KUET_Effervescent Team Notebook struct dsu with rollbacks { 7.6 vector<int> p, rnk; 7.7 Md. Mehrab Hossain Opi, Arnob Sarker, int comps: Sharif Minhazul Islam stack<dsu_save> op; Miller-Rabin-Pollard-Rho [68 lines] - 3e3e5f 16 dsu_with_rollbacks() {} dsu_with_rollbacks(int n) { Contents p.resize(n); No of Digits in n! in base B [7 lines] - 86bfaf 17 rnk.resize(n); for (int i = 0; i < n; i++) { 1 Data Structure p[i] = i;Dsu With Rollback [89 lines] - bc2588 rnk[i] = 0;MO with Update [43 lines] - 2fbf87 8 Misc comps = n;8.1 Persistent Segment Tree [64 lines] - f58bc9 8.2 SQRT Decomposition [96 lines] - a772d3 find_set(p[v]); } Segment Tree [73 lines] - c1fe4f 8.4 bool unite(int v. int u) { 1.7 8.5 v = find set(v): 1.8 8.6 u = find set(u): 1.9 if (v == u) return false; 2 Dynamic Programming if (rnk[v] > rnk[u]) swap(v, u); Divide and Conquer DP [26 lines] - 6d8559 9 String Dvnamic Convex Hull Trick [66 lines] - c283fc 9.1 if (rnk[u] == rnk[v]) rnk[u]++: Knuth Optimization [32 lines] - 911417 return true: LIS O(nlogn) with full path [17 lines] - e7e81f SOS DP [18 lines] - 5063f0 2.5 void rollback() { 2.6 Sibling DP [26 lines] - cfc5ff Palindromic Tree [30 lines] - 9ebc05 if (op.empty()) return; Prefix Function Automaton [21 lines] - b65c0e 19 $dsu_save x = op.top();$ 3 Flow 9.7 Suffix Array [78 lines] - f2f7a0 op.pop(); Blossom [58 lines] - 1b2a6f Suffix Automata [109 lines] - 600ddc 20 comps++; p[x.v] = x.v;3.3 rnk[x.v] = x.rnkv;HopCroftKarp [67 lines] - fac9fc p[x.u] = x.u;rnk[x.u] = x.rnku;3.5 10 Random 3.6 MCMF [116 lines] - 466389 struct query { 4 Game Theory 10.1.2 Stirling Number of the First Kind . . . 21 int v, u; Points to be noted [14 lines] - 6fe124 10.1.3 Stirling Numbers of the Second Kind . . 21 bool united; query(int _v, int _u) : v(_v), u(_u) {} 5 Geometry Geometry [384 lines] - 6bfd7b struct QuervTree { vector<vector<query>> t; dsu_with_rollbacks dsu; int T: 11 6 Graph QueryTree() {} QueryTree(int _T, int n) : T(_T) { 10.2.1 Mobius Function and Inversion 22 dsu = dsu with rollbacks(n): Centroid Decomposition [39 lines] - d5d02b 12 6.3 t.resize(4 * T + 4): 10.2.2 GCD and LCM 6.4 Heavy Light Decomposition [73 lines] - d0e24f 12 10.2.3 Gauss Circle Theorem query& q) { 6.7 if (ul > ur) return; 10.2.5 Formula Cheatsheet 6.8 if $(1 == u1 \&\& r == ur) {$ t[v].push_back(q); 6.9 1 Data Structure return; 1.1 Dsu With Rollback [89 lines] - bc2588 7 Math struct dsu_save { int mid = (1 + r) / 2;int v, rnkv, u, rnku; dsu_save() {} dsu_save(int _v, int _rnkv, int _u, int _rnku)

: v(_v), rnkv(_rnkv), u(_u), rnku(_rnku) { }

```
int find_set(int v) { return (v == p[v]) ? v :
  op.push(dsu_save(v, rnk[v], u, rnk[u]));
void add_to_tree(int v, int l, int r, int ul, int ur,
  add_to_tree(2 * v, 1, mid, ul, min(ur, mid), q);
  add_to_tree(2 * v + 1, mid + 1, r, max(ul, mid +
     1), ur, q);
```

```
while (T > q.T) undo(T--, L, R);
                                                                 while (L > q.L) add(--L);
  void add_query(query q, int 1, int r) {
      add_to_tree(1, 0, T - 1, 1, r, q); }
                                                                 while (R < q.R) add(++R);
                                                                 while (L < q.L) remove(L++);
  void dfs(int v, int 1, int r, vector<int>& ans) {
    for (query& q : t[v]) {
                                                                 while (R > q.R) remove(R--);
      q.united = dsu.unite(q.v, q.u);
                                                                 ans[q.id] = get();
    if (1 == r)
      ans[1] = dsu.comps;
                                                             1.3 MO [28 lines] - bed3e5
    else {
                                                             const int N = 2e5 + 5;
      int mid = (1 + r) / 2;
                                                             const int Q = 2e5 + 5;
      dfs(2 * v, 1, mid, ans);
                                                             const int SZ = sqrt(N) + 1;
      dfs(2 * v + 1, mid + 1, r, ans);
                                                             struct qry {
                                                               int 1, r, id, blk;
    for (query q : t[v]) {
                                                               bool operator<(const qry& p) const {</pre>
      if (q.united) dsu.rollback();
                                                                 return blk == p.blk ? r < p.r : blk < p.blk;
                                                             };
  vector<int> solve() {
                                                             qry query[Q];
    vector<int> ans(T):
                                                             11 ans[0]:
    dfs(1, 0, T - 1, ans);
                                                             void add(int id) {}
    return ans:
                                                             void remove(int id) {}
                                                             11 get() {}
                                                             int n, q;
                                                             void MO() {
1.2 MO with Update [43 lines] - 2fbf87
                                                               sort(query, query + q);
//1 indexed
                                                               int cur_1 = 0, cur_r = -1;
//Complexity:O(S \times Q + Q \times \frac{N^2}{S^2})
                                                               for (int i = 0; i < q; i++) {
//S = (2*n^2)^(1/3)
                                                                 qry q = query[i];
const int block_size = 2720; // 4310 for 2e5
                                                                 while (cur_1 > q.1) add(--cur_1);
const int mx = 1e5 + 5;
                                                                 while (cur_r < q.r) add(++cur_r);</pre>
struct Query {
                                                                 while (cur_l < q.l) remove(cur_l++);</pre>
  int L, R, T, id;
                                                                 while (cur_r > q.r) remove(cur_r--);
  Query() {}
                                                                 ans[q.id] = get();
  Query(int _L, int _R, int _T, int _id) : L(_L),
      R(_R), T(_T), id(_id) {}
  bool operator<(const Query &x) const {</pre>
                                                             /* 0 indexed. */
    if (L / block_size == x.L / block_size) {
                                                             1.4 Persistent Segment Tree [64 lines] - f58bc9
      if (R / block_size == x.R / block_size) return T <</pre>
                                                             const int mxn = 4e5+5;
      return R / block_size < x.R / block_size;</pre>
                                                             int root[mxn], leftchild[25*mxn], rightchild[25*mxn],
                                                                 value[25*mxn], a[mxn];
    return L / block_size < x.L / block_size;
                                                             int now = 0, n, sz = 1;
                                                             int 1, r;
} Q[mx];
struct Update {
                                                             int build(int L, int R){
  int pos;
                                                               int node = ++now:
  int old, cur;
                                                               if(L == R){
  Update(){};
                                                                 //initialize
  Update(int _p, int _o, int _c) : pos(_p), old(_o),
                                                                 //value[node] = a[L]:
      cur( c){}:
                                                                 return node:
} U[mx]:
int ans[mx]:
                                                               int mid = (L+R)>>1;
inline void add(int id) {}
                                                               leftchild[node] = build(L, mid):
inline void remove(int id) {}
                                                               rightchild[node] = build(mid+1, R);
inline void update(int id, int L, int R) {}
inline void undo(int id, int L, int R) {}
                                                               //value[node] = value[leftchild[node]] +
inline int get() {}
                                                                   value[rightchild[node]];
void MO(int nq, int nu) {
                                                               return node;
  sort(Q + 1, Q + nq + 1);
  int L = 1, R = 0, T = nu;
  for (int i = 1; i <= nq; i++) {
                                                             int update(int nownode, int L, int R, int ind, int val){
    Query q = Q[i];
                                                               int node = ++now;
    while (T < q.T) update(++T, L, R);
                                                               if(L == R){
```

```
//value[node] = value[nownode]+val;
    //update value[node]
   return node:
  int mid = (L+R)>>1;
  leftchild[node] = leftchild[nownode];
  rightchild[node] = rightchild[nownode];
  if (mid >= ind) {//change condition as required
   leftchild[node] = update(leftchild[nownode], L,
        mid, ind, val);
   rightchild[node] = update(rightchild[nownode],
        mid+1, R, ind, val);
  //value[node] = value[leftchild[node]] +
      value[rightchild[node]];
  //combine value[node]
 return node:
int query(int nownode, int L, int R){
  if(1 > R \mid \mid r < L) return 0;
 if(L>=1 \&\& r >= R){
   return value[nownode]:
 int mid = (L+R)>>1;
 //change as required
 return query(leftchild[nownode], L, mid) +
      query(rightchild[nownode], mid+1, R);
void persistant(){
 root[0] = build(1, n);
 while(m--){
   if(ck == 2){
      cout << query(root[idx], 1, n) << "\n";</pre>
    else{
          root[sz++] = update(root[idx], 1, n, ind,
}
1.5 SQRT Decomposition [96 lines] - a772d3
struct sqrtDecomposition {
  static const int sz = 320; //sz = sqrt(N);
  int numberofblocks:
  struct node {
   int L. R:
   bool islazy = false;
   11 lazyval=0;
    //extra data needed for different problems
    void ini(int 1, int r) {
     for(int i=1; i<=r; i++)
        //...initialize as need
```

L=1, R=r;

void semiupdate(int 1, int r, 11 val) {

```
if(1>r) return:
    if(islazv){
      for(int i=L; i<=R; i++){
        //...distribute lazy to everyone
      islazy = 0;
      lazyval = 0;
    for(int i=1; i<=r; i++){
      //...do it manually
  void fullupdate(ll val){
    if(islazy){
      //...only update lazyval
    else{
      for(int i=L: i<=R: i++){
        //...everyone are not equal, make them equal
      islazy = 1;
      //update lazyval
  void update(int 1, int r, 11 val){
    if(1<=L && r>=R) fullupdate(val);
    else semiupdate(max(1, L), min(r, R), val);
  11 semiquery(int 1, int r){
    if(1>r) return 0;
    if(islazy){
      for(int i=L; i<=R; i++){
        //...distribute lazy to everyone
      islazv = 0;
      lazyval = 0;
    11 \text{ ret} = 0;
    for(int i=1; i<=r; i++){
      //...take one by one
    return ret;
  11 fullquery(){
    //return stored value:
  11 query(int 1, int r){
    if(1<=L && r>=R) return fullquery();
    else return semiquery(max(1, L), min(r, R));
};
vector<node> blocks:
void init(int n){
  numberofblocks = (n+sz-1)/sz;
  int curL = 1, curR = sz;
  blocks.resize(numberofblocks+5);
  for(int i=1; i<=numberofblocks; i++){</pre>
    curR = min(n, curR);
    blocks[i].ini(curL, curR);
    curL += sz;
    curR += sz;
void update(int 1, int r, ll val){
```

```
int left = (1-1)/sz+1:
    int right = (r-1)/sz+1:
   for(int i=left; i<=right; i++){</pre>
     blocks[i].update(1, r, val);
 11 query(int 1, int r){
   int left = (1-1)/sz+1;
   int right = (r-1)/sz+1;
   11 \text{ ret} = 0;
   for(int i=left; i<=right; i++){</pre>
     ret += blocks[i].query(1, r);
   return ret;
};
1.6 Segment Tree [73 lines] - c1fe4f
/*edit:data,combine,build check datatype*/
template<typename T>
struct SegmentTree {
#define lc (C << 1)
#define rc (C << 1 | 1)
#define M((L+R)>>1)
 struct data {
   T sum:
   data() :sum(0) {};
 };
  vector<data>st:
  vector<bool>isLazv:
  vector<T>lazv:
  int N:
  SegmentTree(int _N) :N(_N) {
    st.resize(4 * N);
   isLazy.resize(4 * N);
    lazy.resize(4 * N);
 void combine(data% cur, data% 1, data% r) {
    cur.sum = 1.sum + r.sum;
 void push(int C, int L, int R) {
   if (!isLazy[C]) return;
   if (L != R) {
     isLazy[lc] = 1;
     isLazy[rc] = 1;
     lazy[lc] += lazy[C];
     lazy[rc] += lazy[C];
    st[C].sum = (R - L + 1) * lazy[C];
   lazy[C] = 0;
   isLazy[C] = false;
  void build(int C, int L, int R) {
   if (L == R) {
      st[C].sum = 0:
     return;
   build(lc, L, M);
   build(rc, M + 1, R);
    combine(st[C], st[lc], st[rc]);
  data Query(int i, int j, int C, int L, int R) {
   push(C, L, R);
    if (j < L \mid | i > R \mid | L > R) return data(); //
```

default val O/INF

```
if (i <= L && R <= j) return st[C];
    data ret:
    data d1 = Query(i, j, lc, L, M);
    data d2 = Query(i, j, rc, M + 1, R);
    combine(ret, d1, d2);
   return ret;
  void Update(int i, int j, T val, int C, int L, int R)
   push(C, L, R);
   if (j < L || i > R || L > R) return;
   if (i <= L && R <= j) {
      isLazy[C] = 1;
      lazy[C] = val;
      push(C, L, R);
      return;
    Update(i, j, val, lc, L, M);
    Update(i, j, val, rc, M + 1, R);
    combine(st[C], st[lc], st[rc]);
  void Update(int i, int j, T val) {
    Update(i, j, val, 1, 1, N);
 T Query(int i, int j) {
    return Query(i, j, 1, 1, N).sum;
};
1.7 Sqrt Tricks [8 lines] - addf19
1. Size of the block is not always Sqrt, adjust it as
    necessary. if o(n/b+b) then take n/b = b and
    calculate b.
   *it is possible to solve a Mo problem without any
```

```
2. MO's Algorithm
       remove operation. For L in one block R only
       increases, for every range we can start L from
       the last of that block
3. Sqrt Decomposition by time of queries.
```

