## CS & IT

ENGINERING

Graph Theory

Matching and Covering



Lecture No.10



TOPICS TO BE COVERED



01 Independent set

...

02 Maximal Independent set

. . .

03 Dominance set

...

04 Minimal dominating set

. . .

05 Domination number

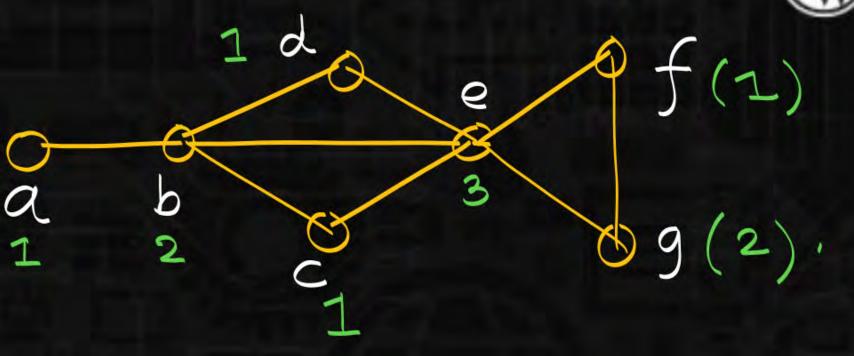


$$x(G) = 3.$$

$$2 \rightarrow \{a, c, d, f\}$$

$$2 \rightarrow \{b, g\}$$

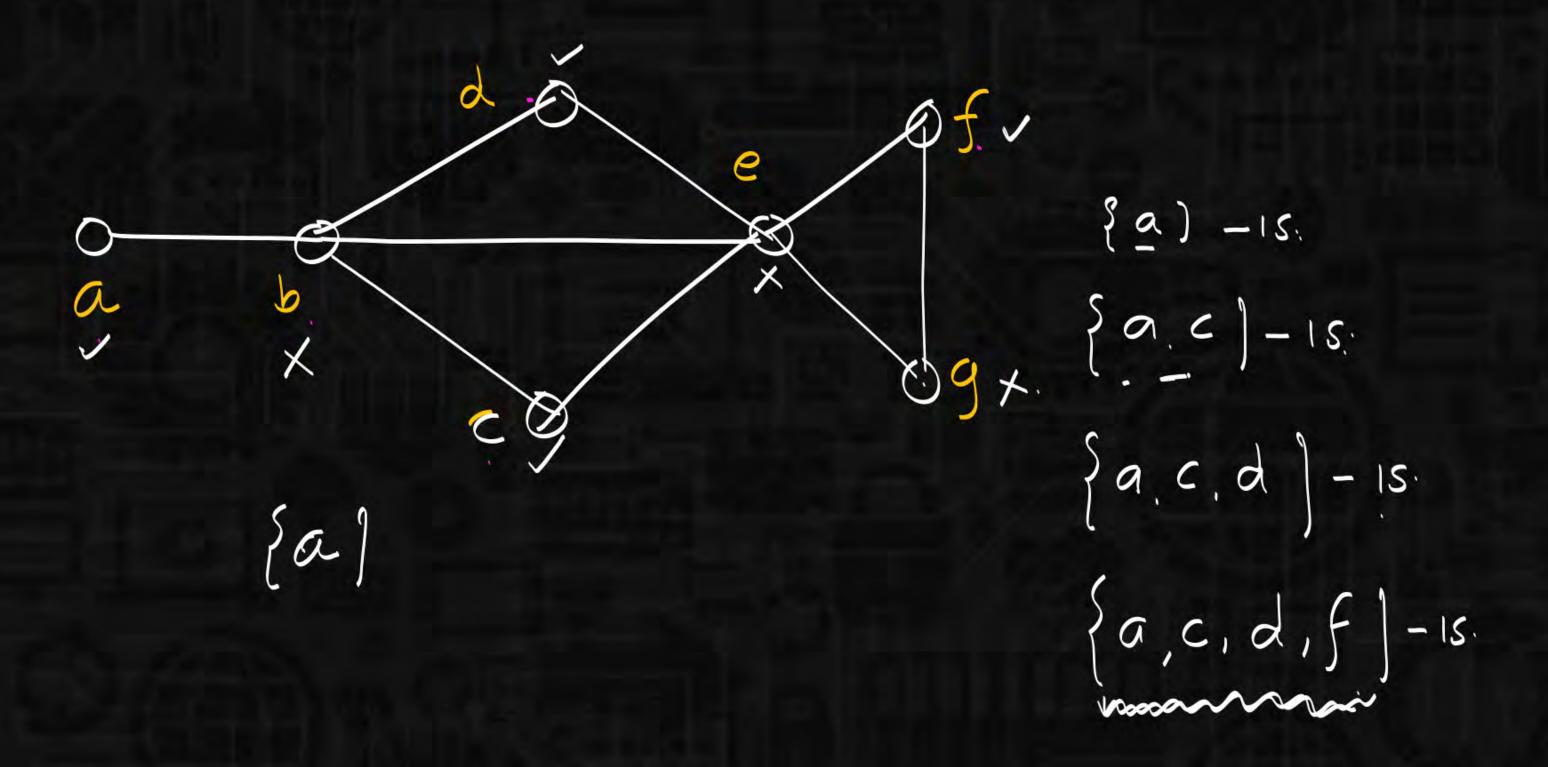
$$3 \rightarrow \{e\}$$



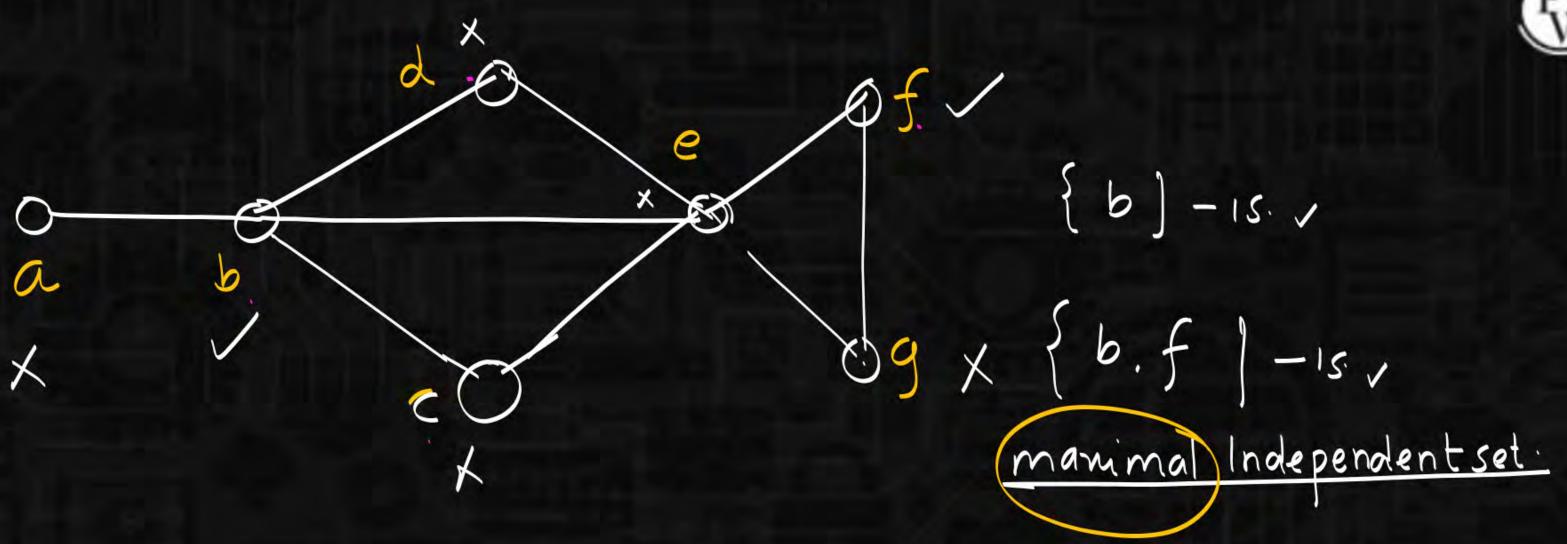


(a,c,d,f) non adjacent vertices. Independent set:

