```
Base type
\sigma ::=
          uint
         bool
\psi ::=
                                         Source type
         \sigma[n]
                                         Constant
c ::=
e ::=
                                         Source expression
        e_1 + e_2
         \mathbf{cond}(e, e_1, e_2)
        e_1 > e_2
       x[e]
     | \quad \mathbf{input}_i^{\sigma}
                                         Source statement
s ::=
        \psi x
         x := e
         for x in n_1 ... n_2 \{s\}
         x[e_1] := e_2
         if(e, s_1, s_2)
         \mathbf{out}\, e
         s_1; s_2
```

Figure 1: Source language

Figure 2: Source runtime

$$\rho \vdash e \downarrow v$$

Figure 3: Source expression evaluation

$$\frac{\operatorname{default}(\psi) = v}{\rho \vdash \psi \ x \downarrow \rho[x \mapsto v];} \quad \operatorname{SC_DECL}$$

$$\frac{\rho \vdash e \downarrow v}{\rho \vdash x := e \downarrow \rho[x \mapsto v];} \quad \operatorname{SC_ASSGN}$$

$$\frac{\rho[x \mapsto n_1] \vdash \operatorname{loop} x \ \operatorname{until} \ n_2 \left\{s\right\} \downarrow \rho_1; O}{\rho \vdash \operatorname{for} x \ \operatorname{in} \ n_1 \dots n_2 \left\{s\right\} \downarrow \rho_1 \setminus x; O} \quad \operatorname{SC_FORT}$$

$$\frac{\rho(x) > n_2}{\rho \vdash \operatorname{loop} x \ \operatorname{until} \ n_2 \left\{s\right\} \downarrow \rho; O} \quad \operatorname{SC_LOOPT}$$

$$\frac{\rho(x) \leq n_2}{\rho \vdash s \downarrow \rho_1; O_1}$$

$$\frac{\rho(p_1|\operatorname{dom}(\rho))[x \mapsto \rho(x) + 1] \vdash \operatorname{loop} x \ \operatorname{until} \ n_2 \left\{s\right\} \downarrow \rho_2; O_2}{\rho \vdash \operatorname{loop} x \ \operatorname{until} \ n_2 \left\{s\right\} \downarrow \rho_2; O_1, O_2} \quad \operatorname{SC_LOOPI}$$

$$\frac{\rho \vdash x \downarrow [c_i]_{n_2}}{\rho \vdash e_1 \downarrow n_1}$$

$$\frac{\rho \vdash e_1 \downarrow n_1}{\rho \vdash e_2 \downarrow c}$$

$$\frac{\rho \vdash x \downarrow [c_i]_{n_2} [n_1 \mapsto c]];}{\rho \vdash x [e_1] := e_2 \downarrow \rho[x \mapsto [c_i]_{n_2} [n_1 \mapsto c]];} \quad \operatorname{SC_AWRITE}$$

$$\frac{\rho \vdash e \downarrow c}{c = \top \Rightarrow i = 1}$$

$$c = \bot \Rightarrow i = 2$$

$$\frac{\rho \vdash s_1 \downarrow \rho_1; O}{\rho \vdash \operatorname{if}(e, s_1, s_2) \downarrow \rho'; O} \quad \operatorname{SC_IF}$$

$$\frac{\rho \vdash e \downarrow v}{\rho \vdash \operatorname{out} e \downarrow \rho; v,} \quad \operatorname{SC_OUT}$$

$$\frac{\rho \vdash s_1 \downarrow \rho_1; O_1}{\rho \vdash s_2 \downarrow \rho_2; O_2} \quad \operatorname{SC_SEQ}$$

