

Wallace Multiplier with Booth Encoding

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What is a multiplier

A circuit that multiply two numbers.

Input:

$$a = a_{n-1} a_{n-2} \cdots a_0 = \sum_{i=0}^{n-1} 2^i a_i$$

$$b = b_{n-1} b_{n-2} \cdots b_0 = \sum_{i=0}^{n-1} 2^i b_i$$

Output:

$$\begin{aligned} z &= z_{2n-1} z_{2n-1} \cdots z_0 \\ &= a \times b \end{aligned}$$

Wallace Multiplier

Directly expand the product:

$$a \times b = \sum_i \sum_j 2^{i+j} a_i b_j \quad (1)$$

Now we have n^2 partial product $2^{i+j} a_i b_j$, $a_i b_j$ is 0 or 1.
Define 2^{i+j} be the “weight” of such a partial product.

Naïve method: sum up all partial products in a binary tree.

Since n -digit addition takes $O(\log n)$ time, this method takes $O(\log^2 n)$ time.

Wallace Multiplier

Optimization: group partial products by their weight (each group called a column). Only add up partial products in the same group.

$$2^i a_1 + 2^i a_2 + 2^i a_3 = 2^{i+1} b_1 + 2^i b_2 \quad (\text{Full adder})$$

Transform 3 partial products to 2 (3-2 compressor)

$$2^i a_1 + 2^i a_2 = 2^{i+1} b_1 + 2^i b_2 \quad (\text{Full adder})$$

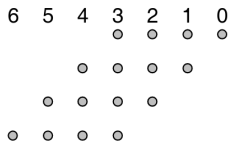
Transform 2 partial products to 2 (2-2 compressor)

In each stage we compress each column as much as possible. Until there are at most two items in each column.

It remains to execute an $2n$ -bit addition.

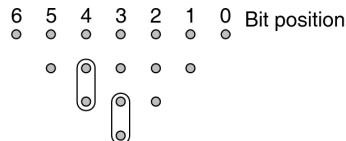
Wallace Multiplier: Example

Partial products



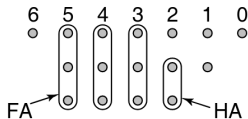
(a)

First stage



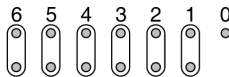
(b)

Second stage



(c)

Final adder



(d)

Notice that in each stage we reduce the number of partial products by roughly $1/3$, there are roughly $\log_{3/2}(n)$ stages.

Optimization Techniques

Fast enough. But may I make it even faster?

A wallace multiplier consists of three parts:

- ▶ **Partial product generation:** generate less partial products
- ▶ **Compression:** compress more in one stage (use 4-3, 5-3 compressor)
- ▶ **Addition:** faster adder (not covered in this topic)

Last two optimizations are easy. The first is tricky.

Booth Recoding

Define $a_{-1} = 0$, $a_n = a_{n-1}$ (sign extension).

$$\begin{aligned} a \times b &= \sum_{i=0}^{n-1} 2^i a_i b = \sum_{i=0}^{n-1} (2^{i+1} - 2^i) a_i b \\ &= \sum_{i=0}^{n-1} 2^i (-a_{i+1} + a_i) b \end{aligned}$$

Example:

$$1100111 = 2^7 - 2^5 + 2^3 - 2^0$$

Now we have

$$1100111 \times b = 2^7 b + 2^5 (-b) + 2^3 b + 2^0 (-b) \quad (2)$$

Not faster in worst case.

Radix-4 Booth Recoding

$$\begin{aligned}
 a &= \sum_{i=0}^{n-1} 2^i a_i \\
 &= \sum_{i=0}^{(n-1)/2} 4^i (-2a_{2i+1} + a_{2i} + a_{2i-1})
 \end{aligned}$$

Look like this:

$$\begin{aligned}
 & \qquad \qquad \qquad (-2a_1 + a_0 + a_{-1}) 4^0 \\
 & \qquad \qquad (-2a_3 + a_2 + a_1) 4^1 \\
 & (-2a_5 + a_4 + a_3) 4^2 \\
 = & \qquad \dots + 8a_3 + 4a_2 + 2a_1 + a_0 + \dots
 \end{aligned}$$

b and $2b$ can be calculated immediately.

$-b = \bar{b} + 1$ (invert all bits, then plus 1).

Radix-4 Booth Recoding: Example

5-bit multiplication: $(-10) \times (-14) \implies 10110_{(2)} \times 10010_{(2)}$.

Booth-recode: $10010_{(2)} = (-1) \times 4^2 + 1 \times 4^1 + (-2) \times 4^0$

0	0	0	0	0	1	0	0	1	0	$+ \overline{10110}_{(2)} \times 2 \times 4^0$
1	1	1	1	0	1	1	0	1	0	$+ 10110_{(2)} \times 1 \times 4^1 + 4^0 \times 2$
0	0	1	0	0	1	0	0			$+ \overline{10110}_{(2)} \times 1 \times 4^2$
				0	1					$+ 4^2$
<hr/>										
0	0	1	0	0	0	1	1	0	0	$= 140$

Reduce the number of partial products by nearly $1/4$.

Radix-8 Booth Recoding

Even more?

$$\begin{aligned} a \times b &= \sum_{i=0}^{n-1} 2^i a_i b \\ &= \sum_{i=0}^{(n-1)/3} 8^i z_i b \quad (\text{corner case ignored}) \end{aligned}$$

Where

$$z_i = -4a_{3i+2} + 2a_{3i+1} + a_{3i} + a_{3i-1} \in \{\pm 4, \pm 3, \pm 2, \pm 1, 0\}.$$

In this case we need to calculate $3b$, which is not cheap.

Tradeoff between the cost of two stages.

Chisel Implementation

Work In Progress

`https://github.com/sequencer/arithmetic/blob/multiplier/arithmetic/
src/multiplier/WallaceMultiplier.scala`

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