# COMPETITIVE PROGRAMMING

## **GEOMETRIC ALGORITHMS**

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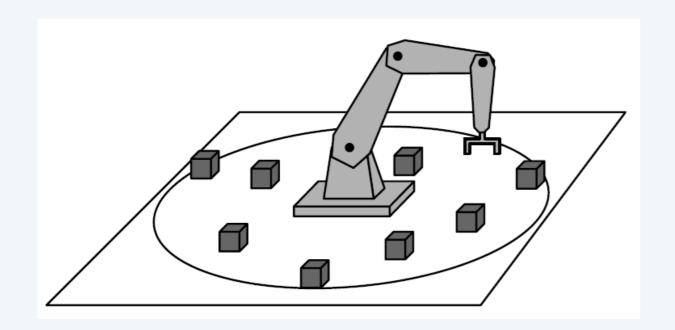


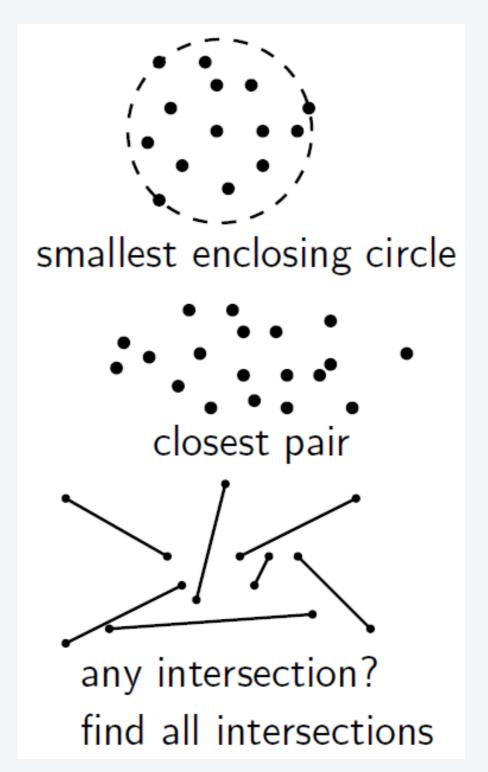
**Sharif University of Technology** 

Spring 2024

## Representative problems

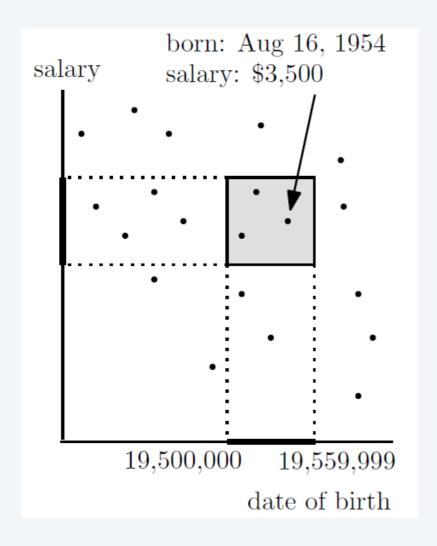
- Smallest enclosing circle
- Closest pair
- Segment intersections

