop
for(Emp e:emp) {
    for(Dept d:dept) {
        System.out.println(e.ename + " -- " + d.dname);
    }
}
```

```
SELECT e.ename, d.dname FROM emps e
CROSS JOIN depts d;

SELECT e.ename, d.dname FROM depts d
CROSS JOIN emps e;
```

```
SELECT e.ename, d.dname FROM emps AS e
CROSS JOIN depts AS d;
-- can use AS keyword for table alias

SELECT emps.ename, depts.dname FROM emps
CROSS JOIN depts;
-- using alias is not mandatory, we can directly use table name

SELECT ename, dname FROM emps
CROSS JOIN depts;
-- if column names are different in both tables, writing alias/tablename is optional

SELECT ename, dname, deptno FROM emps
CROSS JOIN depts;
-- ERROR: Column 'deptno' in field list is ambiguous.

SELECT ename, dname, depts.deptno FROM emps
CROSS JOIN depts;
```

Inner Join

- Joining two tables (Getting data from two tables based on some condition).

- Obviously the table must be related by some way (some column).
 - One DEPT can have Many EMP.
 - Many EMP can be in One DEPT.
 - This relation is established with "deptno" column in depts and in emps.

```
-- display ename and his dname.
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno;

-- display ename and names of dept in which he is not working
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno != d.deptno;
```

Equi-join

- When in any join query condition is of equality, it is referred as "equi-join".

Non-equi-join

- When in any join query condition is of non-equality(<, >, <=, >=, !=), it is referred as "non-equi-join".

Left Outer Join

- Left Join = Intersection + Extra rows from Left table

```
SELECT e.ename, d.dname FROM depts d
LEFT OUTER JOIN emps e ON e.deptno = d.deptno;

SELECT e.ename, d.dname FROM depts d
LEFT JOIN emps e ON e.deptno = d.deptno;
-- OUTER keyword is optional

SELECT e.ename, d.dname FROM emps e
RIGHT OUTER JOIN depts d ON e.deptno = d.deptno;
```

Right Outer Join

- Right Join = Intersection + Extra rows from Right table

```
SELECT e.ename, d.dname FROM depts d
RIGHT OUTER JOIN emps e ON e.deptno = d.deptno;

SELECT e.ename, d.dname FROM depts d
RIGHT JOIN emps e ON e.deptno = d.deptno;
-- OUTER keyword is optional
```

```
SELECT e.ename, d.dname FROM emps e
LEFT JOIN depts d ON e.deptno = d.deptno;
```

Full Outer Join

- Full Join = Intersection + Extra rows from Left table + Extra rows from Right table
- Full Outer Join is not supported in MySQL. It can work well in Oracle, MS-SQL, ...

```
SELECT e.ename, d.dname FROM emps e
FULL OUTER JOIN depts d ON e.deptno = d.deptno;
```

Set Operators

- Used to combine results of two queries (if output contains same number of columns).

```
(SELECT dname AS name FROM depts)
UNION ALL
(SELECT ename FROM emps);
```

```
(SELECT sal FROM emp)
UNION ALL
(SELECT price FROM books);
```

```
(SELECT e.ename, d.dname FROM emps e
LEFT OUTER JOIN depts d ON e.deptno = d.deptno)
UNION ALL
(SELECT e.ename, d.dname FROM emps e
RIGHT OUTER JOIN depts d ON e.deptno = d.deptno);
```

```
(SELECT e.ename, d.dname FROM emps e
LEFT OUTER JOIN depts d ON e.deptno = d.deptno)
UNION
(SELECT e.ename, d.dname FROM emps e
RIGHT OUTER JOIN depts d ON e.deptno = d.deptno);
-- simulation of full outer join in MySQL
```

Self Join

```
-- print ename and his manager name.
SELECT e.ename, m.ename AS mname FROM emps e
```

```
INNER JOIN emps m ON e.mgr = m.empno;

-- print ename and his manager name.
SELECT e.ename, m.ename AS mname FROM emps e
LEFT JOIN emps m ON e.mgr = m.empno;
```

Joins Practice

```
-- display ename, emp's dname and emp's dist.
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno;

SELECT e.ename, d.dname, a.dist FROM emps e
LEFT JOIN depts d ON e.deptno = d.deptno
INNER JOIN addr a ON e.empno = a.empno;
```

```
-- display ename and his meeting topics.
SELECT * FROM emps;

SELECT * FROM meeting;

SELECT * FROM emp_meeting;

SELECT e.ename, m.topic FROM emp_meeting em
INNER JOIN emps e ON em.empno = e.empno
INNER JOIN meeting m ON em.meetno = m.meetno;
```

-- emps are travelling from their home town to attend few meetings. Display name of emp and meeting topic and from where he is travelling.

```
SELECT e.ename, m.topic, a.dist, a.tal
FROM emp_meeting em
INNER JOIN emps e ON em.empno = e.empno
INNER JOIN meeting m ON em.meetno = m.meetno
INNER JOIN addr a ON e.empno = a.empno;
```

-- emps are representing their depts in few meetings. Display name of emp and meeting topic and their dept.

```
SELECT e.ename, m.topic, d.dname
FROM emp_meeting em
INNER JOIN emps e ON em.empno = e.empno
INNER JOIN meeting m ON em.meetno = m.meetno
LEFT JOIN depts d ON e.deptno = d.deptno;
```

```
-- print dname and number (count) of emps in that dept.
SELECT deptno, COUNT(empno) FROM emps
GROUP BY deptno;

SELECT d.dname, COUNT(e.empno) FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno
GROUP BY d.dname;

SELECT d.dname, COUNT(e.empno) FROM emps e
RIGHT JOIN depts d ON e.deptno = d.deptno
GROUP BY d.dname;
```

```
-- display emps and their number of meetings in desc order of meeting count.
SELECT e.ename, m.topic FROM emp_meeting em
INNER JOIN emps e ON em.empno = e.empno
INNER JOIN meeting m ON em.meetno = m.meetno;

SELECT em.empno, COUNT(em.meetno)
FROM emp_meeting em
GROUP BY em.empno;

SELECT e.ename, COUNT(em.meetno)
FROM emp_meeting em
INNER JOIN emps e ON e.empno = em.empno
GROUP BY e.ename;

SELECT e.ename, COUNT(em.meetno)
FROM emp_meeting em
INNER JOIN emps e ON e.empno = em.empno
GROUP BY e.ename
ORDER BY COUNT(em.meetno) DESC;
```

```
-- display all emps in DEV dept.
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno;

SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno
WHERE d.dname = 'DEV';
```

```
SELECT columns FROM table1
xxx JOIN table2 ON condition
xxx JOIN table3 ON condition ...
WHERE condition
GROUP BY column
HAVING condition
```



```
ORDER BY column
LIMIT n;
```

Non-standard joins

```
-- display ename and dname.
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno;

SELECT e.ename, d.dname FROM emps e
JOIN depts d ON e.deptno = d.deptno;
-- by default join is INNER (in MySQL).

SELECT e.ename, d.dname FROM emps e
CROSS JOIN depts d ON e.deptno = d.deptno;
-- In MySQL, we can apply condition on CROSS JOIN

SELECT e.ename, d.dname FROM emps e
CROSS JOIN depts d WHERE e.deptno = d.deptno;
-- You may use WHERE clause with CROSS JOIN
-- However choose INNER JOIN if applicable.

SELECT e.ename, d.dname FROM emps e, depts d
WHERE e.deptno = d.deptno;
-- Join without JOIN keyword is old-style join.

SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d USING (deptno);
-- joined columns from both tables have SAME name
-- the condition can be given using USING keyword
-- USING (colname) --> t1.colname = t2.colname;
-- This is always eqi-join.
-- This works only in MySQL.

SELECT e.ename, d.dname FROM emps e
NATURAL JOIN depts d;
-- num of joined columns = 1 (same name)
-- NATURAL JOIN = Implicit Join Condition
-- Equality Check of Columns with Same Name in Both tables.
-- In this example
-- NATURAL JOIN: ON e.deptno = d.deptno.