*keep overall solution and sqrt(n) updates in a

when the vector size exceeds sqrt(n) you can add these updates with overall solution using o(n) 4. If sum of N positive numbers are S, there are at most sqrt(S) distinct values. 5. Randomization

vector and for a query iterate over all of them,

6. Baby step, gaint step

```
1.8 Treap [166 lines] - 8eef59
struct Treap {
 struct Node {
    int val, priority, cnt; // value, priority, subtree
        size
   Node* 1. * r:
                              // left child, right child
        pointer
   Node() {} //rng from template
   Node(int key) : val(key), priority(rng()),
       1(nullptr), r(nullptr) {}
  typedef Node* node;
 node root;
  Treap() : root(0) {}
  int cnt(node t) { return t ? t->cnt : 0; } // return
      subtree size
```

```
void updateCnt(node t) {
  if (t) t->cnt = 1 + cnt(t->1) + cnt(t->r); //
      update subtree size
void push(node cur) {
 ; // Lazy Propagation
void combine(node& cur, node 1, node r) {
  if (!1) {
    cur = r;
    return;
  if (!r) {
    cur = 1;
    return;
  // Merge Operations like in segment tree
void reset(node& cur) {
  if (!cur) return; // To reset other fields of cur
      except value and cnt
void operation(node& cur) {
  if (!cur) return:
  reset(cur);
  combine(cur, cur->1, cur);
  combine(cur, cur, cur->r);
// Split(T, key): split the tree in two tree. Left
    pointer contains all value
// less than or equal to key. Right pointer contains
void split(node t, node& 1, node& r, int key) {
  if (!t)
    return void(l = r = nullptr);
  push(t);
  if (t->val <= key) {
    split(t->r, t->r, r, key), l = t;
  else {
    split(t->1, 1, t->1, key), r = t;
  updateCnt(t);
  operation(t);
void splitPos(node t, node& 1, node& r, int k, int add
  if (!t) return void(1 = r = 0):
  push(t):
  int idx = add + cnt(t->1):
  if (idx \le k)
    splitPos(t->r, t->r, r, k, idx + 1), l = t;
    splitPos(t->1, 1, t->1, k, add), r = t;
  updateCnt(t);
  operation(t);
// Merge(T1,T2): merges 2 tree into one. The tree with
    root of higher
// priority becomes the new root.
```

```
void merge(node& t, node 1, node r) {
 push(1):
 push(r);
  if (!1 || !r)
   t = 1 ? 1 : r;
  else if (l->priority > r->priority)
   merge(1->r, 1->r, r), t = 1;
   merge(r->1, 1, r->1), t = r;
 updateCnt(t);
 operation(t);
node merge_treap(node 1, node r) {
  if (!1) return r;
  if (!r) return 1;
  if (1->priority < r->priority) swap(1, r);
 node L. R:
 split(r, L, R, 1->val);
 1->r = merge\_treap(1->r, R);
 1->1 = merge\_treap(L, 1->1);
 updateCnt(1);
 operation(1);
 return 1:
// insert creates a set.all unique value.
void insert(int val) {
 if (!root) {
   root = new Node(val):
   return;
 node 1, r, mid, mid2, rr;
  mid = new Node(val);
  split(root, 1, r, val);
  merge(1, 1, mid); // these 3 lines will create
      multiset.
  merge(root, 1, r);
  /*split(root, l, r, val - 1); // l contains all
      small values.
    merge(l, l, mid):
                                  // l contains new val
        too.
    split(r, mid2, rr, val);
                                  // rr contains all
        greater values.
    merge(root, l, rr);*/
// removes all similar values.
void erase(int val) {
 node 1, r, mid;
  /* Removes all similar element*/
  split(root, 1, r, val - 1);
  split(r, mid, r, val);
  merge(root, 1, r);
  /*Removes single instance*/
  /*split(root, l, r, val - 1);
    split(r, mid, r, val);
    merge(mid, mid \rightarrow l, mid \rightarrow r);
    merge(l, l, mid);
    merge(root, l, r);*/
void clear(node cur) {
 if (!cur) return;
  clear(cur->1), clear(cur->r);
 delete cur;
```

```
cout << t->val << ' ';
   inorder(t->r);
  void inorder() {
    inorder(root);
   puts("");
  //1 indexed - xth element after sorting.
  int find_by_order(int x) {
   if (!x) return -1;
    x--;
   node 1, r, mid;
    splitPos(root, 1, r, x - 1);
    splitPos(r, mid, r, 0);
    int ans = -1:
   if (cnt(mid) == 1) ans = mid->val;
    merge(r, mid, r);
    merge(root, 1, r);
  // 1 indexed. index of val in sorted array. -1 if not
      found.
  int order of kev(int val) {
   node 1. r. mid:
    split(root, 1, r, val - 1);
    split(r, mid, r, val);
    int ans = -1;
    if (cnt(mid) == 1) ans = 1 + cnt(1);
    merge(r, mid, r);
    merge(root, 1, r);
    return ans;
};
1.9 Trie Bit [61 lines] - 390174
struct Trie {
 struct node {
   int next[2]:
   int cnt. fin:
   node() :cnt(0), fin(0) {
      for (int i = 0; i < 2; i++) next[i] = -1;
 };
 vector<node>data;
 Trie() {
   data.push_back(node());
 void key_add(int val) {
   int cur = 0:
   for (int i = 30; i >= 0; i--) {
     int id = (val >> i) & 1;
      if (data[cur].next[id] == -1) {
       data[cur].next[id] = data.size();
        data.push_back(node());
      cur = data[cur].next[id];
      data[cur].cnt++;
    data[cur].fin++;
```

void clear() { clear(root); }

void inorder(node t) {

if (!t) return:

inorder(t->1);

```
int key_search(int val) {
    int cur = 0:
    for (int i = 30; ~i; i--) {
      int id = (val >> i) & 1;
      if (data[cur].next[id] == -1) return 0;
      cur = data[cur].next[id];
    return data[cur].fin;
 void key_delete(int val) {
    int cur = 0;
    for (int i = 30; ~i; i--) {
      int id = (val >> i) & 1;
      cur = data[cur].next[id];
      data[cur].cnt--;
    data[cur].fin--:
  bool key_remove(int val) {
    if (key_search(val)) return key_delete(val), 1;
    return 0:
  int maxXor(int x) {
    int cur = 0:
    int ans = 0:
    for (int i = 30: ~i: i--) {
      int b = (x >> i) & 1;
      if (data[cur].next[!b] + 1 &&
          data[data[cur].next[!b]].cnt > 0) {
        ans += (1LL << i);
        cur = data[cur].next[!b];
      else cur = data[cur].next[b];
    return ans;
2 Dynamic Programming
2.1~ Divide and Conquer DP [26 lines] - 6d8559
11 G,L;///total group,cell size
ll dp[8001][801],cum[8001];
11 C[8001];///value of each cell
inline ll cost(ll l,ll r){
 return(cum[r]-cum[l-1])*(r-l+1);
void fn(ll g,ll st,ll ed,ll r1,ll r2){
  if(st>ed)return;
 11 \text{ mid}=(\text{st+ed})/2.\text{pos}=-1:
  dp[mid][g]=inf;
  for(int i=r1;i<=r2;i++){
    11 tcost=cost(i,mid)+dp[i-1][g-1];
    if(tcost<dp[mid][g]){
        dp[mid][g]=tcost,pos=i;
  fn(g,st,mid-1,r1,pos);
  fn(g,mid+1,ed,pos,r2);
int main(){
 for(int i=1;i<=L;i++)
    cum[i]=cum[i-1]+C[i];
  for(int i=1;i<=L;i++)</pre>
```

```
dp[i][1]=cost(1,i);
  for(int i=2; i \le G; i++) fn(i,1,L,1,L);
2.2 Dynamic Convex Hull Trick [66 lines] - c283fc
const int N = 3e5 + 9:
const int mod = 1e9 + 7;
//add lines with -m and -b and return -ans to
//make this code work for minimums. (not -x)
const ll inf = -(1LL \ll 62);
struct line {
 11 m, b;
  mutable function<const line*() > succ;
  bool operator < (const line& rhs) const {</pre>
    if (rhs.b != inf) return m < rhs.m;
    const line* s = succ();
    if (!s) return 0;
    11 x = rhs.m:
    return b - s \rightarrow b < (s \rightarrow m - m) * x:
struct CHT : public multiset<line> {
  bool bad(iterator y) {
    auto z = next(y);
    if (v == begin()) {
      if (z == end()) return 0;
      return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x \rightarrow m \&\& y \rightarrow b
        \langle = x - \rangle b:
    return 1.0 * (x -> b - y -> b) * (z -> m - y -> m)
        >= 1.0 * (y -> b - z -> b) * (y -> m - x -> m);
  void add(ll m, ll b) {
    auto y = insert({ m, b });
    y->succ = [ = ] { return next(y) == end() ? 0 :
        &*next(y); };
    if (bad(y)) {
      erase(v);
      return;
    while (next(y) != end() && bad(next(y)))
        erase(next(v)):
    while (y != begin() && bad(prev(y))) erase(prev(y));
  11 query(11 x) {
    assert(!empty());
    auto 1 = *lower_bound((line) {
      x. inf
    }):
    return 1.m * x + 1.b:
CHT* cht;
ll a[N], b[N];
int32_t main() {
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  int n;
  for(int i = 0; i < n; i++) cin >> a[i];
  for(int i = 0; i < n; i++) cin >> b[i];
```

```
cht = new CHT():
  cht \rightarrow add(-b[0], 0):
  11 \text{ ans} = 0;
  for(int i = 1; i < n; i++) {
    ans = -cht -> query(a[i]);
    cht -> add(-b[i], -ans);
  cout << ans << nl;</pre>
 return 0;
2.3 Knuth Optimization [32 lines] - 911417
/*It is applicable where recurrence is in the form :
dp[i][j] = mini < k < j \{ dp[i][k] + dp[k][j] \} + C[i][j]
condition for applicability is:
A[i, j-1] \leftarrow A[i, j] \leftarrow A[i+1, j]
Where.
A[i][j]-the smallest k that gives optimal answer, like-
dp[i][j] = dp[i-1][k] + C[k][j]
C[i][j]-given cost function
also applicable if: C[i][j]satisfies the following 2
    conditions:
C[a][c]+C[b][d] <= C[a][d]+C[b][c], a <= b <= c <= d
C\lceil b \rceil \lceil c \rceil \le C\lceil a \rceil \lceil d \rceil, a \le b \le c \le d
reduces time complexity from O(n^3) to O(n^2)*/
for(int s=0;s<=k;s++)//s-length(size)of substring</pre>
  for (int l=0; l+s <= k; l++) {//l-left point}
    int r=1+s;//r-right point
    if(s<2){
      res[1][r]=0;//DP base-nothing to break
      mid[1][r]=1;/*mid is equal to left border*/
      continue:
    int mleft=mid[l][r-1];/*Knuth's trick: getting
         bounds on m*/
    int mright=mid[l+1][r];
    res[1][r]=inf;
    for(int m=mleft;m<=mright;m++){/*iterating for m in</pre>
         the bounds only*/
      int64 tres=res[l][m]+res[m][r]+(x[r]-x[l]);
      if(res[1][r]>tres){//relax current solution
        res[1][r]=tres;
        mid[l][r]=m:
int64 answer=res[0][k]:
2.4 LIS O(nlogn) with full path [17 lines] - e7e81f
int num[MX],mem[MX],prev[MX],array[MX],res[MX],maxlen;
void LIS(int SZ.int num∏){
  CLR(mem),CLR(prev),CLR(array),CLR(res);
  int i.k:
  maxlen=1;
  array[0]=-inf;
  RFOR(i,1,SZ+1) array[i]=inf;
  prev[0]=-1,mem[0]=num[0];
  FOR(i,SZ){
    k=lower_bound(array,array+maxlen+1,num[i])-array;
    if(k==1) array[k]=num[i],mem[k]=i,prev[i]=-1;
    else array[k]=num[i],mem[k]=i,prev[i]=mem[k-1];
    if(k>maxlen) maxlen=k;
```

```
for(i=mem[maxlen];i!=-1;i=prev[i])res[k++]=num[i];
2.5 SOS DP [18 lines] - 5063f0
//iterative version
for(int mask = 0; mask < (1<<N); ++mask){</pre>
  dp[mask][-1] = A[mask]; //handle base case separately
      (leaf states)
  for(int i = 0; i < N; ++i){
    if(mask & (1<<i))
      dp[mask][i] = dp[mask][i-1] +
          dp[mask^(1 < i)][i-1];
      dp[mask][i] = dp[mask][i-1];
 F[mask] = dp[mask][N-1];
//memory optimized, super easy to code.
for(int i = 0; i < (1 << N); ++i)
 F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask <
    (1 << N); ++ mask){
  if(mask & (1<<i))
    F[mask] += F[mask^(1<<i)];
2.6 Sibling DP [26 lines] - cfc5ff
/*/dividing tree into min group such that each group
    cost not exceed k*/
ll n,k,dp[mx][mx];
vector<pair<11,11>>adj[mx];///must be rooted tree
ll sibling_dp(ll par,ll idx,ll remk){
 if(remk<0)return inf;</pre>
  if(adj[par].size()<idx+1)return 0;</pre>
 11 u=adj[par][idx].first;
  if(dp[u][remk]!=-1)
    return dp[u][remk];
  11 ret=inf,under=0,sibling=0;
  if(par!=0){//creating new group
    under=1+dfs(u,0,k);
    sibling=dfs(par,idx+1,remk);
    ret=min(ret,under+sibling);
  //divide the current group
 11 temp=remk-adj[par][idx].second;
  for(ll chk=temp;chk>=0;chk--){
    11 siblingk=temp-chk;
    under=0.sibling=0:
    under=dfs(u,0,chk);
    sibling=dfs(par,idx+1,siblingk);
    ret=min(ret,under+sibling);
  return dp[u][remk]=ret;
3 Flow
3.1 Blossom [58 lines] - 1b2a6f
// Finds Maximum matching in General Graph
// Complexity O(NM)
// mate[i] = j means i is paired with j
// source: https://codeforces.com/blog/entry
    /92339?#comment-810242
vector<int> Blossom(vector<vector<int>>& graph) {
```

//mate contains matched edge.

```
int n = graph.size(), timer = -1;
 vector<int> mate(n, -1), label(n), parent(n),
   orig(n), aux(n, -1), q;
 auto lca = [\&](int x, int y) {
   for (timer++; ; swap(x, \dot{y})) {
      if (x == -1) continue;
      if (aux[x] == timer) return x;
      aux[x] = timer;
     x = (mate[x] == -1 ? -1 : orig[parent[mate[x]]]);
 };
 auto blossom = [&](int v, int w, int a) {
   while (orig[v] != a) {
     parent[v] = w; w = mate[v];
     if (label[w] == 1) label[w] = 0, q.push_back(w);
      orig[v] = orig[w] = a; v = parent[w];
 }:
 auto augment = [&](int v) {
   while (v != -1) {
     int pv = parent[v], nv = mate[pv];
     mate[v] = pv; mate[pv] = v; v = nv;
 };
 auto bfs = [&](int root) {
   fill(label.begin(), label.end(), -1);
   iota(orig.begin(), orig.end(), 0);
   q.clear();
   label[root] = 0; q.push_back(root);
   for (int i = 0; i < (int)q.size(); ++i) {
     int v = q[i];
     for (auto x : graph[v]) {
       if (label[x] == -1) {
         label[x] = 1; parent[x] = v;
         if (mate[x] == -1)
            return augment(x), 1;
         label[mate[x]] = 0; q.push_back(mate[x]);
       else if (label[x] == 0 \&\& orig[v] != orig[x]) {
         int a = lca(orig[v], orig[x]);
         blossom(x, v, a); blossom(v, x, a);
     }
   return 0:
 // Time halves if you start with (any) maximal
      matching.
 for (int i = 0; i < n; i++)
   if (mate[i] == -1)
     bfs(i):
 return mate:
3.2 Dinic [72 lines] - a786f1
/*.Complexity: O(V^2 E)
 .Call Dinic with total number of nodes.
 .Nodes start from 0.
 . Capacity is long long data.
 .make graph with create edge(u,v,capacity).
 .Get max flow with maxFlow(src,des).*/
#define eb emplace_back
struct Dinic {
 struct Edge {
```

