1 Formal Development

Correctness theorem that we will aim at:
If:

- $\Gamma \vdash s \rightsquigarrow \widetilde{s} \mid \Gamma'$ (Compilation of source command)
- $\Gamma \sim \rho$ (Source environment, maps free variables to values)
- $\rho \vdash s \longrightarrow \rho'; O$ (Source semantics, no MPC)
- $\Gamma \vdash \rho \hookrightarrow \widetilde{\rho}$ (Compiling source environment to C++ environment, free secret variables are mapped to wire ranges)
- $\Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2$ (Compiling source enviroment to circuit environment, secret variables are converted to shares)

Then:

- $\widetilde{\rho} \vdash \widetilde{s} \Longrightarrow \widetilde{\rho}'; \kappa$ (The C++ ABY program)
- $\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa \longmapsto \widehat{\rho}'_1, \widehat{\rho}'_2; O$ (Circuit evaluation)

m	::=		Secret label
	$\mid$	$\mathcal{A}$	
	$\mid$	$\mathcal{B}$	
ℓ	::=		Label
	$\mid$	$\mathcal{P}$	
	$\mid$	m	
σ	::=		Base type
	$\mid$	uint^{ℓ}	
	$\mid$	bool^{ℓ}	
τ	::=		Type
	$\mid$	σ	
	$\mid$	$\sigma[\,]$	
c	::=		Constant
	$\mid$	n	
	$\mid$	$\top$	
	$\mid$	$\bot$	
e	::=		Source expression
	$\mid$	c	
	$\mid$	x	
	$\mid$	$e_1 + e_2$	
	$\mid$	$\mathbf{cond}(e, e_1, e_2)$	
	$\mid$	$e_1 > e_2$	
	$\mid$	$x[e]$	
s	::=		Source statement
	$\mid$	$\tau\, x$	
	$\mid$	$x := e$	
	$\mid$	$\mathbf{for}(x := n_1; x \leq n_2; x := x + 1)\ s$	
	$\mid$	$x[e_1] := e_2$	
	$\mid$	$\mathbf{if}(e, s_1, s_2)$	
	$\mid$	$\mathbf{out}\, e$	
	$\mid$	$s_1; s_2$	

Figure 1: Source language

$$
\begin{array}{lll}
v & ::= & \text{Source value} \\
 & | & c \\
 & | & [\,\overline{c_i}^{\,i}\,] \\
 & & \\
\rho & ::= & \text{Source runtime environment} \\
 & | & \cdot \\
 & | & \rho[x \mapsto v] \\
 & & \\
O & ::= & \text{Source observation} \\
 & | & \cdot \\
 & | & v \\
 & | & O_1, O_2
\end{array}
$$

Figure 2: Source runtime

$\boxed{\rho \vdash e \downarrow v}$

$$\frac{}{\rho \vdash c \downarrow c}\ \text{SE\_CONST}$$

$$\frac{}{\rho \vdash x \downarrow \rho[x]}\ \text{SE\_VAR}$$

$$\frac{\rho \vdash e_i \downarrow n_i}{\rho \vdash e_1 + e_2 \downarrow n_1 + n_2}\ \text{SE\_ADD}$$

$$\frac{\rho \vdash e \downarrow \top \quad \rho \vdash e_1 \downarrow v}{\rho \vdash \mathbf{cond}(e, e_1, e_2) \downarrow v}\ \text{SE\_CONDT}$$

$$\frac{\rho \vdash e \downarrow \bot \quad \rho \vdash e_2 \downarrow v}{\rho \vdash \mathbf{cond}(e, e_1, e_2) \downarrow v}\ \text{SE\_CONDF}$$

$$\frac{\rho \vdash e_i \downarrow n_i}{\rho \vdash e_1 > e_2 \downarrow n_1 > n_2}\ \text{SE\_GT}$$

$$\frac{\rho \vdash e \downarrow n \quad \rho \vdash x \downarrow [\,\overline{c_i}^{\,i}\,]}{\rho \vdash x[e] \downarrow c_n}\ \text{SE\_AREAD}$$