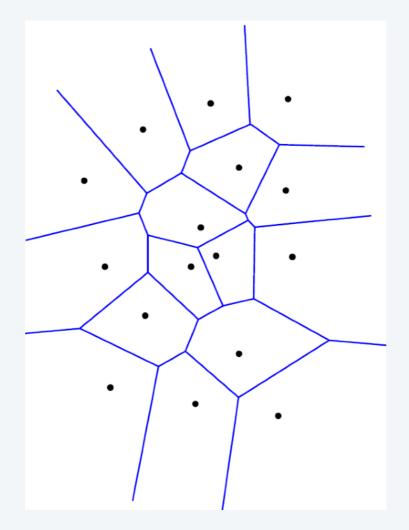


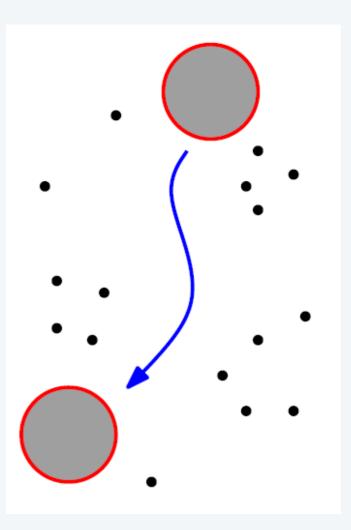


## Representative problems (cont'd)

- Range searching
- Nearest neighbor search
- Motion planning

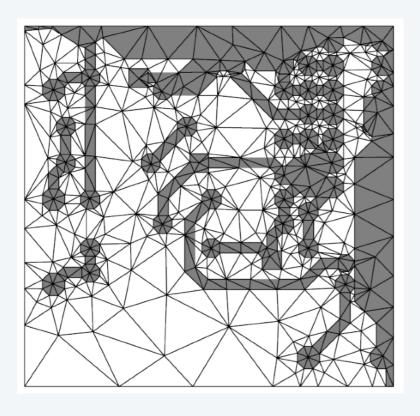


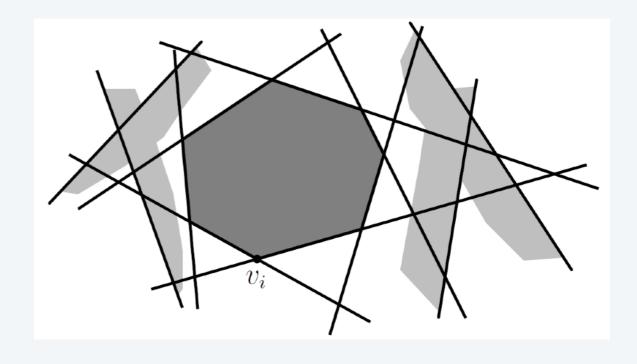


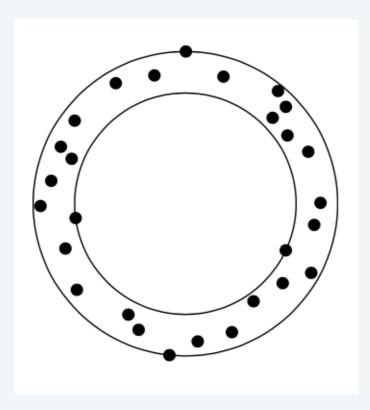


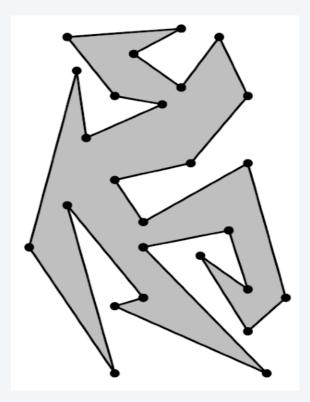
## Geometric objects in 2D

- Point
- Line / Segment
- Circle / Disk
- Polygon / Triangle
- Polygonal Chain
- Half-plane
- ...









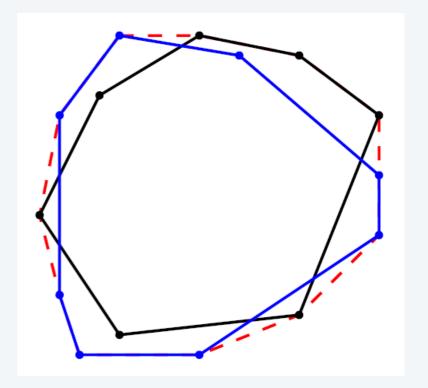
## Geometry problems

## In programming contests:

- Significant impact on the result
- Tricky to implement!

#### Implementation Issues:

- Degenerate cases
- Precision



#### General advices:

- Avoid divisions as much as possible
  - e.g., Instead of if (a/b == c), write if (a == b\*c)
- Avoid using floating-point number whenever you can
- Use long long and double instead of int and float
- Use epsilon in equality tests
  - e.g., Instead of if (x == 0), write if (abs(x) < EPS)

## Geometric primitives

### Dot product:

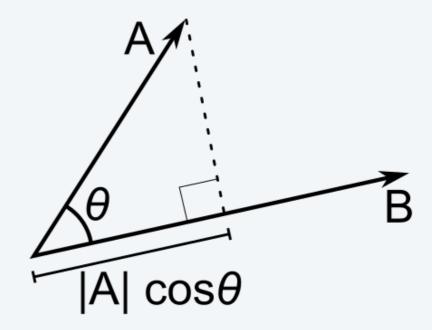
- $A = (a_1, a_2), B = (b_1, b_2)$
- A B =  $a_1b_1 + a_2b_2$

#### Geometric interpretation:

• A • B =  $||A|| ||B|| \cos(\theta)$ 

## Applications:

- Finding vector length:
  - $A \bullet A = ||A||^2$
- Computing projection:
  - $|A| \cos(\theta) = A \bullet B / \|B\|$
- Finding angles:
  - $\theta = \cos^{-1}(A \bullet B / ||A|| ||B||)$