-- display all possible depts for Amit & Nilesh.
SELECT e.ename, d.dname FROM emps e
CROSS JOIN depts d WHERE e.ename IN ('AMIT', 'NILESH');
```

Natural Join

- table1: a, b, c, d
- table2: a, b, x, y

- `table1 NATURAL JOIN table2 --> ON t1.a = t2.a AND t1.b = t2.b;`

MySQL - RDBMS

Agenda

- Sub-queries
- Views

Q & A

```
SELECT deptno, SUM(sal) FROM emp
WHERE job != 'MANAGER'
GROUP BY deptno;
```

```
UPDATE dept d
INNER JOIN emp e ON e.deptno = d.deptno
SET d.dname = 'ACCOUNTING'
WHERE e.ename = 'KING';

SELECT * FROM dept;
```

```
SELECT d.deptno, d.dname, SUM(e.sal) FROM emp e
INNER JOIN dept d ON e.deptno = d.deptno
GROUP BY d.deptno, d.dname;
```

Assignment 6

```
SELECT o.onum, c.cname FROM orders o
INNER JOIN customers c ON o.cnum = c.cnum;
```

```
SELECT o.onum, c.cname, s.sname FROM orders o
INNER JOIN customers c ON o.cnum = c.cnum
INNER JOIN salespeople s ON o.snum = s.snum;
```

```
-- wrong query
SELECT o.onum, c.cname, s.sname FROM orders o
INNER JOIN customers c ON o.cnum = c.cnum
INNER JOIN salespeople s ON c.snum = s.snum;
```

```
SELECT c.cname, s.sname, s.comm FROM customers c
INNER JOIN salespeople s ON c.snum = s.snum
WHERE s.comm > 0.12;
```

```
SELECT o.onum, s.sname, o.amt * s.comm FROM orders o
INNER JOIN customers c ON o.cnum = c.cnum
INNER JOIN salespeople s ON o.snum = s.snum
WHERE c.rating > 100;
```

```
SELECT s1.snum snum1, s1.sname sname1, s2.snum s2, s2.sname sname2 FROM
salespeople s1
INNER JOIN salespeople s2 ON s1.city = s2.city
WHERE s1.snum < s2.snum;
```

RDBMS Preparations

- Interview preparations -- SQL queries -- Joins, Sub-queries
 - Assignments
 - Assessments
 - Hackerrank -- SQL -- Easy & Medium.
- CCEE preparations -- MCQ
 - Moodle MCQ activity
 - Data entry till end of module -- After that read-only
 - Search functionality -- Topicwise, Wordwise.
- Lab exams
 - Lab Assignments & Assessments

Sub-queries

Single row sub-queries

- Return single row from inner query.

```
USE classwork;

-- display emp with max sal.

SELECT * FROM emp ORDER BY sal DESC LIMIT 1;
-- will produce partial output if multiple emps have same max sal.

SELECT * FROM emp WHERE sal = MAX(sal);
-- error: Group fns cannot be used in WHERE clause

SET @maxsal = (SELECT MAX(sal) FROM emp);
```

```

SELECT @maxsal;
SELECT * FROM emp WHERE sal = @maxsal;

SELECT * FROM emp
WHERE sal = (SELECT MAX(sal) FROM emp);

-- display emp with second highest sal
SELECT * FROM emp ORDER BY sal DESC LIMIT 1, 1;

SET @sal2 = (SELECT DISTINCT sal FROM emp ORDER BY sal DESC LIMIT 1,1);
SELECT @sal2;
SELECT * FROM emp WHERE sal = @sal2;

SELECT * FROM emp WHERE sal = (SELECT DISTINCT sal FROM emp ORDER BY sal DESC
LIMIT 1,1);

-- display emp with third highest sal
SELECT * FROM emp WHERE sal = (SELECT DISTINCT sal FROM emp ORDER BY sal DESC
LIMIT 2,1);

-- display emps working in dept of 'KING'
SET @dno = (SELECT deptno FROM emp WHERE ename = 'KING');
SELECT * FROM emp WHERE deptno = @dno;

SELECT * FROM emp WHERE deptno = (SELECT deptno FROM emp WHERE ename = 'KING');

SELECT * FROM emp WHERE deptno = (SELECT deptno FROM emp WHERE ename = 'KING') AND
ename != 'KING';

```

Multi-row sub-queries

- Inner query returns multiple rows.
- They can be compared using ANY, ALL or IN operator.

```

-- find all emps having sal more than all salesman.
SELECT sal FROM emp WHERE job = 'SALESMAN';

SELECT MAX(sal) FROM emp WHERE job = 'SALESMAN';

SELECT * FROM emp WHERE sal > (SELECT MAX(sal) FROM emp WHERE job = 'SALESMAN');

SELECT * FROM emp WHERE sal >
ALL(SELECT sal FROM emp WHERE job = 'SALESMAN');
-- (sal > 1600 AND sal > 1250 AND sal > 1250 AND sal > 1500)

```

```

-- find emp with sal less than sal of "any" emp in deptno=20
SELECT sal FROM emp WHERE deptno = 20;

SELECT * FROM emp WHERE sal < (SELECT MAX(sal) FROM emp WHERE deptno = 20);

```

```
SELECT * FROM emp WHERE sal < ANY(SELECT sal FROM emp WHERE deptno = 20);
-- sal < 800 OR sal < 2975 OR sal < 3000 OR sal < 1100 OR sal < 3000.
```

```
-- display the depts which have emps.
SELECT deptno FROM emp;

SELECT * FROM dept WHERE deptno = ANY(SELECT deptno FROM emp);
-- deptno = 10 OR deptno = 20 OR deptno = 30

SELECT * FROM dept WHERE deptno IN (SELECT deptno FROM emp);
-- deptno = 10 OR deptno = 20 OR deptno = 30
```

- ANY vs IN operator
 - ANY can be used in sub-queries only, while IN can be used with/without sub-queries.
 - ANY can be used for comparison (<, >, <=, >=, =, or !=), while IN can be used only for equality comparison (=).
 - Both operators are logically similar to OR operator.
- ANY vs ALL operator
 - Both can be used for comparison (<, >, <=, >=, =, or !=).
 - Both are usable only with sub-queries.
 - ANY is similar to logical OR, while ALL is similar to logical AND.

```
-- display depts which do not have any emp.

SELECT * FROM dept WHERE deptno NOT IN (SELECT deptno FROM emp);
-- deptno != 10 AND deptno != 20 AND deptno != 30
-- NOT (deptno = 10 OR deptno = 20 OR deptno = 30)

SELECT * FROM dept WHERE deptno != ALL(SELECT deptno FROM emp);
```

Corelated sub-queries

- SELECT ... FROM table WHERE col = (SELECT ...)
- By default inner query is executed for each row of the outer query.
- If no optimization settings are enabled, sub-queries are slower than joins.

```
SELECT * FROM dept WHERE deptno IN (SELECT deptno FROM emp);
-- 10, ACC --> SELECT deptno FROM emp
-- 20, RES --> SELECT deptno FROM emp
-- 30, SAL --> SELECT deptno FROM emp
-- 40, OPS --> SELECT deptno FROM emp
```

- The sub-query execution can be speed-up if inner queries return/process less number rows.

```
SELECT * FROM dept WHERE deptno IN (SELECT DISTINCT deptno FROM emp);
```

- This is typically done by using WHERE clause in inner query.
- The WHERE clause in inner query may depend on current row of the outer query. This kind of query is called as "co-related sub-query".

