

No need to show the parse tree, included for better understanding

Obtained Marks

Name:

ID:

Problem 1: CFG

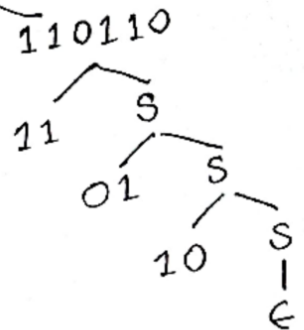
Consider the following two languages.

$L_1 = \{w \in \{0, 1\}^* : \text{The length of } w \text{ is even.}\}$

$L_2 = \{w \in \{0, 1, \#\}^* : w = x\#1^n\#0^{2n+1} \text{ where } x \in L_1 \text{ and } n \geq 0\}$

- a) Design a context-free grammar whose language is L_1 . [Points 6]

Solution 1: $S \rightarrow 00S \mid 01S \mid 10S \mid 11S \mid \epsilon$



Solution 2: $S \rightarrow 0S0 \mid 0S1 \mid 1S0 \mid 1S1 \mid \epsilon$

- b) Design a context-free grammar whose language is L_2 . [Points 4]

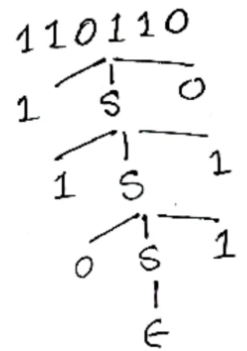
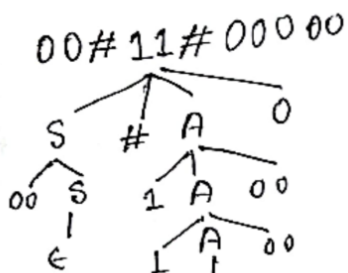
$$w = x\#1^n\#0^{2n+1}$$

$$= x\#1^n\#0^{2n}0$$

$$S' \rightarrow S\#A0$$

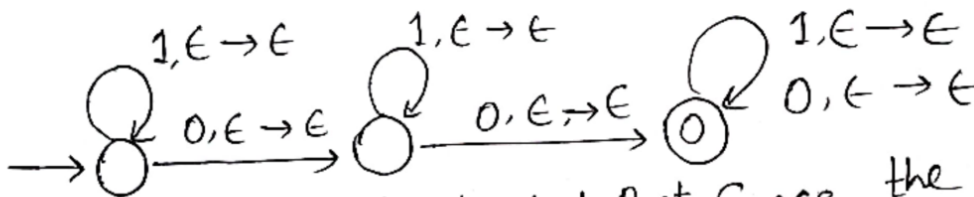
$$S \rightarrow 00S \mid 01S \mid 10S \mid 11S \mid \epsilon$$