## Geometric primitives (cont'd)

## Cross product:

- $A = (a_1, a_2), B = (b_1, b_2)$
- $A \times B = a_1b_2 a_2b_1$

#### Geometric interpretation:

•  $A \times B = ||A|| ||B|| \sin(\theta)$ 

### **Applications:**

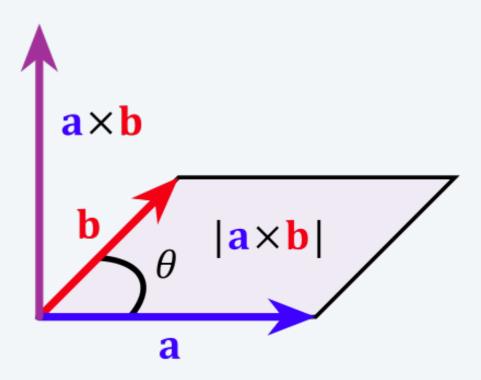
- Computing area of a triangle:
  - $|A \times B| / 2$



• 
$$A \times B = -(B \times A)$$

• Distance of point r to the line through p and q:

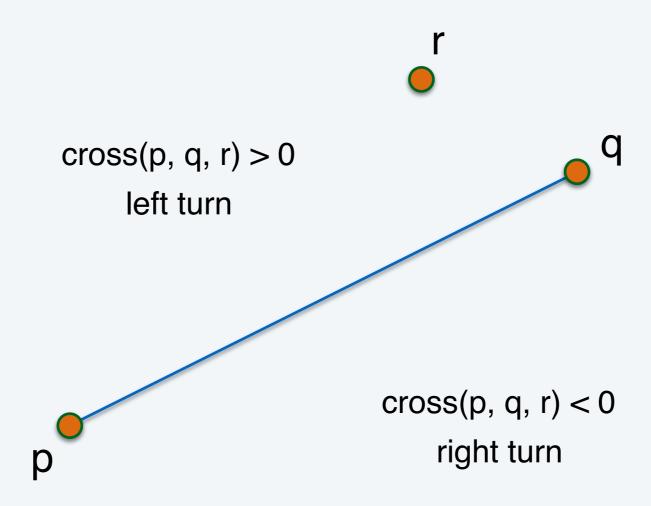
• 
$$(q - p) \times (r - p) / ||q - p||$$



## Orientation of three points

## Define:

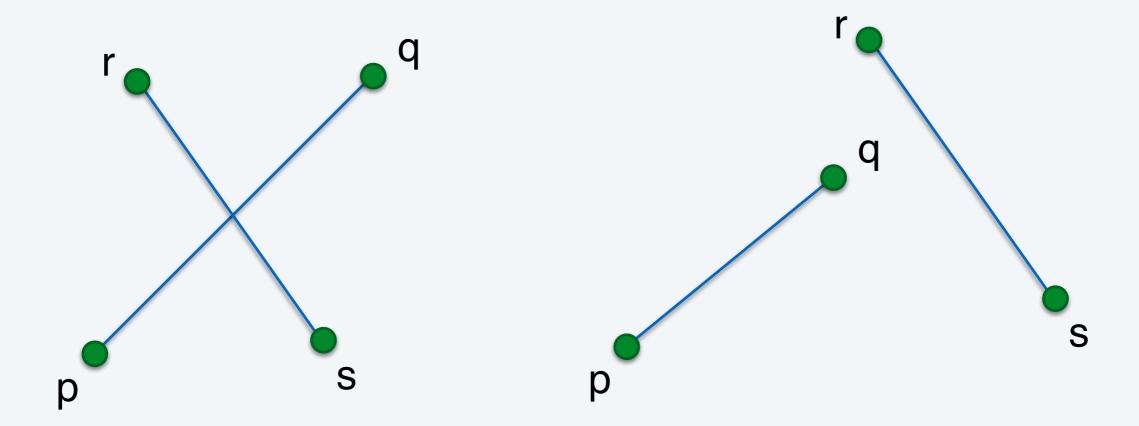
• cross(p, q, r) =  $(q - p) \times (r - p)$ 



## Segment-segment intersection test

## Proper intersection:

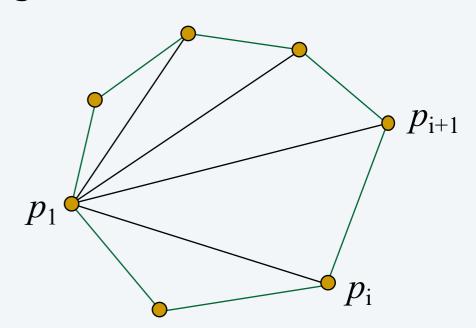
- $cross(p, q, r) \cdot cross(p, q, s) < 0$
- and  $cross(r, s, p) \cdot cross(r, s, q) < 0$