Figure 3: Source expression evaluation

$\boxed{\rho \vdash s \longrightarrow \rho'; O}$

$$\frac{\mathsf{default}(\tau) = v}{\rho \vdash \tau\, x \longrightarrow \rho[x \mapsto v]; \cdot} \quad \text{SC\_DECL}$$

$$\frac{\rho \vdash e \downarrow v}{\rho \vdash x := e \longrightarrow \rho[x \mapsto v]; \cdot} \quad \text{SC\_ASSGN}$$

$$\frac{n_1 > n_2}{\rho \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1)\ s \longrightarrow \rho; \cdot} \quad \text{SC\_FORT}$$

$$\frac{\begin{array}{l} n_2 \geq n_1 \\ \rho[x \mapsto n_1] \vdash s \longrightarrow \rho_1; O_1 \\ \rho_1 \vdash \mathbf{for}(x := n_1 + 1; x \leq n_2; x := x + 1)\ s \longrightarrow \rho_2; O_2 \end{array}}{\rho \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1)\ s \longrightarrow \rho_2; O_1, O_2} \quad \text{SC\_FORI}$$

$$\frac{\begin{array}{c} \rho \vdash x \downarrow [\,\overline{c_i}^{\,i}\,] \\ \rho \vdash e_1 \downarrow n \\ \rho \vdash e_2 \downarrow c \end{array}}{\rho \vdash x[e_1] := e_2 \longrightarrow \rho[x \mapsto [\,\overline{c_i}^{\,i}\,][n \mapsto c]]; O_1, O_2} \quad \text{SC\_AWRITE}$$

$$\frac{\begin{array}{c} \rho \vdash e \downarrow \top \\ \rho \vdash s_1 \longrightarrow \rho'; O \end{array}}{\rho \vdash \mathbf{if}(e, s_1, s_2) \longrightarrow \rho'; O} \quad \text{SC\_IFT}$$

$$\frac{\begin{array}{c} \rho \vdash e \downarrow \bot \\ \rho \vdash s_2 \longrightarrow \rho'; O \end{array}}{\rho \vdash \mathbf{if}(e, s_1, s_2) \longrightarrow \rho'; O} \quad \text{SC\_IFF}$$

$$\frac{\rho \vdash e \downarrow v}{\rho \vdash \mathbf{out}\, e \longrightarrow \rho; v} \quad \text{SC\_OUT}$$

$$\frac{\begin{array}{c} \rho \vdash s_1 \longrightarrow \rho_1; O_1 \\ \rho_1 \vdash s_2 \longrightarrow \rho_2; O_2 \end{array}}{\rho \vdash s_1; s_2 \longrightarrow \rho_2; O_1, O_2} \quad \text{SC\_SEQ}$$

Figure 4: Source command evaluation

$\boxed{v : \tau}$

$$\frac{}{n : \mathsf{uint}^{\mathcal{P}}} \text{ V\_INT}$$

$$\frac{}{\top : \mathsf{bool}^{\mathcal{P}}} \text{ V\_TRUE}$$

$$\frac{}{\bot : \mathsf{bool}^{\mathcal{P}}} \text{ V\_FALSE}$$

$$\frac{c_i : \sigma}{[\,\overline{c_i}^{\,i}\,] : \sigma[\,]} \text{ V\_ARR}$$

Figure 5: Value typing

$\widetilde{e}$	::=		Target expression
	|	c	
	|	x	
	|	$\widetilde{e}_1 +_\ell \widetilde{e}_2$	
	|	$\mathbf{cond}_\ell(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)$	
	|	$\widetilde{e}_1 >_\ell \widetilde{e}_2$	
	|	$x[\widetilde{e}]$	
	|	$\widetilde{e} \rhd m$	
$\widetilde{s}$	::=		Target statement
	|	$\tau\, x$	
	|	$x := \widetilde{e}$	
	|	$\mathbf{for}(x := n_1; x \leq n_2; x := x + 1)\ \widetilde{s}$	
	|	$x[\widetilde{e}_1] := \widetilde{e}_2$	
	|	$\mathbf{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2)$	
	|	$\mathbf{out}\, \widetilde{e}$	
	|	$\widetilde{s}_1; \widetilde{s}_2$	
Γ	::=		Type environment
	|	$\cdot$	
	|	$\Gamma, x : \tau$	

