

Assignment - 2

Q1]

$$w = \{2, 3, 4, 5\}$$

$$m = 5$$

$$p = \{3, 4, 5, 6\}$$

$$S^0 = \{(0, 0)\}$$

$$S_1^0 = \{(0, 0), (3, 2)\}$$

$$S_2^1 = \{(0, 0), (3, 2)\}$$

$$S_2^1 = \{(4, 3), (7, 5)\}$$

$$S_3^2 = \{(0, 0), (3, 2), (4, 3), (7, 5)\}$$

$$S_3^2 = \{(5, 4), (8, 6), (9, 7), (12, 9)\}$$

$$S^3 = \{(0, 0), (3, 2), (4, 3), (7, 5), (5, 4), (8, 8)\}$$

$$S_4^3 = \{(6, 5), (9, 7), (10, 8), (13, 10), (11, 9)\}$$

$$S^4 = \{(0, 0), (3, 2), (4, 3), (6, 5), \boxed{(7, 5)}, (5, 4), \text{max profit}, (6, 5)\}$$

$$(7, 5) - (5, 4) = 2, 1$$

$$(7, 5) - (4, 3) = (3, 2) \quad 0, 2$$

$$(3, 2) - (3, 2) = 0, 0 \quad 0, 1$$

$$\boxed{[1, 1, 0, 0]}$$

Q2]

Dynamic :

- (i) Guarantee optimal solution
- (ii) Subproblems overlap
- (iii) Does more work compared to both
- (iv) No specialized set of feasible solutions
- (v) Employ memorization

Divide and conquer :

- (i) Does not aim for optimal solution
- (ii) Subproblems are independent
- (iii) Slower and inefficient
- (iv) No memorization

Greedy :

- (i) Does not guarantee optimal solution
- (ii) Subproblems do not overlap
- (iii) Does little work
- (iv) Select choice which is locally optimum
- (v) No memorization

Q3] Algorithm.

Sum of subsets (S, k, n)

{ // Find all subsets of $w[1:n]$ that sum to m

$x[k] = 1$

if ($S + w[k] == m$)

write $x[1:k]$

else if ($S + w[k] + w[k+1] \leq m$)

Sum of subsets ($S + w[k], k+1, n - w[k]$)

if ($S + n - w[k] > m$) and ($S + w[k+1] \leq m$)

$x[k] = 0$

Sum of subsets ($S, k+1, n - w[k]$)

eg: $n = 4, w = \{4, 5, 8, 9\}$ & $m = 9$, set $S = 0$

Step 1: $i = 1$, Adding item w_i

~~sum = sum + w_i~~ $S = S + w_i = 0 + 4 = 4$

$S \leq m$, \therefore add to soln set

$x[i] = x[1] = 1 \Rightarrow x = [1, 0, 0, 0]$

Step 2: $i = 2$, Adding item w_i

$S = S + w_i = 4 + 5 = 9$

$S == m$

so solution found & add to solⁿ set

$x[2] = 1 \Rightarrow x = [1, 1, 0, 0]$

P.T.O

Q1] Max $Z = -3x_1 + 2x_2 + 0s_1 + 0s_2 + 0s_3 = 0$
 subject to $-x_1 + 2x_2 + s_1 = 4$
 $3x_1 + 2x_2 + s_2 = 14$
 $x_1 - x_2 + s_3 = 3$

Simplex Iteration No.	Table:	coeff of					RHS Soln	Ratio
	Basic Variable	x_1	x_2	s_1	s_2	s_3		
0	Z	-3	-2	0	0	0	0	
	s_1	-1	2	1	0	0	4	—
	s_2	3	2	0	1	0	14	$\frac{14}{3} = 4.67$
	s_3	1	-1	0	0	1	3	$\frac{3}{1} = 3$
1	Z	0	-5	0	0	3	9	
s_3 exits	s_1	0	1	1	0	1	7	7
x_1 enters	s_2	0	5	0	1	-3	5	1
	x_1	1	-1	0	0	1	3	—
1	Z	0	0	0	1	0	14	
s_2 leaves	s_1	0	0	1	-1/5	8/5	6	
x_2 enters	x_2	0	1	0	1/5	-3/5	1	
	x_1	1	0	0	1/5	2/5	4	

$x_1 = 4, x_2 = 1, \therefore Z_{\max} = 14$

Q5] (a) Rabin Karp Matcher (T, p, d, q)

$n = T.length$

$m = p.length$

$h = d^{m-1} \bmod q$

$p = 0$

$t_0 = 0$

for $i = 1$ to m

$p = (dp + p[i]) \bmod q$

$t_0 = (dt_0 + T[i]) \bmod q$

for $s = 0$ to $n - m$

if $p == t_s$

if $p[1 \dots m] == T[s+1 \dots s+m]$

print ("Pattern occurs with shift" s)

if $s < n - m$

$t_{s+1} = (d(t_s - T[s+1])h + T[s+m+1]) \bmod q$

(b) Compute Prefix Function (p)

$m = p.length$

let $\pi[1 \dots m]$ be a new array

$\pi[1] = 0$

$k = 0$

for $q = 2$ to m

while $k > 0 \wedge p[k+1] \neq p[q]$

$k = \pi[k]$

if $p[k+1] == p[q]$

$k = k + 1$

$\pi[q] = k$

return π

KMP-matcher (T, p)

$n = T.length$

$m = p.length$

$\pi = \text{Compute Prefix Function}(p)$

$q = 0$

for $i = 1$ to n

while $q > 0 \wedge p[q+1] \neq p[i]$

$q = \pi[q]$

if $p[q+1] == p[i]$

$q = q + 1$

if $q == m$

Print ("Pattern occurs with shift" i)

$q = \pi[q]$