```
SELECT * FROM dept d WHERE d.deptno IN
(SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno);
-- 10, ACC --> SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno --> 3 rows
(10,10,10)
-- 20, RES --> SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno --> 5 rows
(20,20,20,20,20)
-- 30, SAL --> SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno --> 6 rows
(30,30,30,30,30,30)
-- 40, OPS --> SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno --> 0 rows
```

```
SELECT * FROM dept d WHERE EXISTS
(SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno);
-- EXISTS check if sub-query return row(s).
```

```
-- display all depts which do not contain any emp.
SELECT * FROM dept d WHERE NOT EXISTS
(SELECT e.deptno FROM emp e WHERE e.deptno = d.deptno);
```

Sub-query in projection

```
-- display number of emps in each dept along with total number of emps.
SELECT deptno, COUNT(empno) FROM emp
GROUP BY deptno;
```

```
SELECT COUNT(empno) FROM emp;
```

```
(SELECT deptno, COUNT(empno) FROM emp
GROUP BY deptno)
```

```
UNION
```

```
(SELECT NULL, COUNT(empno) FROM emp);
```

```
-- +-----+-----+
-- | deptno | COUNT(empno) |
-- +-----+-----+
-- |      20 |           5 |
```

```
-- |      30 |      6 |
-- |      10 |      3 |
-- |     NULL |     14 |
-- +-----+-----+

SELECT deptno,
COUNT(empno) AS deptcnt,
(SELECT COUNT(empno) FROM emp) AS totalcnt
FROM emp
GROUP BY deptno;

-- +-----+-----+-----+
-- | deptno | deptcnt | totalcnt |
-- +-----+-----+-----+
-- |      20 |      5 |      14 |
-- |      30 |      6 |      14 |
-- |      10 |      3 |      14 |
-- +-----+-----+-----+
```

Sub-query in FROM clause

- Inner-query can be written in FROM clause of SELECT statement. The output of the inner query is treated as a table (MUST give table alias) and outer query execute on that table.
- This is called as "Derived table" or "Inline view".

```
-- display empno, ename, sal and category of emp.
-- < 1500 -> POOR, 1500-2500 -> MID, > 2500 -> RICH
SELECT empno, ename, sal, CASE
WHEN sal < 1500 THEN 'POOR'
WHEN sal BETWEEN 1500 AND 2500 THEN 'MIDDLE'
ELSE 'RICH'
END AS category
FROM emp;
```

```
-- display number of emps in each category.
SELECT category, COUNT(empno) FROM
(SELECT empno, ename, sal, CASE
WHEN sal < 1500 THEN 'POOR'
WHEN sal BETWEEN 1500 AND 2500 THEN 'MIDDLE'
ELSE 'RICH'
END AS category
FROM emp) AS emp_category
GROUP BY category;
```

Sub-query in DML


```
-- INSERT new emp with name 'JOHN' and sal=2000 in dept 'OPERATIONS'.
SELECT deptno FROM dept WHERE dname = 'OPERATIONS';

INSERT INTO emp(ename, sal, deptno) VALUES ('JOHN', 2000, (SELECT deptno FROM dept
WHERE dname = 'OPERATIONS'));

SELECT * FROM emp;
```

```
-- give comm 100 to all emps in OPERATIONS dept.
UPDATE emp SET comm = 100 WHERE deptno =
(SELECT deptno FROM dept WHERE dname = 'OPERATIONS');

-- UPDATE emp SET comm = 100 WHERE deptno = 40;

SELECT * FROM emp;
```

```
-- delete emps from OPERATIONS dept.
DELETE FROM emp WHERE deptno =
(SELECT deptno FROM dept WHERE dname = 'OPERATIONS');

-- DELETE FROM emp WHERE deptno = 40;
```

- In MySQL, DML cannot be performed on the table from which inner query is selecting.

```
INSERT INTO emp(empno,ename,sal) VALUES(1000, 'JACK', 6000);

-- delete emp with max sal.
DELETE FROM emp
WHERE sal = (SELECT MAX(sal) FROM emp);
-- error: not allowed in MySQL

SET @maxsal = (SELECT MAX(sal) FROM emp);
DELETE FROM emp WHERE sal = @maxsal;
```

SQL Performance

- Modern RDBMS implement lot of optimization mechanisms (internally based on data structures and algorithms) like caching (materialization), semijoins or hashjoins, etc.
- These optimization logic will differ from RDBMS to RDBMS.

```
SELECT @@optimizer_switch;
```

- In MySQL, query Performance is measured in terms of query cost. Lower the cost, better is performance.
- The cost of query depends on data in the table(s), MySQL version, server machine config, optimizer settings, etc.

```
EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE sal < (SELECT MAX(sal) FROM emp WHERE deptno = 20);
-- 1.65

EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE sal < ANY(SELECT sal FROM emp WHERE deptno = 20);
-- 1.65
```

Views

- View is projection of the data.
- CREATE VIEW viewname AS SELECT ...;
- In MySQL views are non-materialized i.e. output of SELECT statement (of view) is not saved in server disk. Only SELECT query is saved in server disk.

```
SELECT empno, ename, sal, CASE
WHEN sal < 1500 THEN 'POOR'
WHEN sal BETWEEN 1500 AND 2500 THEN 'MIDDLE'
ELSE 'RICH'
END AS category
FROM emp;

CREATE VIEW v_empcategory AS
SELECT empno, ename, sal, CASE
WHEN sal < 1500 THEN 'POOR'
WHEN sal BETWEEN 1500 AND 2500 THEN 'MIDDLE'
ELSE 'RICH'
END AS category
FROM emp;

SHOW TABLES;

SHOW FULL TABLES;

SELECT * FROM v_empcategory;

SELECT category, COUNT(empno) FROM v_empcategory
GROUP BY category;

EXPLAIN FORMAT=JSON
SELECT category, COUNT(empno) FROM v_empcategory
GROUP BY category;
```

```
SHOW CREATE VIEW v_empcategory;

SELECT * FROM v_empcategory;

ALTER VIEW v_empcategory AS
SELECT empno, sal, CASE
WHEN sal < 1500 THEN 'POOR'
WHEN sal BETWEEN 1500 AND 2500 THEN 'MIDDLE'
ELSE 'RICH'
END AS category
FROM emp;

SELECT * FROM v_empcategory;

DROP VIEW v_empcategory;

SHOW FULL TABLES;
```

```
CREATE VIEW v_empsal AS
SELECT empno, ename, sal FROM emp;

SELECT * FROM v_empsal;

CREATE VIEW v_richemp AS
SELECT * FROM emp WHERE sal > 2500;

SELECT * FROM v_richemp;

CREATE VIEW v_empincome AS
SELECT empno, ename, sal, comm, sal + IFNULL(comm,0.0) income FROM emp;

SELECT * FROM v_empincome;

CREATE VIEW v_empjobsummary AS
SELECT job, SUM(sal) salsum, AVG(sal) salavg, MAX(sal) salmax, MIN(sal) salmin
FROM emp
GROUP BY job;

SELECT * FROM v_empjobsummary;

INSERT INTO emp(empno,ename,job,sal,deptno) VALUES (1000, 'JILL', 'DEV', 2000,
40);

SELECT * FROM v_empjobsummary;
```

```
SELECT * FROM v_empsal;

INSERT INTO v_empsal VALUES (1001, 'JOY', 2300);
```

```
-- DML ops are allowed on Simple Views.
```

```
SELECT * FROM emp;
```

```
SELECT * FROM v_empjobsummary;
```

```
DELETE FROM v_empjobsummary WHERE job = 'DEV';
```

```
-- DML ops are NOT allowed on Complex Views.
```

MySQL - RDBMS

Agenda

- Views
- Security
- Transactions
- Row locking
- Indexes

Views

- Projection of data -- SELECT.
- CREATE VIEW viewname AS SELECT ...
- MySQL views are non-materialized.
- When query is executed on view, data is read from target table.
- Simple view: DQL + DML
- Complex view: DQL

```
SELECT USER(), DATABASE();
-- | sunbeam@localhost | classwork |

SHOW FULL TABLES;

SHOW CREATE VIEW v_richemp;

SELECT empno, ename, sal FROM v_richemp;

INSERT INTO v_richemp(empno, ename, sal) VALUES(1000, 'JAMES', 2600);