Figure 6: Target language

$\boxed{\Gamma \vdash e : \tau \rightsquigarrow \widetilde{e}}$

$$\frac{}{\Gamma \vdash n : \mathsf{uint}^{\mathcal{P}} \rightsquigarrow n} \text{ S\_CONST}$$

$$\frac{}{\Gamma \vdash \top : \mathsf{bool}^{\mathcal{P}} \rightsquigarrow \top} \text{ S\_TRUE}$$

$$\frac{}{\Gamma \vdash \bot : \mathsf{bool}^{\mathcal{P}} \rightsquigarrow \bot} \text{ S\_FALSE}$$

$$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau \rightsquigarrow x} \text{ S\_VAR}$$

$$\frac{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \rightsquigarrow \widetilde{e}_i}{\Gamma \vdash e_1 + e_2 : \mathsf{uint}^{\mathcal{P}} \rightsquigarrow \widetilde{e}_1 +_{\mathcal{P}} \widetilde{e}_2} \text{ S\_PADD}$$

$$\frac{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{A}} \rightsquigarrow \widetilde{e}_i}{\Gamma \vdash e_1 + e_2 : \mathsf{uint}^{\mathcal{A}} \rightsquigarrow \widetilde{e}_1 +_{\mathcal{A}} \widetilde{e}_2} \text{ S\_SADD}$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \mathsf{bool}^{\mathcal{P}} \rightsquigarrow \widetilde{e} \\ \Gamma \vdash e_i : \sigma \rightsquigarrow \widetilde{e}_i\end{array}}{\Gamma \vdash \mathbf{cond}(e, e_1, e_2) : \sigma \rightsquigarrow \mathbf{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)} \text{ S\_PCOND}$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \mathsf{bool}^{\mathcal{B}} \rightsquigarrow \widetilde{e} \\ \Gamma \vdash e_i : \sigma \rightsquigarrow \widetilde{e}_i \\ \mathsf{label}(\sigma) = \mathcal{B}\end{array}}{\Gamma \vdash \mathbf{cond}(e, e_1, e_2) : \sigma \rightsquigarrow \mathbf{cond}_{\mathcal{B}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2)} \text{ S\_SCOND}$$

$$\frac{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{P}} \rightsquigarrow \widetilde{e}_i}{\Gamma \vdash e_1 > e_2 : \mathsf{bool}^{\mathcal{P}} \rightsquigarrow \widetilde{e}_1 >_{\mathcal{P}} \widetilde{e}_2} \text{ S\_PGT}$$

$$\frac{\Gamma \vdash e_i : \mathsf{uint}^{\mathcal{B}} \rightsquigarrow \widetilde{e}_i}{\Gamma \vdash e_1 > e_2 : \mathsf{bool}^{\mathcal{B}} \rightsquigarrow \widetilde{e}_1 >_{\mathcal{B}} \widetilde{e}_2} \text{ S\_SGT}$$

$$\frac{\begin{array}{c}\Gamma \vdash x : \sigma[\,] \rightsquigarrow x \\ \Gamma \vdash e : \mathsf{uint}^{\mathcal{P}} \rightsquigarrow \widetilde{e}\end{array}}{\Gamma \vdash x[e] : \sigma \rightsquigarrow x[\widetilde{e}]} \text{ S\_AREAD}$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \sigma_1 \rightsquigarrow \widetilde{e} \\ \mathsf{base}(\sigma_1) = \mathsf{base}(\sigma_2) \\ \mathsf{label}(\sigma_2) = m\end{array}}{\Gamma \vdash e : \sigma_2 \rightsquigarrow \widetilde{e} \rhd m} \text{ S\_SUB}$$

Figure 7: Expression compilation

$\boxed{\Gamma \vdash s \rightsquigarrow \widetilde{s} \mid \Gamma'}$

$$\frac{}{\Gamma \vdash \tau\, x \rightsquigarrow \tau\, x \mid \Gamma, x : \tau} \quad \text{C\_DECL}$$

$$\frac{\begin{array}{c}\Gamma(x) = \sigma \\ \Gamma \vdash e : \sigma \rightsquigarrow \widetilde{e}\end{array}}{\Gamma \vdash x := e \rightsquigarrow x := \widetilde{e} \mid \Gamma} \quad \text{C\_VASSGN}$$

