```
Secret label
m
           ::=
                     \mathcal{A}
                     \mathcal{B}
                                                                                  Label
                     \mathcal{P}
                     m
                                                                                  Base type
                     \mathsf{uint}^\ell
                     \mathsf{bool}^\ell
            ::=
                                                                                  Type
                     \sigma
                    \sigma[\ ]
            ::=
                                                                                  Constant
                     n
                     \top
                                                                                  {\bf Source\ expression}
                     \boldsymbol{x}
                    e_1 + e_2
                    \mathbf{cond}(e, e_1, e_2)
                    e_1 > e_2
                     x[e]
                                                                                  Source statement
s
                     \tau x
                    for(x := n_1; x \le n_2; x := x + 1) s
                    x[e_1] := e_2

if (e, s_1, s_2)
                     \mathbf{out}\, e
                     s_1; s_2
```

Figure 1: Source language

Figure 2: Source runtime

Figure 3: Source expression evaluation

$$\begin{array}{c} \frac{\text{default}(\tau) = v}{\rho \vdash \tau \, x \longrightarrow \rho[x \mapsto v];} \quad \text{SC\_DECL} \\ \frac{\rho \vdash e \downarrow v}{\rho \vdash x := e \longrightarrow \rho[x \mapsto v];} \quad \text{SC\_ASSGN} \\ \frac{n_1 > n_2}{\rho \vdash \text{for}(x := n_1; x \leq n_2; x := x + 1) \, s \longrightarrow \rho;} \quad \text{SC\_FORT} \\ \frac{n_2 \geq n_1}{\rho \vdash \text{for}(x := n_1 + 1; x \leq n_2; x := x + 1) \, s \longrightarrow \rho_2; O_2} \\ \frac{\rho \vdash \text{for}(x := n_1 + 1; x \leq n_2; x := x + 1) \, s \longrightarrow \rho_2; O_2}{\rho \vdash \text{for}(x := n_1 ; x \leq n_2; x := x + 1) \, s \longrightarrow \rho_2; O_1, O_2} \quad \text{SC\_FORI} \\ \frac{\rho \vdash x \downarrow [\overline{c_i}^i]}{\rho \vdash e_1 \downarrow n} \\ \frac{\rho \vdash e_1 \downarrow n}{\rho \vdash e_2 \downarrow c} \\ \frac{\rho \vdash e_1 \downarrow n}{\rho \vdash e_1 \downarrow n} \\ \frac{\rho \vdash e \downarrow \downarrow}{\rho \vdash s_1 \longrightarrow \rho'; O} \quad \text{SC\_AWRITE} \\ \frac{\rho \vdash e \downarrow \bot}{\rho \vdash s_2 \longrightarrow \rho'; O} \\ \frac{\rho \vdash e \downarrow \bot}{\rho \vdash \text{if}(e, s_1, s_2) \longrightarrow \rho'; O} \quad \text{SC\_IFT} \\ \frac{\rho \vdash e \downarrow v}{\rho \vdash \text{out} \, e \longrightarrow \rho; v} \quad \text{SC\_OUT} \\ \frac{\rho \vdash s_1 \longrightarrow \rho_1; O_1}{\rho \vdash s_1 \longrightarrow \rho_1; O_1} \\ \frac{\rho \vdash s_1 \longrightarrow \rho_2; O_2}{\rho \vdash s_1; s_2 \longrightarrow \rho_2; O_2} \quad \text{SC\_SEQ} \end{array}$$

Figure 4: Source command evaluation

Figure 5: Target language

## $\Gamma \vdash e : \tau \leadsto \widetilde{e}$

Figure 6: Expression compilation

$$\begin{array}{c|c} \hline \Gamma \vdash s \leadsto \widetilde{s} \mid \Gamma' \\ \hline \hline \Gamma \vdash \tau x \leadsto \tau x \mid \Gamma, x : \tau \\ \hline \Gamma(x) = \sigma \\ \hline \Gamma \vdash e : \sigma \leadsto \widetilde{e} \\ \hline \Gamma \vdash x := e \leadsto x := \widetilde{e} \mid \Gamma \\ \hline \Gamma, x : \operatorname{uint}^{\mathcal{P}} \vdash s \leadsto \widetilde{s} \mid - \\ \hline x \not\in \operatorname{modifies}(s) \\ \hline \Gamma \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) s \leadsto \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \widetilde{s} \mid \Gamma \\ \hline \Gamma \vdash x : \sigma[ \ ] \leadsto x \\ \hline \Gamma \vdash a : \operatorname{uint}^{\mathcal{P}} \leadsto \widetilde{e}_1 \\ \hline \Gamma \vdash e_2 : \sigma \leadsto \widetilde{e}_2 \\ \hline \hline \Gamma \vdash x [e_1] := e_2 \leadsto x [\widetilde{e}_1] := \widetilde{e}_2 \mid \Gamma \\ \hline \Gamma \vdash e : \operatorname{bool}^{\mathcal{P}} \leadsto \widetilde{e} \\ \hline \Gamma \vdash s_i \leadsto \widetilde{s}_i \mid - \\ \hline \Gamma \vdash \operatorname{if}(e, s_1, s_2) \leadsto \operatorname{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \mid \Gamma \\ \hline \Gamma \vdash \operatorname{out} e \leadsto \operatorname{out} \widetilde{e} \mid \Gamma \\ \hline \Gamma \vdash \operatorname{out} e \leadsto \operatorname{out} \widetilde{e} \mid \Gamma \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \Gamma \vdash s_1 : s_2 \leadsto \widetilde{s}_1 \mid \widetilde{s}_2 \mid \Gamma' \\ \hline \end{array}$$