SELECT empno, ename, sal FROM v_richemp;

INSERT INTO v_richemp(empno, ename, sal) VALUES(1001, 'HARRY', 2200);

SELECT empno, ename, sal FROM v_richemp;

SELECT empno, ename, sal FROM emp;

EXPLAIN FORMAT=JSON
SELECT empno, ename, sal FROM v_richemp;

INSERT INTO v_richemp(empno, ename, sal) VALUES(1002, 'MERRY', 2000);

ALTER VIEW v_richemp AS
SELECT * FROM emp WHERE sal > 2500
WITH CHECK OPTION;

INSERT INTO v_richemp(empno, ename, sal) VALUES(1003, 'ADAM', 1500);
-- error: CHECK OPTION failed -- sal <= 2500.
```

```
INSERT INTO v_richemp(empno, ename, sal) VALUES(1004, 'EVE', 2800);
```

```
CREATE VIEW v_richemp2 AS
SELECT empno, ename, sal, comm FROM v_richemp;

SELECT * FROM v_richemp2;

SHOW CREATE VIEW v_richemp2;

DROP VIEW v_richemp;

SELECT * FROM v_richemp2;
-- error: invalid table/view.

DROP VIEW v_richemp2;
```

```
SELECT e.empno, e.ename, d.deptno, d.dname FROM emp e INNER JOIN dept d ON
e.deptno = d.deptno;

CREATE VIEW v_empdept AS
SELECT e.empno, e.ename, d.deptno, d.dname FROM emp e INNER JOIN dept d ON
e.deptno = d.deptno;

SELECT * FROM v_empdept;

-- display all emps in ACCOUNTING dept.
SELECT e.empno, e.ename, d.deptno, d.dname FROM emp e INNER JOIN dept d ON
e.deptno = d.deptno
WHERE d.dname = 'ACCOUNTING';

SELECT * FROM v_empdept
WHERE dname = 'ACCOUNTING';

EXPLAIN FORMAT=JSON
SELECT e.empno, e.ename, d.deptno, d.dname FROM emp e INNER JOIN dept d ON
e.deptno = d.deptno
WHERE d.dname = 'ACCOUNTING';

EXPLAIN FORMAT=JSON
SELECT * FROM v_empdept
WHERE dname = 'ACCOUNTING';
```

- Assign: sales database
 - create view on all three tables joined together with all columns.
 - solve join assignments on that view.

JSON = Java Script Object Notation

- Data format to represent the data
- Another data format example is XML, CSV, ...
- JSON --> C struct like -- key-value pairs

```
"table": {  
  "table_name": "d",  
  "access_type": "ALL",  
  "rows_examined_per_scan": 4,  
  "rows_produced_per_join": 1,  
}
```

- { ... } --> object
- [...] --> array
- "..." --> string
- 123.34, true, null --> number, bool

DCL

```
sunbeam> SHOW DATABASES;  
-- classwork, sales, hr  
  
root> SHOW GRANTS FOR sunbeam@localhost;  
  
root> REVOKE ALL PRIVILEGES ON hr.* FROM sunbeam@localhost;  
  
sunbeam> SHOW DATABASES;  
-- classwork, sales
```

```
root> CREATE USER 'mgr'@'%' IDENTIFIED BY 'mgr';  
  
root> CREATE USER 'teamlead'@'localhost' IDENTIFIED BY 'teamlead';  
  
root> CREATE USER 'dev1'@'localhost' IDENTIFIED BY 'dev1';  
  
root> CREATE USER 'dev2'@'localhost' IDENTIFIED BY 'dev2';  
  
root> SELECT user, host FROM mysql.user;  
  
root> SHOW GRANTS FOR 'mgr'@'%';  
-- USAGE -- no permissions  
  
mgr> SHOW DATABASES;  
  
root> GRANT ALL ON classwork.* TO 'mgr'@'%' WITH GRANT OPTION;  
  
mgr> SHOW DATABASES;
```

```
mgr> GRANT INSERT, UPDATE, DELETE, SELECT ON classwork.emp TO
'teamlead'@'localhost';

mgr> GRANT INSERT, UPDATE, DELETE, SELECT ON classwork.dept TO
'teamlead'@'localhost';

mgr> GRANT INSERT, SELECT ON classwork.books TO 'teamlead'@'localhost';

root> SHOW GRANTS FOR 'teamlead'@'localhost';
```

```
teamlead> SHOW DATABASES;

teamlead> USE classwork;

teamlead> SHOW TABLES;

teamlead> DELETE FROM books;
-- error: Access denied

teamlead> INSERT INTO books VALUES (1, 'Atlas Shrugged', 'Ayn Rand', 'Novell',
727.29);

teamlead> SELECT * FROM books;
```

```
mgr> USE classwork;

mgr> CREATE VIEW v_empsummary AS
SELECT job, SUM(sal) salsum, AVG(sal) salavg FROM emp GROUP BY job;

mgr> GRANT SELECT ON classwork.v_empsummary TO 'dev1'@'localhost';

mgr> CREATE VIEW v_deptsummary AS
SELECT d.dname, SUM(e.sal) salsum, AVG(e.sal) salavg FROM emp e RIGHT JOIN dept d
ON e.deptno = d.deptno GROUP BY d.dname;

mgr> GRANT SELECT ON classwork.v_deptsummary TO 'dev1'@'localhost';

dev1> SHOW DATABASES;

dev1> USE classwork;

dev1> SHOW FULL TABLES;

dev1> SELECT * FROM v_deptsummary;

dev1> DESCRIBE v_deptsummary;
```



```
dev1> SHOW CREATE VIEW v_deptsummary;
-- error: Access denied.
```

```
root> SHOW GRANTS FOR 'teamlead'@'localhost';

root> REVOKE SELECT ON classwork.books FROM 'teamlead'@'localhost';

root> SHOW GRANTS FOR 'teamlead'@'localhost';

teamlead> SELECT * FROM books;
-- error: Access denied.
```

Transaction

```
-- sunbeam@localhost, classwork

CREATE TABLE accounts(id INT PRIMARY KEY, type CHAR(20),balance DECIMAL(9,2));

INSERT INTO accounts VALUES
(1, 'Saving', 20000.00),
(2, 'Current', 60000.00),
(3, 'Saving', 5000.00),
(4, 'Saving', 3000.00),
(5, 'Current', 50000.00),
(6, 'Saving', 10000.00);

SELECT * FROM accounts;

-- transfer Rs. 2000 from acc 1 to acc 3.
UPDATE accounts SET balance = balance - 2000
WHERE id = 1;

UPDATE accounts SET balance = balance + 2000
WHERE id = 3;

SELECT * FROM accounts;
```

- In MySQL, by default each DML operation is executed as a transaction with the single query and it is auto-committed.
- In MySQL, transaction is explicitly started using START TRANSACTION command. It is finalized using COMMIT command or disarded using ROLLBACK command.
- TCL commands
 - START TRANSACTION
 - COMMIT
 - ROLLBACK

```
-- transfer Rs. 5000 from acc 6 to acc 4.  
START TRANSACTION;  
  
UPDATE accounts SET balance = balance - 5000  
WHERE id = 6;  
  
SELECT * FROM accounts;  
  
UPDATE accounts SET balance = balance + 5000  
WHERE id = 4;  
  
SELECT * FROM accounts;  
  
COMMIT;
```

```
-- transfer Rs. 2000 from acc 6 to acc 4.  
START TRANSACTION;  
  
UPDATE accounts SET balance = balance - 2000  
WHERE id = 6;  
  
SELECT * FROM accounts;  
  
UPDATE accounts SET balance = balance + 2000  
WHERE id = 4;  
  
SELECT * FROM accounts;  
  
ROLLBACK;  
  
SELECT * FROM accounts;
```

```
START TRANSACTION;  
  
DELETE FROM accounts;  
  
SELECT * FROM accounts;  
  
ROLLBACK;  
  
SELECT * FROM accounts;
```

```
DELETE FROM accounts;  
-- delete without transaction  
-- single dml query tx --> auto-committed.  
-- START TX ++ DELETE ++ COMMIT.
```

```
-- all rows are deleted permanently.  
  
