

# 1 Formal Development

Correctness theorem that we will aim at:

If:

- $\Gamma \vdash s \rightsquigarrow \tilde{s} \mid \Gamma'$
- $\Gamma \sim \rho$
- $\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2$

Then:

- $\rho \vdash s \downarrow \rho'; O$
- $\tilde{\rho} \vdash \tilde{s} \Longrightarrow \tilde{\rho}'; \kappa$
- $\hat{\rho}_1, \hat{\rho}_2 \vdash \kappa \longmapsto \tilde{\rho}'_1, \tilde{\rho}'_2; O$

$m$	$::=$   $\mathcal{A}$   $\mathcal{B}$	Secret label
$\ell$	$::=$   $\mathcal{P}$   $m$	Label
$\sigma$	$::=$   $\text{uint}^\ell$   $\text{bool}^\ell$	Base type
$\tau$	$::=$   $\sigma$   $\sigma[]$	Type
$c$	$::=$   $n$   $\top$   $\perp$	Constant
$e$	$::=$   $c$   $x$   $e_1 + e_2$   $\mathbf{cond}(e, e_1, e_2)$   $e_1 > e_2$   $x[e]$	Source expression
$s$	$::=$   $\tau x$   $x := e$   $\mathbf{for}(x := n_1; x \leq n_2; x := x + 1) s$   $x[e_1] := e_2$   $\mathbf{if}(e, s_1, s_2)$   $\mathbf{out} e$   $s_1; s_2$	Source statement

Figure 1: Source language

$v$	$::=$	Source value
	$\begin{array}{ c} c \\ \hline [\overline{c_i}^i] \end{array}$	
$\rho$	$::=$	Source runtime environment
	$\begin{array}{ c} \cdot \\ \hline \rho[x \mapsto v] \end{array}$	
$O$	$::=$	Source observation
	$\begin{array}{ c} \cdot \\ \hline v \\ \hline O_1, O_2 \end{array}$	

Figure 2: Source runtime

$$\boxed{\rho \vdash e \downarrow v}$$

$$\begin{array}{c}
\frac{}{\rho \vdash c \downarrow c} \quad \text{SE\_CONST} \\
\\
\frac{}{\rho \vdash x \downarrow \rho(x)} \quad \text{SE\_VAR} \\
\\
\frac{\rho \vdash e_i \downarrow n_i}{\rho \vdash e_1 + e_2 \downarrow n_1 + n_2} \quad \text{SE\_ADD} \\
\\
\frac{\rho \vdash e \downarrow \top \quad \rho \vdash e_1 \downarrow v}{\rho \vdash \mathbf{cond}(e, e_1, e_2) \downarrow v} \quad \text{SE\_CONDT} \\
\\
\frac{\rho \vdash e \downarrow \perp \quad \rho \vdash e_2 \downarrow v}{\rho \vdash \mathbf{cond}(e, e_1, e_2) \downarrow v} \quad \text{SE\_CONDF} \\
\\
\frac{\rho \vdash e_i \downarrow n_i}{\rho \vdash e_1 > e_2 \downarrow n_1 > n_2} \quad \text{SE\_GT} \\
\\
\frac{\rho \vdash e \downarrow n \quad \rho \vdash x \downarrow [\overline{c_i}^i]}{\rho \vdash x[e] \downarrow c_n} \quad \text{SE\_AREAD}
\end{array}$$