int u, v;

```
11 cap, flow = 0;
   Edge() {}
   Edge(int u, int v, ll cap) :u(u), v(v), cap(cap) {}
 int N;
  vector<Edge>edge;
  vector<vector<int>>adj;
  vector<int>d, pt;
  Dinic(int N) : N(N), edge(0), adj(N), d(N), pt(N) {}
  void addEdge(int u, int v, ll cap) {
   if (u == v) return;
    edge.eb(u, v, cap);
    adj[u].eb(edge.size() - 1);
    edge.eb(v, u, 0);
    adj[v].eb(edge.size() - 1);
  bool bfs(int s, int t) {
    queue<int>q({ s }):
    fill(d.begin(), d.end(), N + 1);
    d[s] = 0:
    while (!q.empty()) {
      int u = q.front();q.pop();
      if (u == t) break;
      for (int k : adj[u]) {
        Edge& e = edge[k]:
        if (e.flow<e.cap && d[e.v]>d[e.u] + 1) {
          d[e.v] = d[e.u] + 1:
          q.emplace(e.v);
   return d[t] != N + 1;
 ll dfs(int u, int T, ll flow = -1) {
    if (u == T \mid | flow == 0) return flow;
    for (int& i = pt[u];i < adj[u].size();i++) {</pre>
      Edge& e = edge[adj[u][i]];
      Edge& oe = edge[adj[u][i] ^ 1];
      if (d[e.v] == d[e.u] + 1) {
       11 amt = e.cap - e.flow;
        if (flow != -1 && amt > flow) amt = flow;
        if (ll pushed = dfs(e.v, T, amt)) {
          e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
   }
    return 0:
 11 maxFlow(int s. int t) {
   11 total = 0:
    while (bfs(s, t)) {
      fill(pt.begin(), pt.end(), 0);
      while (ll flow = dfs(s, t)) {
       total += flow;
    return total;
};
```

```
Covering Problems:
> Maximum Independent Set(Bipartite): Largest set of
    nodes which do not have any edge between them. sol:
    V-(MaxMatching)
> Minimum Vertex Cover(Bipartite): -Smallest set of
    nodes to cover all the edges -sol: MaxMatching
> Minimum Edge Cover(General graph): -Smallest set of
    edges to cover all the nodes -sol: V-(MaxMatching)
    (if edge cover exists, does not exit for isolated
> Minimum Path Cover(Vertex disjoint) DAG: -Minimum
    number of vertex disjoint paths that visit all the
    nodes -sol: make a bipartite graph using same nodes
    in two sides, one side is "from" other is "to", add
    edges from "from" to "to", then ans is
    V-(MaxMatching)
> Minimum Path Cover(Vertex Not Disjoint) General graph:
    -Minimum number of paths that visit all the nodes
    -sol: consider cycles as nodes then it will become a
    path cover problem with vertex disjoint on DAG
3.4 HopCroftKarp [67 lines] - fac9fc
/*. Finds Maximum Matching In a bipartite graph
  . Complexity O(E\sqrt{V})
  .1-indexed
  .No default constructor
  .add single edge for (u, v)*/
  static const int inf = 1e9;
  vector<int>matchL, matchR, dist;
  //matchL contains value of matched node for L part.
  vector<vector<int>>adj;
  HK(int n) : n(n), matchL(n + 1),
  matchR(n + 1), dist(n + 1), adj(n + 1) {
  void addEdge(int u, int v) {
    adj[u].push_back(v);
  bool bfs() {
    queue<int>q;
    for (int u = 1; u \le n; u++) {
      if (!matchL[u]) {
        dist[u] = 0;
        q.push(u);
      else dist[u] = inf;
    dist[0] = inf;///unmatched node matches with 0.
    while (!a.emptv()) {
      int u = q.front();
      q.pop();
      for (auto v : adj[u]) {
        if (dist[matchR[v]] == inf) {
          dist[matchR[v]] = dist[u] + 1;
          q.push(matchR[v]);
    return dist[0] != inf;
```

3.3 Flow [6 lines] - 6ebca7

```
bool dfs(int u) {
    if (!u) return true:
    for (auto v : adj[u]) {
      if (dist[matchR[v]] == dist[u] + 1
          && dfs(matchR[v])) {
        matchL[u] = v;
        matchR[v] = u;
        return true;
    dist[u] = inf;
    return false;
  int max_match() {
    int matching = 0;
    while (bfs()) {
      for (int u = 1; u \le n; u++) {
        if (!matchL[u])
          if (dfs(u))
            matching++;
    return matching;
};
3.5 Hungarian [116 lines] - 64902f
/* Complexity: O(n^3) but optimized
   It finds minimum cost maximum matching.
   For finding maximum cost maximum matching
   add -cost and return -matching()
   1-indexed */
struct Hungarian {
  long long c[N][N], fx[N], fy[N], d[N];
  int 1[N], r[N], arg[N], trace[N];
  queue<int> q;
  int start, finish, n;
  const long long inf = 1e18;
  Hungarian() {}
  Hungarian(int n1, int n2) : n(max(n1, n2)) {
   for (int i = 1; i <= n; ++i) {
      fv[i] = 1[i] = r[i] = 0;
      for (int j = 1; j \le n; ++j) c[i][j] = inf;
  void add_edge(int u, int v, long long cost) {
    c[u][v] = min(c[u][v], cost);
  inline long long getC(int u, int v) {
    return c[u][v] - fx[u] - fv[v]:
  void initBFS() {
    while (!q.empty()) q.pop();
    q.push(start);
    for (int i = 0; i <= n; ++i) trace[i] = 0;
    for (int v = 1; v \le n; ++v) {
      d[v] = getC(start, v);
      arg[v] = start;
    finish = 0;
  void findAugPath() {
    while (!q.empty()) {
      int u = q.front();
```

```
for (int v = 1; v \le n; ++v) if (!trace[v]) {
      long long w = getC(u, v);
      if (!w) {
        trace[v] = u;
        if (!r[v]) {
          finish = v;
          return;
        q.push(r[v]);
      if (d[v] > w) {
        d[v] = w;
        arg[v] = u;
 }
}
void subX_addY() {
 long long delta = inf;
  for (int v = 1; v \le n; ++v) if (trace[v] == 0 &&
      d[v] < delta) {</pre>
    delta = d[v];
  // Rotate
  fx[start] += delta:
  for (int v = 1: v \le n: ++v) if (trace[v]) {
    int u = r[v]:
    fv[v] -= delta:
    fx[u] += delta;
  else d[v] -= delta;
  for (int v = 1; v \le n; ++v) if (!trace[v] && !d[v])
    trace[v] = arg[v];
    if (!r[v]) {
      finish = v;
      return;
    q.push(r[v]);
void Enlarge() {
  do {
    int u = trace[finish]:
    int nxt = l[u]:
    l[u] = finish:
    r[finish] = u:
    finish = nxt:
  } while (finish);
long long maximum_matching() {
  for (int u = 1; u <= n; ++u) {
    fx[u] = c[u][1]:
    for (int v = 1; v \le n; ++v) {
      fx[u] = min(fx[u], c[u][v]);
 for (int v = 1; v \le n; ++v) {
    fv[v] = c[1][v] - fx[1];
    for (int u = 1; u <= n; ++u) {
      fy[v] = min(fy[v], c[u][v] - fx[u]);
```

```
for (int u = 1: u <= n: ++u) {
     start = u:
     initBFS():
     while (!finish) {
       findAugPath();
       if (!finish) subX_addY();
     Enlarge();
    long long ans = 0;
    for (int i = 1; i <= n; ++i) {
     if (c[i][l[i]] != inf) ans += c[i][l[i]];
     else l[i] = 0;
   return ans;
3.6 MCMF [116 lines] - 466389
/*Credit: ShahjalalShohaq
  .Works for both directed, undirected and with negative
      cost too
  .doesn't work for negative cycles
  .for undirected edges just make the directed flag
  . Complexity: O(\min(E^2 *V \log V, E \log V * flow))*/
using T = long long;
const T inf = 1LL << 61;
struct MCMF {
 struct edge {
   int u, v;
   T cap, cost;
    edge(int _u, int _v, T _cap, T _cost, int _id) {
     v = v;
     cap = _cap;
     cost = _cost;
     id = _id;
 };
  int n, s, t, mxid;
 T flow, cost;
 vector<vector<int>> g;
 vector<edge> e;
 vector<T> d, potential, flow_through;
 vector<int> par;
 bool neg;
 MCMF() {}
 MCMF(int _n) { // O-based indexing
   n = n + 10:
   g.assign(n, vector<int>());
   neg = false;
   mxid = 0:
  void add_edge(int u, int v, T cap, T cost, int id =
     -1, bool directed = true) {
    if (cost < 0) neg = true;
    g[u].push_back(e.size());
    e.push_back(edge(u, v, cap, cost, id));
    g[v].push_back(e.size());
    e.push_back(edge(v, u, 0, -cost, -1));
    mxid = max(mxid, id);
    if (!directed) add_edge(v, u, cap, cost, -1, true);
```

};

```
bool dijkstra() {
  par.assign(n, -1);
  d.assign(n, inf);
  priority_queue<pair<T, T>, vector<pair<T, T>>,
      greater<pair<T, T>> > q;
  d[s] = 0;
  q.push(pair<T, T>(0, s));
  while (!q.empty()) {
   int u = q.top().second;
   T nw = q.top().first;
    q.pop();
    if (nw != d[u]) continue;
    for (int i = 0; i < (int)g[u].size(); i++) {
      int id = g[u][i];
      int v = e[id].v;
      T cap = e[id].cap;
      T w = e[id].cost + potential[u] - potential[v];
      if (d[u] + w < d[v] && cap > 0) {
        d[v] = d[u] + w:
        par[v] = id;
        q.push(pair<T, T>(d[v], v));
  for (int i = 0; i < n; i++) { // update potential
    if (d[i] < inf) potential[i] += d[i];</pre>
  return d[t] != inf:
T send_flow(int v, T cur) {
  if (par[v] == -1) return cur;
  int id = par[v];
  int u = e[id].u;
  T w = e[id].cost;
  T f = send_flow(u, min(cur, e[id].cap));
  cost += f * w:
  e[id].cap -= f;
  e[id ^1].cap += f;
  return f;
//returns {maxflow, mincost}
pair<T, T> solve(int _s, int _t, T goal = inf) {
 s = s:
 t = t:
  flow = 0, cost = 0;
  potential.assign(n, 0);
  if (neg) {
    // run Bellman-Ford to find starting potential
    d.assign(n, inf);
    for (int i = 0, relax = true; i < n \&\& relax; i++)
      for (int u = 0; u < n; u++) {
        for (int k = 0; k < (int)g[u].size(); k++) {</pre>
          int id = g[u][k];
          int v = e[id].v;
          T cap = e[id].cap, w = e[id].cost;
          if (d[v] > d[u] + w && cap > 0) {
            d[v] = d[u] + w;
            relax = true;
```

```
for (int i = 0; i < n; i++) if (d[i] < inf)
          potential[i] = d[i]:
    while (flow < goal && dijkstra()) flow +=
        send_flow(t, goal - flow);
    flow_through.assign(mxid + 10, 0);
    for (int u = 0; u < n; u++) {
      for (auto v : g[u]) {
        if (e[v].id >= 0) flow_through[e[v].id] = e[v
    return make_pair(flow, cost);
};
4 Game Theory
4.1 Points to be noted [14 lines] - 6fe124
>[First Write a Brute Force solution]
>Nim = all xor
>Misere Nim = Nim + corner case: if all piles are 1,
    reverse(nim)
>Bogus Nim = Nim
>Staircase Nim = Odd indexed pile Nim (Even indexed pile
    doesnt matter, as one player can give bogus moves to
    drop all even piles to ground)
>Sprague Grundy: [Every impartial game under the normal
    play convention is equivalent to a one-heap game of
Every tree = one nim pile = tree root value; tree leaf
    value = 0; tree node value = mex of all child nodes.
[Careful: one tree node can become multiple new tree
    roots(multiple elements in one node), then the value
    of that node = xor of all those root values]
>Hackenbush(Given a rooted tree; cut an edge in one
    move; subtree under that edge gets removed; last
    player to cut wins):
Colon: //G(u) = (G(v1) + 1) \oplus (G(v2) + 1) \oplus \cdots [v1, v2, \cdots]
    are childs of ul
For multiple trees ans is their xor
>Hackenbush on graph (instead of tree given an rooted
fusion: All edges in a cycle can be fused to get a tree
    structure; build a super node, connect some single
    nodes with that super node, number of single nodes
    is the number of edges in the cycle.
Sol: [Bridge component tree] mark all bridges, a group
    of edges that are not bridges, becomes one component
    and contributes number of edges to the hackenbush.
    (even number of edges contributes 0, odd number of
    edges contributes 1)
5 Geometry
5.1 Geometry [384 lines] - 6bfd7b
namespace Geometry
  #define M_PI(acos(-1.0))
  double eps=1e-8;
  typedef double T; //coordinate point type
  struct pt //Point
```

T x, y;