$$\frac{\begin{array}{c}\Gamma, x : \mathsf{uint}^{\mathcal{P}} \vdash s \rightsquigarrow \widetilde{s} \mid \_ \\ x \notin \mathsf{modifies}(s)\end{array}}{\Gamma \vdash \mathbf{for}(x := n_1; x \leq n_2; x := x + 1)\ s \rightsquigarrow \mathbf{for}(x := n_1; x \leq n_2; x := x + 1)\ \widetilde{s} \mid \Gamma} \quad \text{C\_FOR}$$

$$\frac{\begin{array}{c}\Gamma \vdash x : \sigma[\,] \rightsquigarrow x \\ \Gamma \vdash e_1 : \mathsf{uint}^{\mathcal{P}} \rightsquigarrow \widetilde{e}_1 \\ \Gamma \vdash e_2 : \sigma \rightsquigarrow \widetilde{e}_2\end{array}}{\Gamma \vdash x[e_1] := e_2 \rightsquigarrow x[\widetilde{e}_1] := \widetilde{e}_2 \mid \Gamma} \quad \text{C\_AWRITE}$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \mathsf{bool}^{\mathcal{P}} \rightsquigarrow \widetilde{e} \\ \Gamma \vdash s_i \rightsquigarrow \widetilde{s}_i \mid \_\end{array}}{\Gamma \vdash \mathbf{if}(e, s_1, s_2) \rightsquigarrow \mathbf{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \mid \Gamma} \quad \text{C\_IF}$$

$$\frac{\begin{array}{c}\Gamma \vdash e : \sigma \rightsquigarrow \widetilde{e} \\ \mathsf{label}(\sigma) = m\end{array}}{\Gamma \vdash \mathbf{out}\, e \rightsquigarrow \mathbf{out}\, \widetilde{e} \mid \Gamma} \quad \text{C\_OUT}$$

$$\frac{\begin{array}{c}\Gamma \vdash s_1 \rightsquigarrow \widetilde{s}_1 \mid \Gamma_1 \\ \Gamma_1 \vdash s_2 \rightsquigarrow \widetilde{s}_2 \mid \Gamma'\end{array}}{\Gamma \vdash s_1; s_2 \rightsquigarrow \widetilde{s}_1; \widetilde{s}_2 \mid \Gamma'} \quad \text{C\_SEQ}$$

Figure 8: Command compilation

r	::=		Wire id range
$\widetilde{w}$	::=		Compiled base value
	|	c	
	|	r	
$\widetilde{v}$	::=		Compiled value
	|	$\widetilde{w}$	
	|	$[\overline{\widetilde{w}_i}^{\,i}]$	
κ	::=		Circuit
	|	$\cdot$	
	|	$\oplus(r_1, r_2, r_3)$	
	|	$\mathsf{Mux}(r_1, r_2, r_3, r_4)$	
	|	$\mathsf{Gt}(r_1, r_2, r_3)$	
	|	$r_1 \rhd_m r_2$	
	|	$\mathsf{Out}(r)$	
	|	κ_1, κ_2	
$\widetilde{\rho}$	::=		Runtime environment
	|	$\cdot$	
	|	$\widetilde{\rho}[x \mapsto \widetilde{v}]$	

Figure 9: Target runtime

$\boxed{\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa}$

$$\frac{}{\widetilde{\rho} \vdash c \Downarrow c; \cdot} \quad \text{EE\_CONST}$$

$$\frac{}{\widetilde{\rho} \vdash x \Downarrow \widetilde{\rho}[x]; \cdot} \quad \text{EE\_VAR}$$