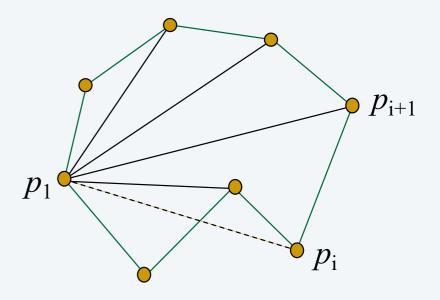


## Area of a simple polygon

#### Problem:

- Given  $p_1, p_2, ..., p_n$  around perimeter of a polygon P
- If P is convex, we can decompose it into triangles:
  - 2 area =  $|\Sigma_{i=2..n-1} (p_i p_1) \times (p_{i+1} p_1)|$
- It also works for non-convex polygons!
  - sum of "signed areas"
- Alternative formula:
  - let  $pi = (x_i y_i)$
  - 2 area =  $|\Sigma_{i=1..n} (x_i y_{i+1}) \times (x_{i+1} y_i)|$

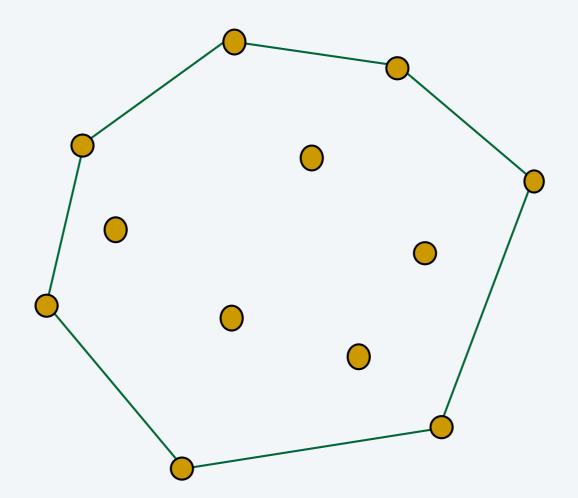




## Convex hull problem

## Problem:

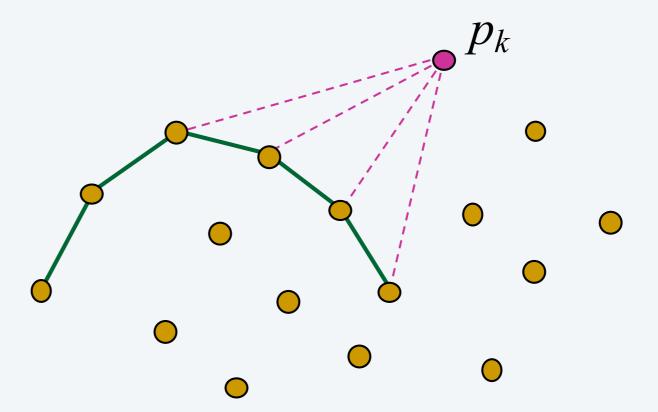
• Given n points in the plane, find the smallest convex polygon containing all the points.



## Graham scan

## Algorithm:

- sort points in x-coordinates: p<sub>1</sub>, p<sub>2</sub>,..., p<sub>n</sub>.
- start with an empty chain
- **for** k = 1 to n:
  - while the last two points of the chain and  $p_k$  make a left turn:
    - remove the last point from the chain
  - add p<sub>k</sub> to the chain
- return the chain



#### Graham scan

#### Algorithm:

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## Analysis:

- · Each point is added to chain exactly once, and is deleted at most once.
  - So the scan takes O(n) time.
- But we need to sort the points at first.
  - Overall runtime is O(n log n).

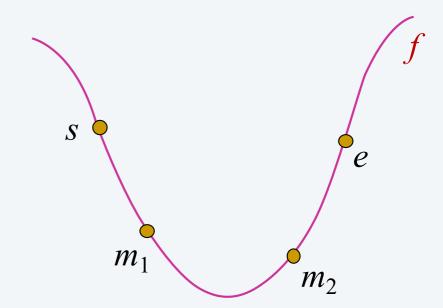
## Ternary search

#### Problem:

• Find the minimum of a convex function f

#### Solution:

- start with search interval [s, e]
- while |e s| is not small enough:
  - let  $m_1 = s + (e s)/3$
  - let  $m_2 = e (e s)/3$
  - if  $f(m_1) < f(m_2)$ , then set e to  $m_2$
  - otherwise set s to m<sub>1</sub>



## Team reference document



KTH Royal Institute of Technology

# Omogen Heap

Simon Lindholm, Johan Sannemo, Mårten Wiman

https://github.com/kth-competitive-programming/kactl/

#### References

- J. Park, Introduction to Programming Contests, Stanford, Winter 2012.
- M. de Berg, O. Cheong, M. van Kreveld, M. Overmars, Computational Geometry: Algorithms and Applications, 3rd edition, Springer, 2008.
- J. O'Rourke, Computational Geometry in C, 2nd edition, Cambridge University Press, 1998.