$$A \rightarrow 1A00 \mid \#$$



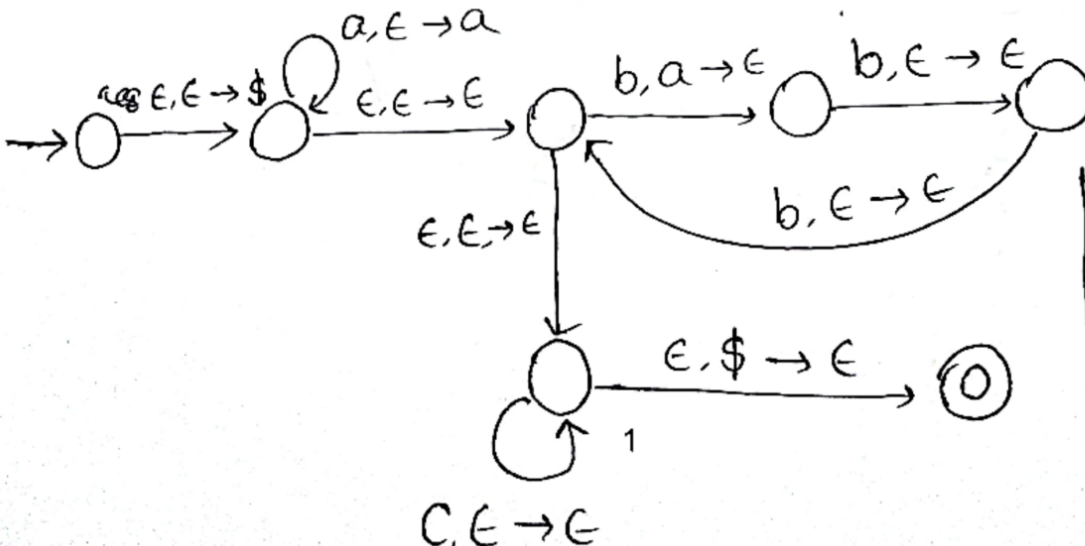
Problem 2: PDF

- a) Give a PDA for the language $L = \{w \in \{0, 1\}^* : w \text{ contains at least two 0s.}\}$ [4 Points]



You may push in the stack. But, Since the language is regular, no need to use the stack.

- b) Give a PDA for the language $L = \{w \in \{a, b, c\}^* : w = a^n b^{3n} c^m \text{ where } n, m \geq 0\}$ [6 Points]



Solution 2: You may also push three 'a' by reading one 'a' in the stack.

Problem 3 : Derivations, Parse Trees, and Ambiguity

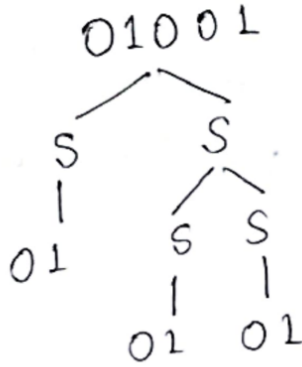
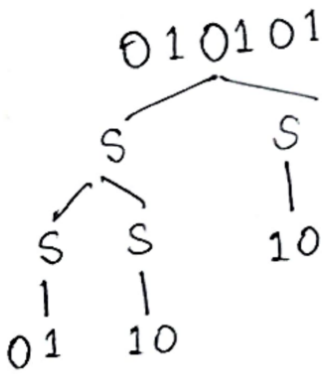
a) Consider the following context free grammar:

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10$$

i) Give a left derivation for the string 011010. [2 Points]

$$\begin{aligned} S &\rightarrow SS \\ &\rightarrow 01^S \\ &\rightarrow 01SS \\ &\rightarrow 0110S \rightarrow 011010 \end{aligned}$$

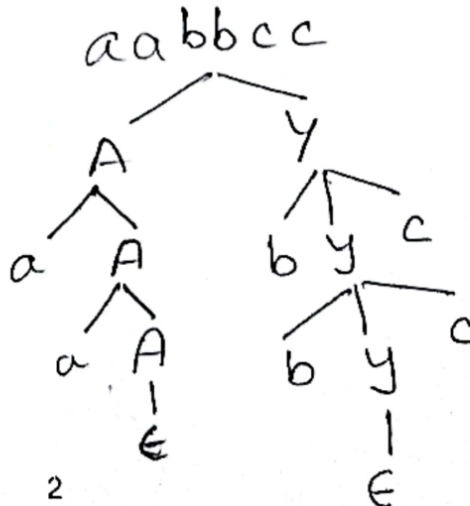
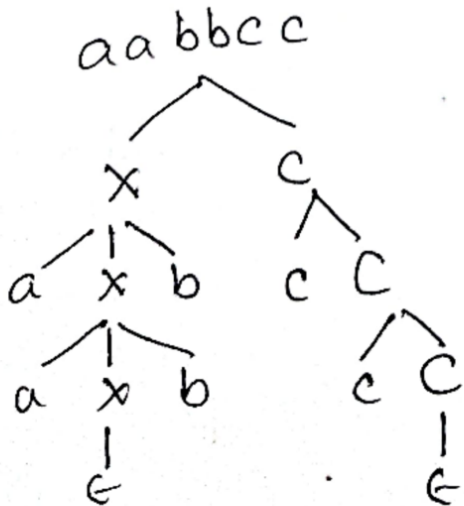
ii) Show that the grammar above is ambiguous by demonstrating two different parse trees for 010101. [3 Points]



b) Consider the following context free grammar:

$$\begin{aligned} S &\rightarrow XC \mid AY \\ X &\rightarrow aXb \mid \varepsilon \\ Y &\rightarrow bYc \mid \varepsilon \\ A &\rightarrow aA \mid \varepsilon \\ C &\rightarrow cC \mid \varepsilon \end{aligned}$$

