```
Secret label
           ::=
                   \mathcal{A}
                                                               Label
                                                               Base type
           ::=
                    uint
                    bool
                                                               Type
           ::=
                   \sigma^\ell
                   \mathsf{uint}^\ell[\ ]
                                                               Expression
                    e_1 \oplus e_2
                   e_1 ? e_2 : e_3
                   e_1 \oplus_s e_2
                    \mathbf{mux}_s \ e \ e_1 \ e_2
                    e_1 >_s e_2
                    e \rhd s
                                                               Command
                    \tau x = e
                   \mathbf{for}\;x\;\in\;[n\ldots m]\;\mathbf{do}\;c
                   x[e_1] := e_2
                    \mathbf{if}\ e\ c_1\ c_2
                    \mathbf{out}\; e
                    c_1; c_2
Γ
                                                               Type environment
                   \Gamma, x:\tau
```

$$\begin{array}{l} \Gamma \vdash e : \tau \leadsto e' \\ \hline \Gamma \vdash n : \mathsf{uint}^{\mathcal{P}} \leadsto n & \mathsf{S\_CONST} \\ \hline \Gamma \vdash n : \mathsf{uint}^{\mathcal{P}} \leadsto n & \mathsf{S\_VAR} \\ \hline \Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_PBINOP} \\ \hline \Gamma \vdash e_1 \oplus e_2 : \mathsf{uint}^{\mathcal{P}} \leadsto e_1' \oplus e_2' & \mathsf{S\_PBINOP} \\ \hline \Gamma \vdash e_i \oplus e_2 : \mathsf{uint}^{\mathcal{A}} \leadsto e_1' \oplus e_2' & \mathsf{S\_SBINOP} \\ \hline \Gamma \vdash e_i \oplus \mathsf{bool}^{\mathcal{P}} \leadsto e' & \mathsf{\Gamma} \vdash e_i : \tau \leadsto e_i' & \mathsf{S\_PCOND} \\ \hline \Gamma \vdash e : \mathsf{bool}^{\mathcal{B}} \leadsto e' & \mathsf{\Gamma} \vdash e_i : \tau \leadsto e_i' & \mathsf{S\_PCOND} \\ \hline \Gamma \vdash e : \mathsf{bool}^{\mathcal{B}} \leadsto e' & \mathsf{\Gamma} \vdash e_i : \tau \leadsto e_i' & \mathsf{S\_SCOND} \\ \hline \Gamma \vdash e : \mathsf{e} : \mathsf{init}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_PGT} \\ \hline \Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_PGT} \\ \hline \Gamma \vdash e_i : \mathsf{uint}^{\mathcal{B}} \leadsto e_i' & \mathsf{S\_SGT} \\ \hline \Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_SGT} \\ \hline \Gamma \vdash x : \mathsf{uint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_SGT} \\ \hline \Gamma \vdash x : \mathsf{uint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_AREAD} \\ \hline \Gamma \vdash x : \mathsf{uint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_AREAD} \\ \hline \Gamma \vdash x : \mathsf{lint}^{\mathcal{P}} \leadsto e_i' & \mathsf{S\_AREAD} \\ \hline \end{array}$$