pt(){}

```
pt(T_x,T_y):x(_x),y(_y){}
  pt operator+(pt p){
    return{x+p.x,y+p.y};
  pt operator-(pt p){
    return{x-p.x,y-p.y};
  pt operator*(T d){
    return{x*d, y*d};
  pt operator*(pt d){/*I added for General linear
       transformation, not sure about that function*/
    return{x*d.x,y*d.y};
  pt operator/(T d){
    return{x/d,y/d};/*only for floating point*/
  pt operator/(pt d){/*I added for General linear
       transformation, not sure about that function*/
    return\{x/d.x,y/d.y\};
  bool operator<(const pt% p)const {</pre>
    if(x!=p.x)
      return x<p.x;
    return y<p.y;
  bool operator==(pt b){
    return x==b.x \&\& y==b.y;
  bool operator!=(pt b){
    return!(*(this)==b);
  friend ostream& operator << (ostream& os, const pt p) {
    return os<<"("<<p.x<<","<<p.y<<")";
  friend istream& operator>>(istream& is,pt &p){
    is>>p.x>>p.y;
    return is;
};
T sq(pt p){
  return p.x*p.x+p.y*p.y;
double Abs(pt p){
  return sqrtl(sq(p));
pt translate(pt v,pt p){ /*To translate an object by a
    vector v*/
  return p+v;
pt scale(pt c,double factor,pt p){/*To scale an object
    by a certain ratio factor around a center*/
  return c+(p-c)*factor;
pt rot(pt p,double a) {/*To rotate a point by angle
  return{p.x*cos(a)-p.y*sin(a),p.x*sin(a)+p.y*
       cos(a)};
pt perp(pt p){/*To rotate a point 90 degree*/
  return{-p.y,p.x};
pt linearTransfo(pt p,pt q,pt r,pt fp,pt fq){/*so far
     don't know about that function*/
```

```
return fp+(r-p)*(fq-fp)/(q-p);
T dot(pt v,pt w){
  return v.x*w.x+v.y*w.y;
bool isPerp(pt v,pt w){
  return dot(v,w)==0;
double angle(pt v,pt w){/*Find the smallest angle of
     two vector*/
  double cosTheta=dot(v,w)/Abs(v)/Abs(w);
  return acos(max(-1.0,min(1.0,cosTheta)));
T cross(pt v,pt w){
  return v.x*w.y-v.y*w.x;
T orient(pt a,pt b,pt c){
  return cross(b-a,c-a); /*if c is left side+ve,c is
      right side-ve.on line 0*/
bool inAngle(pt a,pt b,pt c,pt p){/*if p is in the
  assert(orient(a,b,c)!=0);
   if(orient(a,b,c)<0)</pre>
     swap(b,c):
  return orient(a,b,p)>=0 && orient(a,c,p)<=0;
double orientedAngle(pt a,pt b,pt c){/*the actual
     angle from ab to ac*/
   if(orient(a,b,c)>=0)
      return angle(b-a,c-a);
   else
      return 2*M_PI-angle(b-a,c-a);
///line
struct line{
  pt v;
  Tc;
  line(){}
  line(pt p,pt q){/*From points P and Q*/
    v=(q-p), this->c=cross(v,p);
  line(T a,T b,T c){/*From equation ax+by=c*/
     v=pt(b,-a),this->c=c;
  line(pt v,T c){/*From direction vector v and offset
    this->v=v,this->c=c;
  double getY(double x){/*self made, not sure if it is
      okau*/
     assert(v.x!=0):
     double ret=(double)(c+v.y*x)/v.x;
    return ret:
  double getX(double y){/*self made, not sure if it is
       okay*/
     assert(v.v!=0);
     double ret=(double)(c-v.x*y)/-v.y;
    return ret;
  T side(pt p){/*which side a point is*/
      return cross(v,p)-c;
```

```
double dist(pt p){/*point to line dist*/
     return abs(side(p))/Abs(v);
   double sqDist(pt p){/*square dist*/
     return side(p)*side(p)/(double)sq(v);
   line perpThrough(pt p){/*perpendicular line with
       point p*/
       return line(p,p+perp(v));
   bool cmpProj(pt p,pt q){/*compare function to sort
       points on a line*/
       return dot(v,p)<dot(v,q);
   line translate(pt t){/*translate with vector t*/
       return line(v,c+cross(v,t));
   line shiftLeft(double dist){/*translate with
       distance dist*/
       return line(v,c+dist*Abs(v));
   }
   pt proj(pt p){
       return p-perp(v)*side(p)/sq(v);
   pt refl(pt p){
       return p-perp(v)*2*side(p)/sq(v);
};
 bool areParallel(line 11.line 12){
   return(l1.v.x*l2.v.y==l1.v.y*l2.v.x);
 bool areSame(line 11,line 12){
   return areParallel(11,12) and(11.v.x*12.c==12.v.x*
       11.c) and (11.v.v*12.c==12.v.v*11.c);
 bool inter(line 11,line 12,pt& out){
  T d=cross(l1.v,l2.v);
   if(d==0)return false;
   out=(12.v*11.c-11.v*12.c)/d;
   return true;
 line intBisector(line 11,line 12,bool interior){/*if
     change sign then returns the other one*/
   assert(cross(11.v.12.v)!=0):
   double sign=interior?1:-1:
   return line(12.v/Abs(12.v)+11.v*sign/Abs(11.v),
           12.c/Abs(12.v)+11.c*sign/Abs(11.v));
 //segment
 bool inDisk(pt a,pt b,pt p){/*check weather point p is
     in diameter AB*/
   return dot(a-p,b-p) \le 0;
 bool onSegment(pt a,pt b,pt p){/*check weather point p
     is in segment AB*/
   return orient(a,b,p)==0 and inDisk(a,b,p);
 bool properInter(pt a,pt b,pt c,pt d,pt% i){
   double oa=orient(c,d,a),
          ob=orient(c,d,b),
          oc=orient(a,b,c),
          od=orient(a,b,d);
 //Proper intersection exists iff opposite signs
   if (oa*ob<0) and oc*od<0)
```

```
i=(a*ob-b*oa)/(ob-oa):
    return 1:
  return 0;
/*To create sets of points we need a comparison
    function*/
 struct cmpX{
  bool operator()(pt a,pt b){
       return make_pair(a.x,a.y) < make_pair(b.x,b.y);
};
set<pt,cmpX>inters(pt a,pt b,pt c,pt d){
   if(properInter(a,b,c,d,out))
    return{out};
   set<pt,cmpX>s;
   if(onSegment(c,d,a))s.insert(a);
   if(onSegment(c,d,b))s.insert(b);
   if(onSegment(a,b,c))s.insert(c);
   if(onSegment(a,b,d))s.insert(d);
   return s:
 bool LineSegInter(line 1,pt a,pt b,pt& out){
   if(l.side(a)*l.side(b)>eps)return 0;
  return inter(1,line(a,b),out);
 double segPoint(pt a,pt b,pt p){/*returns distance
     from a point p to segment AB*/
   if(a!=b){}
       line l(a,b);
       if(l.cmpProj(a,p)and l.cmpProj(p,b))
         return l.dist(p);
   return min(Abs(p-a),Abs(p-b));
double segSeg(pt a,pt b,pt c,pt d){/*returns distance
     from a segment AB to segment CD*/
   if(properInter(a,b,c,d,dummy))return 0;
   return min(min(min(segPoint(a,b,c),segPoint(a,b,
       d)),segPoint(c,d,a)),segPoint(c,d,b));
/*int latticePoints(pt a,pt b){
  // requires int representation
  return = qcd(abs(a.x-b.x), abs(a.y-b.y))+1;
 \frac{1}{A} = i + (b/2) - 1: here
     A=area, i=pointsinside, b=pointsonline
 Polygon*/
bool isConvex(vector<pt>&p){
   bool hasPos=0,hasNeg=0;
   for(int i=0,n=p.size();i<n;i++){</pre>
    int o=orient(p[i],p[(i+1)%n],p[(i+2)%n]);
    if(o>0)hasPos=1:
    if(o<0)hasNeg=true;</pre>
   return! (hasPos and hasNeg);
 double areaTriangle(pt a,pt b,pt c){
   return abs(cross(b-a,c-a))/2.0;
 double areaPolygon(const vector<pt>&p){
  double area=0.0;
   for(int i=0,n=p.size();i<n;i++){</pre>
```

```
area+=cross(p[i],p[(i+1)\%n]);
   return fabs(area)/2.0;
 bool pointInPolygon(const vector<pt>&p,pt q){/*returns
      true if pt q is in polygon p*/
   bool c=false;
   for(int i=0,n=p.size();i<n;i++){</pre>
      int j=(i+1)%p.size();
      if((p[i].y \le q.y \text{ and } q.y \le p[j].y \text{ or } p[j].y \le q.y \text{ and}
          q.y < p[i].y) and
        q.x < p[i].x+(p[j].x-p[i].x)*(q.y-p[i].y)/
            (p[j].y-p[i].y))
   }
   return c;
 ll is_point_in_convex(vector<pt>& p, pt &x) { // O(log
     11 n = p.size(); /*this function from
          YouKnowWho*/
     if (n < 3) return 1;
     ll a =orient(p[0], p[1], x), b = orient(p[0], p[n]
          -1], x);
      if (a < 0 | | b > 0) return 1:
     11 1 = 1, r = n - 1;
      while (1 + 1 < r) {
          int mid = 1 + r \gg 1;
          if (p[0], p[mid], x) >= 0) 1 = mid;
          else r = mid;
     11 k = orient(p[1], p[r], x);
     if (k \le 0) return -k;
     if (1 == 1 && a == 0) return 0;
     if (r == n - 1 \&\& b == 0) return 0;
      return -1;
 pt centroidPolygon(vector<pt>&p){/*from rezaul, i don't
      know about that*/
   pt c(0,0);
   double scale=6.0*areaPolygon(p);
// if(scale<eps)return c;</pre>
   for(int i=0,n=p.size();i<n;i++){</pre>
     int j=(i+1)%n;
     c=c+(p[i]+p[j])*cross(p[i],p[j]);
   return c/scale;
///Circle
 pt circumCenter(pt a,pt b,pt c){/*return the center of
      the circle go through point a,b,c*/
   b=b-a.c=c-a:
   assert(cross(b,c)!=0);
   return a+perp(b*sq(c)-c*sq(b))/cross(b,c)/2;
 bool circle2PtsRad(pt p1,pt p2,double r,pt& c){
   double d2=sq(p1-p2);
   double det=r*r/d2-0.25;
   if(det<0.0)return false;
   double h=sqrt(det);
   c.x=(p1.x+p2.x)*0.5+(p1.y-p2.y)*h;
   c.y=(p1.y+p2.y)*0.5+(p2.x-p1.x)*h;
   return true;
```

```
int circleLine(pt c,double r,line l,pair<pt,pt>&
     out) {/*circle line intersection*/
   double h2=r*r-l.sqDist(c);
   if(h2<0)return 0; /*the line doesn't touch the
       circle:*/
   pt p=1.proj(c);
  pt h=1.v*sqrt(h2)/Abs(1.v);
   out=make_pair(p-h,p+h);
  return 1+(h2>0);
int circleCircle(pt c1,double r1,pt c2,double
     r2,pair<pt,pt>& out){/*circle circle
     intersection*/
   pt d=c2-c1;
   double d2=sq(d);
   if (d2==0) {//concentric circles
     assert(r1!=r2);
     return 0:
   double pd=(d2+r1*r1-r2*r2)/2;
   double h2=r1*r1-pd*pd/d2;//h ^ 2
   if(h2<0)return 0:
   pt p=c1+d*pd/d2, h=perp(d)*sqrt(h2/d2);
   out=make_pair(p-h,p+h);
  return 1+h2>0:
 int tangents(pt c1,double r1,pt c2,double r2,bool
     inner,vector<pair<pt,pt>>&out){
   if(inner)r2=-r2;/*returns tangent(the line which
       touch a circle in one point) of two circle*/
  pt d=c2-c1;/*the same code can be used to find the
       tangent to a circle passing through a point by
       setting r2 to 0*/
   double dr=r1-r2,d2=sq(d),h2=d2-dr*dr;
   if(d2==0 \text{ or } h2<0){}
     assert(h2!=0);
     return 0;
  for(int sign :{-1,1}){
       pt v=pt(d*dr+perp(d)*sqrt(h2)*sign)/d2;
       out.push_back(make_pair(c1+v*r1,c2+v*r2));
  return 1+(h2>0);
//Convex Hull-Monotone Chain
pt H[100000+5]:
vector<pt>monotoneChain(vector<pt>&points){
   sort(points.begin(),points.end());
   vector<pt>ret;
  ret.clear();
  for(int i=0,sz=points.size();i<sz;i++){</pre>
     while(st>=2 and
         orient(H[st-2],H[st-1],points[i])<0)st--;</pre>
     H[st++]=points[i];
   int taken=st-1;
  for(int i=points.size()-2;i>=0;i--){
     while(st>=taken+2 and
         orient(H[st-2],H[st-1],points[i])<0)st--;
     H[st++]=points[i];
  for(int i=0;i<st;i++)ret.push_back(H[i]);</pre>
  return ret;
```

```
KUET_Effervescent Team Notebook - September 11, 2024
```

```
rMat[1][3]=0.0:
                                                              rMat[2][0] = (u*w*(1-cos(ang))-v*sqrt(L)*sin(ang))/L;
  //Convex Hull-Monotone Chain from you_know_who
                                                              rMat[2][1] = (v*w*(1-cos(ang)) + u*sqrt(L)*sin(ang))/L;
  vector<pt> monotoneChain(vector<pt> &v) {
                                                              rMat[2][2]=(w2 + (u2 + v2) * cos(ang)) / L;
      if(v.size()==1) return v;
                                                              rMat[2][3]=0.0; rMat[3][0]=0.0; rMat[3][1]=0.0;
      sort(v.begin(), v.end());
                                                              rMat[3][2]=0.0; rMat[3][3]=1.0;
      vector<pt> up(2*v.size()+2), down(2*v.size()+2);
      int szup=0, szdw=0;
                                                            /*double ang:
                                                              double u, v, w; //points = the point to be rotated
      for(int i=0;i<v.size();i++) {
          while(szup>1 && orient(up[szup-2],
                                                              Point point, rotated; //u,v,w=unit vector of line
              up[szup-1], v[i])>=0)
                                                              inMat[0][0] = points.x; inMat[1][0] = points.y;
              szup--;
                                                              inMat[2][0] = points.z; inMat[3][0] = 1.0;
          while(szdw>1 && orient(down[szdw-2],
                                                              setMat(ang, u, v, w); mulMat();
                                                              rotated.x = outMat[0][0]; rotated.y = outMat[1][0];
              down[szdw-1], v[i]) <= 0
                                                              rotated.z = outMat[2][0]:*/
              szdw--;
          up[szup++]=v[i];
          down[szdw++]=v[i];
                                                            6 Graph
                                                            6.1 2SAT [92 lines] - 5289ec
      if(szdw>1) szdw--:
                                                            struct TwoSat {
      reverse(up.begin(), up.begin()+szup);
                                                              vector<bool>vis:
      for(int i=0;i<szup-1;i++) down[szdw++] = up[i];</pre>
                                                              vector<vector<int>>adj, radj;
      if(szdw==2 && down[0].x==down[1].x &&
                                                              vector<int>dfs_t, ord, par;
          down[0].y==down[1].y)
                                                              int n, intime; //For n node there will be 2*n node in
          szdw--:
                                                                  SAT.
      sz = szdw:
                                                              void init(int N) {
      return down:
                                                               n = N:
                                                                intime = 0:
 double cosA(double a,double b,double c){
                                                                vis.assign(N * 2 + 1, false);
      double val=b*b+c*c-a*a:
                                                                adj.assign(N * 2 + 1, vector<int>());
      val/=(2*b*c);
                                                                radj.assign(N * 2 + 1, vector\langle int \rangle());
      return acos(val);
                                                                dfs_t.resize(N * 2 + 1);
                                                                ord.resize(N * 2 + 1);
  double triangle(double a,double b,double c){
                                                                par.resize(N * 2 + 1);
      double s=(a+b+c)/2;
      return sqrtl(s*(s-a)*(s-b)*(s-c));
                                                              inline int neg(int x) {
                                                                return x \le n ? x + n : x - n;
using namespace Geometry;
                                                              inline void add_implication(int a, int b) {
5.2 Rotation Matrix [39 lines] - f97f03
                                                                if (a < 0) a = n - a;
struct { double x; double v; double z; } Point;
                                                                if (b < 0) b = n - b;
double rMat[4][4];
                                                                adj[a].push_back(b);
double inMat[4][1] = {0.0, 0.0, 0.0, 0.0};
                                                                radj[b].push_back(a);
double outMat[4][1] = {0.0, 0.0, 0.0, 0.0};
void mulMat() {
                                                              inline void add_or(int a, int b) {
 for(int i = 0; i < 4; i++){
                                                                add_implication(-a, b);
    for(int j = 0; j < 1; j++){
                                                                add_implication(-b, a);
      outMat[i][j] = 0;
      for(int k = 0; k < 4; k++)
                                                              inline void add_xor(int a, int b) {
        outMat[i][j] += rMat[i][k] * inMat[k][j];
                                                                add_or(a, b);
                                                                add or (-a, -b):
                                                              inline void add and(int a. int b) {
void setMat(double ang, double u, double v, double w){
                                                                add or(a, b):
  double L = (u * u + v * v + w * w):
                                                                add_or(a, -b);
  ang = ang * PI / 180.0; /*converting to radian
                                                                add_or(-a, b);
      value*/
  double u2 = u*u; double v2 = v*v; double w2 = w*w;
                                                              inline void force_true(int x) {
  rMat[0][0]=(u2+(v2+w2)*cos(ang))/L;
                                                                if (x < 0) x = n - x;
 rMat[0][1]=(u*v*(1-cos(ang))-w*sqrt(L)*sin(ang))/L;
                                                                add_implication(neg(x), x);
 rMat[0][2]=(u*w*(1-cos(ang))+v*sqrt(L)*sin(ang))/L;
                                                              inline void add_xnor(int a, int b) {
 rMat[0][3]=0.0;
 rMat[1][0] = (u*v*(1-cos(ang))+w*sqrt(L)*sin(ang))/L;
                                                                add_or(a, -b);
 rMat[1][1]=(v2+(u2+w2)*cos(ang))/L;
                                                                add_or(-a, b);
  rMat[1][2]=(v*w*(1-cos(ang))-u*sqrt(L)*sin(ang))/L;
```

```
inline void add_nand(int a, int b) {
    add or (-a, -b):
  inline void add_nor(int a, int b) {
    add_and(-a, -b);
  inline void force_false(int x) {
    if (x < 0) x = n - x;
    add_implication(x, neg(x));
  inline void topsort(int u) {
    vis[u] = 1:
    for (int v : radj[u]) if (!vis[v]) topsort(v);
    dfs_t[u] = ++intime;
  inline void dfs(int u, int p) {
    par[u] = p, vis[u] = 1;
    for (int v : adj[u]) if (!vis[v]) dfs(v, p);
  void build() {
    int i. x:
    for (i = n * 2, intime = 0; i >= 1; i--) {
      if (!vis[i]) topsort(i);
      ord[dfs_t[i]] = i;
    vis.assign(n * 2 + 1, 0);
    for (i = n * 2:i > 0:i--) {
      x = ord[i]:
      if (!vis[x]) dfs(x, x);
  bool satisfy(vector<int>& ret)//ret contains the value
      that are true if the graph is satisfiable.
    build();
    vis.assign(n * 2 + 1, 0);
    for (int i = 1; i \le n * 2; i++) {
      int x = ord[i];
      if (par[x] == par[neg(x)]) return 0;
      if (!vis[par[x]]) {
        vis[par[x]] = 1;
        vis[par[neg(x)]] = 0;
    for (int i = 1;i <= n;i++) if (vis[par[i]])
        ret.push_back(i);
    return 1:
};
6.2 BridgeTree [66 lines] - f8e197
int N, M, timer, compid;
vector<pair<int, int>> g[mx];
bool used[mx], isBridge[mx];
int comp[mx], tin[mx], minAncestor[mx];
vector<int> Tree[mx]; // Store 2-edge-connected
    component tree. (Bridge tree).
void markBridge(int v, int p) {
  tin[v] = minAncestor[v] = ++timer;
  used[v] = 1;
  for (auto& e : g[v]) {
    int to, id;
```

tie(to, id) = e;