Figure 3: Source expression evaluation

$$\boxed{\rho \vdash s \downarrow \rho'; O}$$

$$\begin{array}{c}
\frac{\text{default}(\tau) = v}{\rho \vdash \tau x \downarrow \rho[x \mapsto v]; \cdot} \quad \text{SC\_DECL} \\
\\
\frac{\rho \vdash e \downarrow v}{\rho \vdash x := e \downarrow \rho[x \mapsto v]; \cdot} \quad \text{SC\_ASSGN} \\
\\
\frac{n_1 > n_2}{\rho \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) s \downarrow \rho; \cdot} \quad \text{SC\_FORT} \\
\\
\frac{\begin{array}{l} n_2 \geq n_1 \\ \rho[x \mapsto n_1] \vdash s \downarrow \rho_1; O_1 \\ \rho_1 \vdash \mathbf{for}(x := n_1 + 1; x \leq n_2; x := x + 1) s \downarrow \rho_2; O_2 \end{array}}{\rho \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) s \downarrow \rho_2; O_1, O_2} \quad \text{SC\_FORI} \\
\\
\frac{\begin{array}{l} \rho \vdash x \downarrow [\overline{c_i}^i] \\ \rho \vdash e_1 \downarrow n \\ \rho \vdash e_2 \downarrow c \end{array}}{\rho \vdash x[e_1] := e_2 \downarrow \rho[x \mapsto [\overline{c_i}^i][n \mapsto c]]; O_1, O_2} \quad \text{SC\_AWRITE} \\
\\
\frac{\begin{array}{l} \rho \vdash e \downarrow \top \\ \rho \vdash s_1 \downarrow \rho'; O \end{array}}{\rho \vdash \mathbf{if}(e, s_1, s_2) \downarrow \rho'; O} \quad \text{SC\_IFT} \\
\\
\frac{\begin{array}{l} \rho \vdash e \downarrow \perp \\ \rho \vdash s_2 \downarrow \rho'; O \end{array}}{\rho \vdash \mathbf{if}(e, s_1, s_2) \downarrow \rho'; O} \quad \text{SC\_IFF} \\
\\
\frac{\rho \vdash e \downarrow v}{\rho \vdash \mathbf{out} e \downarrow \rho; v} \quad \text{SC\_OUT} \\
\\
\frac{\begin{array}{l} \rho \vdash s_1 \downarrow \rho_1; O_1 \\ \rho_1 \vdash s_2 \downarrow \rho_2; O_2 \end{array}}{\rho \vdash s_1; s_2 \downarrow \rho_2; O_1, O_2} \quad \text{SC\_SEQ}
\end{array}$$

Figure 4: Source command evaluation

$$\boxed{v : \tau}$$

$$\begin{array}{c} \frac{}{n : \text{uint}^{\mathcal{P}}} \quad \text{V\_INT} \\ \frac{}{\top : \text{bool}^{\mathcal{P}}} \quad \text{V\_TRUE} \\ \frac{}{\perp : \text{bool}^{\mathcal{P}}} \quad \text{V\_FALSE} \\ \frac{c_i : \sigma}{[\overline{c_i}^i] : \sigma[]} \quad \text{V\_ARR} \end{array}$$

Figure 5: Value typing

$\tilde{e}$	::=	$\begin{array}{ l} c \\ x \\ \tilde{e}_1 +_{\ell} \tilde{e}_2 \\ \mathbf{cond}_{\ell}(\tilde{e}, \tilde{e}_1, \tilde{e}_2) \\ \tilde{e}_1 >_{\ell} \tilde{e}_2 \\ x[\tilde{e}] \\ \tilde{e} \triangleright m \end{array}$	Target expression
$\tilde{s}$	::=	$\begin{array}{ l} \tau x \\ x := \tilde{e} \\ \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) \tilde{s} \\ x[\tilde{e}_1] := \tilde{e}_2 \\ \mathbf{if}(\tilde{e}, \tilde{s}_1, \tilde{s}_2) \\ \mathbf{out} \tilde{e} \\ \tilde{s}_1; \tilde{s}_2 \end{array}$	Target statement
$\Gamma$	::=	$\begin{array}{ l} . \\ \Gamma, x : \tau \end{array}$	Type environment

