

1. (15 points) True or False

Determine whether the following statements are true or false.

- (a) (3') Merge sort is not an in place sorting since it requires $\Theta(\log n)$ space. ☐ True ☐ False
- (b) (3') When an array contains distinct elements and is arranged in descending order, the time complexity of quick sort (always choose the last element as pivot) is $\Theta(n^2)$. ☐ True ☐ False
- (c) (3') There exists a comparison-base sort algorithm that needs $O(1)$ extra space and takes $o(n \log n)$ time. ☐ True ☐ False
- (d) (3') For an array $\{a_n\}$ with distinct elements, for fixed i, j , if $\forall a_k \neq a_i, a_k \neq a_j, (a_k - a_i)(a_k - a_j) > 0$, then a_i and a_j will be compared in any case when using randomized quick-sort to make $\{a_n\}$ sorted. ☐ True ☐ False
- (e) (3') When performing flagged bubble sort on an array of n elements, the minimum number of comparisons is $n - 1$. ☐ True ☐ False

2. (15 points) Fill in the blanks

- (a) (3') If we use randomized quick-sort (i.e., randomly choosing pivots) to sort the array $[3, 4, 6, 2, 1, 5, 8, 0]$, the probability that 2 and 5 are compared is _____.
- (b) (3') The worst case runtime of insertion sort is $\Theta(\text{_____})$, and the best case runtime is $\Theta(\text{_____})$. (The array has n entries.)
- (c) (3') In average case, the space complexity of quick sort on n elements is $\Theta(\text{_____})$.
- (d) (3') The following is the process of sorting an array $\{7, 8, 5, 2, 4, 6, 3\}$ by a sorting algorithm $\{7, 8, 5, 2, 4, 6, 3\} \rightarrow \{7, 5, 2, 4, 6, 3, 8\} \rightarrow \{5, 2, 4, 6, 3, 7, 8\} \rightarrow \{2, 4, 5, 3, 6, 7, 8\} \rightarrow \{2, 4, 3, 5, 6, 7, 8\} \rightarrow \{2, 3, 4, 5, 6, 7, 8\} \rightarrow \{2, 3, 4, 5, 6, 7, 8\}$
Which sorting algorithm is it? _____.
- (e) (3') The array $\{8, 9, 10, 4, 5, 6, 20, 1, 2\}$ is the result of _____ sort after two passes of the outer loop. (fill in the blank with "bubble" or "insertion")

3. (20 points) Alternating bubble sort

```
template <typename Type>
void bubble(Type* const array, int n) {
    int lower = 0;
    int upper = n - 1;
    while (true) {
        int new_upper = lower;
        for (int i = lower; i < upper; ++i) {
            if (array[i] > array[i + 1]) {
                Type tmp = array[i];
                array[i] = array[i + 1];
                array[i + 1] = tmp;
                new_upper = i;
            }
        }
    }
}
```

```

    upper = new_upper;
    if (lower == upper) { break; }
    int new_lower = upper;
    for (int i = upper; i > lower; --i) {
        if (array[i - 1] > array[i]) {
            Type tmp = array[i];
            array[i] = array[i - 1];
            array[i - 1] = tmp;
            new_lower = i;}
    }
    lower = new_lower;
    if (lower == upper) { break; }
}

```

Suppose we apply this algorithm to an array with n entries.

- (a) (3') Compared with the standard one-way bubble sort, the main optimization of the given code is: _____.
 - A. Using binary search to locate swap positions
 - B. Using a bidirectional pass to push both the maximum and minimum toward the boundaries, shrinking the unsorted range
 - C. Using a faster swap trick to reduce assignments
- (b) (3') In the worst case (array in reverse order), the time complexity of this algorithm is Θ (_____).
- (c) (6') After the first for loop (from lower to upper) in the first iteration of the while loop completes, the statement “upper = new_upper” updates the right boundary of the unsorted range to the last swap position i , skipping the already sorted suffix. Thus, the unsorted range of the array shrinks to [_____, _____] (fill in the name of the variable). If no swap occurs in this pass, then *new_upper* will be equal to _____.
- (d) (3') For this code to sort a custom **Type** correctly, Type must define the operator $>$. Additionally, swapping via a temporary requires that the type supports copying or moving. ☐ True ☐ False
- (e) (5') Sort the array {3, 7, 4, 2, 6, 5, 1, 0} using the algorithm above. Write down the array after each **for** loop.