12

```
if (to == p) continue;
    if (used[to]) minAncestor[v] = min(minAncestor[v].
        tin[to]):
    else {
      markBridge(to, v);
      minAncestor[v] = min(minAncestor[v],
          minAncestor[to]);
      if (minAncestor[to] > tin[v]) isBridge[id] = true;
      // if (tin[u] \leq minAncestor[v]) ap [u] = 1;
void markComp(int v, int p) {
 used[v] = 1;
  comp[v] = compid;
  for (auto& e : g[v]) {
    int to, id;
    tie(to. id) = e:
    if (isBridge[id]) continue;
    if (used[to]) continue;
    markComp(to, v);
}
vector<pair<int, int>> edges;
void addEdge(int from, int to, int id) {
  g[from].push_back({ to, id });
  g[to].push_back({ from, id });
  edges[id] = { from, to };
void initB() {
  for (int i = 0; i <= compid; ++i) Tree[i].clear();</pre>
  for (int i = 1; i <= N; ++i) used[i] = false;
  for (int i = 1; i <= M; ++i) isBridge[i] = false;</pre>
  timer = compid = 0;
void bridge_tree() {
  initB():
  markBridge(1, -1); //Assuming graph is connected.
  for (int i = 1; i \le N; ++i) used[i] = 0;
  for (int i = 1; i \le N; ++i) {
    if (!used[i]) {
      markComp(i, -1);
      ++compid;
  for (int i = 1; i <= M; ++i) {
    if (isBridge[i]) {
      int u, v;
      tie(u, v) = edges[i];
      // connect two componets using edge.
      Tree[comp[u]].push_back(comp[v]);
      Tree[comp[v]].push_back(comp[u]);
      int x = comp[u];
      int y = comp[v];
6.3 Centroid Decomposition [39 lines] - d5d02b
11 n,subsize[mx];
vector<int>adj[mx];
bool b[mx];
int cpar[mx];
vector<int>ctree[mx];
```

```
void calculatesize(ll u,ll par){
  subsize[u]=1:
  for(ll i=0;i<(ll)adj[u].size();i++){
    ll v=adi[u][i]:
    if(v==par or b[v]==true)continue;
    calculatesize(v,u);
    subsize[u]+=subsize[v];
11 getcentroid(ll u,ll par,ll n){
 ll ret=u;
  for(ll i=0;i<(ll)adj[u].size();i++){
    11 v=adj[u][i];
    if(v==par or b[v]==true)continue;
    if(subsize[v]>(n/2)){
      ret=getcentroid(v,u,n);
      break;
 return ret;
void decompose(ll u, int p){
  calculatesize(u,-1);
  11 c=getcentroid(u,-1,subsize[u]);
  b[c]=true:
  cpar[c] = p;
  //if(p != -1)ctree[p].push_back(c);
  for(ll i=0;i<(ll)adj[c].size();i++){
    11 v=adj[c][i];
    if(b[v] == true) continue;
    decompose(v, c);
6.4 DSU on Tree [56 lines] - 391fb6
//extra data you need
vector<int> adj[mxn];
vector<int> *dsu[mxn];
void call(int u, int p=-1){
  sz[u] = 1;
  for(auto v: adj[u]){
   if(v != p){
  dep[v] = dep[u]+1;
      call(v, u);
      sz[u] += sz[v];
void dfs(int u, int p = -1, int isb = 1){
  int mx=-1. big=-1:
 for(auto v: adj[u]){
    if(v != p \&\& sz[v]>mx){
      mx = sz[v];
      big = v;
  for(auto v: adj[u]){
    if(v != p && v != big){
      dfs(v, u, 0);
  if(big != -1){
    dfs(big, u, 1);
    dsu[u] = dsu[big];
```

```
else{
    dsu[u] = new vector<int>();
  dsu[u]->push_back(u);
  //calculation
  for(auto v: adi[u]){
   if (v == p \mid | v == big) continue;
   for(auto x: *dsu[v]){
      dsu[u]->push_back(x);
      //calculation
  //calculate ans for node u
 if(isb == 0){
   for(auto x: *dsu[u]){
      //reverse calculation
int main() {
 //input graph
 dep[1] = 1;
 call(1);
 dfs(1):
6.5 Heavy Light Decomposition [73 lines] - d0e24f
/*Heavy Light Decomposition
Build Complexity O(n)
Query Complexity O(lq^2 n)
Call init() with number of nodes
It's probably for the best to not do"using namespace
namespace hld {
 //N is the maximum number of nodes
 /*par, lev, size corresponds to
      parent, depth, subtree-size*/
  //head[u] is the starting node of the chain u is in
  //in[u] to out[u] keeps the subtree indices
  const int N=100000+7;
  vector<int>g[N]:
  int par[N],lev[N],head[N],size[N],in[N],out[N];
 int cur_pos,n;
  //returns the size of subtree rooted at u
  /*maintains the child with the largest subtree at the
      front of q[u]*/
  //WARNING: Don't change anything here specially with
      size[]if Jon Snow
  int dfs(int u,int p){
    size[u]=1,par[u]=p;
   lev[u] = lev[p] + 1;
   for(auto &v : g[u]){
      if (v==p) continue;
      size[u] += dfs(v,u);
      if(size[v]>size[g[u].front()]){
        swap(v,g[u].front());
   return size[u];
  //decomposed the tree in an array
  //note that there is no physical array here
```

```
void decompose(int u,int p){
                                                                 for(i=0;i<deg[u.v];i++){
    in[u]=++cur pos:
                                                                   v.v=adi[u.v][i].v:
    for(auto &v : g[u]){
                                                                   int cost=adj[u.v][i].w+u.w;
      if (v==p) continue;
                                                                   for(v.k=u.k;v.k<K;v.k++){
      head[v] = (v = g[u].front()? head[u]: v);
                                                                     if(cost==inf)break;
      decompose(v,u);
                                                                     if(val[v.v][v.k]>cost){
                                                                       swap(cost,val[v.v][v.k]);
    out[u]=cur_pos;
                                                                       v.w=val[v.v][v.k];
                                                                       Q.push(v);
  //initializes the structure with _n nodes
                                                                       break;
  void init(int _n,int root=1){
    n=_n;
    cur_pos=0;
                                                                   for(v.k++;v.k<K;v.k++){
    dfs(root,0);
                                                                     if(cost==inf)break;
    head[root]=root;
    decompose(root,0);
                                                              }
  //checks whether p is an ancestor of u
 bool isances(int p,int u){
    return in[p] <= in[u] and out[u] <= out[p];
                                                             6.7 LCA [46 lines] - 9de12b
                                                             const int Lg = 22;
  //Returns the maximum node value in the path u-v
                                                             vector<int>adj[mx];
 11 query(int u,int v){
                                                             int level[mx];
    11 ret=-INF:
                                                             int dp[Lg][mx];
    while(!isances(head[u].v)){
                                                             void dfs(int u) {
      ret=max(ret,seg.query(1,1,n,in[head[u]],in[u]));
                                                               for (int i = 1; i < Lg; i++)
      u=par[head[u]]:
                                                               for (int v : adj[u]) {
    swap(u,v);
                                                                 if (dp[0][u] == v)continue;
    while(!isances(head[u],v)){
                                                                 level[v] = level[u] + 1;
      ret=max(ret,seg.query(1,1,n,in[head[u]],in[u]));
                                                                 dp[0][v] = u;
      u=par[head[u]];
                                                                 dfs(v);
    if(in[v]<in[u])swap(u,v);</pre>
    ret=max(ret,seg.query(1,1,n,in[u],in[v]));
                                                             int lca(int u, int v) {
    return ret;
                                                               int diff = level[v] - level[u];
  //Adds val to subtree of u
                                                               for (int i = 0;i < Lg;i++)</pre>
 void update(int u,ll val){
                                                                 if (diff & (1 << i))
    seg.update(1,1,n,in[u],out[u],val);
                                                                   v = dp[i][v];
                                                               for (int i = Lg - 1; i >= 0; i--)
                                                                 if (dp[i][u] != dp[i][v])
                                                                   u = dp[i][u], v = dp[i][v];
6.6 K'th Shortest path [40 lines] - 9f3788
int m,n,deg[MM],source,sink,K,val[MM][12];
                                                              return u == v ? u : dp[0][u];
struct edge{
                                                             int kth(int u, int k) {
 int v,w;
}adj[MM] [500];
                                                              for (int i = Lg - 1; i >= 0; i--)
struct info{
                                                                 if (k & (1 << i))
                                                                   u = dp[i][u];
  int v,w,k;
 bool operator<(const info &b)const{</pre>
                                                              return u:
    return w>b.w:
                                                             //kth node from u to v. Oth is u.
                                                             int go(int u, int v, int k) {
                                                              int 1 = lca(u, v);
priority_queue<info, vector<info>>Q;
void kthBestShortestPath(){
                                                               assert(k <= d);</pre>
  int i,j;
 info u,v;
  for(i=0;i<n;i++)
                                                              k -= level[u] - level[l];
    for(j=0;j<K;j++)val[i][j]=inf;
  u.v=source,u.k=0,u.w=0;
  Q.push(u);
  while(!Q.empty()){
                                                                LCA(u,v) with root r:
    u=Q.top();
                                                                lca(u,v)^{l}ca(u,r)^{l}ca(v,r)
    Q.pop();
                                                                Distance between u,v:
```

```
level(u) + level(v) - 2*level(lca(u,v))
                                                          6.8 SCC [43 lines] - 4da431
                                                          /*components: number of SCC.
                                                          sz: size of each SCC.
                                                          comp: component number of each node.
                                                          Create reverse graph.
                                                          Run find_scc() to find SCC.
                                                          Might need to create condensation graph by
                                                              create_condensed().
                                                          Think about indeg/outdeg
                                                          for multiple test cases- clear
                                                               adj/radj/comp/vis/sz/topo/condensed.*/
      if(val[v.v][v.k]>cost)swap(cost, val[v.v][v.k]);
                                                          vector<int>adj[mx], radj[mx];
                                                          int comp[mx], vis[mx], sz[mx], components;
                                                          vector<int>topo:
                                                          void dfs(int u) {
                                                            vis[u] = 1:
                                                            for (int v : adj[u])
                                                              if (!vis[v]) dfs(v);
                                                            topo.push_back(u);
                                                          void dfs2(int u. int val) {
                                                            comp[u] = val:
  dp[i][u] = dp[i - 1][dp[i - 1][u]];
                                                            sz[val]++:
                                                            for (int v : radj[u])
                                                              if (comp[v] == -1)
                                                                dfs2(v, val);
                                                          void find_scc(int n) {
                                                            memset(vis, 0, sizeof vis);
                                                            memset(comp, -1, sizeof comp);
                                                            for (int i = 1;i <= n;i++)
if (level[v] < level[u])swap(u, v);</pre>
                                                              if (!vis[i])
                                                                dfs(i);
                                                            reverse(topo.begin(), topo.end());
                                                            for (int u : topo)
                                                              if (comp[u] = -1)
                                                                dfs2(u, ++components);
                                                          vector<int>condensed[mx]:
                                                          void create condensed(int n) {
                                                            for (int i = 1:i \le n:i++)
                                                              for (int v : adj[i])
                                                                if (comp[i] != comp[v])
                                                                  condensed[comp[i]].push_back(comp[v]);
                                                          }
                                                          6.9 kuhn [31 lines] - 30d06c
                                                          int n. k:
                                                          vector<vector<int>> g;
int d = level[u] + level[v] - (level[1] << 1);</pre>
                                                          vector<int> mt;
                                                          vector<bool> used;
if (level[1] + k <= level[u]) return kth(u, k);</pre>
                                                          bool try_kuhn(int v) {
return kth(v, level[v] - level[l] - k);
                                                              if (used[v])
                                                                  return false;
                                                              used[v] = true;
                                                              for (int to : g[v]) {
                                                                  if (mt[to] == -1 || try_kuhn(mt[to])) {
                                                                      mt[to] = v;
```

```
return true;
    return false;
int main() {
    //... reading the graph ...
    mt.assign(k, -1);
    for (int v = 0; v < n; ++v) {
        used.assign(n, false);
        try_kuhn(v);
    for (int i = 0; i < k; ++i)
        if (mt[i] != -1)
            printf("\frac{d}{d}", mt[i] + 1, i + 1);
7 Math
7.1 Big Sum [13 lines] - 8d9520
11 bigsum(11 a, 11 b, 11 m) {
  if (b == 0) return 0;
  ll sum: a %= m:
  if (b & 1) {
    sum = bigsum((a * a) % m, (b - 1) / 2, m);
    sum = (sum + (a * sum) % m) % m;
    sum = (1 + (a * sum) % m) % m;
    sum = bigsum((a * a) % m, b / 2, m);
    sum = (sum + (a * sum) % m) % m;
  return sum:
7.2 CRT [52 lines] - 59a568
11 ext_gcd(11 A, 11 B, 11* X, 11* Y) {
 11 x2, y2, x1, y1, x, y, r2, r1, q, r;
  x2 = 1; v2 = 0;
  x1 = 0; v1 = 1;
  for (r2 = A, r1 = B; r1 != 0; r2 = r1, r1 = r, x2 =
      x1, y2 = y1, x1 = x, y1 = y) {
    q = r2 / r1;
   r = r2 \% r1;
    x = x2 - (q * x1);
    y = y2 - (q * y1);
  *X = x2; *Y = y2;
  return r2:
/*----*/
class ChineseRemainderTheorem {
  typedef long long vlong;
  typedef pair<vlong, vlong> pll;
  /** CRT Equations stored as pairs of vector. See
      addEqation()*/
  vector<pll> equations;
  public:
  void clear() {
    equations.clear();
  /** Add equation of the form x = r \pmod{m}*/
  void addEquation(vlong r, vlong m) {
```

```
equations.push_back({ r, m });
 pll solve() {
   if (equations.size() == 0) return \{-1,-1\}; /// No
        equations to solve
   vlong a1 = equations[0].first;
   vlong m1 = equations[0].second;
   /** Initially x = a_0 \pmod{m_0}*/
   /** Merge the solution with remaining equations */
   for (int i = 1; i < equations.size(); i++) {</pre>
      vlong a2 = equations[i].first;
     vlong m2 = equations[i].second;
     vlong g = \_gcd(m1, m2);
      if (a1 % g != a2 % g) return { -1,-1 }; ///
          Conflict in equations
      /** Merge the two equations*/
      vlong p, q;
      ext_gcd(m1 / g, m2 / g, &p, &q);
      vlong mod = m1 / g * m2;
      vlong x = ((__int128)a1 * (m2 / g) \% mod * q \% mod
         + (__int128)a2 * (m1 / g) % mod * p % mod) %
      /** Merged equation*/
     a1 = x:
     if (a1 < 0) a1 += mod:
     m1 = mod;
   return { a1, m1 };
7.3 FFT [85 lines] - 4ca8f0
template<typename float_t>
struct mycomplex {
 float_t x, v;
 mycomplex<float_t>(float_t _x = 0, float_t _y = 0) :
      x(_x), y(_y) {}
 float_t real() const { return x; }
 float_t imag() const { return y; }
 void real(float_t _x) { x = _x; }
 void imag(float_t _y) { y = _y; }
 mycomplex<float_t>& operator+=(const
      mycomplex<float_t> &other) { x += other.x; y +=
      other.y; return *this; }
 mycomplex<float_t>& operator==(const
     mycomplex<float_t> &other) { x -= other.x; y -=
      other.y; return *this; }
 mycomplex<float_t> operator+(const mycomplex<float_t>
      &other) const { return mycomplex < float_t > (*this)
      += other: }
 mycomplex<float_t> operator-(const mycomplex<float_t>
     &other) const { return mycomplex<float_t>(*this)
      -= other: }
 mycomplex<float_t> operator*(const mycomplex<float_t>
      &other) const {
   return {x * other.x - y * other.y, x * other.y +
       other.x * y};
 mycomplex<float_t> operator*(float_t mult) const {
   return {x * mult, y * mult};
 friend mycomplex<float_t> conj(const
     mycomplex<float_t> &c) {
   return {c.x, -c.y};
```

```
friend ostream& operator << (ostream & stream, const
      mycomplex<float_t> &c) {
   return stream << '(' << c.x << ", " << c.y << ')';
};
using cd = mycomplex<double>;
void fft(vector<cd> & a, bool invert) {
  int n = a.size();
  for (int i = 1, j = 0; i < n; i++) {
    int bit = n >> 1;
   for (; j & bit; bit >>= 1)
      j ^= bit;
    j ~= bit:
    if (i < j)
      swap(a[i], a[j]);
  for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 : 1);
    cd wlen(cos(ang), sin(ang));
    for (int i = 0; i < n; i += len) {
      cd w(1);
      for (int j = 0; j < len / 2; j++) {
        cd u = a[i+j], v = a[i+j+len/2] * w;
       a[i+i] = u + v:
       a[i+j+len/2] = u - v;
        w = w*wlen:
 if (invert) {
   for (cd & x : a){
      double z = n;
      z=1/z;
      x = x*z;
    // x /= n;
void multiply (const vector<bool> & a, const
    vector<bool> & b, vector<bool> & res) {//change all
    the bool to your type needed
  vector<cd> fa (a.begin(), a.end()), fb (b.begin(),
      b.end()):
  size t n = 1:
  while (n < max (a.size(), b.size())) n <<= 1;
 n <<= 1:
  fa.resize (n), fb.resize (n);
  fft (fa, false), fft (fb, false);
  for (size_t i=0; i<n; ++i)
   fa[i] =fa[i] * fb[i];
  fft (fa. true):
  res.resize (n):
  for (size_t i=0; i<n; ++i)</pre>
    res[i] = round(fa[i].real());
  while(res.back()==0) res.pop_back();
void pow(const vector<bool> &a, vector<bool> &res, long
    long int k){
  vector<bool> po=a;
 res.resize(1);
 res[0] = 1;
  while(k){
```

```
if(k&1){
      multiply(po, res, res);
   multiply(po, po, po);
   k/=2;
7.4 GaussElimination [39 lines] - aa53e0
template<typename ld>
int gauss(vector<vector<ld>>& a, vector<ld>& ans) {
 const ld EPS = 1e-9;
 int n = a.size();//number of equations
 int m = a[0].size() - 1;///number of variables
 vector<int>where(m, -1);///indicates which row
      contains the solution
 for (col = 0, row = 0; col < m && row < n; ++col) {
    int sel = row://which row contains the maximum
   for (int i = row + 1; i < n; i++)
     if (abs(a[i][col]) > abs(a[sel][col]))
    if (abs(a[sel][col]) < EPS) continue; ///it's
        basically 0.
    a[sel].swap(a[row]);///taking the max row up
    where [col] = row:
   ld t = a[row][col];
    for (int i = col; i <= m; i++) a[row][i] /= t;
    for (int i = 0; i < n; i++) {
     if (i != row) {
       ld c = a[i][col];
       for (int j = col; j <= m; j++)
          a[i][i] -= a[row][i] * c;
   row++;
 ans.assign(m, 0);
 for (int i = 0; i < m; i++)
    if (where[i] != -1)
      ans[i] = a[where[i]][m] / a[where[i]][i];
  for (int i = 0; i < n; i++) {
   ld sum = 0:
   for (int j = 0; j < m; j++)
      sum += ans[j] * a[i][j];
    if (abs(sum - a[i][m]) > EPS) ///L.H.S!=R.H.S
      ans.clear()://No solution
 return row;
7.5 \text{ GaussMod2} [44 lines] - e8fae4
template<typename T>
struct Gauss {
 int bits = 60:
 vector<T>table;
 Gauss() {
   table = vector<T>(bits, 0);
  //call with constructor to define bit size.
 Gauss(int _bits) {
   bits = _bits;