ROLLBACK;  
-- useless -- no current tx.  
  
SELECT * FROM accounts;
```

- Transaction

```
START TRANSACTION;  
  
dml1;  
dml2;  
dml3;  
  
COMMIT;  
-- changes of dml 1,2,3 are permanent.  
-- transaction is completed.
```

```
START TRANSACTION;  
  
dml1;  
dml2;  
dml3;  
  
ROLLBACK;  
-- changes of dml 1,2,3 are discarded.  
-- transaction is completed.
```

- In MySQL, all DML operations are auto-committed (for each query).
- However MySQL can be configured to start transaction automatically after commit/rollback of previous transaction.

```
SELECT @@autocommit;  
  
SET @@autocommit=0;  
  
SELECT @@autocommit;  
  
SELECT * FROM books;  
  
DELETE FROM books;  
  
ROLLBACK;  
  
SELECT * FROM books;
```

```
SELECT * FROM dept;

DELETE FROM dept;

SELECT * FROM dept;

ROLLBACK;

SELECT * FROM dept;

SELECT * FROM accounts;

INSERT INTO accounts VALUES
(1, 'Saving', 20000.00),
(2, 'Current', 60000.00),
(3, 'Saving', 5000.00),
(4, 'Saving', 3000.00),
(5, 'Current', 50000.00),
(6, 'Saving', 10000.00);

SELECT * FROM accounts;

COMMIT;

SET @@autocommit=1;

SELECT @@autocommit;
```

- Transaction is set of DML queries executed as a single unit.
- Transaction is limited to same RDBMS server.

```
START TRANSACTION;

SELECT * FROM emp;

DELETE FROM emp WHERE empno < 1100;
-- changed in tx -- not permanent

SELECT * FROM emp;

TRUNCATE TABLE books;
-- ddl command -- current tx is auto committed.
-- DELETE FROM emp ... changes are permanent

SELECT * FROM books;

ROLLBACK;
-- has no effect ... tx is already committed.

SELECT * FROM emp;
```

```
START TRANSACTION;  
  
SELECT * FROM dept;  
  
DELETE FROM dept;  
  
SELECT * FROM dept;  
  
EXIT;  
-- auto rollback current tx.
```

```
SELECT * FROM dept;
```

MySQL - RDBMS

Agenda

- Transactions
- Locking
- Indexes
- Constraints
- ALTER TABLE

Transactions

- In applications transactions are done programmatically.
 - Example 1: Funds transfer
 - accounts
 - Example 2: Online Ordering System
 - orders
 - order_items
 - payments
- JDBC -- Java database connectivity

```
try {  
    // create connection  
    con = DriverManager.getConnection(...);  
    // start transaction  
    con.setAutoCommit(false);  
    // create statements  
    stmt1 = con.prepareStatement("INSERT ...");  
    stmt1.executeUpdate();  
    stmt2 = con.prepareStatement("INSERT ...");  
    stmt2.executeUpdate();  
    stmt3 = con.prepareStatement("INSERT ...");  
    stmt3.executeUpdate();  
    // commit transaction  
    con.commit();  
}  
catch(Exception e) {  
    // rollback transaction  
    con.rollback();  
}
```

- MySQL prompt

```
START TRANSACTION;
```

```
INSERT INTO orders VALUES(...);

INSERT INTO order_items VALUES(...);
INSERT INTO order_items VALUES(...);
INSERT INTO order_items VALUES(...);

INSERT INTO payments VALUES(...);

COMMIT; -- OR ROLLBACK;
```

Savepoints

- Savepoint is state of database within a transaction.
- User can rollback to any of the savepoint using
 - ROLLBACK TO sa;
- In this case all changes after savepoint "sa" are discarded.
- ROLLBACK TO "sa" do not ROLLBACK the whole transaction. The same transaction can be continued further (can make more DML queries).
- Transaction is completed only when COMMIT or ROLLBACK is done. All savepoint memory is released.
- COMMIT is not allowed upto a savepoint. We can only commit whole transaction.

```
START TRANSACTION;

INSERT INTO orders VALUES(...);
SAVEPOINT sa1;

INSERT INTO order_items VALUES(...);
INSERT INTO order_items VALUES(...);
INSERT INTO order_items VALUES(...);
SAVEPOINT sa2;

INSERT INTO payments VALUES(...);

ROLLBACK TO sa2;
-- revert db state back to sa2.
-- only payments query will be rolledback.
-- transaction is not yet completed

INSERT INTO payments VALUES(...);
-- continue ops in same transaction

COMMIT; -- or ROLLBACK
-- transaction is completed.
```

```
START TRANSACTION;

INSERT INTO orders VALUES(...);
SAVEPOINT sa1;
```

```
INSERT INTO order_items VALUES(...);
INSERT INTO order_items VALUES(...);
INSERT INTO order_items VALUES(...);
SAVEPOINT sa2;

INSERT INTO payments VALUES(...);

ROLLBACK TO sa1;
-- revert db state back to sa1.
-- order_items & payment queries are rolled back.

INSERT INTO order_items VALUES(...);
INSERT INTO payments VALUES(...);

COMMIT; -- or ROLLBACK
```

Transaction properties/characteristics

- Atomicity: All DML queries in transaction will be successful or failed/discarded. Partial transaction never committed.
- Consistent: At the end of transaction same state is visible to all the users.
- Isolation: Each transaction is isolated from each other. All transactions are added in a transaction queue at server side and process sequentially.
- Durability: At the end of transaction, final state is always saved (on server side).

Transaction Internals

```
root> SELECT * FROM accounts;

sunbeam> SELECT * FROM accounts;

sunbeam> START TRANSACTION;

sunbeam> DELETE FROM accounts WHERE id = 6;

sunbeam> SELECT * FROM accounts;
-- changes visible in current transaction

root> SELECT * FROM accounts;
-- changes not visible in other transactions

sunbeam> COMMIT;

sunbeam> SELECT * FROM accounts;

root> SELECT * FROM accounts;
-- changes are visible to other users after commit is done
```



```
sunbeam> START TRANSACTION;

sunbeam> DELETE FROM accounts WHERE id = 5;

sunbeam> SELECT * FROM accounts;

root> SELECT * FROM accounts;

sunbeam> ROLLBACK;

sunbeam> SELECT * FROM accounts;

root> SELECT * FROM accounts;
```

```
sunbeam> START TRANSACTION;

sunbeam> DELETE FROM accounts WHERE id = 5;

sunbeam> SELECT * FROM accounts;

root> SELECT * FROM accounts;

root> DELETE FROM accounts WHERE id = 4;
-- single dml tx -- autocommitted

root> SELECT * FROM accounts;

sunbeam> SELECT * FROM accounts;

sunbeam> COMMIT;

sunbeam> SELECT * FROM accounts;

root> SELECT * FROM accounts;
```

- When an user is in a transaction, changes done by the user are saved in a temp table. These changes are visible to that user.
- However this temp table is not accessible/visible to other users and hence changes under progress in a transaction are not visible to other users.
- When an user is in a transaction, changes committed by other users are not visible to him. Because he is dealing with temp data.

Row locking