$$\frac{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}{\widetilde{\rho} \vdash \widetilde{e}_1 +_{\mathcal{P}} \widetilde{e}_2 \Downarrow n_1 + n_2; \kappa_1, \kappa_2} \quad \text{EE\_PADD}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow r_i; \kappa_i \\ r_3 = \mathsf{next\_range}()\end{array}}{\widetilde{\rho} \vdash \widetilde{e}_1 +_{\mathcal{A}} \widetilde{e}_2 \Downarrow r_3; \kappa_1, \kappa_2, \oplus(r_1, r_2, r_3)} \quad \text{EE\_SADD}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e} \Downarrow \top; \kappa \\ \widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow \widetilde{v}; \kappa_1\end{array}}{\widetilde{\rho} \vdash \mathbf{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow \widetilde{v}; \kappa, \kappa_1} \quad \text{EE\_PCONDT}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e} \Downarrow \bot; \kappa \\ \widetilde{\rho} \vdash \widetilde{e}_2 \Downarrow \widetilde{v}; \kappa_2\end{array}}{\widetilde{\rho} \vdash \mathbf{cond}_{\mathcal{P}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow \widetilde{v}; \kappa, \kappa_2} \quad \text{EE\_PCONDF}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e} \Downarrow r; \kappa \\ \widetilde{\rho} \vdash \widetilde{e}_i \Downarrow r_i; \kappa_i \\ r_3 = \mathsf{next\_range}()\end{array}}{\widetilde{\rho} \vdash \mathbf{cond}_{\mathcal{B}}(\widetilde{e}, \widetilde{e}_1, \widetilde{e}_2) \Downarrow r_3; \kappa, \kappa_1, \kappa_2, \mathsf{Mux}(r, r_1, r_2, r_3)} \quad \text{EE\_SCOND}$$

$$\frac{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}{\widetilde{\rho} \vdash \widetilde{e}_1 >_{\mathcal{P}} \widetilde{e}_2 \Downarrow n_1 > n_2; \kappa_1, \kappa_2} \quad \text{EE\_PGT}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow r_i; \kappa_i \\ r_3 = \mathsf{next\_range}()\end{array}}{\widetilde{\rho} \vdash \widetilde{e}_1 >_{\mathcal{B}} \widetilde{e}_2 \Downarrow r_3; \kappa_1, \kappa_2, \mathsf{Gt}(r_1, r_2, r_3)} \quad \text{EE\_SGT}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e} \Downarrow n; \kappa_1 \\ \widetilde{\rho} \vdash x \Downarrow [\overline{\widetilde{w}_i}^{\,i}]; \kappa_2\end{array}}{\widetilde{\rho} \vdash x[\widetilde{e}] \Downarrow \widetilde{w}_n; \kappa_1, \kappa_2} \quad \text{EE\_AREAD}$$

$$\frac{\begin{array}{c}\widetilde{\rho} \vdash \widetilde{e} \Downarrow r; \kappa \\ r' = \mathsf{next\_range}()\end{array}}{\widetilde{\rho} \vdash \widetilde{e} \rhd m \Downarrow r'; \kappa, r \rhd_m r'} \quad \text{EE\_COERCE}$$

Figure 10: Target expression evaluation

$\boxed{\widetilde{\rho} \vdash \widetilde{s} \Longrightarrow \widetilde{\rho}'; \kappa}$

$$\frac{\mathsf{default}(\tau) = \widetilde{v}; \kappa}{\widetilde{\rho} \vdash \tau\, x \Longrightarrow \widetilde{\rho}[x \mapsto \widetilde{v}]; \kappa} \quad \text{EC\_DECL}$$

$$\frac{\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa}{\widetilde{\rho} \vdash x := \widetilde{e} \Longrightarrow \widetilde{\rho}[x \mapsto \widetilde{v}]; \kappa} \quad \text{EC\_ASSGN}$$

$$\frac{n_1 > n_2}{\widetilde{\rho} \vdash \textbf{for}(x := n_1; x \leq n_2; x := x + 1)\ \widetilde{s} \Longrightarrow \widetilde{\rho}; \cdot} \quad \text{EC\_FORT}$$

$$\frac{\begin{array}{l} n_2 \geq n_1 \\ \widetilde{\rho}[x \mapsto n_1] \vdash \widetilde{s} \Longrightarrow \widetilde{\rho}_1; \kappa_1 \\ \widetilde{\rho}_1 \vdash \textbf{for}(x := n_1 + 1; x \leq n_2; x := x + 1)\ \widetilde{s} \Longrightarrow \widetilde{\rho}_2; \kappa_2 \end{array}}{\widetilde{\rho} \vdash \textbf{for}(x := n_1; x \leq n_2; x := x + 1)\ \widetilde{s} \Longrightarrow \widetilde{\rho}_2; \kappa_1, \kappa_2} \quad \text{EC\_FORI}$$