    table = vector<T>(bits, 0);
```

```
int basis()//return rank/size of basis
    int ans = 0:
    for (int i = 0:i < bits:i++)
     if (table[i])
        ans++;
   return ans;
  bool can(T x)//can x be obtained from the basis
    for (int i = bits - 1; i >= 0; i--) x = min(x, x)
        table[i]):
    return x == 0;
  void add(T x) {
    for (int i = bits - 1; i >= 0 \&\& x; i--) {
      if (table[i] == 0) {
        table[i] = x:
        x = 0:
      else x = min(x, x ^ table[i]);
  T getBest() {
    T x = 0:
    for (int i = bits - 1:i >= 0:i--)
      x = max(x, x \hat{table[i]}):
   return x:
  void Merge(Gauss& other) {
    for (int i = bits - 1; i >= 0; i--)
        add(other.table[i]);
};
7.6 Karatsuba Idea [5 lines] - 6944e1
Three subproblems:
a = xH yH
d = xL vL
e = (xH + xL)(yH + yL) - a - d
Then xv = a rn + e rn/2 + d
7.7 Linear Diophatine [19 lines] - 7c6f05
int extended_gcd(ll a, ll b, ll& x, ll& y) {
 if (b == 0)\{x = 1; y = 0; return a; \}
  11 d = extended_gcd(b, a % b, x1, y1);
  x = y1; y = x1 - y1 * (a / b);
  return d;
/*x'=x+(k*B/q), y'=y-(k*A/q); infinite soln
if A=B=0,C must equal 0 and any x,y is solution;
if A/B=0, (x,y)=(C/A,k)/(k,C/B)*/
bool LDE(11 A,11 B,11 C,11 &x,11 &y){
  int g=gcd(A,B);
  if(C%g!=0)return false;
  int a=A/g, b=B/g, c=C/g;
  extended_gcd(a,b,x,y); //ax+by=1
  if(g<0)\{a*=-1;b*=-1;c*=-1;\}//Ensure\ qcd(a,b)=1
  x*=c;y*=c;//ax+by=c
  return true; //Solution Exists
7.8 Matrix [100 lines] - a33f18
template<typename T>
struct Matrix {
```

```
T MOD = 1e9 + 7; ///change if necessary
T add(T a. T b) const {
 T res = a + b:
  if (res >= MOD) return res - MOD;
  return res;
T sub(T a, T b) const {
 T res = a - b;
  if (res < 0) return res + MOD;
  return res;
T mul(T a, T b) const {
 T res = a * b;
  if (res >= MOD) return res % MOD;
 return res;
int R, C;
vector<vector<T>>mat:
Matrix(int _R = 0, int _C = 0) {
 R = R, C = C;
 mat.resize(R);
 for (auto& v : mat) v.assign(C, 0);
void print() {
 for (int i = 0:i < R:i++)
    for (int j = 0; j < C; j++)
      cout << mat[i][i] << " \n"[i == C - 1]:
void createIdentitv() {
 for (int i = 0; i < R; i++)
    for (int j = 0; j < C; j++)
     mat[i][j] = (i == j);
Matrix operator+(const Matrix& o) const {
 Matrix res(R, C);
 for (int i = 0; i < R; i++)
   for (int j = 0; j < C; j++)
     res[i][j] = add(mat[i][j] + o.mat[i][j]);
Matrix operator-(const Matrix& o) const {
  Matrix res(R, C);
 for (int i = 0; i < R; i++)
    for (int j = 0; j < C; j++)
      res[i][j] = sub(mat[i][j] + o.mat[i][j]);
Matrix operator*(const Matrix& o) const {
 Matrix res(R. o.C):
 for (int i = 0; i < R; i++)
    for (int j = 0; j < o.C; j++)
     for (int k = 0; k < C; k++)
        res.mat[i][j] = add(res.mat[i][j],
            mul(mat[i][k], o.mat[k][j]));
 return res:
Matrix pow(long long x) {
 Matrix res(R, C);
  res.createIdentity();
 Matrix<T> o = *this;
  while (x) {
   if (x \& 1) res = res * o;
    o = o * o;
    x >>= 1:
```

```
return res:
  Matrix inverse()///Only square matrix & non-zero
      determinant
    Matrix res(R, R + R);
    for (int i = 0; i < R; i++) {
      for (int j = 0; j < R; j++)
       res.mat[i][j] = mat[i][j];
      res.mat[i][R + i] = 1;
    for (int i = 0;i < R;i++) {
      ///find row 'r' with highest value at [r][i]
      for (int j = i + 1; j < R; j++)
        if (abs(res.mat[j][i]) > abs(res.mat[tr][i]))
      ///swap the row
      res.mat[tr].swap(res.mat[i]);
      ///make 1 at [i][i]
      T val = res.mat[i][i];
      for (int j = 0; j < R + R; j++) res.mat[i][j] /=
      ///eliminate [r][i] from every row except i.
      for (int j = 0; j < R; j++) {
        if (j == i) continue;
        for (int k = R + R - 1; k >= i; k--) {
          res.mat[j][k] -= res.mat[i][k] * res.mat[j][i]
              / res.mat[i][i];
      }
    Matrix ans(R, R);
    for (int i = 0; i < R; i++)
      for (int j = 0; j < R; j++)
        ans.mat[i][j] = res.mat[i][R + j];
    return ans;
7.9 Miller-Rabin-Pollard-Rho [68 lines] - 3e3e5f
ll powmod(ll a, ll p, ll m) \{///(a^p \% m)
 ll result = 1:
 a %= m:
 while (p) {
   if (p & 1)
      result = (vll)result * a % m;
   a = (vll)a * a % m;
   p >>= 1;
 return result:
bool check_composite(ll n, ll a, ll d, int s) {
 ll x = powmod(a, d, n);
 if (x = 1 | x = n - 1)
   return false;
  for (int r = 1; r < s; r++) {
   x = (v11)x * x % n;
   if (x == n - 1)
      return false;
 return true;
bool MillerRabin(ll n) {
```

```
if (n < 2) return false;
  int r = 0:
  11 d = n - 1;
  while ((d \& 1) == 0) {
   d >>= 1;
  for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
    if (n == a) return true;
    if (check_composite(n, a, d, r))
      return false;
 return true;
11 mult(11 a, 11 b, 11 mod) {
 return (vll)a * b % mod;
11 f(11 x, 11 c, 11 mod) {
 return (mult(x, x, mod) + c) % mod;
11 rho(11 n) {
 if (n \% 2 == 0) return 2;
  ll x = myrand() \% n + 1, y = x, c = myrand() \% n + 1,
      g = 1;
  while (g == 1) {
   x = f(x, c, n);
    y = f(y, c, n);
    y = f(y, c, n);
    g = \_gcd(abs(x - y), n);
 return g;
set<ll>prime;
void prime_factorization(ll n) {
 if (n == 1) return;
 if (MillerRabin(n)) {
    prime.insert(n);
   return;
  11 x = n;
  while (x == n) x = rho(n);
  prime_factorization(x);
 prime_factorization(n / x);
//call prime_factorization(n) for prime factors.
//call MillerRabin(n) to check if prime.
7.10 Mod Inverse [5 lines] - 772679
int modInv(int a, int m) {
    int x, y; //if q==1 Inverse doesn't exist
    int g = gcdExt(a, m, x, y);
    return (x % m + m) % m;
7.11 NTT [96 lines] - 6faca3
ll power(ll a, ll p, ll mod) {
 if (p==0) return 1;
 11 ans = power(a, p/2, mod);
  ans = (ans * ans) \% mod;
  if(p%2)
            ans = (ans * a) \% mod;
 return ans;
int primitive_root(int p) {
 vector<int> factor;
```

```
int phi = p-1, n = phi;
  for (int i=2; i*i<=n; i++) {
   if (n%i) continue;
   factor.push_back(i);
    while (n\%i==0) n/=i;
 if (n>1) factor.push_back(n);
  for (int res =2; res<=p; res++) {
    bool ok = true;
    for (int i=0; i<factor.size() && ok; i++)
      ok &= power(res, phi/factor[i], p) != 1;
   if (ok) return res;
 return -1;
int nttdata(int mod, int &root, int &inv, int &pw) {
  int c = 0, n = mod-1;
  while (n\%2==0) c++. n/=2:
 pw = (mod-1)/n:
  int g = primitive_root(mod);
 root = power(g, n, mod);
  inv = power(root, mod-2, mod);
 return c;
const int M = 786433:
struct NTT {
 int N:
 vector<int> perm;
  int mod, root, inv, pw;
 NTT(){}
 NTT(int mod, int root, int inv, int pw) : mod(mod),
      root(root), inv(inv), pw(pw) {}
  void precalculate() {
   perm.resize(N);
    perm[0] = 0;
   for (int k=1; k<N; k<<=1) {
     for (int i=0; i<k; i++) {
       perm[i] <<= 1;
       perm[i+k] = 1 + perm[i];
  void fft(vector<11> &v, bool invert = false) {
    if (v.size() != perm.size()) {
      N = v.size();
      assert(N && (N&(N-1)) == 0);
      precalculate();
    for (int i=0; i<N; i++)
      if (i < perm[i])</pre>
        swap(v[i], v[perm[i]]);
    for (int len = 2; len <= N; len <<=1) {
      11 factor = invert ? inv: root;
      for (int i=len; i<pw; i<<=1)
       factor = (factor * factor) % mod;
      for (int i=0; i<N; i+=len) {
       11 w = 1;
        for (int j=0; j<len/2; j++) {
         11 x = v[i+j], y = (w*v[i+j+len/2]) \mod;
          v[i+j] = (x+y) \% mod;
         v[i+j+len/2] = (x-y+mod) \%mod;
          w = (w*factor)\%mod;
```

```
if (invert) {
      ll n1 = power(N, mod-2, mod);
      for (11 \& x: v) x = (x*n1) \% mod;
  vector<ll> multiply(vector<ll> a, vector<ll> &b) {
    while (a.size()) && a.back() == 0) a.pop_back();
    while (b.size() \&\& b.back() == 0) b.pop_back();
    int n = 1:
    while (n < a.size() + b.size()) n <<=1;
    a.resize(n);
    b.resize(n);
    fft(a);
    fft(b);
    for (int i=0; i<n; i++) a[i] = (a[i] * b[i]) M;
    fft(a. true):
    while (a.size() && a.back() == 0) a.pop_back();
 }
  //
         int mod=786433, root, inv, pw;
         nttdata(mod, root, inv, pw);
  //
         NTT nn = NTT(mod, root, inv, pw);
7.12 No of Digits in n! in base B [7 lines] - 86bfaf
11 NoOfDigitInNFactInBaseB(11 N,11 B){
 11 i;
 double ans=0;
 for(i=1;i<=N;i++)ans+=log(i);
  ans=ans/log(B),ans=ans+1;
 return(11)ans;
7.13 SOD Upto N [16 lines] - d8aa2c
11 SOD UpTo N(11 N){
 ll i, j, ans=0; ///upto\ N\ in\ Sqrt(N)
 for(i=1;i*i<=N;i++){
    ans+=((j*(j+1))/2)-(((i-1)*i)/2);
    ans+=((j-i)*i);
 return ans;
11 SODUptoN(11 N){
 11 res=0,u=sqrt(N);
 for(ll i=1;i<=u;i++)
   res+=(N/i)-i;
  res*=2,res+=u;
  return res;
7.14 Sieve Phi Mobius [26 lines] - 353c39
const int N = 1e7:
vector<int>pr;
int mu[N + 1], phi[N + 1], lp[N + 1];
void sieve() {
 phi[1] = 1, mu[1] = 1;
  for (int i = 2; i <= N; i++) {
    if (lp[i] == 0) {
      lp[\bar{i}] = i;
      phi[i] = i - 1;
      pr.push_back(i);
```

```
for (int j = 0; j < pr.size() && i * pr[j] <= N;
        i++) {
      lp[i * pr[j]] = pr[j];
      if (i % pr[j] == 0) {
        phi[i * pr[j]] = phi[i] * pr[j];
        phi[i * pr[j]] = phi[i] * phi[pr[j]];
  for (int i = 2; i \le N; i++) {
    if (lp[i / lp[i]] == lp[i]) mu[i] = 0;
    else mu[i] = -1 * mu[i / lp[i]];
8 Misc
8.1 Bit hacks [12 lines] - dd22ef
\# x & -x is the least bit in x.
# iterate over all the subsets of the mask
for (int s=m; ; s=(s-1)\&m) {
... you can use s ...
if (s==0) break;
# c = x\&-x, r = x+c; (((r^x) >> 2)/c) | r is the
next number after x with the same number of bits set.
# __builtin_popcount(x) //number of ones in binary
  __builtin_popcountll(x) // for long long
# __builtin_clz(x) // number of leading zeros
  __builtin_ctz(x) // number of trailing zeros, they
      also have long long version
8.2 Bitset C++ [13 lines] - a6a7a4
bitset<17>BS;
BS[1] = BS[7] = 1;
cout<<BS._Find_first()<<endl; // prints 1</pre>
bs._Find_next(idx). This function returns first set bit
    after index idx.for example:
bitset<17>BS:
BS[1] = BS[7] = 1:
cout<<BS._Find_next(1)<<','<<BS._Find_next(3)<<endl; //</pre>
    prints 7.7
So this code will print all of the set bits of BS:
for(int i=BS._Find_first();i< BS.size();i =</pre>
    BS._Find_next(i))
    cout << i << endl;
//Note that there isn't any set bit after idx,
    BS._Find_next(idx) will return BS.size(); same as
    calling BS._Find_first() when bitset is clear;
8.3 Template [33 lines] - 7aea62
// #pragma GCC optimize("03,unroll-loops")
// #pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <typename A, typename B> ostream&
    operator << (ostream& os, const pair <A, B>& p) {
    return os << '(' << p.first << ", " << p.second <<
    ')'; }
```

```
template <typename T_container, typename T = typename
    enable_if<!is_same<T_container, string>::value,
    typename T_container::value_type>::type> ostream&
    operator << (ostream& os, const T_container& v) { os
    << '{'; string sep; for (const T& x : v) os << sep
    << x, sep = ", "; return os << '}'; }
void dbg_out() { cerr << endl; }</pre>
template <typename Head, typename... Tail> void
    dbg_out(Head H, Tail...T) { cerr << " " << H;</pre>
    dbg_out(T...); }
#ifdef SMIE
#define debug(args...) cerr << "(" << #args << "):",
    dbq_out(args)
#define debug(args...)
#endif
template <typename T> inline T gcd(T a, T b) { T c; while
    (b) { c = b; b = a \% b; a = c; } return a; } // better
    than __qcd
ll powmod(ll a, ll b, ll MOD) { ll res = 1;a %=
    MOD; assert(b >= 0); for (; b; b >>= 1) { if (b & 
    1)res = res * a % MOD; a = a * a % MOD; }return res;
template <typename T>using orderedSet = tree<T,</pre>
    null_type, less_equal<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
//order_of_key(k) - number of element strictly less than
//find_by_order(k) - k'th element in set.(0
    indexed)(iterator)
mt19937
rng(chrono::steady_clock::now().time_since_epoch()
//uniform_int_distribution<int>(0, i)(rng)
int main(int argc, char* argv[]) {
  ios_base::sync_with_stdio(false);//DON'T CC++
  cin.tie(NULL);//DON'T use for interactive
  int seed = atoi(argv[1]);
8.4 build [2 lines] - 801989
#!/bin/bash
>&2 echo -e "Making [$2]\t: $1.cpp" && g++ -std=gnu++17
    -Wshadow -Wall -Wextra -Wno-unused-result -02 -g
    -fsanitize=undefined -fsanitize=address $2 "$1.cpp"
8.5 check [15 lines] - 478053
#!/bin/bash
build $1
TESTNO=0
for INP in $1.in*; do
 printf "\n======\n"
  printf "INPUT %d" $TESTNO
 printf "\n=====\n"
    cat $INP
  printf "\n=====\n"
  printf "OUTPUT %d" $TESTNO
  printf "\n=====\n"
    ./$1 < $INP
    mv $INP $1.in$TESTNO 2>/dev/null
    TESTNO=$((TESTNO+1))
done
```

```
8.6 debug [3 lines] - 859f78
#!/bin/bash
build "$1" -DSMIE && >&2 echo -e "Running\t\t:
    $1\n-----" && "./$1"
8.7 stress [15 lines] - 62e61a
#!/bin/bash
build $1 $2 && build $1_gen $2 && build $1_brute $2 &&
for((i = 1; ; ++i)); do
   echo -e "\nTest Case "$i
    ./$1 gen $i > inp
    ./$1 < inp > out1
    ./$1_brute < inp > out2
    diff -w out1 out2 || break
echo -e "=======\nINPUT\n----"
echo -e "\nOUTPUT\n----"
cat out1
echo -e "\nEXPECTED\n----"
cat out2
8.8 vimrc [14 lines] - ffdf4e
filetype plugin indent on
set rnu wfw hls is ar aw wrap mouse=a
let mapleader=' '
im jk <esc>
tno jk <c-w>N
no <leader>d "_d
im {<cr> {<cr>}<esc>0
nn ff :let @+ = expand("%:p") < cr >
nn cd :cd %:h<cr>
au BufNewFile *.cpp -r ./template.cpp | 14
ca hash w !cpp -dD -P -fpreprocessed \| tr -d
    '[:space:]' \| md5sum \| cut -c-6
9 String
9.1 Aho-Corasick [124 lines] - 2d8d6c
const int NODE=3000500;//Maximum Nodes
const int LGN=30;
                    ///Maximum Number of Tries
const int MXCHR=53: ///Maximum Characters
const int MXP=5005:
struct node {
 int val:
 int child[MXCHR];
 vector<int>graph;
 void clear(){
   CLR(child.0):
   val=0:
    graph.clear();
}Trie[NODE+10];
int maxNodeId,fail[NODE+10],par[NODE+10];
int nodeSt[NODE+10],nodeEd[NODE+10];
vlong csum[NODE+10],pLoc[MXP];
void resetTrie(){
 maxNodeId=0;
int getNode(){
 int curNodeId=++maxNodeId;
 Trie[curNodeId].clear();
```

```
return curNodeId:
inline void upd(vlong pos){
 csum[pos]++:
inline vlong qry(vlong pos){
 vlong res=csum[pos];
 return res;
struct AhoCorasick {
 int root, size, euler;
 void clear(){
   root=getNode();
   size=euler=0;
 inline int getname(char ch){
   if(ch=='-')return 52;
   else if(ch \ge A' \&\& ch \le Z') return 26 + (ch - A'):
   else return(ch-'a'):
 void addToTrie(string &s,int id){
  //Add string s to the Trie in general way
    int len=SZ(s),cur=root;
   FOR(i,0,len-1){
     int c=getname(s[i]);
      if(Trie[cur].child[c]==0){
       int curNodeId=getNode();
       Trie[curNodeId].val=c;
        Trie[cur].child[c]=curNodeId;
      cur=Trie[cur].child[c];
   pLoc[id]=cur;
    size++;
  void calcFailFunction(){
    aueue<int>Q;
    Q.push(root);
    while(!Q.empty()){
     int s=0.front();
      Q.pop();
    //Add all the children to the queue:
      FOR(i,0,MXCHR-1){
       int t=Trie[s].child[i]:
       if(t!=0){
          Q.push(t);
          par[t]=s;
      if(s==root){/*Handle special case when s is
        fail[s]=par[s]=root:
       continue:
//Find fall back of s:
      int p=par[s],f=fail[p];;
      int val=Trie[s].val;
/*Fall back till you found a node who has got val as a
     while(f!=root && Trie[f].child[val]==0){
       f=fail[f];
     fail[s]=(Trie[f].child[val]==0)? root :
          Trie[f].child[val];
```

```
//Self fall back not allowed
     if(s==fail[s]){
       fail[s]=root;
      Trie[fail[s]].graph.push_back(s);
  void dfs(int pos){
    ++euler;
    nodeSt[pos]=euler;
   for(auto x: Trie[pos].graph){
     dfs(x);
   nodeEd[pos]=euler;
 //Returns the next state
 int goTo(int state,int c){
    if(Trie[state].child[c]!=0){/*No need to fall
        back*/
     return Trie[state].child[c]:
  //Fall back now:
    int f=fail[state];
    while(f!=root && Trie[f].child[c]==0){
     f=fail[f]:
    int res=(Trie[f].child[c]==0)?
       root:Trie[f].child[c];
 /*Iterate through the whole text and find all the
     matchings*/
  void findmatching(string &s){
   int cur=root,idx=0;
    int len=SZ(s);
    while(idx<len){
     int c=getname(s[idx]);
      cur=goTo(cur,c);
     upd(nodeSt[cur]);
     idx++;
}acorasick;
9.2 Double Hasing [50 lines] - 1a70c1
struct SimpleHash {
    int len:
    long long base, mod;
    vector<int> P. H. R:
    SimpleHash() {}
    SimpleHash(string str, long long b, long long m) {
       base = b, mod = m, len = str.size();
        P.resize(len + 4, 1), H.resize(len + 3, 0),
           R.resize(len + 3, 0);
       for (int i = 1; i <= len + 3; i++)
           P[i] = (P[i - 1] * base) \% mod;
        for (int i = 1; i <= len; i++)
           H[i] = (H[i - 1] * base + str[i - 1] + 1007)
                % mod;
        for (int i = len; i >= 1; i--)
           R[i] = (R[i + 1] * base + str[i - 1] + 1007)
   }
```

```
KUET_Effervescent Team Notebook - September 11, 2024
```

```
9.4 Manacher [16 lines] - 2b3cab
                                                    vector<int> manacher_odd(string s) {
                                                      int n = s.size():
                                                      s = "$" + s + "^";
                                                      vector < int > p(n + 2);
                                                      int 1 = 1, r = 1;
                                                      for(int i = 1; i <= n; i++) {
                                                        p[i] = max(0, min(r - i, p[1 + (r - i)]));
                                                        while(s[i - p[i]] == s[i + p[i]]) {
                                                         p[i]++;
                                                        if(i + p[i] > r) {
                                                         l = i - p[i], r = i + p[i];
                                                      return vector<int>(begin(p) + 1, end(p) - 1);
                                                    9.5 Palindromic Tree [30 lines] - 9ebc05
                                                    struct PalindromicTree{
                                                      int n.idx.t:
                                                      vector<vector<int>> tree:
                                                      vector<int> len.link:
                                                      string s: // 1-indexed
                                                      PalindromicTree(string str){
                                                        s="$"+str:
                                                        n=s.size();
                                                        len.assign(n+5,0);
                                                        link.assign(n+5,0);
                                                        tree.assign(n+5, vector<int>(26,0));
                                                      void extend(int p){
                                                        while (s[p-len[t]-1]!=s[p]) t=link[t];
return ((long long)sh1.range_hash(l, r) << 32) ^
                                                        int x=link[t],c=s[p]-'a';
                                                        while(s[p-len[x]-1]!=s[p]) x=link[x];
                                                        if(!tree[t][c]){
                                                          tree[t][c]=++idx;
return ((long long)sh1.reverse_hash(1, r) << 32)
                                                          len[idx]=len[t]+2;
                                                          link[idx]=len[idx]==1?2:tree[x][c];
                                                        t=tree[t][c];
                                                      void build(){
                                                        len[1]=-1.link[1]=1:
                                                        len[2]=0, link[2]=1;
                                                        idx=t=2;
                                                        for(int i=1;i<n;i++) extend(i);</pre>
                                                    9.6 Prefix Function Automaton [21 lines] - b65c0e
                                                    /* create prefix function array in 26n.*/
                                                    int aut[mxn] [26];
                                                    int lps[mxn];
                                                    void automaton(string &s){
                                                      int n = s.size();
                                                      aut[0][s[0] - 'a'] = 1;
                                                      for(int i = 1; i < n; i++){
                                                        for(int j = 0; j < 26; j++){
                                                         if(j == s[i] - 'a'){
                                                            aut[i][j] = i + 1;
                                                            lps[i + 1] = aut[lps[i]][j];
```