```
sunbeam> SELECT * FROM accounts;

sunbeam> START TRANSACTION;
```

```
sunbeam> DELETE FROM accounts WHERE id = 1;
-- row locked

sunbeam> SELECT * FROM accounts;

root> SELECT * FROM accounts;

root> UPDATE accounts SET balance = 10000 WHERE id = 1;
-- blocked

sunbeam> COMMIT;
-- root user unblocked

root> SELECT * FROM accounts;
```

```
sunbeam> SELECT * FROM accounts;

sunbeam> START TRANSACTION;

sunbeam> DELETE FROM accounts WHERE id = 2;
-- row locked

root> UPDATE accounts SET balance = 40000 WHERE id = 2;
-- blocked

sunbeam> ROLLBACK;
-- root user unblocked

root> SELECT * FROM accounts;

sunbeam> SELECT * FROM accounts;
```

```
sunbeam> SELECT * FROM accounts;

sunbeam> START TRANSACTION;

sunbeam> DELETE FROM accounts WHERE id = 2;
-- row locked

root> UPDATE accounts SET balance = 30000 WHERE id = 2;
-- root user is blocked
-- auto unblocked after some time if other user has not done COMMIT/ROLLBACK.

sunbeam> ROLLBACK;

sunbeam> SELECT * FROM accounts;
```

Pessimistic Locking

```
sunbeam> SELECT * FROM accounts;

sunbeam> START TRANSACTION;

sunbeam> SELECT * FROM accounts WHERE id = 2 FOR UPDATE;
-- row locked

root> SELECT * FROM accounts;

root> SELECT * FROM accounts WHERE id = 2 FOR UPDATE;
-- blocked

sunbeam> DELETE FROM accounts WHERE id = 2;

sunbeam> ROLLBACK;
-- root user unblocked
```

Table locking

```
sunbeam> SELECT * FROM depts;

sunbeam> START TRANSACTION;

sunbeam> DELETE FROM depts WHERE deptno = 40;
-- whole table is locked (bcoz no primary key)

root> DELETE FROM depts WHERE deptno = 30;
-- blocked

sunbeam> COMMIT;
-- root user unblocked

sunbeam> SELECT * FROM depts;

root> SELECT * FROM depts;
```

```
DESCRIBE accounts;

DESCRIBE depts;
```

Indexes

- Faster searching

Simple Index

```
SELECT * FROM books;

EXPLAIN FORMAT=JSON
SELECT * FROM books WHERE subject = 'C Programming';
-- 1.55

CREATE INDEX idx_books_subject ON books(subject);

EXPLAIN FORMAT=JSON
SELECT * FROM books WHERE subject = 'C Programming';
-- 0.90

DESCRIBE books;

SHOW INDEXES FROM books;

CREATE INDEX idx_books_author ON books(author DESC);

DESCRIBE books;

SHOW INDEXES FROM books;
```

```
EXPLAIN FORMAT=JSON
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno;
-- 1.70

CREATE INDEX idx1 ON emps(deptno);
CREATE INDEX idx2 ON depts(deptno);

EXPLAIN FORMAT=JSON
SELECT e.ename, d.dname FROM emps e
INNER JOIN depts d ON e.deptno = d.deptno;
-- 1.62
```

Unique Index

- Duplicate values are not allowed.

```
CREATE UNIQUE INDEX idx3 ON emps(ename);

DESCRIBE emps;

SELECT * FROM emps;

SELECT * FROM emps WHERE ename = 'Nitin';
```

```
INSERT INTO emps VALUES (6, 'Rahul', 70, 1);
-- error: Duplicate entry

INSERT INTO emps VALUES (7, NULL, 50, 5);
-- (multiple) NULL value is allowed, but duplicate is not allowed.

CREATE UNIQUE INDEX idx4 ON emps(mgr);
-- error
```

Composite Index

```
SELECT * FROM emp;

DESCRIBE emp;

EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE deptno = 20 AND job = 'ANALYST';
-- 1.65

CREATE INDEX idx_dj ON emp(deptno ASC, job ASC);

DESCRIBE emp;

SHOW INDEXES FROM emp;

EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE deptno = 20 AND job = 'ANALYST';
-- 0.70

EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE sal = 5000;

EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE deptno = 20;
-- 1.00

EXPLAIN FORMAT=JSON
SELECT * FROM emp WHERE job = 'ANALYST';
-- 1.65
```

```
CREATE TABLE students(std INT, roll INT, name CHAR(30), marks DECIMAL(5,2));

INSERT INTO students VALUES (1, 1, 'Soham', 99);
INSERT INTO students VALUES (1, 2, 'Sakshi', 96);
INSERT INTO students VALUES (1, 3, 'Prisha', 98);
INSERT INTO students VALUES (2, 1, 'Madhu', 97);
INSERT INTO students VALUES (2, 2, 'Om', 95);
```

```
CREATE UNIQUE INDEX idx ON students(std,roll);

INSERT INTO students VALUES (1, 2, 'Ram', 99);
-- error: duplicate combination of std+roll not allowed

SELECT * FROM students;

INSERT INTO students VALUES (3, 1, 'Ram', 99);

DESCRIBE students;
```

Clustered Index

- Clustered index is auto-created on Primary key.
- It is internally a unique index that is used to lookup rows quickly on server disk.
- If Primary key is not present in the table, then a hidden (synthetic) column is created by RDBMS and Clustered index is created on it.

Drop Index

```
SHOW INDEXES FROM books;

DROP INDEX idx_books_author ON books;

DESCRIBE books;
```

Constraints

- Five Constraints
 - NOT NULL
 - Unique
 - Primary key
 - Foreign key
 - Check
- Types of constraints (based on syntax)
 - Column level

```
CREATE TABLE customers(
    id INT PRIMARY KEY,
    name CHAR(30) NOT NULL,
    email CHAR(40) UNIQUE NOT NULL,
    mobile CHAR(12) UNIQUE,
    addr VARCHAR(100)
);
```

- NOT NULL, UNIQUE, PRIMARY, Foreign, CHECK

- Table level

```
CREATE TABLE customers(
  id INT,
  name CHAR(30) NOT NULL,
  email CHAR(40) NOT NULL,
  mobile CHAR(12),
  addr VARCHAR(100),
  PRIMARY KEY(id),
  UNIQUE(email),
  UNIQUE(mobile)
);
```

- UNIQUE, PRIMARY, Foreign, CHECK

NOT NULL

- NULL value is not allowed in the column.
- Can be given at column level only.

```
CREATE TABLE temp1(c1 INT, c2 INT, c3 INT NOT NULL);

DESCRIBE temp1;

INSERT INTO temp1 VALUES (1, 1, 1);
INSERT INTO temp1 VALUES (NULL, 2, 2);
INSERT INTO temp1(c1,c3) VALUES (3,3);
SELECT * FROM temp1;

INSERT INTO temp1 VALUES (4, 4, NULL);
-- error: c3 cannot be NULL
INSERT INTO temp1(c1,c2) VALUES (5,5);
-- error: c3 cannot be NULL

SHOW INDEXES FROM temp1;
```

Unique

- Cannot have duplicate value in the column.
- However NULL value(s) allowed.
- Unique constraint internally creates unique index.
- Unique constraint on combination of multiple columns internally creates Composite Unique index. Must be at table level -- UNIQUE(c1,c2).

```
CREATE TABLE temp2(c1 INT, c2 INT, UNIQUE(c1));
-- CREATE TABLE temp2(c1 INT UNIQUE, c2 INT);

INSERT INTO temp2 VALUES (1, 1);
INSERT INTO temp2 VALUES (2, 2);
INSERT INTO temp2 VALUES (3, 3);
INSERT INTO temp2 VALUES (4, 2);

INSERT INTO temp2 VALUES (3, 5);
-- error: c1 cannot be duplicated

INSERT INTO temp2 VALUES (NULL, 5);
INSERT INTO temp2 VALUES (NULL, 6);

SELECT * FROM temp2;

SHOW INDEXES FROM temp2;
```

```
DROP TABLE IF EXISTS students;

CREATE TABLE students(std INT, roll INT, name CHAR(30), marks DECIMAL(5,2),
UNIQUE(std,roll));

INSERT INTO students VALUES (1, 1, 'Soham', 99);
INSERT INTO students VALUES (1, 2, 'Sakshi', 96);
INSERT INTO students VALUES (1, 3, 'Prisha', 98);
INSERT INTO students VALUES (2, 1, 'Madhu', 97);
INSERT INTO students VALUES (2, 2, 'Om', 95);

INSERT INTO students VALUES (1, 2, 'Ram', 99);
-- error: duplicate combination of std+roll not allowed
```

Primary key

- Primary key --> Column(s)
- "Identity" of each row/record.
- Primary key is "like" Unique constraint (cannot be duplicated) + NOT NULL constraint (cannot be NULL).
- In a table there is single primary key, but a table can have multiple unique key (constraints).