$$\frac{\begin{array}{c} \widetilde{\rho} \vdash x \Downarrow [\,\overline{\widetilde{w}_i}^{\,i}\,]; \cdot \\ \widetilde{\rho} \vdash \widetilde{e}_1 \Downarrow n; \kappa_1 \\ \widetilde{\rho} \vdash \widetilde{e}_2 \Downarrow \widetilde{w}; \kappa_2 \end{array}}{\widetilde{\rho} \vdash x[\widetilde{e}_1] := \widetilde{e}_2 \Longrightarrow \widetilde{\rho}[x \mapsto [\,\overline{\widetilde{w}_i}^{\,i}\,][n \mapsto \widetilde{w}]]; \kappa_1, \kappa_2} \quad \text{EC\_AWRITE}$$

$$\frac{\begin{array}{c} \widetilde{\rho} \vdash \widetilde{e} \Downarrow \top; \kappa_1 \\ \widetilde{\rho} \vdash \widetilde{s}_1 \Longrightarrow \widetilde{\rho}'; \kappa_2 \end{array}}{\widetilde{\rho} \vdash \textbf{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \Longrightarrow \widetilde{\rho}'; \kappa_1, \kappa_2} \quad \text{EC\_IFT}$$

$$\frac{\begin{array}{c} \widetilde{\rho} \vdash \widetilde{e} \Downarrow \bot; \kappa_1 \\ \widetilde{\rho} \vdash \widetilde{s}_2 \Longrightarrow \widetilde{\rho}'; \kappa_2 \end{array}}{\widetilde{\rho} \vdash \textbf{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \Longrightarrow \widetilde{\rho}'; \kappa_1, \kappa_2} \quad \text{EC\_IFF}$$

$$\frac{\widetilde{\rho} \vdash \widetilde{e} \Downarrow r; \kappa}{\widetilde{\rho} \vdash \textbf{out}\, \widetilde{e} \Longrightarrow \widetilde{\rho}; \kappa, \mathsf{Out}(r)} \quad \text{EC\_OUT}$$

$$\frac{\begin{array}{c} \widetilde{\rho} \vdash \widetilde{s}_1 \Longrightarrow \widetilde{\rho}_1; \kappa_1 \\ \widetilde{\rho}_1 \vdash \widetilde{s}_2 \Longrightarrow \widetilde{\rho}_2; \kappa_2 \end{array}}{\widetilde{\rho} \vdash \widetilde{s}_1; \widetilde{s}_2 \Longrightarrow \widetilde{\rho}_2; \kappa_1, \kappa_2} \quad \text{EC\_SEQ}$$

Figure 11: Target command evaluation

$$
\begin{array}{lcll}
\widehat{c} & ::= & & \text{Circuit value (share)} \\
 & | & \mathcal{E}_1^m(c) & \\
 & | & \mathcal{E}_2^m(c) & \\
 & & & \\
\widehat{\rho} & ::= & & \text{Circuit environment} \\
 & | & \cdot & \\
 & | & \widehat{\rho}[r \mapsto \widehat{c}] &
\end{array}
$$

Figure 12: Circuit runtime

$\boxed{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa \longmapsto \widehat{\rho}'_1, \widehat{\rho}'_2; O}$