inline int range_hash(int 1, int r) {

inline int reverse_hash(int 1, int r) {

1] * R[r + 2] % mod);

1] * H[1] % mod);

}

struct DoubleHash {

DoubleHash() {}

SimpleHash sh1, sh2;

9.3 KMP [23 lines] - 99c570

while(j>=0 and P[i]!=P[j])

while($j \ge 0$ and T[i]!=P[j])

//pattern found at index i-j

char P[maxn],T[maxn];

void kmpPreprocess(){

int b[maxn].n.m:

int i=0, j=-1;

j=b[j];

b[i]=j;

void kmpSearch(){

i=b[i];

i++; j++;

 $if(j==m){$

int i=0, j=0; while(i<n){

i++; j++;

while(i<m){

b[0]=-1:

DoubleHash(string str) {

int 12 . int r2) {

};

int hashval = H[r + 1] - ((long long)P[r - 1 +

return (hashval < 0 ? hashval + mod : hashval):

int hashval = R[1 + 1] - ((long long)P[r - 1 +

return (hashval < 0 ? hashval + mod : hashval);</pre>

sh1 = Simple Hash(str, 1949313259, 2091573227);

sh2 = SimpleHash(str, 1997293877, 2117566807);

long long concate(DoubleHash& B , int l1 , int r1 ,

int len1 = r1 - 11+1, len2 = r2 - 12+1;

long long x1 = sh1.range_hash(l1, r1) ,

x1 = (x1 * B.sh1.P[len2]) % 2091573227;

x1 = (x1 * B.sh2.P[len2]) % 2117566807;

inline long long range_hash(int 1, int r) {

inline long long reverse_hash(int 1, int r) {

sh2.reverse_hash(1, r);

long long newx2 = (x1 + x2) % 2117566807;

long long newx1 = (x1 + x2) % 2091573227;

