

# 1 Formal Development

**Lemma 1** (Value inversion). *Inversion lemma for values:*

1. If  $v : \text{uint}$ , then  $v = n$
2. If  $v : \text{bool}$ , then  $v = \top$  or  $v = \perp$
3. If  $v : \sigma[n]$ , then  $v = [c_i]_n$  and  $c_i : \sigma$

**Lemma 2** (Consistency of source type and target type). *If  $\psi \sim \tau$ , then one of the following holds:*

1.  $\psi = \sigma$  and  $\tau = \sigma^\ell$ .
2.  $\psi = \sigma[n]$  and  $\tau = \sigma^\ell[n]$ .

**Lemma 3** (Compilation of source environment). *If  $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$  and  $\Gamma(x) = \tau$ , then one of the following holds:*

1.  $\tau = \sigma^{\mathcal{P}}$ ,  $\rho(x) = c$ ,  $c : \sigma$ , and  $\tilde{\rho}(x) = c$ .
2.  $\tau = \sigma^m$ ,  $\rho(x) = c$ ,  $c : \sigma$ ,  $\tilde{\rho}(x) = r$ , and  $(\hat{\rho}_1[r], \hat{\rho}_2[r]) = \mathcal{E}_m(c)$ .
3.  $\tau = \sigma^{\mathcal{P}}[n]$ ,  $\rho(x) = [c_i]_n$ ,  $\forall i \in \{0, n-1\}$ .  $c_i : \sigma$ , and  $\tilde{\rho}(x) = [c_i]_n$ .
4.  $\tau = \sigma^m[n]$ ,  $\rho(x) = [c_i]_n$ ,  $\forall i \in \{0, n-1\}$ .  $c_i : \sigma$ ,  $\tilde{\rho}(x) = [r_i]_n$ ,  $\forall i \in \{0, n-1\}$ .  $(\hat{\rho}_1[r_i], \hat{\rho}_2[r_i]) = \mathcal{E}_m(c_i)$

**Lemma 4** (Soundness of public, scalar expressions). *If*

1.  $\Gamma \vdash e : \sigma^{\mathcal{P}} \rightsquigarrow \tilde{e}$
2.  $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$

*Then*

- (a)  $\rho \vdash e \Downarrow c$
- (b)  $c : \sigma$
- (c)  $\tilde{\rho} \vdash \tilde{e} \Downarrow c$

**Lemma 5** (Soundness of public, array expressions). *If*

1.  $\Gamma \vdash e : \sigma^{\mathcal{P}}[n] \rightsquigarrow \tilde{e}$
2.  $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$

*Then*

- (a)  $\rho \vdash e \Downarrow [c_i]_n$
- (b)  $\forall i \in \{0, n-1\}$ .  $c_i : \sigma$

$$(c) \tilde{\rho} \vdash \tilde{e} \Downarrow [c_i]_n$$

**Lemma 6** (Soundness of secret, scalar expressions). *If*

1.  $\Gamma \vdash e : \sigma^m \rightsquigarrow \tilde{e}$
2.  $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$

*Then*

- (a)  $\rho \vdash e \Downarrow c$
- (b)  $c : \sigma$
- (c)  $\tilde{\rho} \vdash \tilde{e} \Downarrow \kappa^e$
- (d)  $\hat{\rho}_1, \hat{\rho}_2 \vdash \kappa^e \Downarrow b_1, b_2$
- (e)  $c = \mathcal{D}_m(b_1, b_2)$

**Lemma 7** (Soundness of secret, array expressions). *If*

1.  $\Gamma \vdash e : \sigma^m[n] \rightsquigarrow \tilde{e}$
2.  $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$

*Then*

- (a)  $\rho \vdash e \Downarrow [c_i]_n$
- (b)  $\forall i \in \{0, n-1\}. c_i : \sigma$
- (c)  $\tilde{\rho} \vdash \tilde{e} \Downarrow [\kappa_i^e]_n$
- (d)  $\forall i \in \{0, n-1\}. \hat{\rho}_1, \hat{\rho}_2 \vdash \kappa_i^e \Downarrow b_{1\ i}, b_{2\ i}, \text{ s.t. } c_i = \mathcal{D}_m(b_{1\ i}, b_{2\ i})$

**Lemma 8** (Target semantics correspondence). *If:*

1.  $\Gamma \vdash s \rightsquigarrow \tilde{s} \mid \Gamma'$
2.  $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$
3.  $\rho \vdash s \Downarrow \rho'; O$

*then:*

- (a)  $\tilde{\rho} \vdash \tilde{s} \Downarrow \tilde{\rho}'; \kappa^s$
- (b)  $\hat{\rho}_1, \hat{\rho}_2 \vdash \kappa^s \Downarrow \tilde{\rho}'_1, \tilde{\rho}'_2; O$
- (c)  $\Gamma' \vdash \rho' \hookrightarrow \tilde{\rho}'; \hat{\rho}'_1, \hat{\rho}'_2$

**Lemma 9** (Soundness of source semantics). *If:*

1.  $\Gamma \vdash s \rightsquigarrow \tilde{s} \mid \Gamma'$

$$2. \Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$$

*then:*

$$(a) \rho \vdash s \Downarrow \rho'; O$$

$$(b) \Gamma' \vdash \rho' \hookrightarrow \tilde{\rho}'; \tilde{\rho}'_1, \tilde{\rho}'_2$$

**Theorem 10** (Soundness). *If:*

$$1. \Gamma \vdash s \rightsquigarrow \tilde{s} \mid \Gamma'$$

$$2. \Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$$

*then:*

$$(a) \rho \vdash s \Downarrow \rho'; O$$

$$(b) \tilde{\rho} \vdash \tilde{s} \Downarrow \tilde{\rho}'; \kappa^s$$

$$(c) \hat{\rho}_1, \hat{\rho}_2 \vdash \kappa^s \Downarrow \tilde{\rho}'_1, \tilde{\rho}'_2; O$$

$$(d) \Gamma' \vdash \rho' \hookrightarrow \tilde{\rho}'; \tilde{\rho}'_1, \tilde{\rho}'_2$$