```
CREATE TABLE cdac_students(
    prn CHAR(16) PRIMARY KEY,
    name CHAR(40) NOT NULL,
    email CHAR(30) UNIQUE NOT NULL,
    mobile CHAR(12) UNIQUE NOT NULL,
```



```
addr VARCHAR(100)
);
```

- The primary key can be combination multiple columns. It is called as Composite Primary Key.

```
DROP TABLE IF EXISTS students;

CREATE TABLE students(std INT, roll INT, name CHAR(30), marks DECIMAL(5,2),
PRIMARY KEY(std,roll));

INSERT INTO students VALUES (1, 1, 'Soham', 99);
INSERT INTO students VALUES (1, 2, 'Sakshi', 96);
INSERT INTO students VALUES (1, 3, 'Prisha', 98);
INSERT INTO students VALUES (2, 1, 'Madhu', 97);
INSERT INTO students VALUES (2, 2, 'Om', 95);

INSERT INTO students VALUES (1, 2, 'Ram', 99);
-- error: duplicate combination of std+roll not allowed

DESCRIBE students;
```

- The Primary key internally creates UNIQUE index by name "PRIMARY".

```
SHOW INDEXES FROM cdac_students;

SHOW INDEXES FROM students;
```

```
CREATE TABLE t (c1 INT DEFAULT 0, ...)
```

MySQL - RDBMS

Agenda

- Constraints
 - Surrogate Primary Key
 - Foreign Key
 - Check
- ALTER table
- PSM / PL-SQL
 - Stored procedure
 - Functions
 - Triggers

Constraints

- Restrict values in the column.
- Constraints are checked/verified while performing DML operations. It slows down DML operations.
- Constraints ensure valid data is entered in the database.
- Unique key, Primary key and Foreign key constraints internally create indexes. It will help searching faster on these keys/columns.

Primary Key

- Primary key --> Identity of row/record/tuple.
- Natural primary key

```
CREATE TABLE customers(  
    email CHAR(40) PRIMARY KEY,  
    password CHAR(40),  
    name CHAR(40),  
    addr CHAR(100),  
    birth DATE  
);
```

- Composite primary key

```
CREATE TABLE students(  
    email CHAR(40),  
    password CHAR(40),  
    name CHAR(40),  
    grade CHAR(2),  
    course_code CHAR(20),  
    PRIMARY KEY(course_code, email)  
);
```

- Surrogate primary key
 - Usually auto-generated.
 - Oracle/Pg-SQL --> Sequences
 - MS-SQL --> Identity
 - MySQL --> AUTO_INCREMENT

```
CREATE TABLE students(  
    regno INT AUTO_INCREMENT,  
    email CHAR(40),  
    password CHAR(40),  
    name CHAR(40),  
    grade CHAR(2),  
    course_code CHAR(20),  
    PRIMARY KEY(regno)  
);
```

```
CREATE TABLE items(  
    id INT PRIMARY KEY AUTO_INCREMENT,  
    name CHAR(30),  
    price DECIMAL(5,2)  
);  
  
INSERT INTO items(name,price) VALUES('A', 10);  
INSERT INTO items(name,price) VALUES('B', 15);  
INSERT INTO items(name,price) VALUES('C', 20);  
SELECT * FROM items;  
  
ALTER TABLE items AUTO_INCREMENT = 100;  
  
INSERT INTO items(name,price) VALUES('X', 50);  
INSERT INTO items(name,price) VALUES('Y', 55);  
SELECT * FROM items;  
  
INSERT INTO items(id,name,price) VALUES (1000, 'P', 30);  
SELECT * FROM items;  
INSERT INTO items(name,price) VALUES('Q', 60);  
SELECT * FROM items;
```

Foreign Key

```
DESCRIBE emps;  
  
DESCRIBE depts;  
  
SELECT * FROM depts;
```

```

SELECT * FROM emps;

DROP TABLE emps;

DROP TABLE depts;

CREATE TABLE depts (deptno INT, dname VARCHAR(20), PRIMARY KEY(deptno));

INSERT INTO depts VALUES (10, 'DEV');
INSERT INTO depts VALUES (20, 'QA');
INSERT INTO depts VALUES (30, 'OPS');
INSERT INTO depts VALUES (40, 'ACC');

DESCRIBE depts;

CREATE TABLE emps (empno INT, ename VARCHAR(20), deptno INT, mgr INT, FOREIGN KEY
(deptno) REFERENCES depts(deptno));

INSERT INTO emps VALUES (1, 'Amit', 10, 4);
INSERT INTO emps VALUES (2, 'Rahul', 10, 3);
INSERT INTO emps VALUES (3, 'Nilesh', 20, 4);

INSERT INTO emps VALUES (4, 'Nitin', 50, 5);
-- error: a foreign key constraint fails
INSERT INTO emps VALUES (5, 'Sarang', 50, NULL);
-- error: a foreign key constraint fails

INSERT INTO emps VALUES (4, 'Nitin', 30, 5);
INSERT INTO emps VALUES (5, 'Sarang', 30, NULL);

SELECT * FROM depts;
SELECT * FROM emps;

INSERT INTO emps VALUES (6, 'Vishal', NULL, 3);
-- FK can be NULL
SELECT * FROM emps;

SELECT * FROM depts;
DELETE FROM depts WHERE deptno=40;

DELETE FROM depts WHERE deptno=30;
-- error: a foreign key constraint fails

DROP TABLE depts;
-- error: Cannot drop table 'depts' referenced by a foreign key constraint

```

- depts "1" ---- "*" emps
 - Parent-child relationship
 - Parent = depts table
 - Child = emps table
- Foreign key

- Cannot add/update in child row, if corresponding row is absent in parent table.
- Cannot delete parent row, if corresponding rows are present in child table.

```

DROP TABLE emps;
DROP TABLE depts;

CREATE TABLE depts (deptno INT, dname VARCHAR(20), PRIMARY KEY(deptno));

INSERT INTO depts VALUES (10, 'DEV');
INSERT INTO depts VALUES (20, 'QA');
INSERT INTO depts VALUES (30, 'OPS');
INSERT INTO depts VALUES (40, 'ACC');

DESCRIBE depts;

CREATE TABLE emps (empno INT, ename VARCHAR(20), deptno INT, mgr INT, FOREIGN KEY
(deptno) REFERENCES depts(deptno) ON DELETE CASCADE ON UPDATE CASCADE);

INSERT INTO emps VALUES (1, 'Amit', 10, 4);
INSERT INTO emps VALUES (2, 'Rahul', 10, 3);
INSERT INTO emps VALUES (3, 'Nilesh', 20, 4);

SELECT * FROM depts;
SELECT * FROM emps;

DELETE FROM depts WHERE deptno = 20;
-- ON DELETE CASCADE: If parent row is deleted, corresponding child rows will be
deleted automatically.

SELECT * FROM depts;
SELECT * FROM emps;

UPDATE depts SET deptno=100 WHERE dname='DEV';
-- ON UPDATE CASCADE: If parent row (primary key) is updated, corresponding child
rows (foreign key) will be updated automatically.

SELECT * FROM depts;
SELECT * FROM emps;

DROP TABLE depts;
-- error: Cannot drop table 'depts' referenced by a foreign key constraint

```

- Foreign is mapped to the primary key of other table.
- If PK is Composite primary key, the Foreign key can be Composite key.