Figure 4: Source command evaluation

 $v:\psi$

$$\frac{c:\delta(c)}{c:\sigma} \quad \text{V_CONS}$$

$$\frac{c_i:\sigma}{[c_i]_n:\sigma[n]} \quad \text{V_ARR}$$

Figure 5: Value typing

```
Secret label
m ::=
         | A
| B
\ell ::=
                                                                                                              Label
                   \mathcal{P}
                   m
                                                                                                              Type
\tau \quad ::=
         \mid \quad \sigma^\ell
          \sigma^{\ell}[n]
\widetilde{e} ::=
                                                                                                              Target expression
                  c
            |\widetilde{e}_1 +_{\ell} \widetilde{e}_2|
                   \operatorname{\mathbf{cond}}_{\ell}(\widetilde{e},\widetilde{e}_1,\widetilde{e}_2)
                   \widetilde{e}_1 >_{\ell} \widetilde{e}_2
                   x[\widetilde{e}]
                   \mathbf{input}_i^{(\sigma^m)}
                    \tilde{e} \triangleright m
                    [\widetilde{e}_i]_n
\widetilde{s} ::=
                                                                                                              Target statement
                   \tau\,x=\widetilde{e}
                   for(x := n_1; x \le n_2; x := x + 1) \tilde{s}
                   x[\widetilde{e}_1] := \widetilde{e}_2
                   \mathbf{if}(\widetilde{e},\widetilde{s}_1,\widetilde{s}_2)
                   \mathbf{out}\,\widetilde{e}
                   \widetilde{s}_1; \widetilde{s}_2
\Gamma ::=
                                                                                                              Type environment
         \begin{array}{c|c} | & \cdot \\ & \Gamma, x : \tau \end{array}
```

Figure 6: Target language

$\Gamma \vdash e : \tau \leadsto \widetilde{e}$

$$\begin{array}{c} \overline{\Gamma \vdash c : \delta(c)^{\mathcal{P}} \leadsto c} & \text{S_CONS} \\ \\ \frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau \leadsto x} & \text{S_VAR} \\ \\ \frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i}{\Gamma \vdash e_1 + e_2 : \text{uint}^{\mathcal{A}} \leadsto \widetilde{e}_1 +_{\mathcal{P}} \widetilde{e}_2} & \text{S_PADD} \\ \\ \frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{A}} \leadsto \widetilde{e}_i}{\Gamma \vdash e_1 + e_2 : \text{uint}^{\mathcal{A}} \leadsto \widetilde{e}_1 +_{\mathcal{A}} \widetilde{e}_2} & \text{S_SADD} \\ \\ \Gamma \vdash e_i : bool^{\mathcal{P}} \leadsto \widetilde{e}_i \\ \hline{\Gamma \vdash e_i : o^{\ell} \leadsto \widetilde{e}_i} & \text{S_PCOND} \\ \\ \Gamma \vdash e : bool^{\mathcal{B}} \leadsto \widetilde{e}_i \\ \hline{\Gamma \vdash e_i : o^{\mathcal{B}} \leadsto \widetilde{e}_i} & \text{S_SCOND} \\ \\ \frac{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i}{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_PGT} \\ \hline{\Gamma \vdash e_i : \text{uint}^{\mathcal{B}} \leadsto \widetilde{e}_i} & \text{S_PGT} \\ \hline{\Gamma \vdash e_i : \text{uint}^{\mathcal{B}} \leadsto \widetilde{e}_i} & \text{S_SGT} \\ \hline{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_SGT} \\ \hline{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_SGT} \\ \hline{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_SGT} \\ \hline{\Gamma \vdash e_i : \text{uint}^{\mathcal{P}} \leadsto \widetilde{e}_i} & \text{S_AREAD} \\ \hline{\Gamma \vdash input}_i^{\sigma} : \sigma^m \leadsto input_i^{(\sigma^m)}} & \text{S_INP} \\ \hline{\Gamma \vdash e : \sigma^m \leadsto \widetilde{e} \bowtie m} & \text{S_SUB} \\ \hline \end{array}$$

