

An algorithm for computing the k -numerical range

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Visual examples of the k -numerical range

$$A_1 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & \omega & 0 & 0 & 0 \\ 0 & 0 & \omega^2 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}, \quad A_2 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & \omega & 0 & 0 & 0 \\ 0 & 0 & \omega^2 & 0 & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2} \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

where $\omega = \frac{-1+i\sqrt{3}}{2}$

Introduction

The k -numerical range is defined as:

$$W_k(A) = \left\{ \frac{1}{k} \sum_{i=1}^k x_i^* A x_i : x_i^* x_j = \delta_{ij} \right\}$$

The k -numerical range:

- is convex
- is compact
- contains averages of k eigenvalues (counting multiplicity)

The real part of the k -numerical range

For $A \in M_n(\mathbb{C})$, $\text{Re}(W_k(A)) = W_k(\text{Re}(A))$.

Linear Transform Property

For $z, w \in \mathbb{C}$, $A \in M_n(\mathbb{C})$, $W_k(zA + wI) = zW_k(A) + w$.

Technical Lemma

For $\lambda_1 \geq \dots \geq \lambda_n$, $0 \leq c_i \leq 1$, $\sum_{i=1}^n c_i = k \in \mathbb{N}$. Then $\sum_{i=1}^n c_i \lambda_i \leq \sum_{i=1}^k \lambda_i$.

Convex Combination Lemma

For $A = A^*$, any $z \in W_k(A)$ has the form $z = \frac{1}{k} \sum_{i=1}^n c_i \lambda_i$, where $\lambda_1 \geq \dots \geq \lambda_n$ are eigenvalues of A , $0 \leq c_i \leq 1$, $\sum_{i=1}^n c_i = k$.

Maximum Real Part Theorem

Suppose $A \in M_n(\mathbb{C})$, and v_1, \dots, v_n are eigenvectors of $\text{Re}(A)$ for eigenvalues $\lambda_1 \geq \dots \geq \lambda_n$. Then $\frac{1}{k} \sum_{i=1}^k v_i^* A v_i \in W_k(A)$ has maximum real part.

These results are the key tools necessary to generalize an algorithm due to Carl Cowen in the case when $k = 1$.

Explanation of the algorithm

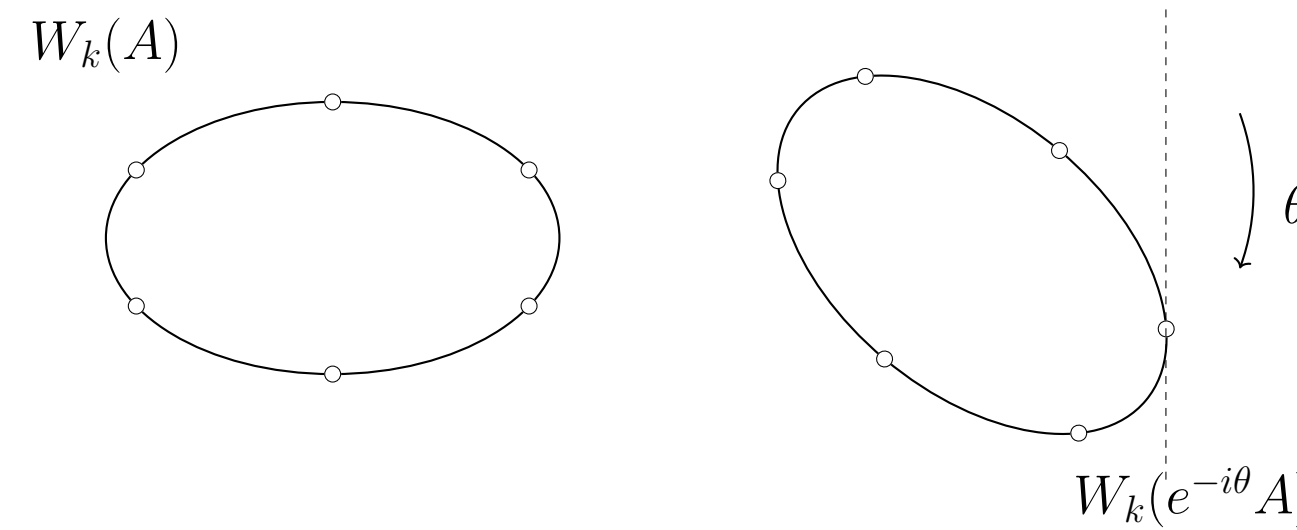


Figure 1: The point on $W_k(e^{-i\theta}A)$ with maximum real part is achieved by the k -largest eigenvectors of $\text{Re}(e^{-i\theta}A)$. So, it suffices to find the eigenvectors corresponding to the k -largest eigenvalues of $\text{Re}(e^{-i\theta}A)$ for various θ .

