1 Formal Development

Correctness theorem that we will aim at:

If:

- $\Gamma \vdash s \leadsto \widetilde{s} \mid \Gamma'$
- $\Gamma \sim \rho$
- $\bullet \ \rho \vdash s \longrightarrow \rho'; O$
- $\bullet \ \Gamma \vdash \rho \hookrightarrow \widetilde{\rho}$
- $\Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2$

Then:

- $\bullet \ \widetilde{\rho} \vdash \widetilde{s} \Longrightarrow \widetilde{\rho}'; \kappa$
- $\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa \longmapsto \widehat{\rho}'_1, \widehat{\rho}'_2; O$

```
Secret label
m
                      \mathcal{A}
                      \mathcal{B}
\ell
                                        Label
                      \mathcal{P}
                                        Base type
            ::=
                      \mathsf{uint}^\ell
                      \mathsf{bool}^\ell
                                        Type
            ::=
                      \sigma[\ ]
c
            ::=
                                        Constant
                      n
                      \top
                                        Source expression
                    \begin{aligned} e_1 + e_2 \\ \mathbf{cond}(e, e_1, e_2) \end{aligned}
                    e_1 > e_2
x[e]
           ::=
                                                                                      Source statement
s
                    \tau x
                    x := e
                    \mathbf{for}(x := n_1; x \le n_2; x := x + 1) \ s
                    x[e_1] := e_2
                    \mathbf{if}(e,s_1,s_2)
                    \mathbf{out}\, e
                    s_1; s_2
```

Figure 1: Source language

Figure 2: Source runtime

Figure 3: Source expression evaluation

$$\begin{array}{c} \operatorname{default}(\tau) = v \\ \hline \rho \vdash \tau \: x \longrightarrow \rho[x \mapsto v]; \end{array} \quad \text{SC_DECL} \\ \hline \rho \vdash e \downarrow v \\ \hline \rho \vdash x := e \longrightarrow \rho[x \mapsto v]; \end{array} \quad \text{SC_ASSGN} \\ \hline \begin{array}{c} n_1 > n_2 \\ \hline \rho \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \: s \longrightarrow \rho; \end{array} \quad \text{SC_FORT} \\ \hline \begin{array}{c} n_2 \geq n_1 \\ \hline \rho_1 \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \: s \longrightarrow \rho_2; O_2 \\ \hline \rho \vdash \operatorname{for}(x := n_1 + 1; x \leq n_2; x := x + 1) \: s \longrightarrow \rho_2; O_2 \\ \hline \rho \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \: s \longrightarrow \rho_2; O_1, O_2 \end{array} \quad \text{SC_FORI} \\ \hline \begin{array}{c} \rho \vdash x \downarrow [\overline{c_i}^i] \\ \rho \vdash e_1 \downarrow n \\ \rho \vdash e_2 \downarrow c \\ \hline \hline \rho \vdash x [e_1] := e_2 \longrightarrow \rho[x \mapsto [\overline{c_i}^i][n \mapsto c]]; O_1, O_2 \end{array} \quad \text{SC_AWRITE} \\ \hline \begin{array}{c} \rho \vdash e \downarrow \top \\ \rho \vdash s_1 \longrightarrow \rho'; O \\ \hline \rho \vdash \operatorname{if}(e, s_1, s_2) \longrightarrow \rho'; O \end{array} \quad \text{SC_IFT} \\ \hline \begin{array}{c} \rho \vdash e \downarrow \bot \\ \rho \vdash s_2 \longrightarrow \rho'; O \\ \hline \rho \vdash \operatorname{if}(e, s_1, s_2) \longrightarrow \rho'; O \end{array} \quad \text{SC_IFF} \\ \hline \begin{array}{c} \rho \vdash e \downarrow v \\ \hline \rho \vdash \operatorname{out} e \longrightarrow \rho; v \end{array} \quad \text{SC_OUT} \\ \hline \begin{array}{c} \rho \vdash s_1 \longrightarrow \rho_1; O_1 \\ \rho_1 \vdash s_2 \longrightarrow \rho_2; O_2 \\ \hline \rho \vdash s_1; s_2 \longrightarrow \rho_2; O_2 \end{array} \quad \text{SC_SEQ} \end{array}$$