```

CREATE TABLE students(
    email CHAR(40),
    password CHAR(40),
    name CHAR(40),
    grade CHAR(2),

```

```

        course_code CHAR(20),
        PRIMARY KEY(course_code, email)
    );

CREATE TABLE marks(
    id INT,
    subject CHAR(20),
    marks INT,
    course_id CHAR(20),
    email CHAR(40),
    FOREIGN KEY (course_id,email) REFERENCES students(course_code, email);
);

```

- Foreign key internally creates index on the table. It also helps in faster searching.

```

DESCRIBE emps;

SHOW INDEXES FROM emps;

```

- Foreign key constraint can be disabled temporarily in some cases (like backup/restore).

```

SELECT @@foreign_key_checks;

CREATE TABLE dept_backup(deptno INT, dname CHAR(40), loc CHAR(40), PRIMARY
KEY(deptno));

CREATE TABLE emp_backup(empno INT, ename CHAR(40), sal DECIMAL(8,2), deptno
INT,
PRIMARY KEY(empno), FOREIGN KEY (deptno) REFERENCES dept_backup(deptno));

INSERT INTO dept_backup SELECT * FROM dept;
SELECT * FROM dept_backup;

SET @@foreign_key_checks=0;

INSERT INTO emp_backup(empno,ename,sal,deptno) SELECT empno,ename,sal,deptno
FROM emp;
-- insert is fast, bcoz FK is disabled.

INSERT INTO emp_backup VALUES(1000, 'JOHN', 2000, 60);
-- allowed, bcoz FK checks are disabled -- but wrong

SET @@foreign_key_checks=1;
-- FK check is enabled -- further DML ops.

INSERT INTO emp_backup VALUES(1001, 'JACK', 2200, 60);
-- error: FK checks are enabled

SELECT * FROM emp_backup;

```

- Foreign key can be for the same table. It is called as "self-referencing" FK.

```
CREATE TABLE emps(  
  empno INT,  
  ename CHAR(40),  
  mgr INT,  
  PRIMARY KEY(empno),  
  FOREIGN KEY(mgr) REFERENCES emps(empno)  
);
```

Check

- Arbitrary conditions (application specific) to be applied on the column.
- Do not work in MySQL version <= 8.0.15

```
CREATE TABLE employees(  
  id INT PRIMARY KEY,  
  ename CHAR(40) CHECK (LENGTH(ename) > 1),  
  age INT NOT NULL CHECK (age > 18),  
  sal DECIMAL(7,2) CHECK (sal > 1000),  
  comm DECIMAL(7,2),  
  CHECK((sal + IFNULL(comm,0)) > 1200)  
);
```

```
INSERT INTO employees VALUES (1, 'A', 20, 2000, NULL);  
-- error: LENGTH(ename) > 1  
INSERT INTO employees VALUES (1, 'Om', 16, 2000, NULL);  
-- error: age > 18  
INSERT INTO employees VALUES (1, 'Om', 20, 900, NULL);  
-- error: sal > 1000  
INSERT INTO employees VALUES (1, 'Om', 20, 1100, NULL);  
-- error: (sal + IFNULL(comm,0)) > 1200  
INSERT INTO employees VALUES (1, 'Om', 20, 1100, 200);  
-- okay
```

Constraint names

```
CREATE TABLE employees(  
  id INT,  
  ename CHAR(40),  
  age INT NOT NULL,  
  sal DECIMAL(7,2),  
  comm DECIMAL(7,2),
```

```

deptno INT,
PRIMARY KEY(id),
FOREIGN KEY(deptno) REFERENCES departments(deptno),
UNIQUE(ename),
CHECK((sal + IFNULL(comm,0)) > 1200)
);
-- names of constraints are given auto by db

```

```

CREATE TABLE employees(
  id INT,
  ename CHAR(40),
  age INT NOT NULL,
  sal DECIMAL(7,2),
  comm DECIMAL(7,2),
  deptno INT,
  CONSTRAINT pk_employees PRIMARY KEY(id),
  CONSTRAINT fk_dept FOREIGN KEY(deptno) REFERENCES departments(deptno),
  CONSTRAINT uk_ename UNIQUE(ename),
  CONSTRAINT chk_income CHECK((sal + IFNULL(comm,0)) > 1200)
);

```

Show Constraints

```

SHOW CREATE TABLE emps;

SELECT TABLE_NAME,
       COLUMN_NAME,
       CONSTRAINT_NAME,
       REFERENCED_TABLE_NAME,
       REFERENCED_COLUMN_NAME
FROM INFORMATION_SCHEMA.KEY_COLUMN_USAGE
WHERE TABLE_SCHEMA = 'classwork'
  AND TABLE_NAME = 'emps'
  AND REFERENCED_COLUMN_NAME IS NOT NULL;

```

ALTER Table

- CREATE TABLE -- Table Structure (Metadata)
- DML operations -- Table Data
- ALTER TABLE -- Change table structure/metadata
 - Add column, Remove column, Change column data type/name, Add/Remove constraint, ...
 - Not recommended in production database.
 - After alteration table storage become inefficient.

```
DESCRIBE emp_backup;
```



```
ALTER TABLE emp_backup ADD COLUMN job CHAR(20);

SELECT * FROM emp_backup;

UPDATE emp_backup e SET e.job = (SELECT job FROM emp WHERE empno = e.empno);

SELECT * FROM emp_backup;

DESCRIBE emp_backup;

ALTER TABLE emp_backup MODIFY job VARCHAR(40);
-- can change data type to compatible data type

DESCRIBE emp_backup;

ALTER TABLE emp_backup MODIFY job INT;
-- error: cannot change data type to incompatible.

ALTER TABLE emp_backup CHANGE ename name CHAR(30);

DESCRIBE emp_backup;

ALTER TABLE emp_backup DROP COLUMN sal;

DESCRIBE emp_backup;
```

```
ALTER TABLE emp ADD PRIMARY KEY (empno);

ALTER TABLE emp ADD UNIQUE(ename);

SHOW CREATE TABLE emp;

ALTER TABLE dept ADD PRIMARY KEY (deptno);

ALTER TABLE emp ADD FOREIGN KEY (deptno) REFERENCES dept (deptno);
```

```
ALTER TABLE emp DROP PRIMARY KEY;

SHOW CREATE TABLE emp;

ALTER TABLE emp DROP CONSTRAINT ename;

ALTER TABLE emp DROP CONSTRAINT emp_ibfk_1;
```

PSM / PL-SQL

Stored procedure

- Default DELIMITER is semicolon.
- When ; is found, client submit the code/query to the server.
- It should be changed temporarily to implement stored procedure using DELIMITER keyword.

Steps of Stored Procedure programming.

- step 1: Create a .sql file (like psm01.sql).
- step 2: Use SOURCE command on mysql CLI to execute it.

```
SOURCE D:/sep21/DAC/dbt/day09/psm01.sql
```

- step 3: Call the procedure.

```
CALL sp_hello1();
```

Stored Procedure Result into Table

```
CREATE TABLE result(id INT, val CHAR(100));
```

Stored Procedure Params

```
// arg n --> input to function --> in param
int sqr(int n) {
    return n * n;
}
void main() {
    // ...
    res = sqr(5);
    // ...
}
```

```
// arg n --> input to function --> in param
// arg r --> output from function --> out param
void sqr(int n, int *r) {
    *r = n * n;
}
void main() {
    // ...
    sqr(5, &res);
    // ...
}
```

```
// arg n --> input to fn & output from fn --> in-out param
void sqr(int *n) {
    *n = (*n) * (*n);
}
void main() {
    // ...
    res = 5;
    res = sqr(&res);
    // ...
}
```

```
CREATE PROCEDURE sp_sqr1(IN p_n INT, OUT p_r INT)
BEGIN
    SET p_r = p_n * p_n;
END
```

```
CALL sp_sqr1(5, @res1)
SELECT @res1;
```

```
CREATE PROCEDURE sp_sqr2(INOUT p_n INT)
BEGIN
    SET p_n = p_n * p_n;
END
```

```
SET @res2 = 5;
CALL sp_sqr2(@res2);
SELECT @res2;
```