Figure 6: Target language

$$\boxed{\Gamma \vdash e : \tau \rightsquigarrow \tilde{e}}$$

$$\begin{array}{c}
\frac{}{\Gamma \vdash n : \text{uint}^{\mathcal{P}} \rightsquigarrow n} \quad \text{S\_CONST} \\
\\
\frac{}{\Gamma \vdash \top : \text{bool}^{\mathcal{P}} \rightsquigarrow \top} \quad \text{S\_TRUE} \\
\\
\frac{}{\Gamma \vdash \perp : \text{bool}^{\mathcal{P}} \rightsquigarrow \perp} \quad \text{S\_FALSE} \\
\\
\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau \rightsquigarrow x} \quad \text{S\_VAR} \\
\\
\frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \rightsquigarrow \tilde{e}_i}{\Gamma \vdash e_1 + e_2 : \text{uint}^{\mathcal{P}} \rightsquigarrow \tilde{e}_1 +_{\mathcal{P}} \tilde{e}_2} \quad \text{S\_PADD} \\
\\
\frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{A}} \rightsquigarrow \tilde{e}_i}{\Gamma \vdash e_1 + e_2 : \text{uint}^{\mathcal{A}} \rightsquigarrow \tilde{e}_1 +_{\mathcal{A}} \tilde{e}_2} \quad \text{S\_SADD} \\
\\
\frac{\Gamma \vdash e : \text{bool}^{\mathcal{P}} \rightsquigarrow \tilde{e} \quad \Gamma \vdash e_i : \sigma \rightsquigarrow \tilde{e}_i}{\Gamma \vdash \mathbf{cond}(e, e_1, e_2) : \sigma \rightsquigarrow \mathbf{cond}_{\mathcal{P}}(\tilde{e}, \tilde{e}_1, \tilde{e}_2)} \quad \text{S\_PCOND} \\
\\
\frac{\Gamma \vdash e : \text{bool}^{\mathcal{B}} \rightsquigarrow \tilde{e} \quad \Gamma \vdash e_i : \sigma \rightsquigarrow \tilde{e}_i \quad \text{label}(\sigma) = \mathcal{B}}{\Gamma \vdash \mathbf{cond}(e, e_1, e_2) : \sigma \rightsquigarrow \mathbf{cond}_{\mathcal{B}}(\tilde{e}, \tilde{e}_1, \tilde{e}_2)} \quad \text{S\_SCOND} \\
\\
\frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \rightsquigarrow \tilde{e}_i}{\Gamma \vdash e_1 > e_2 : \text{bool}^{\mathcal{P}} \rightsquigarrow \tilde{e}_1 >_{\mathcal{P}} \tilde{e}_2} \quad \text{S\_PGT} \\
\\
\frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{B}} \rightsquigarrow \tilde{e}_i}{\Gamma \vdash e_1 > e_2 : \text{bool}^{\mathcal{B}} \rightsquigarrow \tilde{e}_1 >_{\mathcal{B}} \tilde{e}_2} \quad \text{S\_SGT} \\
\\
\frac{\Gamma \vdash x : \sigma[\ ] \rightsquigarrow x \quad \Gamma \vdash e : \text{uint}^{\mathcal{P}} \rightsquigarrow \tilde{e}}{\Gamma \vdash x[e] : \sigma \rightsquigarrow x[\tilde{e}]} \quad \text{S\_AREAD} \\
\\
\frac{\Gamma \vdash e : \sigma_1 \rightsquigarrow \tilde{e} \quad \text{base}(\sigma_1) = \text{base}(\sigma_2) \quad \text{label}(\sigma_2) = m}{\Gamma \vdash e : \sigma_2 \rightsquigarrow \tilde{e} \triangleright m} \quad \text{S\_SUB}
\end{array}$$

Figure 7: Expression compilation

$$\boxed{\Gamma \vdash s \rightsquigarrow \tilde{s} \mid \Gamma'}$$

$$\begin{array}{c}
\overline{\Gamma \vdash \tau x \rightsquigarrow \tau x \mid \Gamma, x : \tau} \quad \text{C\_DECL} \\
\\
\frac{\Gamma(x) = \sigma \quad \Gamma \vdash e : \sigma \rightsquigarrow \tilde{e}}{\Gamma \vdash x := e \rightsquigarrow x := \tilde{e} \mid \Gamma} \quad \text{C\_VASSGN} \\
\\
\frac{\Gamma, x : \text{uint}^{\mathcal{P}} \vdash s \rightsquigarrow \tilde{s} \mid - \quad x \notin \text{modifies}(s)}{\Gamma \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) s \rightsquigarrow \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) \tilde{s} \mid \Gamma} \quad \text{C\_FOR} \\
\\
\frac{\Gamma \vdash x : \sigma[\ ] \rightsquigarrow x \quad \Gamma \vdash e_1 : \text{uint}^{\mathcal{P}} \rightsquigarrow \tilde{e}_1 \quad \Gamma \vdash e_2 : \sigma \rightsquigarrow \tilde{e}_2}{\Gamma \vdash x[e_1] := e_2 \rightsquigarrow x[\tilde{e}_1] := \tilde{e}_2 \mid \Gamma} \quad \text{C\_AWRITE} \\
\\
\frac{\Gamma \vdash e : \text{bool}^{\mathcal{P}} \rightsquigarrow \tilde{e} \quad \Gamma \vdash s_i \rightsquigarrow \tilde{s}_i \mid -}{\Gamma \vdash \mathbf{if}(e, s_1, s_2) \rightsquigarrow \mathbf{if}(\tilde{e}, \tilde{s}_1, \tilde{s}_2) \mid \Gamma} \quad \text{C\_IF} \\
\\
\frac{\Gamma \vdash e : \sigma \rightsquigarrow \tilde{e} \quad \text{label}(\sigma) = m}{\Gamma \vdash \mathbf{out} e \rightsquigarrow \mathbf{out} \tilde{e} \mid \Gamma} \quad \text{C\_OUT} \\
\\
\frac{\Gamma \vdash s_1 \rightsquigarrow \tilde{s}_1 \mid \Gamma_1 \quad \Gamma_1 \vdash s_2 \rightsquigarrow \tilde{s}_2 \mid \Gamma'}{\Gamma \vdash s_1; s_2 \rightsquigarrow \tilde{s}_1; \tilde{s}_2 \mid \Gamma'} \quad \text{C\_SEQ}
\end{array}$$