Figure 4: Source command evaluation

 $v:\tau$

Figure 5: Value typing

Figure 6: Target language

$\Gamma \vdash e : \tau \leadsto \widetilde{e}$

$$\begin{array}{c|c} \hline{\Gamma \vdash n : \mathsf{uint}^{\mathcal{P}} \leadsto n} & \text{S_CONST} \\ \hline \hline{\Gamma \vdash T : \mathsf{bool}^{\mathcal{P}} \leadsto \top} & \text{S_TRUE} \\ \hline \hline{\Gamma \vdash T : \mathsf{bool}^{\mathcal{P}} \leadsto \bot} & \text{S_FALSE} \\ \hline \hline{\Gamma \vdash L : \mathsf{bool}^{\mathcal{P}} \leadsto \bot} & \text{S_VAR} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_PADD} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_PADD} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{A}} \leadsto \widetilde{e}_i} & \text{S_SADD} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{bool}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_SADD} \\ \hline \hline{\Gamma \vdash e : \mathsf{bool}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_PCOND} \\ \hline \hline{\Gamma \vdash \mathsf{cond}(e, e_1, e_2) : \sigma \leadsto \mathsf{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)} & \text{S_PCOND} \\ \hline \hline{\Gamma \vdash \mathsf{cond}(e, e_1, e_2) : \sigma \leadsto \mathsf{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)} & \text{S_PCOND} \\ \hline \hline{\Gamma \vdash \mathsf{cond}(e, e_1, e_2) : \sigma \leadsto \mathsf{cond}_{\mathcal{B}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)} & \text{S_PCOND} \\ \hline \hline{\Gamma \vdash \mathsf{cond}(e, e_1, e_2) : \sigma \leadsto \mathsf{cond}_{\mathcal{B}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)} & \text{S_PGT} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_PGT} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{B}} \leadsto \widetilde{e}_i} & \text{S_PGT} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{B}} \leadsto \widetilde{e}_i} & \text{S_SGT} \\ \hline \hline{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_AREAD} \\ \hline \hline{\Gamma \vdash e : \mathsf{uint}^{\mathcal{P}} \leadsto \widetilde{e}} & \text{S_AREAD} \\ \hline \hline{\Gamma \vdash e : \sigma_1 \leadsto \widetilde{e}} & \text{base}(\sigma_1) = \mathsf{base}(\sigma_2) \\ \hline \hline \mathsf{label}(\sigma_2) = m \\ \hline \hline{\Gamma \vdash e : \sigma_2 \leadsto \widetilde{e} \vDash m} & \text{S_SUB} \\ \hline \hline \hline \hline \end{array}$$

Figure 7: Expression compilation

$$\begin{array}{c|c} \hline \Gamma \vdash s \leadsto \widetilde{s} \mid \Gamma' \\ \hline \hline \Gamma \vdash \tau x \leadsto \tau x \mid \Gamma, x : \tau \\ \hline \Gamma(x) = \sigma \\ \hline \Gamma \vdash e : \sigma \leadsto \widetilde{e} \\ \hline \Gamma \vdash x := e \leadsto x := \widetilde{e} \mid \Gamma \\ \hline \Gamma, x : \operatorname{uint}^P \vdash s \leadsto \widetilde{s} \mid - \\ x \not\in \operatorname{modifies}(s) \\ \hline \Gamma \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) s \leadsto \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \widetilde{s} \mid \Gamma \\ \hline \Gamma \vdash x : \sigma[\] \leadsto x \\ \Gamma \vdash e_1 : \operatorname{uint}^P \leadsto \widetilde{e}_1 \\ \hline \Gamma \vdash e_2 : \sigma \leadsto \widetilde{e}_2 \\ \hline \hline \Gamma \vdash x \mid e_1] := e_2 \leadsto x \mid \widetilde{e}_1] := \widetilde{e}_2 \mid \Gamma \\ \hline \Gamma \vdash e : \operatorname{bool}^P \leadsto \widetilde{e} \\ \hline \Gamma \vdash s_i \leadsto \widetilde{s}_i \mid - \\ \hline \Gamma \vdash \operatorname{if}(e, s_1, s_2) \leadsto \operatorname{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \mid \Gamma \\ \hline \Gamma \vdash \operatorname{out} e \leadsto \operatorname{out} \widetilde{e} \mid \Gamma \\ \hline \Gamma \vdash \operatorname{s}_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \mid \widetilde{s}_2 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\ \hline \Gamma \vdash s_1 \bowtie \widetilde{s}_1 \bowtie \widetilde{s}_1 \mid \Gamma' \\$$