To compute a boundary point of $W_k(A)$ in the direction θ , we rotate the matrix: $e^{-i\theta}A$. The maximum part theorem shows that the point in $W_k(e^{-i\theta}A)$ with maximum real part is achieved by the k eigenvectors corresponding to the k -largest eigenvalues of $\text{Re}(e^{-i\theta}A)$.

By evaluating this over many angles θ , we trace the boundary of $W_k(A)$.

How it works: Rotate $A \rightarrow$ Take real part $\text{Re}(e^{-i\theta}A) \rightarrow$ Use k eigenvectors with largest eigenvalues \rightarrow Average \rightarrow Plot the result.

To approximate the k -numerical range of $A_0 = \text{diag}(i, -i, 1, \dots, 1) \in M_7(\mathbb{C})$, we compute x^*Ax for 1000 random unit vectors x . The vertices of the range lie at i , $-i$, and 1. Since most diagonal entries are 1, the values cluster near 1 on the real axis. The resulting approximation for $k = 1$ is shown below.

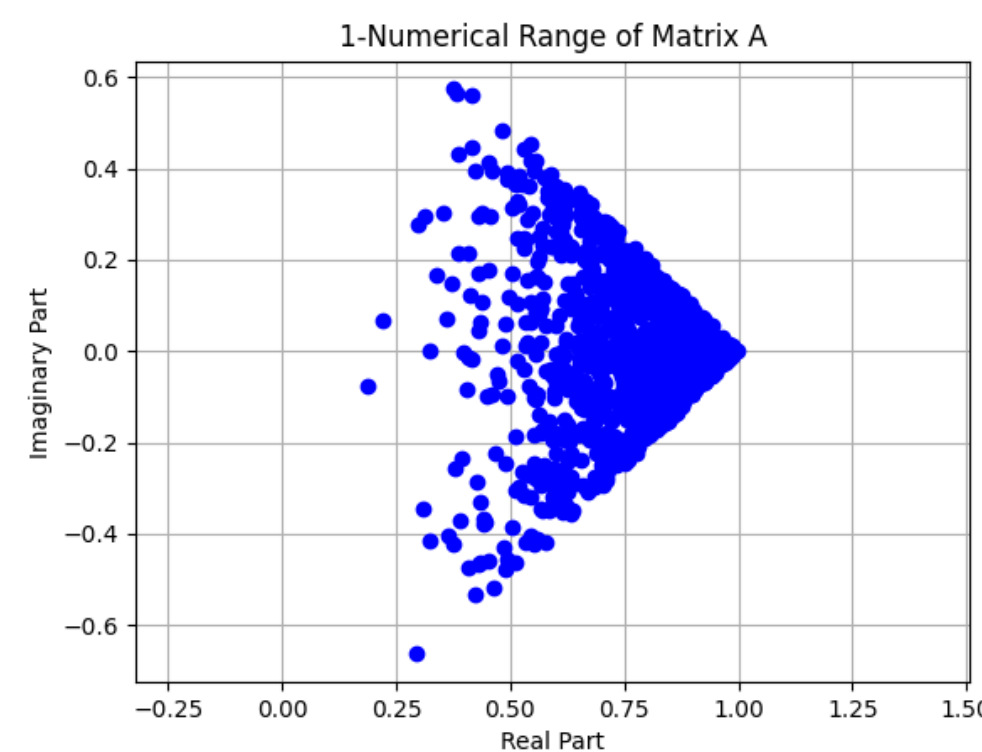


Figure 2: Approximation of the numerical range of $\text{diag}(i, -i, 1, \dots, 1) \in M_7(\mathbb{C})$; computed x^*Ax for 1000 randomly chosen unit vectors x