Figure 7: Command compilation

```
Wire id range
                    ::=
\widetilde{w}
                                                                                                    Compiled base value
                     ::=
                                     c
                                                                                                   Compiled value
\widetilde{v}
                    ::=
                                     \widetilde{w}
                                    [\,\overline{\widetilde{w}_i}^{\,\,i}\,]
                                                                                                    Circuit
                                   egin{array}{l} \oplus (r_1, r_2, r_3) \ \otimes (r_1, r_2, r_3) \ \mathrm{Mux}(r_1, r_2, r_3, r_4) \ \mathrm{Gt}(r_1, r_2, r_3) \end{array}
                                    r_1 \rhd_m r_2
\mathsf{Out}(r)
                                     \kappa_1, \kappa_2
\widetilde{
ho}
                                                                                                    Runtime environment
                                    \overset{\cdot}{\widetilde{\rho}}[x\mapsto \widetilde{v}]
```

Figure 8: Target runtime

$$\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa$$

Figure 9: Target expression evaluation

$$\begin{array}{c} \frac{\operatorname{default}(\tau) = \widetilde{v}; \kappa}{\widetilde{\rho} \vdash \tau \, x \Longrightarrow \widetilde{\rho}[x \mapsto \widetilde{v}]; \kappa} & \text{EC\_DECL} \\ \\ \frac{\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa}{\widetilde{\rho} \vdash \kappa : = \widetilde{e} \Longrightarrow \widetilde{\rho}[x \mapsto \widetilde{v}]; \kappa} & \text{EC\_ASSGN} \\ \\ \frac{n_1 > n_2}{\widetilde{\rho} \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \ \widetilde{s} \Longrightarrow \widetilde{\rho};} & \text{EC\_FORT} \\ \\ \frac{n_2 \geq n_1}{\widetilde{\rho}[x \mapsto n_1] \vdash \widetilde{s} \Longrightarrow \widetilde{\rho}_1; \kappa_1} \\ \widetilde{\rho}_1 \vdash \operatorname{for}(x := n_1 + 1; x \leq n_2; x := x + 1) \ \widetilde{s} \Longrightarrow \widetilde{\rho}_2; \kappa_2}{\widetilde{\rho} \vdash \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \ \widetilde{s} \Longrightarrow \widetilde{\rho}_2; \kappa_1, \kappa_2} & \text{EC\_FORI} \\ \\ \\ \frac{\widetilde{\rho} \vdash \kappa_1 \Downarrow n_1 \vdash \kappa_1}{\widetilde{\rho} \vdash \kappa_1 \vdash \kappa_1 \vdash \kappa_1} \\ \widetilde{\rho} \vdash \kappa_1 \Downarrow n_1 \vdash \kappa_1} \\ \widetilde{\rho} \vdash \kappa_1 \Downarrow n_1 \vdash \kappa_1} \\ \widetilde{\rho} \vdash \kappa_1 \Longrightarrow \widetilde{\rho}'; \kappa_2} \\ \widetilde{\rho} \vdash \operatorname{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \Longrightarrow \widetilde{\rho}'; \kappa_1, \kappa_2} & \text{EC\_AWRITE} \\ \\ \\ \frac{\widetilde{\rho} \vdash \widetilde{e} \Downarrow \vdash \kappa_1}{\widetilde{\rho} \vdash \widetilde{s}_1 \Longrightarrow \widetilde{\rho}'; \kappa_2} \Longrightarrow \widetilde{\rho}'; \kappa_1, \kappa_2} \\ \widetilde{\rho} \vdash \operatorname{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \Longrightarrow \widetilde{\rho}'; \kappa_1, \kappa_2} & \operatorname{EC\_IFT} \\ \\ \\ \widetilde{\rho} \vdash \widetilde{e} \Downarrow \vdash r; \kappa} \\ \widetilde{\rho} \vdash \operatorname{out} \widetilde{e} \Longrightarrow \widetilde{\rho}; \kappa, \operatorname{Out}(r) & \operatorname{EC\_OUT} \\ \\ \widetilde{\rho} \vdash \widetilde{s}_1 \varinjlim \widetilde{\rho}_1; \kappa_1} \\ \widetilde{\rho} \vdash \widetilde{s}_2 \Longrightarrow \widetilde{\rho}_2; \kappa_2} \\ \widetilde{\rho} \vdash \widetilde{s}_1; \widetilde{s}_2 \Longrightarrow \widetilde{\rho}_2; \kappa_2} & \operatorname{EC\_SEQ} \\ \end{array}$$

Figure 10: Target command evaluation