 $\frac{\Gamma \vdash e : \sigma^{\ell} \leadsto e'}{\Gamma \vdash e : \sigma^{s} \leadsto e' \rhd s} \quad \text{S\_SUB}$ 

```
\Gamma \vdash c \leadsto c' \mid \Gamma'
\frac{\Gamma \vdash e : \tau \leadsto e'}{\Gamma \vdash \tau \: x = e \leadsto \tau \: x = e' \mid \Gamma, x : \tau} \quad \text{C-DECL}
    \Gamma(x) = \tau
\frac{\Gamma \vdash e : \tau \leadsto e'}{\Gamma \vdash x := e \leadsto x := e' \mid \Gamma} \quad \text{C-VASSGN}
    \Gamma, x: \mathsf{uint}^{\mathcal{P}} \vdash c \leadsto c' \mid \_
\frac{x \notin \mathsf{modifies}(c)}{\Gamma \vdash \mathbf{for} \ x \in [n \dots m] \ \mathbf{do} \ c \leadsto \mathbf{for} \ x \in [n \dots m] \ \mathbf{do} \ c' \mid \Gamma}
\begin{split} &\Gamma \vdash x : \mathsf{uint}^{\ell}[\;] \leadsto x \\ &\Gamma \vdash e_1 : \mathsf{uint}^{\mathcal{P}} \leadsto e_1' \\ &\frac{\Gamma \vdash e_2 : \mathsf{uint}^{\ell} \leadsto e_2'}{\Gamma \vdash x[e_1] := e_2 \leadsto x[e_1'] := e_2' \mid \Gamma} \quad \text{C_AWRITE} \end{split}
     \Gamma \vdash e : \mathsf{bool}^{\mathcal{P}} \leadsto e'

\frac{\Gamma \vdash c_1 \leadsto c_1' \mid \bot}{\Gamma \vdash c_2 \leadsto c_2' \mid \bot}

\frac{\Gamma \vdash c_2 \leadsto c_2' \mid \bot}{\Gamma \vdash \text{if } e \ c_1 \ c_2 \leadsto \text{if } e' \ c_1' \ c_2' \mid \Gamma}

C_IF
\frac{\Gamma \vdash e : \tau \leadsto e'}{\Gamma \vdash \mathbf{out} \; e \leadsto \mathbf{out} \; e' \mid \Gamma} \quad \text{C-OUT}
\frac{\Gamma \vdash c_1 \leadsto c_1' \mid \Gamma_1}{\Gamma_1 \vdash c_2 \leadsto c_2' \mid \Gamma'} \\ \frac{\Gamma_1 \vdash c_2 \leadsto c_2' \mid \Gamma'}{\Gamma \vdash c_1; c_2 \leadsto c_1'; c_2' \mid \Gamma'} \quad \text{C\_SEQ}
                                                                                                             Runtime base values
    w
                                                                                                             Runtime values
                               \begin{array}{cc} | & w \\ | & [\overline{v_i}^i] \end{array}
   \begin{array}{cccc} \rho & & ::= & & \\ & | & \cdot & \\ & & | & \rho[x \mapsto v] \end{array}
                                                                                                             Runtime environment
```

$$\begin{array}{c} \boxed{\rho_1,\rho_2 \vdash e \Downarrow v_1,v_2} \\ \hline \rho_1,\rho_2 \vdash n \Downarrow n,n \\ \hline \\ \rho_1[x] = \rho_2[x] = v \\ \hline \rho_1,\rho_2 \vdash x \Downarrow v,v \\ \hline \\ \rho_1[x] = w^{s,1} \\ \rho_2[x] = w^{s,2} \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow w_1, u_1 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _2 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e _1 \Downarrow v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2, n_1 > n_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \downarrow \vdash v_1, v_2 \\ \hline \rho_1,\rho_2 \vdash e \vdash v_1, v_2 \\ \hline \rho_1,\rho_$$

$$\begin{split} \boxed{\rho_1, \rho_2 \vdash c \longrightarrow \rho_1', \rho_2'} \\ &\frac{\rho_1, \rho_2 \vdash e \Downarrow v_1, v_2}{\rho_1, \rho_2 \vdash \tau \: x = e \longrightarrow \rho_1[x \mapsto v_1], \rho_2[x \mapsto v_2]} \quad \text{EC\_DECL} \\ &\frac{\rho_1, \rho_2 \vdash e \Downarrow v_1, v_2}{\rho_1, \rho_2 \vdash x := e \longrightarrow \rho_1[x \mapsto v_1], \rho_2[x \mapsto v_2]} \quad \text{EC\_ASSGN} \end{split}$$