Figure 7: Expression compilation

$$\begin{array}{c} \varphi = \sigma \Rightarrow \tau = \sigma^{\ell} \\ \psi = \sigma[n] \Rightarrow \tau = \sigma^{\ell}[n] \\ \widetilde{e} = \operatorname{default}(\tau) \\ \overline{\Gamma \vdash \psi x \leadsto \tau x = \widetilde{e} \mid \Gamma, x : \tau} \end{array} \quad \text{C_DECL} \\ \Gamma(x) = \sigma^{\ell} \\ \overline{\Gamma \vdash e : \sigma^{\ell} \leadsto \widetilde{e}} \\ \overline{\Gamma \vdash x := e \leadsto x := \widetilde{e} \mid \Gamma} \quad \text{C_VASSGN} \\ \hline \Gamma, x : \operatorname{uint}^{\mathcal{P}} \vdash s \leadsto \widetilde{s} \mid - \\ x \not\in \operatorname{modifies}(s) \\ \hline \Gamma \vdash \operatorname{for} x \text{ in } n_1 \ldots n_2 \left\{ s \right\} \leadsto \operatorname{for}(x := n_1; x \leq n_2; x := x + 1) \ \widetilde{s} \mid \Gamma \\ \hline \Gamma \vdash x : \sigma^{\ell}[n] \leadsto x \\ \Gamma \vdash e_1 : \operatorname{uint}^{\mathcal{P}} \leadsto \widetilde{e}_1 \\ \Gamma \vdash e_2 : \sigma^{\ell} \leadsto \widetilde{e}_2 \\ \hline \Gamma \models e_1 < n \\ \hline \Gamma \vdash x[e_1] := e_2 \leadsto x[\widetilde{e}_1] := \widetilde{e}_2 \mid \Gamma \\ \hline \Gamma \vdash e : \operatorname{bool}^{\mathcal{P}} \leadsto \widetilde{e} \\ \Gamma \vdash s_i \leadsto \widetilde{s}_i \mid - \\ \hline \Gamma \vdash \operatorname{if}(e, s_1, s_2) \leadsto \operatorname{if}(\widetilde{e}, \widetilde{s}_1, \widetilde{s}_2) \mid \Gamma \\ \hline \Gamma \vdash \operatorname{e} : \sigma^m \leadsto \widetilde{e} \\ \hline \Gamma \vdash \operatorname{out} e \leadsto \operatorname{out} \widetilde{e} \mid \Gamma \\ \hline \Gamma \vdash s_1 \leadsto \widetilde{s}_1 \mid \Gamma_1 \\ \Gamma \vdash s_2 \leadsto \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_2 \leadsto \widetilde{s}_1 \colon \widetilde{s}_2 \mid \Gamma_2 \\ \hline \Gamma \vdash s_1 \colon s_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \\ \hline \Gamma \vdash s_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \\ \hline \Gamma \vdash s_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \\ \hline \Gamma \vdash s_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \\ \hline \Gamma \vdash s_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \mapsto \widetilde{s}_1 \\ \hline \Gamma \vdash s_1 \mapsto \widetilde{s}_1 \mapsto$$

Figure 8: Command compilation

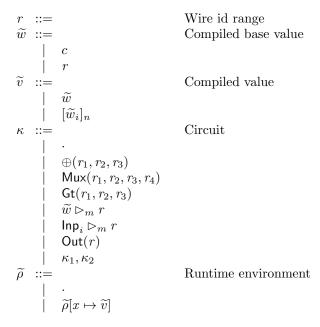


Figure 9: Target runtime

$$\overline{\widetilde{\rho} \vdash \widetilde{e} \Downarrow \widetilde{v}; \kappa}$$
 EE_CONST
$$\overline{\widetilde{\rho} \vdash c \Downarrow c; \cdot}$$
 EE_VAR
$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}$$
 EE_PADD
$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}$$
 EE_PADD
$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}$$
 EE_SADD
$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow n_i; \kappa_i}$$
 Ta = next_range()
$$\overline{\widetilde{\rho} \vdash \widetilde{e}_i + A} \ \widetilde{e}_i \Downarrow r_i; \kappa_i, \kappa_i + C_i + C_i, r_i + C_i + C_i + C_i, \kappa_i + C_i + C$$

Figure 10: Target expression evaluation

 $\frac{\widetilde{\rho} \vdash \widetilde{e}_i \Downarrow \widetilde{w}_i; \kappa_i}{\widetilde{\rho} \vdash [\widetilde{e}_i]_n \Downarrow [\widetilde{w}_i]_n; \kappa_i} \quad \text{EE_ARR}$

Figure 11: Target command evaluation

$$\begin{array}{ll} b & ::= & \text{Share (byte string)} \\ \widehat{\rho} & ::= & \text{Circuit environment} \\ & | & \widehat{\rho}[r \mapsto b] \end{array}$$