Figure 8: Command compilation

$r$	$::=$	Wire id range
$\tilde{w}$	$::=$	Compiled base value
	$\begin{array}{ c} c \\ r \end{array}$	
$\tilde{v}$	$::=$	Compiled value
	$\begin{array}{ c} \tilde{w} \\ [\widetilde{w_i}^i] \end{array}$	
$\kappa$	$::=$	Circuit
	$\begin{array}{ c} \cdot \\ \oplus(r_1, r_2, r_3) \\ \mathbf{Mux}(r_1, r_2, r_3, r_4) \\ \mathbf{Gt}(r_1, r_2, r_3) \\ r_1 \triangleright_m r_2 \\ \mathbf{Out}(r) \\ \kappa_1, \kappa_2 \end{array}$	
$\tilde{\rho}$	$::=$	Runtime environment
	$\begin{array}{ c} \cdot \\ \tilde{\rho}[x \mapsto \tilde{v}] \end{array}$	

Figure 9: Target runtime



$$\boxed{\tilde{\rho} \vdash \tilde{e} \Downarrow \tilde{v}; \kappa}$$

$$\begin{array}{c}
\frac{}{\tilde{\rho} \vdash c \Downarrow c; \cdot} \text{EE\_CONST} \\
\\
\frac{}{\tilde{\rho} \vdash x \Downarrow \tilde{\rho}[x]; \cdot} \text{EE\_VAR} \\
\\
\frac{\tilde{\rho} \vdash \tilde{e}_i \Downarrow n_i; \kappa_i}{\tilde{\rho} \vdash \tilde{e}_1 +_{\mathcal{P}} \tilde{e}_2 \Downarrow n_1 + n_2; \kappa_1, \kappa_2} \text{EE\_PADD} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e}_i \Downarrow r_i; \kappa_i \\ r_3 = \text{next\_range}() \end{array}}{\tilde{\rho} \vdash \tilde{e}_1 +_{\mathcal{A}} \tilde{e}_2 \Downarrow r_3; \kappa_1, \kappa_2, \oplus(r_1, r_2, r_3)} \text{EE\_SADD} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e} \Downarrow \top; \kappa \\ \tilde{\rho} \vdash \tilde{e}_1 \Downarrow \tilde{v}; \kappa_1 \end{array}}{\tilde{\rho} \vdash \mathbf{cond}_{\mathcal{P}}(\tilde{e}, \tilde{e}_1, \tilde{e}_2) \Downarrow \tilde{v}; \kappa, \kappa_1} \text{EE\_PCONDT} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e} \Downarrow \perp; \kappa \\ \tilde{\rho} \vdash \tilde{e}_2 \Downarrow \tilde{v}; \kappa_2 \end{array}}{\tilde{\rho} \vdash \mathbf{cond}_{\mathcal{P}}(\tilde{e}, \tilde{e}_1, \tilde{e}_2) \Downarrow \tilde{v}; \kappa, \kappa_2} \text{EE\_PCONDF} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e} \Downarrow r; \kappa \\ \tilde{\rho} \vdash \tilde{e}_i \Downarrow r_i; \kappa_i \\ r_3 = \text{next\_range}() \end{array}}{\tilde{\rho} \vdash \mathbf{cond}_{\mathcal{B}}(\tilde{e}, \tilde{e}_1, \tilde{e}_2) \Downarrow r_3; \kappa, \kappa_1, \kappa_2, \text{Mux}(r, r_1, r_2, r_3)} \text{EE\_SCOND} \\
\\
\frac{\tilde{\rho} \vdash \tilde{e}_i \Downarrow n_i; \kappa_i}{\tilde{\rho} \vdash \tilde{e}_1 >_{\mathcal{P}} \tilde{e}_2 \Downarrow n_1 > n_2; \kappa_1, \kappa_2} \text{EE\_PGT} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e}_i \Downarrow r_i; \kappa_i \\ r_3 = \text{next\_range}() \end{array}}{\tilde{\rho} \vdash \tilde{e}_1 >_{\mathcal{B}} \tilde{e}_2 \Downarrow r_3; \kappa_1, \kappa_2, \text{Gt}(r_1, r_2, r_3)} \text{EE\_SGT} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e} \Downarrow n; \kappa_1 \\ \tilde{\rho} \vdash x \Downarrow [\widetilde{w}_i^i]; \kappa_2 \end{array}}{\tilde{\rho} \vdash x[\tilde{e}] \Downarrow \widetilde{w}_n; \kappa_1, \kappa_2} \text{EE\_AREAD} \\
\\
\frac{\begin{array}{c} \tilde{\rho} \vdash \tilde{e} \Downarrow r; \kappa \\ r' = \text{next\_range}() \end{array}}{\tilde{\rho} \vdash \tilde{e} \triangleright m \Downarrow r'; \kappa, r \triangleright_m r'} \text{EE\_COERCE}
\end{array}$$