 $x2 = B.sh1.range_hash(12, r2);$

 $x1 = sh2.range_hash(l1, r1);$

 $x2 = B.sh2.range_hash(12, r2);$

return (newx1 << 32) ^ newx2;

sh2.range_hash(1, r);

```
aut[i][j] = aut[lps[i]][j];
    cout << lps[i + 1] << endl;</pre>
9.7 Suffix Array [78 lines] - f2f7a0
struct SuffixArray {
  vector<int> p, c, rank, lcp;
  vector<vector<int>> st;
  SuffixArray(string const& s) {
   build_suffix(s + char(1));
   build_rank(p.size());
   build_lcp(s + char(1));
    build_sparse_table(lcp.size());
  void build_suffix(string const& s) {
   int n = s.size():
    const int MX_ASCII = 256;
   vector<int> cnt(max(MX_ASCII, n), 0);
   p.resize(n): c.resize(n):
   for (int i = 0; i < n; i++) cnt[s[i]]++;
   for (int i=1; i<MX_ASCII; i++) cnt[i]+=cnt[i-1];</pre>
   for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int classes = 1;
   for (int i = 1; i < n; i++) {
      if (s[p[i]] != s[p[i-1]]) classes++;
      c[p[i]] = classes - 1;
    vector<int> pn(n), cn(n);
   for (int h = 0; (1 << h) < n; ++h) {
     for (int i = 0; i < n; i++) {
       pn[i] = p[i] - (1 << h);
       if (pn[i] < 0) pn[i] += n;
      fill(cnt.begin(), cnt.begin() + classes, 0);
      for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
      for (int i=1; i < classes; i++) cnt[i]+=cnt[i-1];
      for (int i=n-1;i>=0;i--) p[--cnt[c[pn[i]]]]=pn[i];
      cn[p[0]] = 0; classes = 1;
      for (int i = 1; i < n; i++) {
       pair<int, int> cur = {c[p[i]], c[(p[i] + (1 <<
           h)) % n]};
       << h)) % n]}:
       if (cur != prev) ++classes;
       cn[p[i]] = classes - 1:
      c.swap(cn);
  void build_rank(int n) {
   rank.resize(n, 0);
   for (int i = 0; i < n; i++) rank[p[i]] = i;
  void build_lcp(string const& s) {
   int n = s.size(), k = 0;
   lcp.resize(n - 1, 0);
                                                        19
   for (int i = 0; i < n; i++) {
```

```
if (rank[i] == n - 1) {
        k = 0:
        continue:
      int j = p[rank[i] + 1];
      while (i + k < n \&\& j + k < n \&\& s[i+k] == s[j+k])
      lcp[rank[i]] = k;
      if (k) k--;
  void build_sparse_table(int n) {
    int lim = __lg(n);
    st.resize(lim + 1, vector<int>(n)); st[0] = lcp;
    for (int k = 1; k \le \lim_{k \to +} k + 1)
      for (int i = 0; i + (1 << k) <= n; i++)
        st[k][i] = min(st[k-1][i], st[k-1][i+(1 <<
  int get_lcp(int i) { return lcp[i]; }
  int get_lcp(int i, int j) {
    if (j < i) swap(i, j);
    j--; /*for lcp from i to j we don't need last lcp*/
    int K = _{-}lg(j - i + 1);
    return min(st[K][i]. st[K][i - (1 << K) + 1]):
};
9.8 Suffix Automata [109 lines] - 600ddc
const int mxc = 26;
/*
  + link
             - longest suffix belonging to another
      endpos-equivalent class.
             - largest string length ending in current
  + len
  + firstPos - first occurance of substring ending at
      current state.
  + adj
             - suffix link tree.
             - number of states.
  + 52
            - number of times state occured in string.
          - number of distinct substring.
  + cnt & SA - for count sorting the nodes.
struct SuffixAutomata{
  struct state{
    int link, len, firstPos;
    int next[mxc]:
    bool is_clone;
    state(){}
    state(int 1){
      len = 1, link = -1;
      is clone = false:
      for(int i=0;i < mxc;i++)next[i] = -1;
  };
  vector<state>t;
  int sz, last;
  vector<ll>cnt,dist, occ,SA;
  vector<vector<int>> adj;
  SuffixAutomata(){
    t.pb(state(0));
    occ.pb(0);
    last = sz = 0;
```

```
int getID(char c){ return c - 'a';}
void extend(char c){
  int idx = ++sz, p = last, id = getID(c);
  t.pb(state(t[last].len + 1));
  t[idx].firstPos = t[idx].len - 1;
  occ.pb(1);
  while (p! = -1 \text{ and } t[p] \cdot next[id] == -1)
    t[p].next[id] = idx;
   p = t[p].link;
  if(p==-1) t[idx].link = 0;
    int q = t[p].next[id];
    if(t[p].len+1 == t[q].len) t[idx].link = q;
      int clone = ++sz;
      state x = t[q];
      x.len = t[p].len+1:
      t.pb(x):
      t[clone].firstPos = t[q].firstPos;
      t[clone].is_clone = true;
      occ.pb(0);
      while (p!=-1 \text{ and } t[p].next[id]==q)
        t[p].next[id] = clone;
        p = t[p].link;
      t[idx].link = t[q].link = clone;
  }
  last = idx;
void build(string &s){
  for(char c:s) extend(c);
  cnt = dist = SA = vector<11>(sz+1);
  adj.resize(sz+1);
  for(int i=0;i<=sz;i++)cnt[t[i].len]++;
  for(int i=1;i<=sz;i++)cnt[i]+=cnt[i-1];
  for(int i=0;i<=sz;i++) SA[--cnt[t[i].len]] = i;</pre>
  for(int i=sz;i>0;i--){
    int idx = SA[i];
    occ[t[idx].link]+=occ[idx];
    adj[t[idx].link].pb(idx);
    dist[idx] = 1;
    for(int j=0; j < mxc; j++) {</pre>
      if(t[idx].next[j]+1){
        dist[idx]+=dist[t[idx].next[j]];
  for(int i=0;i<mxc;i++){</pre>
    if(t[0].next[i]+1) dist[0]+=dist[t[0].next[i]]:
pair<int,int> LCS( string& s){
  int mxlen = 0, bestpos = -1, pos = 0, len = 0;
  int u = 0;
  for(char c:s){
    int v = getID(c);
    while (u and t[u].next[v]!=-1){
      u = t[u].link;
      len = t[u].len;
    if(t[u].next[v]+1){
```

```
u = t[u].next[v]:
      if(len>mxlen){
        mxlen = len;
        bestpos = pos;
      pos++;
    return {bestpos - mxlen + 1, mxlen};
  state &operator[](int index) { return t[index];}
9.9 Trie [28 lines] - 408ef5
const int maxn=100005;
struct Trie{
  int next[27][maxn];
  int endmark[maxn].sz:
  bool created[maxn];
  void insertTrie(string& s){
    int v=0:
    for(int i=0;i<(int)s.size();i++){</pre>
      int c=s[i]-'a';
      if(!created[next[c][v]]){
        next[c][v]=++sz:
        created[sz]=true;
      v=next[c][v];
    endmark[v]++;
  bool searchTrie(string& s){
    for(int i=0;i<(int)s.size();i++){</pre>
      int c=s[i]-'a';
      if(!created[next[c][v]])
        return false;
      v=next[c][v];
    return(endmark[v]>0);
};
9.10 Z-Algorithm [19 lines] - e04285
void compute_z_function(const char*S,int N){
  int L=0.R=0:
  for(int i=1;i<N;++i){</pre>
    if(i>R){
      while (R < N \&\& S[R-L] == S[R]) ++ R;
      Z[i]=R-L.--R:
    else{
      int k=i-L;
      if(Z[k]<R-i+1)Z[i]=Z[k];
      else{
        while (R < N \&\& S[R-k] == S[R]) ++R;
        Z[i]=R-L,--R;
```

10 Random 10.1 Combinatorics

$$\bullet \sum_{k=0}^{n} {n-k \choose k} = Fib_{n+1}$$

$$\bullet \ \binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$$

$$\bullet \ k\binom{n}{k} = n\binom{n-1}{k-1}$$

- Number of binary sequences of length n such that no two 0's are adjacent = Fib_{n+1}
- Number of non-negative solution of $x_1 + x_2 + x_3 + ... + x_k = n$ is $\binom{n+k-1}{n}$

10.1.1 Catalan Number

•
$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

•
$$C_0 = 1, C_1 = 1, C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

- 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786
- Number of correct bracket sequences consisting of n opening brackets.
- Number of ways to completely parenthesize n+1 factors.
- The number of triangulations of a convex polygon with +2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).
- The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.
- The number of monotonic lattice paths from point (0,0) to point (n,n) in a square lattice of size $n \times n$, which do not pass above the main diagonal
- Number of permutation of length n that can be stack sorted.
- The number of non-crossing partitions of a set of n elements.
- The number of rooted full binary tree with n+1 leaves.
- The number of Dyck words of length 2n. A string consisting of n X's and n Y's such that no string prefix has more Y's than X's.
- Number of permutation of length n with no three-term increasing subsequence.
- Number of ways to tile a stairstep shape of height n with n rectangle.

- $C_n^k = \frac{k+1}{n+1} \binom{2n-k}{n-k}$ denote the number of bracket sequences of size 2n with the first k elements being (.
- $N(n,k) = \frac{1}{n} \binom{n}{k} \binom{n}{k-1}$
- The number of expressions containing n pairs of correct parentheses, which contain k distinct nestings. N(4,2) = 6 ()((()),(())(()),(()(())),((()())),((())(())))
- The number of paths from (0,0) to (2n, 0) with steps only northeast and southeast, not staying below the x-axis with k peaks. And sum of all number of peaks is Catalan number.

10.1.2 Stirling Number of the First Kind

- Count permutation according to their number of cycles.
- S(n,k) count the number of permutation of n elements with k disjoint cycles.
- $S(n,k) = (n-1) \times S(n-1,k) + S(n-1,k-1), S(0,0) = 1, S(n,0) = S(0,n) = 0$
- S(n,1) = (n-1)!
- $S(n,n-1) = \binom{n}{2}$
- $\bullet \ \sum_{k=0}^{n} S(n,k) = n!$

10.1.3 Stirling Numbers of the Second Kind

- Number of ways to partition a set of n objects into k non-empty subsets.
- S(n,k) = k * S(n-1,k) + S(n-1,k-1), S(0,0) = 1, S(n,0) = S(0,n) = 0
- $S(n,2) = 2^{n-1} 1$
- $S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$
- S(n,k) * k! = number of ways to color n nodes using colors from 1 to k such that each color is used at least once.

10.1.4 Bell Number

- Counts the number of partitions of a set.
- $\bullet \ B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} * B_k$
- $B_n = \sum_{k=0}^n S(n,k)$, where S is Stirling number of second kind.
- The number of multiplicative partitions of a square free number with i prime factors is the i-th Bell number.

- $B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$
- If a deck is shuffled by removing and reinserting the top card n times, there are n^n possible shuffles. The number of shuffles that return the deck to its original order is B_n , so the probability of returning to the original order is B_n/n^n .

10.1.5 Lucas Theorem

- If p is prime then $\binom{p^a}{k} \equiv 0 \mod p$
- For non-negative integers m and n and a prime p:

$$\binom{m}{n} = \prod_{i=0}^{n} \binom{m_i}{n_i} \pmod{p}$$
 where
$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0 \ n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$
 are the base p expansion.

10.1.6 Derangement

- A permutation such that no element appears in its original position.
- d(n) = (n-1)*(d(n-1)+d(n-2)), d(0) = 1, d(1) = 0
- $d(n) = nd(n-1) + (-1)^n = \lfloor \frac{n!}{n!} \rfloor, n \ge 1$

10.1.7 Burnside Lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

where X^g are the elements fixed by g(g.x = x) If f(n) counts "configurations" of some sort of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k)$$

10.1.8 Eulerian Number

- E(n,k) is the number of permutations of the numbers 1 to n in which exactly k elements are greater than the previous element.
- E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k), E(n,0) = E(n,n-1) = 1

•
$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

•
$$E(n,k) = E(n,n-1-k)$$

•
$$E(0,k) = [k=0]$$

•
$$E(n,1) = 2^n - n - 1$$

|21|

10.2 Number Theory

10.2.1 Mobius Function and Inversion

• define $\mu(n)$ as the sum of the primitive nth roots of unity depending on the factorization of n into prime factors:

$$\mu(x) = \begin{cases} 0 & \text{n is not square free} \\ 1 & \text{n has even number of prime factors} \\ -1 & \text{n has odd number of prime factors} \end{cases}$$

• Mobius Inversion:

$$g(n) = \sum_{d|n} f(d) \leftrightarrow f(n) = \sum_{d|n} \mu(d)g(n/d)$$

- $\sum_{d|n} \mu(d) = [n=1]$
- $\phi(n) = \sum_{d|n} \mu(d) \cdot \frac{n}{d} = n \sum_{d|n} \frac{\mu(d)}{d} = \sum_{d|n} d \cdot \mu(\frac{n}{d})$
- $a|b \to \phi(a)|\phi(b)$
- $\phi(mn) = \phi(m).\phi(n).\frac{d}{\phi(d)}$ where $d = \gcd(m, n)$
- $\bullet \sum_{i=1}^{n} [gcd(i,n) = k] = \phi(\frac{n}{k})$
- $\sum_{i=1}^{n} gcd(i,n) = \sum_{d|n} d.\phi(\frac{n}{d})$
- $\bullet \sum_{i=1}^{n} \frac{1}{\gcd(i,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi(\frac{n}{d}) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$
- $\sum_{i=1}^{n} \frac{i}{\gcd(i,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi(\frac{n}{d}) = \frac{n}{2} \cdot \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$
- $\sum_{i=1}^{n} \frac{n}{\gcd(i,n)} = 2 \cdot \sum_{i=1}^{n} \frac{i}{\gcd(i,n)} 1$

10.2.2 GCD and LCM

- $\bullet \ \gcd(a,b) = \gcd(b, \ a \ mod \ b)$
- If a|b.c, and gcd(a,b) = d, then (a/d)|c.
- GCD is a multiplicative function.
- gcd(a, lcm(b,c)) = lcm(gcd(a,b), gcd(a,c))
- $gcd(n^a 1, n^b 1) = n^{gcd(a,b)} 1$

10.2.3 Gauss Circle Theorem

- Determine the number of lattice points in a circle centered at the origin with radius r.
- number of pairs (m,n) such that $m^2 + n^2 \le r^2$

•
$$N(r) = 1 + 4 \sum_{i=0}^{\infty} (\lfloor \frac{r^2}{4i+1} \rfloor - \lfloor \frac{r^2}{4i+3} \rfloor)$$

10.2.4 Pick's Theorem

According to Pick's Theorem We can calculate the area of any polygon by just counting the number of Interior and Boundary lattice points of that polygon. If number of interior points are I and number of boundary lattice points are B then Area (A) of polygon will be:

$$Area = I + B/2 - 1$$

where I is the number of points in the interior shape, B stands for the number of points on the boundary of the shape.

10.2.5 Formula Cheatsheet

- $\sum_{i=1}^{n} = \frac{1}{m+1}[(n+1)^{m+1} 1 \sum_{i=1}^{n}((i+1)^{m+1} i^{m+1} (m+1)i^{m})]$
- $\sum_{i=0}^{n} c^i = \frac{c^{n+1}-1}{c-1}, c \neq 1$
- $\sum_{i=0}^{\infty} c^i = \frac{1}{1-c}, \sum_{i=1}^{\infty} c^i = \frac{c}{1-c}, |c| < 1$
- $H_n = \sum_{i=1}^n \frac{1}{n}, \sum_{i=1}^n iH_i = \frac{n(n+1)}{2}H_n \frac{n(n-1)}{4}$
- $\bullet \sum_{k=0}^{n} {r+k \choose k} = {r+n+1 \choose n}$