Figure 12: Circuit runtime

$$\begin{split} \widehat{\rho_1}, \widehat{\rho_2} \vdash \kappa \longmapsto \widehat{\rho_1}, \widehat{\rho_2}; O \\ \\ & \widehat{\rho_1}, \widehat{\rho_2} \vdash \cdot \longmapsto \widehat{\rho_1}, \widehat{\rho_2}; \\ & n_1 = \mathcal{D}_A(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_1]) \\ & n_2 = \mathcal{D}_A(\widehat{\rho_1}[r_2], \widehat{\rho_2}[r_2]) \\ & (b_1, b_2) = \mathcal{E}_A(n_1 + n_2) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \oplus (r_1, r_2, r_3) \longmapsto \widehat{\rho_1}[r_3 \mapsto b_1], \widehat{\rho_2}[r_3 \mapsto b_2]; \\ \\ & \top = \mathcal{D}_B(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_1]) \\ & c = \mathcal{D}_B(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_2]) \\ & (b_1, b_2) = \mathcal{E}_B(c) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Mux}(r_1, r_2, r_3, r_4) \longmapsto \widehat{\rho_1}[r_4 \mapsto b_1], \widehat{\rho_2}[r_4 \mapsto b_2]; \\ \\ & \bot = \mathcal{D}_B(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_1]) \\ & c = \mathcal{D}_B(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_3]) \\ & (b_1, b_2) = \mathcal{E}_B(c) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Mux}(r_1, r_2, r_3, r_4) \longmapsto \widehat{\rho_1}[r_4 \mapsto b_1], \widehat{\rho_2}[r_4 \mapsto b_2]; \\ \\ & n_1 = \mathcal{D}_B(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_3]) \\ & (b_1, b_2) = \mathcal{E}_B(c) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Gt}(r_1, r_2, r_3) \longmapsto \widehat{\rho_1}[r_3 \mapsto b_1], \widehat{\rho_2}[r_3 \mapsto b_2]; \\ \hline & (b_1, b_2) = \mathcal{E}_B(n_1 > n_2) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Gt}(r_1, r_2, r_3) \longmapsto \widehat{\rho_1}[r_3 \mapsto b_1], \widehat{\rho_2}[r_3 \mapsto b_2]; \\ \hline & (b_1, b_2) = \mathcal{E}_m(c) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash c \rhd_m r_2 \longmapsto \widehat{\rho_1}[r_2 \mapsto b_1], \widehat{\rho_2}[r_2 \mapsto b_2]; \\ \hline & c = \mathcal{D}_{m_1}(\widehat{\rho_1}[r_1], \widehat{\rho_2}[r_1]) \\ & (b_1, b_2) = \mathcal{E}_m(c) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash r_1 \rhd_m r_2 \longmapsto \widehat{\rho_1}[r_2 \mapsto b_1], \widehat{\rho_2}[r_2 \mapsto b_2]; \\ \hline & (b_1, b_2) = \mathcal{E}_m(\mathsf{get_input}(i)) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Inp}_i \rhd_m r \longmapsto \widehat{\rho_1}[r_2 \mapsto b_1], \widehat{\rho_2}[r_2 \mapsto b_2]; \\ \hline & c = \mathcal{D}_m(\widehat{\rho_1}[r], \widehat{\rho_2}[r]) \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Nut}(r) \longmapsto \widehat{\rho_1}, \widehat{\rho_2}; c, \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Nut}(r_1) \mapsto \widehat{\rho_1}, \widehat{\rho_2}; c, \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Nut}(r_1) \mapsto \widehat{\rho_1}, \widehat{\rho_2}; c, \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Nut}(r_1) \mapsto \widehat{\rho_1}, \widehat{\rho_2}; C, \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \mathsf{Nut}, \kappa_2 \mapsto \widehat{\rho_1}', \widehat{\rho_2}'; O_1 \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \kappa_1, \kappa_2 \mapsto \widehat{\rho_1}', \widehat{\rho_2}'; O_2 \\ \hline & \widehat{\rho_1}, \widehat{\rho_2} \vdash \kappa_1, \kappa_2 \mapsto \widehat{\rho_1}', \widehat{\rho_2}'; O_1, O_2 \\ \hline \end{pmatrix} \quad \text{CKT_SEQ}$$

Figure 13: Circuit evaluation

$$\psi \sim \tau$$

$$\frac{\overline{\sigma \sim \sigma^{\ell}}}{\sigma[n] \sim \sigma^{\ell}[n]} \quad \text{ST_ARR}$$

Figure 14: Source type and target type consistency

$$\Gamma \sim \rho$$

$$\begin{array}{cc} \overline{\cdot \sim} & \text{SEN_EMP} \\ v: \psi \\ \psi \sim \tau \\ \overline{\Gamma \sim \rho} \\ \overline{\Gamma, x: \tau \sim \rho[x \mapsto v]} & \text{SEN_BND} \end{array}$$

Figure 15: Source environment and type environment consistency

Figure 16: Source environment to target environment compilation