Figure 10: Target expression evaluation

$$\boxed{\tilde{\rho} \vdash \tilde{s} \Longrightarrow \tilde{\rho}'; \kappa}$$

$$\begin{array}{c}
\frac{\text{default}(\tau) = \tilde{v}; \kappa}{\tilde{\rho} \vdash \tau x \Longrightarrow \tilde{\rho}[x \mapsto \tilde{v}]; \kappa} \text{ EC\_DECL} \\
\\
\frac{\tilde{\rho} \vdash \tilde{e} \Downarrow \tilde{v}; \kappa}{\tilde{\rho} \vdash x := \tilde{e} \Longrightarrow \tilde{\rho}[x \mapsto \tilde{v}]; \kappa} \text{ EC\_ASSGN} \\
\\
\frac{n_1 > n_2}{\tilde{\rho} \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) \tilde{s} \Longrightarrow \tilde{\rho}; \cdot} \text{ EC\_FORT} \\
\\
\frac{\begin{array}{l} n_2 \geq n_1 \\ \tilde{\rho}[x \mapsto n_1] \vdash \tilde{s} \Longrightarrow \tilde{\rho}_1; \kappa_1 \\ \tilde{\rho}_1 \vdash \mathbf{for}(x := n_1 + 1; x \leq n_2; x := x + 1) \tilde{s} \Longrightarrow \tilde{\rho}_2; \kappa_2 \end{array}}{\tilde{\rho} \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1) \tilde{s} \Longrightarrow \tilde{\rho}_2; \kappa_1, \kappa_2} \text{ EC\_FORI} \\
\\
\frac{\begin{array}{l} \tilde{\rho} \vdash x \Downarrow [\tilde{w}_i^i]; \cdot \\ \tilde{\rho} \vdash \tilde{e}_1 \Downarrow n; \kappa_1 \\ \tilde{\rho} \vdash \tilde{e}_2 \Downarrow \tilde{w}; \kappa_2 \end{array}}{\tilde{\rho} \vdash x[\tilde{e}_1] := \tilde{e}_2 \Longrightarrow \tilde{\rho}[x \mapsto [\tilde{w}_i^i][n \mapsto \tilde{w}]]; \kappa_1, \kappa_2} \text{ EC\_AWRITE} \\
\\
\frac{\begin{array}{l} \tilde{\rho} \vdash \tilde{e} \Downarrow \top; \kappa_1 \\ \tilde{\rho} \vdash \tilde{s}_1 \Longrightarrow \tilde{\rho}'; \kappa_2 \end{array}}{\tilde{\rho} \vdash \mathbf{if}(\tilde{e}, \tilde{s}_1, \tilde{s}_2) \Longrightarrow \tilde{\rho}'; \kappa_1, \kappa_2} \text{ EC\_IFT} \\
\\
\frac{\begin{array}{l} \tilde{\rho} \vdash \tilde{e} \Downarrow \perp; \kappa_1 \\ \tilde{\rho} \vdash \tilde{s}_2 \Longrightarrow \tilde{\rho}'; \kappa_2 \end{array}}{\tilde{\rho} \vdash \mathbf{if}(\tilde{e}, \tilde{s}_1, \tilde{s}_2) \Longrightarrow \tilde{\rho}'; \kappa_1, \kappa_2} \text{ EC\_IFF} \\
\\
\frac{\tilde{\rho} \vdash \tilde{e} \Downarrow r; \kappa}{\tilde{\rho} \vdash \mathbf{out} \tilde{e} \Longrightarrow \tilde{\rho}; \kappa, \text{Out}(r)} \text{ EC\_OUT} \\
\\
\frac{\begin{array}{l} \tilde{\rho} \vdash \tilde{s}_1 \Longrightarrow \tilde{\rho}_1; \kappa_1 \\ \tilde{\rho}_1 \vdash \tilde{s}_2 \Longrightarrow \tilde{\rho}_2; \kappa_2 \end{array}}{\tilde{\rho} \vdash \tilde{s}_1; \tilde{s}_2 \Longrightarrow \tilde{\rho}_2; \kappa_1, \kappa_2} \text{ EC\_SEQ}
\end{array}$$