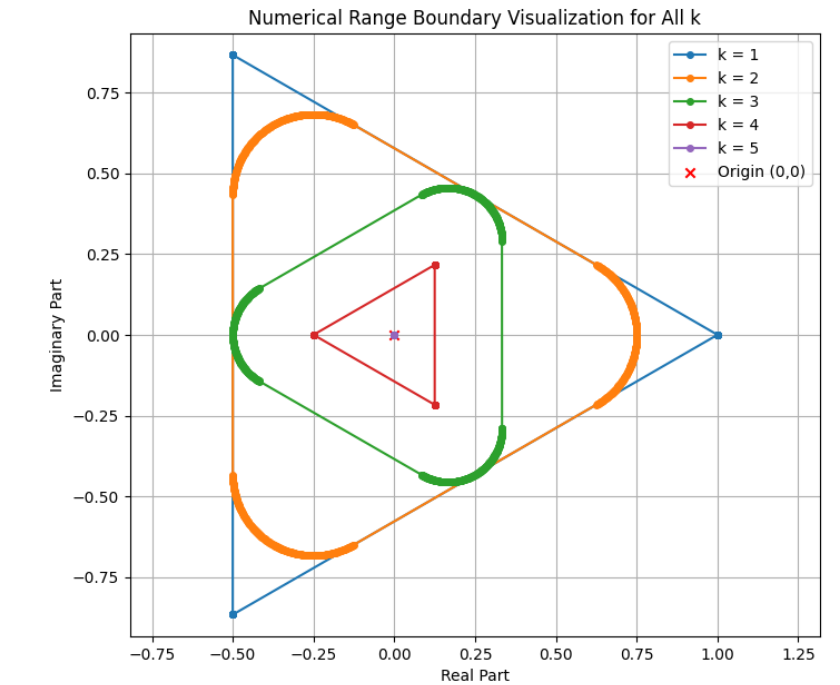


Figure 3: k -numerical range of A_1

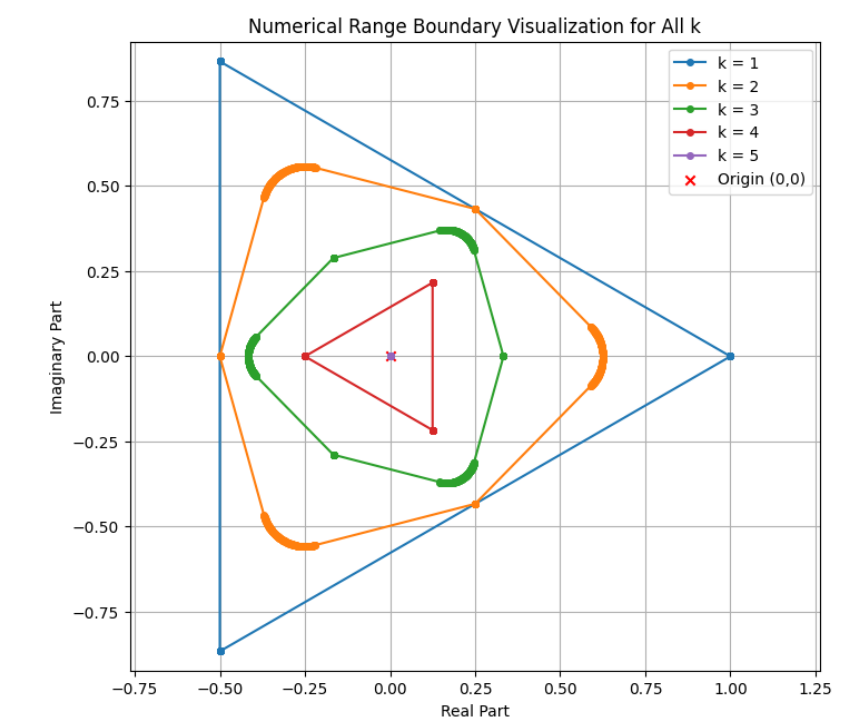


Figure 4: k -numerical range of A_2

References

- [1] Carl C. Cowen, *An effective algorithm for computing the numerical range*, Aug 1995, Unpublished Manuscript.
- [2] Panayiotis J. Psarrakos and Michael J. Tsatsomeros, *Numerical range: (in) a matrix nutshell*, <https://www.math.wsu.edu/math/faculty/tsat/files/short.pdf>, Aug 2002.