Figure 8: Command compilation

Figure 9: Target runtime

$$\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa$$

$$\overline{\widetilde{\rho} \vdash c \Downarrow c;} \cdot \text{EE_CONST}$$

$$\overline{\widetilde{\rho} \vdash x \Downarrow \widetilde{\rho}[x];} \cdot \text{EE_VAR}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow r_i; \kappa_i} \quad \text{EE_PADD}$$

$$\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow r_i; \kappa_i \quad \text{EE_SADD}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 + A \widetilde{e}_2 \Downarrow r_3; \kappa_1, \kappa_2, \oplus (r_1, r_2, r_3)} \quad \text{EE_SADD}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow \widetilde{v}; \kappa_1} \quad \text{EE_PCONDT}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow \widetilde{v}; \kappa_1} \quad \text{EE_PCONDT}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow \widetilde{v}; \kappa, \kappa_1} \quad \text{EE_PCONDF}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow \widetilde{v}; \kappa, \kappa_2} \quad \text{EE_PCONDF}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow \widetilde{v}; \kappa, \kappa_2} \quad \text{EE_PCONDF}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow \widetilde{v}; \kappa, \kappa_2} \quad \text{EE_PCONDF}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow r_3; \kappa, \kappa_1, \kappa_2, \operatorname{Mux}(r, r_1, r_2, r_3)} \quad \text{EE_SCOND}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow r_3; \kappa, \kappa_1, \kappa_2, \operatorname{Mux}(r, r_1, r_2, r_3)} \quad \text{EE_SCOND}$$

$$\overline{\widetilde{\rho} \vdash \operatorname{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow r_3; \kappa, \kappa_1, \kappa_2, \operatorname{Mux}(r, r_1, r_2, r_3)} \quad \text{EE_SCOND}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow r_i; \kappa_i} \quad \text{EE_PGT}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow r_i; \kappa_i} \quad \text{EE_PGT}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow r_i; \kappa_i} \quad \text{EE_SGT}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}_1 \gg r_i; \kappa_1} \quad \text{EE_AREAD}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}} \Downarrow r_i; \kappa_1} \quad \text{EE_AREAD}$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}} \Downarrow r_i; \kappa$$

$$\overline{\widetilde{\rho} \vdash \widetilde{e}} \bowtie r_i; \kappa$$

$$\overline{\widetilde{\rho}} \vdash \widetilde{e}} \vdash \widetilde{e}} \bowtie r_i; \kappa$$

$$\overline{\widetilde{\rho}} \vdash \widetilde{e}} \vdash \widetilde{e}}$$

Figure 10: Target expression evaluation

$$\begin{array}{c} \frac{\operatorname{default}(\tau) = \widetilde{v}; \kappa}{\widetilde{\rho} \vdash \tau \, x \Longrightarrow \widetilde{\rho}[x \mapsto \widetilde{v}]; \kappa} & \text{EC_DECL} \\ \\ \frac{\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa}{\widetilde{\rho} \vdash \kappa : = \widetilde{e} \Longrightarrow \widetilde{\rho}[x \mapsto \widetilde{v}]; \kappa} & \text{EC_ASSGN} \\ \\ \frac{n_1 > n_2}{\widetilde{\rho} \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \ \widetilde{s} \Longrightarrow \widetilde{\rho};} & \text{EC_FORT} \\ \\ \frac{n_2 \geq n_1}{\widetilde{\rho}[x \mapsto n_1] \vdash \widetilde{s} \Longrightarrow \widetilde{\rho}_1; \kappa_1} \\ \widetilde{\rho}_1 \vdash \operatorname{for}(x := n_1 + 1; x \leq n_2; x := x + 1) \ \widetilde{s} \Longrightarrow \widetilde{\rho}_2; \kappa_2}{\widetilde{\rho} \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \ \widetilde{s} \Longrightarrow \widetilde{\rho}_2; \kappa_1, \kappa_2} & \text{EC_FORI} \\ \\ \\ \frac{\widetilde{\rho} \vdash \kappa_1 \Downarrow n_1 + n_2}{\widetilde{\rho} \vdash \kappa_1 \vdash \kappa_1 + n_1} \\ \widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow n_1; \kappa_1} \\ \widetilde{\rho} \vdash \widetilde$$

Figure 11: Target command evaluation

$$\begin{array}{cccc} \widehat{c} & & ::= & & \text{Circuit value (share)} \\ & | & \mathcal{E}_1^m(c) & & & \\ & | & \mathcal{E}_2^m(c) & & & \\ \\ \widehat{\rho} & & ::= & & \text{Circuit environment} \\ & | & \widehat{\rho}[r \mapsto \widehat{c}] & & & \\ \end{array}$$