Figure 11: Target command evaluation

$$\begin{array}{ll}
b & ::= \text{Share (byte string)} \\
\\
\hat{\rho} & ::= \text{Circuit environment} \\
& \quad | \cdot \\
& \quad | \hat{\rho}[r \mapsto b]
\end{array}$$

Figure 12: Circuit runtime

$$\boxed{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa \mapsto \widehat{\rho}'_1, \widehat{\rho}'_2; O}$$

$$\begin{array}{c}
\frac{}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \cdot \mapsto \widehat{\rho}_1, \widehat{\rho}_2; \cdot} \text{CKT\_EMP} \\
\\
\frac{
\begin{array}{l}
n_1 = \mathcal{D}_{\mathcal{A}}(\widehat{\rho}_1[r_1], \widehat{\rho}_2[r_1]) \\
n_2 = \mathcal{D}_{\mathcal{A}}(\widehat{\rho}_1[r_2], \widehat{\rho}_2[r_2]) \\
(b_1, b_2) = \mathcal{E}_{\mathcal{A}}(n_1 + n_2)
\end{array}
}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \oplus(r_1, r_2, r_3) \mapsto \widehat{\rho}_1[r_3 \mapsto b_1], \widehat{\rho}_2[r_3 \mapsto b_2]; \cdot} \text{CKT\_ADD} \\
\\
\frac{
\begin{array}{l}
\top = \mathcal{D}_{\mathcal{B}}(\widehat{\rho}_1[r_1], \widehat{\rho}_2[r_1]) \\
c = \mathcal{D}_{\mathcal{B}}(\widehat{\rho}_1[r_2], \widehat{\rho}_2[r_2]) \\
(b_1, b_2) = \mathcal{E}_{\mathcal{B}}(c)
\end{array}
}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \text{Mux}(r_1, r_2, r_3, r_4) \mapsto \widehat{\rho}_1[r_4 \mapsto b_1], \widehat{\rho}_2[r_4 \mapsto b_2]; \cdot} \text{CKT\_MUXT} \\
\\
\frac{
\begin{array}{l}
\perp = \mathcal{D}_{\mathcal{B}}(\widehat{\rho}_1[r_1], \widehat{\rho}_2[r_1]) \\
c = \mathcal{D}_{\mathcal{B}}(\widehat{\rho}_1[r_3], \widehat{\rho}_2[r_3]) \\
(b_1, b_2) = \mathcal{E}_{\mathcal{B}}(c)
\end{array}
}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \text{Mux}(r_1, r_2, r_3, r_4) \mapsto \widehat{\rho}_1[r_4 \mapsto b_1], \widehat{\rho}_2[r_4 \mapsto b_2]; \cdot} \text{CKT\_MUXF} \\
\\
\frac{
\begin{array}{l}
n_1 = \mathcal{D}_{\mathcal{B}}(\widehat{\rho}_1[r_1], \widehat{\rho}_2[r_1]) \\
n_2 = \mathcal{D}_{\mathcal{B}}(\widehat{\rho}_1[r_2], \widehat{\rho}_2[r_2]) \\
(b_1, b_2) = \mathcal{E}_{\mathcal{B}}(n_1 > n_2)
\end{array}
}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \text{Gt}(r_1, r_2, r_3) \mapsto \widehat{\rho}_1[r_3 \mapsto b_1], \widehat{\rho}_2[r_3 \mapsto b_2]; \cdot} \text{CKT\_GT} \\
\\
\frac{
\begin{array}{l}
c = \mathcal{D}_{m_1}(\widehat{\rho}_1[r_1], \widehat{\rho}_2[r_1]) \\
(b_1, b_2) = \mathcal{E}_m(c)
\end{array}
}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash r_1 \triangleright_m r_2 \mapsto \widehat{\rho}_1[r_2 \mapsto b_1], \widehat{\rho}_2[r_2 \mapsto b_2]; \cdot} \text{CKT\_COERCE} \\
\\
\frac{c = \mathcal{D}_m(\widehat{\rho}_1[r], \widehat{\rho}_2[r])}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \text{Out}(r) \mapsto \widehat{\rho}_1, \widehat{\rho}_2; c} \text{CKT\_OUT} \\
\\
\frac{
\begin{array}{l}
\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa_1 \mapsto \widehat{\rho}'_1, \widehat{\rho}'_2; O_1 \\
\widehat{\rho}'_1, \widehat{\rho}'_2 \vdash \kappa_2 \mapsto \widehat{\rho}''_1, \widehat{\rho}''_2; O_2
\end{array}
}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa_1, \kappa_2 \mapsto \widehat{\rho}'_1, \widehat{\rho}''_2; O_1, O_2} \text{CKT\_SEQ}
\end{array}$$