Show that the grammar above is ambiguous by finding a length 6 string with two parse trees. [5 Points]



Question 03

a) w starts with '0' and the length of w is even.

$$S \rightarrow 0AB$$

$$A \rightarrow 00A \mid 01A \mid 10A \mid 11A \mid \epsilon$$

$$B \rightarrow 0 \mid 1$$

$$0(zz)^*(0+1)$$

b) Every second letter in w is b

$$((a+b)b)^*(\epsilon + a + b)$$

$$S \rightarrow AB$$

$$A \rightarrow XA \mid \epsilon$$

$$X \rightarrow PQ$$

$$P \rightarrow a \mid b$$

$$Q \rightarrow b$$

$$B \rightarrow a \mid b \mid \epsilon$$

Solution 2:

$$S \rightarrow abS \mid bbS \mid A$$

$$_b_b_b_$$

$$A \rightarrow a \mid b \mid \epsilon$$

c) the length of w is divisible by ~~2~~ three.

$$(\Sigma\Sigma\Sigma)^*$$

$$\rightarrow ((a+b)(a+b)(a+b))^*$$

} No need to write ~~it~~ in the answer script (optional)

$$S \rightarrow AS|E$$

$$A \rightarrow BBB$$

$$B \rightarrow a|b$$

d) w starts and ends with different letters.

$$a(a+b)^*b + b(a+b)^*a$$

$$S \rightarrow aAb | bAa$$

$$A \rightarrow aA | bA | E$$

e) the number of a is at least the number of b in w

$$S \rightarrow \underbrace{aSbs | bSas}_{\text{ensuring \# 'b's equal to \# 'a's, means the amount of 'b's won't exceed the amount of 'a's'}} | \overset{\text{more 'a's'}}{aS} | E$$

ensuring # 'b's equal to # 'a's, means the amount of 'b's won't exceed the amount of 'a's

Similar Qs:

w contain more ~~0's~~ 'a's than ~~1's~~ 'b's

$$S \rightarrow XaX$$

$$X \rightarrow aXbX | bXaX | \cancel{aX} | E$$

f) $w = 0^n 1^n$, where n is odd.

wrong \leftarrow

$S \rightarrow ABX$
$A \rightarrow 0X$
$X \rightarrow 00X1 \in$
$B \rightarrow 1Y$
$Y \rightarrow 11Y1 \in$

Solution, since this doesn't ensure equal numbers of 0s & 1s.

$S \rightarrow 00S11|A$

$A \rightarrow 01$

g) $w = 0^i 1^j 2^k$, where $j \geq 2i + 3k$

$S \rightarrow ABC$

$A \rightarrow 0A11| \in$

$B \rightarrow 1B| \in$

$C \rightarrow 111C2| \in$

* first try to solve $j = 2i + 3k$

$0^i 1^j 2^k$

$\rightarrow 0^i 1^{2i+3k} 2^k$

$\rightarrow \underbrace{0^i 1^{2i}}_{\text{}} \cdot \underbrace{1^{3k} 2^k}_{\text{}}$

the additional 1s ($j >$)
will occur in this place.

J) $A = w$ contains at least two 0s

$$L = \{w : w = 0^{3i} \vee 1^{2i}, \forall i \in \mathbb{N}, i \geq 0\}$$

$$S \rightarrow 000S11 \mid A$$

$$A \rightarrow X0X0X$$

$$X \rightarrow 0X \mid 1X \mid \epsilon$$