$$\frac{}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \cdot \longmapsto \widehat{\rho}_1, \widehat{\rho}_2; \cdot} \quad \text{CKT\_EMP}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1[r_1] = \mathcal{E}_1^{\mathcal{A}}(n_1) \\ \widehat{\rho}_2[r_1] = \mathcal{E}_2^{\mathcal{A}}(n_1) \\ \widehat{\rho}_1[r_2] = \mathcal{E}_1^{\mathcal{A}}(n_2) \\ \widehat{\rho}_2[r_2] = \mathcal{E}_2^{\mathcal{A}}(n_2)\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \oplus(r_1, r_2, r_3) \longmapsto \widehat{\rho}_1[r_3 \mapsto \mathcal{E}_1^{\mathcal{A}}((n_1 + n_2))], \widehat{\rho}_2[r_3 \mapsto \mathcal{E}_2^{\mathcal{A}}((n_1 + n_2))]; \cdot} \quad \text{CKT\_ADD}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1[r_1] = \mathcal{E}_1^{\mathcal{B}}(\top) \\ \widehat{\rho}_2[r_1] = \mathcal{E}_2^{\mathcal{B}}(\top) \\ \widehat{\rho}_1[r_2] = \mathcal{E}_1^{\mathcal{B}}(c) \\ \widehat{\rho}_2[r_2] = \mathcal{E}_2^{\mathcal{B}}(c)\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \mathsf{Mux}(r_1, r_2, r_3, r_4) \longmapsto \widehat{\rho}_1[r_4 \mapsto \mathcal{E}_1^{\mathcal{B}}(c)], \widehat{\rho}_2[r_4 \mapsto \mathcal{E}_2^{\mathcal{B}}(c)]; \cdot} \quad \text{CKT\_MUXT}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1[r_1] = \mathcal{E}_1^{\mathcal{B}}(\bot) \\ \widehat{\rho}_2[r_1] = \mathcal{E}_2^{\mathcal{B}}(\bot) \\ \widehat{\rho}_1[r_3] = \mathcal{E}_1^{\mathcal{B}}(c) \\ \widehat{\rho}_2[r_3] = \mathcal{E}_2^{\mathcal{B}}(c)\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \mathsf{Mux}(r_1, r_2, r_3, r_4) \longmapsto \widehat{\rho}_1[r_4 \mapsto \mathcal{E}_1^{\mathcal{B}}(c)], \widehat{\rho}_2[r_4 \mapsto \mathcal{E}_2^{\mathcal{B}}(c)]; \cdot} \quad \text{CKT\_MUXF}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1[r_1] = \mathcal{E}_1^{\mathcal{B}}(n_1) \\ \widehat{\rho}_2[r_1] = \mathcal{E}_2^{\mathcal{B}}(n_1) \\ \widehat{\rho}_1[r_2] = \mathcal{E}_1^{\mathcal{B}}(n_2) \\ \widehat{\rho}_2[r_2] = \mathcal{E}_2^{\mathcal{B}}(n_2) \\ c = n_1 > n_2\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \mathsf{Gt}(r_1, r_2, r_3) \longmapsto \widehat{\rho}_1[r_3 \mapsto \mathcal{E}_1^{\mathcal{B}}(c)], \widehat{\rho}_2[r_3 \mapsto \mathcal{E}_2^{\mathcal{B}}(c)]; \cdot} \quad \text{CKT\_GT}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1[r_1] = \mathcal{E}_1^{m_1}(c) \\ \widehat{\rho}_2[r_1] = \mathcal{E}_2^{m_1}(c)\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash r_1 \rhd_m r_2 \longmapsto \widehat{\rho}_1[r_2 \mapsto \mathcal{E}_1^{m}(c)], \widehat{\rho}_2[r_2 \mapsto \mathcal{E}_2^{m}(c)]; \cdot} \quad \text{CKT\_COERCE}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1[r] = \mathcal{E}_1^{m}(c) \\ \widehat{\rho}_2[r] = \mathcal{E}_2^{m}(c)\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \mathsf{Out}(r) \longmapsto \widehat{\rho}_1, \widehat{\rho}_2; c} \quad \text{CKT\_OUT}$$

$$\frac{\begin{array}{c}\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa_1 \longmapsto \widehat{\rho}'_1, \widehat{\rho}'_2; O_1 \\ \widehat{\rho}'_1, \widehat{\rho}'_2 \vdash \kappa_2 \longmapsto \widehat{\rho}''_1, \widehat{\rho}''_2; O_2\end{array}}{\widehat{\rho}_1, \widehat{\rho}_2 \vdash \kappa_1, \kappa_2 \longmapsto \widehat{\rho}''_1, \widehat{\rho}''_2; O_1, O_2} \quad \text{CKT\_SEQ}$$

Figure 13: Circuit evaluation

$\boxed{\Gamma \sim \rho}$

$$\frac{}{\cdot \sim \cdot}\ \text{SEN\_EMP}$$

$$\frac{\begin{array}{c} v : [\tau]_{\mathcal{P}} \\ \Gamma \sim \rho \end{array}}{\Gamma, x : \tau \sim \rho[x \mapsto v]}\ \text{SEN\_BND}$$