Figure 13: Circuit evaluation

$$\boxed{\Gamma \sim \rho}$$

$$\begin{array}{c}
\frac{}{\cdot \sim \cdot} \text{SEN\_EMP} \\
\\
\frac{
\begin{array}{l}
v : [\tau]_{\mathcal{P}} \\
\Gamma \sim \rho
\end{array}
}{\Gamma, x : \tau \sim \rho[x \mapsto v]} \text{SEN\_BND}
\end{array}$$

Figure 14: Source environment and type environment consistency

$$\boxed{\Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2}$$

$$\frac{}{\Gamma \vdash \cdot \hookrightarrow \cdot; \cdot, \cdot} \text{EN\_EMP}$$

$$\frac{\begin{array}{l} \Gamma(x) = \sigma \\ \text{label}(\sigma) = \mathcal{P} \\ \Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2 \end{array}}{\Gamma \vdash \rho[x \mapsto c] \hookrightarrow \tilde{\rho}[x \mapsto c]; \hat{\rho}_1, \hat{\rho}_2} \text{EN\_PBT}$$

$$\frac{\begin{array}{l} \Gamma(x) = \sigma \\ \text{label}(\sigma) = m \\ r = \text{next\_range}() \\ (b_1, b_2) = \mathcal{E}_m(c) \\ \Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2 \end{array}}{\Gamma \vdash \rho[x \mapsto c] \hookrightarrow \tilde{\rho}[x \mapsto r]; \hat{\rho}_1[r \mapsto b_1], \hat{\rho}_2[r \mapsto b_2]} \text{EN\_SBT}$$

$$\frac{\begin{array}{l} \Gamma(x) = \sigma[ ] \\ \text{label}(\sigma) = \mathcal{P} \\ \Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2 \end{array}}{\Gamma \vdash \rho[x \mapsto [\bar{c}_i^i]] \hookrightarrow \tilde{\rho}[x \mapsto [\bar{c}_i^i]]; \hat{\rho}_1, \hat{\rho}_2} \text{EN\_PARR}$$

$$\frac{\begin{array}{l} \Gamma(x) = \sigma[ ] \\ \text{label}(\sigma) = m \\ r_i = \text{next\_range}() \\ (b_{1\ i}, b_{2\ i}) = \mathcal{E}_m(c_i) \\ \Gamma \vdash \rho \hookrightarrow \tilde{\rho}; \hat{\rho}_1, \hat{\rho}_2 \end{array}}{\Gamma \vdash \rho[x \mapsto [\bar{c}_i^i]] \hookrightarrow \tilde{\rho}[x \mapsto [\bar{r}_i^i]]; \hat{\rho}_1[r_i \mapsto b_{1\ i}], \hat{\rho}_2[r_i \mapsto b_{2\ i}]} \text{EN\_SARR}$$

Figure 15: Source environment to target environment compilation