Figure 12: Circuit runtime

$$\widehat{\rho_1}, \widehat{\rho_2} \vdash \kappa \longmapsto \widehat{\rho_1}, \widehat{\rho_2}; O$$

$$\widehat{\rho_1}, \widehat{\rho_2} \vdash \iota \longmapsto \widehat{\rho_1}, \widehat{\rho_2}; CKT_EMP$$

$$\widehat{\rho_1}[r_1] = \mathcal{E}_1^A(n_1)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^A(n_2)$$

$$\widehat{\rho_1}[r_2] = \mathcal{E}_1^A(n_2)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^A(n_2)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^A(n_2)$$

$$\widehat{\rho_1}, \widehat{\rho_2} \vdash \oplus (r_1, r_2, r_3) \longmapsto \widehat{\rho_1}[r_3 \mapsto \mathcal{E}_1^A((n_1 + n_2))], \widehat{\rho_2}[r_3 \mapsto \mathcal{E}_2^A((n_1 + n_2))]; CKT_ADD$$

$$\widehat{\rho_1}[r_1] = \mathcal{E}_2^B(T)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^B(T)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_1}[r_2] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_1}[r_2] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_1}[r_1] = \mathcal{E}_1^B(\bot)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^B(\bot)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_1}[r_3] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_1}[r_3] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_2}[r_3] = \mathcal{E}_2^B(c)$$

$$\widehat{\rho_1}[r_3] = \mathcal{E}_2^B(n_1)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^B(n_1)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^B(n_2)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^B(n_2)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^B(n_2)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^B(n_2)$$

$$\widehat{\rho_2}[r_2] = \mathcal{E}_2^B(n_2)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^{B(1)}(c)$$

$$\widehat{\rho_1}[r_2] \in \mathcal{E}_1^{B(1)}(c)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^{B(1)}(c)$$

$$\widehat{\rho_2}[r_1] = \mathcal{E}_2^{B(1)}(c)$$

$$\widehat{\rho_1}[r_2] \in \mathcal{E}_2^{B(1)}(c)$$

$$\widehat{\rho_1}[r_2} \in \mathcal$$

Figure 13: Circuit evaluation

$$\frac{\Gamma \sim \rho}{\cdots} \quad \text{SEN_EMP}$$

$$\frac{v: [\tau]_{\mathcal{P}}}{\Gamma \sim \rho}$$

$$\frac{\Gamma \sim \rho}{\Gamma, x: \tau \sim \rho[x \mapsto v]} \quad \text{SEN_BND}$$

Figure 14: Source environment and type environment consistency

Figure 15: Source environment to target environment compilation

$$\begin{array}{c} \overline{\Gamma;\widetilde{\rho}\vdash\rho\hookrightarrow\widehat{\rho}_{1},\widehat{\rho}_{2}} \\ \hline \\ \overline{\Gamma;\widetilde{\rho}\vdash\cdot\hookrightarrow\cdot,\cdot} & \text{CEN_EMP} \\ \hline \\ \Gamma(x) = \sigma \\ \text{label}(\sigma) = \mathcal{P} \\ \overline{\Gamma;\widetilde{\rho}\vdash\rho\hookrightarrow\widehat{\rho}_{1},\widehat{\rho}_{2}} \\ \hline \overline{\Gamma;\widetilde{\rho}\vdash\rho[x\mapsto c]\hookrightarrow\widehat{\rho}_{1},\widehat{\rho}_{2}} & \text{CEN_PBT} \\ \hline \\ \Gamma(x) = \sigma \\ \text{label}(\sigma) = m \\ \widetilde{\rho}[x] = r \\ \overline{\Gamma;\widetilde{\rho}\vdash\rho[x\mapsto c]\hookrightarrow\widehat{\rho}_{1}[r\mapsto\mathcal{E}_{1}^{m}(c)],\widehat{\rho}_{2}[r\mapsto\mathcal{E}_{2}^{m}(c)]} & \text{CEN_SBT} \\ \hline \\ \Gamma(x) = \sigma[] \\ \text{label}(\sigma) = \mathcal{P} \\ \overline{\Gamma;\widetilde{\rho}\vdash\rho[x\mapsto\overline{\rho}_{1},\widehat{\rho}_{2}]} & \text{CEN_PARR} \\ \hline \\ \Gamma(x) = \sigma[] \\ \text{label}(\sigma) = m \\ \overline{\rho}[x] = [\overline{r_{i}}^{i}] \\ \overline{\Gamma;\widetilde{\rho}\vdash\rho[x\mapsto[\overline{c_{i}}^{i}]]\hookrightarrow\widehat{\rho}_{1},\widehat{\rho}_{2}} & \text{CEN_PARR} \\ \hline \\ \Gamma(x) = \sigma[] \\ \text{label}(\sigma) = m \\ \overline{\rho}[x] = [\overline{r_{i}}^{i}] \\ \overline{\Gamma;\widetilde{\rho}\vdash\rho[x\mapsto[\overline{c_{i}}^{i}]]\hookrightarrow\widehat{\rho}_{1},\widehat{\rho}_{2}} & \text{CEN_SARR} \\ \hline \end{array}$$

Figure 16: Source environment to circuit environment compilation