Figure 14: Source environment and type environment consistency

$\boxed{\Gamma \vdash \rho \hookrightarrow \widetilde{\rho}}$

$$\frac{}{\Gamma \vdash \cdot \hookrightarrow \cdot}\ \text{TEN\_EMP}$$

$$\frac{\begin{array}{c} \Gamma(x) = \sigma \\ \mathsf{label}(\sigma) = \mathcal{P} \\ \Gamma \vdash \rho \hookrightarrow \widetilde{\rho} \end{array}}{\Gamma \vdash \rho[x \mapsto c] \hookrightarrow \widetilde{\rho}[x \mapsto c]}\ \text{TEN\_PBT}$$

$$\frac{\begin{array}{c} \Gamma(x) = \sigma \\ \mathsf{label}(\sigma) = m \\ r = \mathsf{next\_range}() \\ \Gamma \vdash \rho \hookrightarrow \widetilde{\rho} \end{array}}{\Gamma \vdash \rho[x \mapsto c] \hookrightarrow \widetilde{\rho}[x \mapsto r]}\ \text{TEN\_SBT}$$

$$\frac{\begin{array}{c} \Gamma(x) = \sigma[\,] \\ \mathsf{label}(\sigma) = \mathcal{P} \\ \Gamma \vdash \rho \hookrightarrow \widetilde{\rho} \end{array}}{\Gamma \vdash \rho[x \mapsto [\,\overline{c_i}^{\,i}\,]] \hookrightarrow \widetilde{\rho}[x \mapsto [\,\overline{c_i}^{\,i}\,]]}\ \text{TEN\_PARR}$$

$$\frac{\begin{array}{c} \Gamma(x) = \sigma[\,] \\ \mathsf{label}(\sigma) = m \\ r_i = \mathsf{next\_range}() \\ \Gamma \vdash \rho \hookrightarrow \widetilde{\rho} \end{array}}{\Gamma \vdash \rho[x \mapsto [\,\overline{c_i}^{\,i}\,]] \hookrightarrow \widetilde{\rho}[x \mapsto [\,\overline{r_i}^{\,i}\,]]}\ \text{TEN\_SARR}$$

Figure 15: Source environment to target environment compilation

$\boxed{\Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2}$

$$\frac{}{\Gamma; \widetilde{\rho} \vdash \cdot \hookrightarrow \cdot, \cdot} \text{ CEN\_EMP}$$

$$\frac{\begin{array}{c}\Gamma(x) = \sigma \\ \mathsf{label}(\sigma) = \mathcal{P} \\ \Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2\end{array}}{\Gamma; \widetilde{\rho} \vdash \rho[x \mapsto c] \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2} \text{ CEN\_PBT}$$

$$\frac{\begin{array}{c}\Gamma(x) = \sigma \\ \mathsf{label}(\sigma) = m \\ \widetilde{\rho}[x] = r \\ \Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2\end{array}}{\Gamma; \widetilde{\rho} \vdash \rho[x \mapsto c] \hookrightarrow \widehat{\rho}_1[r \mapsto \mathcal{E}_1^m(c)], \widehat{\rho}_2[r \mapsto \mathcal{E}_2^m(c)]} \text{ CEN\_SBT}$$

$$\frac{\begin{array}{c}\Gamma(x) = \sigma[\,] \\ \mathsf{label}(\sigma) = \mathcal{P} \\ \Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2\end{array}}{\Gamma; \widetilde{\rho} \vdash \rho[x \mapsto [\,\overline{c_i}^{\,i}\,]] \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2} \text{ CEN\_PARR}$$

$$\frac{\begin{array}{c}\Gamma(x) = \sigma[\,] \\ \mathsf{label}(\sigma) = m \\ \widetilde{\rho}[x] = [\,\overline{r_i}^{\,i}\,] \\ \Gamma; \widetilde{\rho} \vdash \rho \hookrightarrow \widehat{\rho}_1, \widehat{\rho}_2\end{array}}{\Gamma; \widetilde{\rho} \vdash \rho[x \mapsto [\,\overline{c_i}^{\,i}\,]] \hookrightarrow \widehat{\rho}_1[r_i \mapsto \mathcal{E}_1^m(c_i)], \widehat{\rho}_2[r_i \mapsto \mathcal{E}_2^m(c_i)]} \text{ CEN\_SARR}$$

Figure 16: Source environment to circuit environment compilation