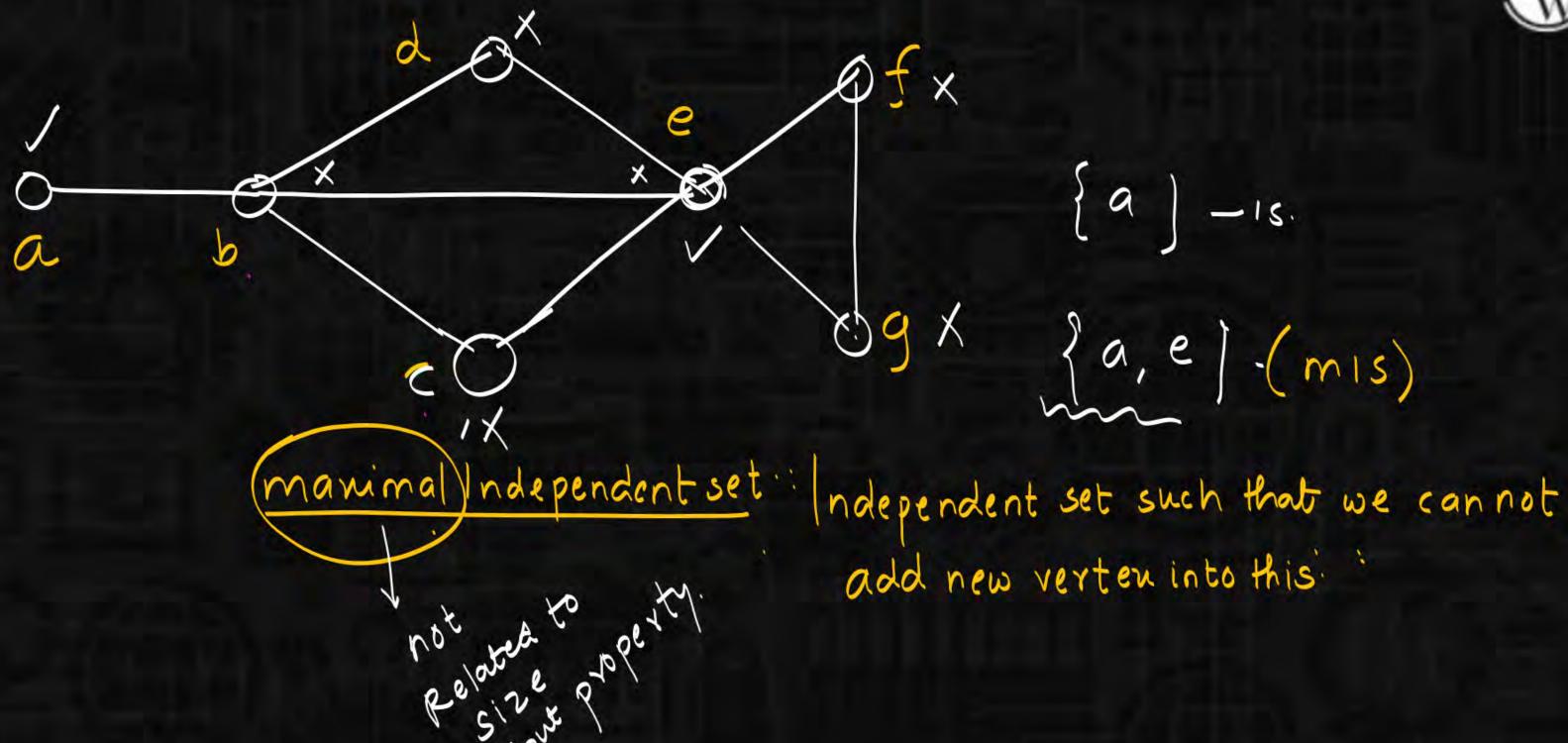


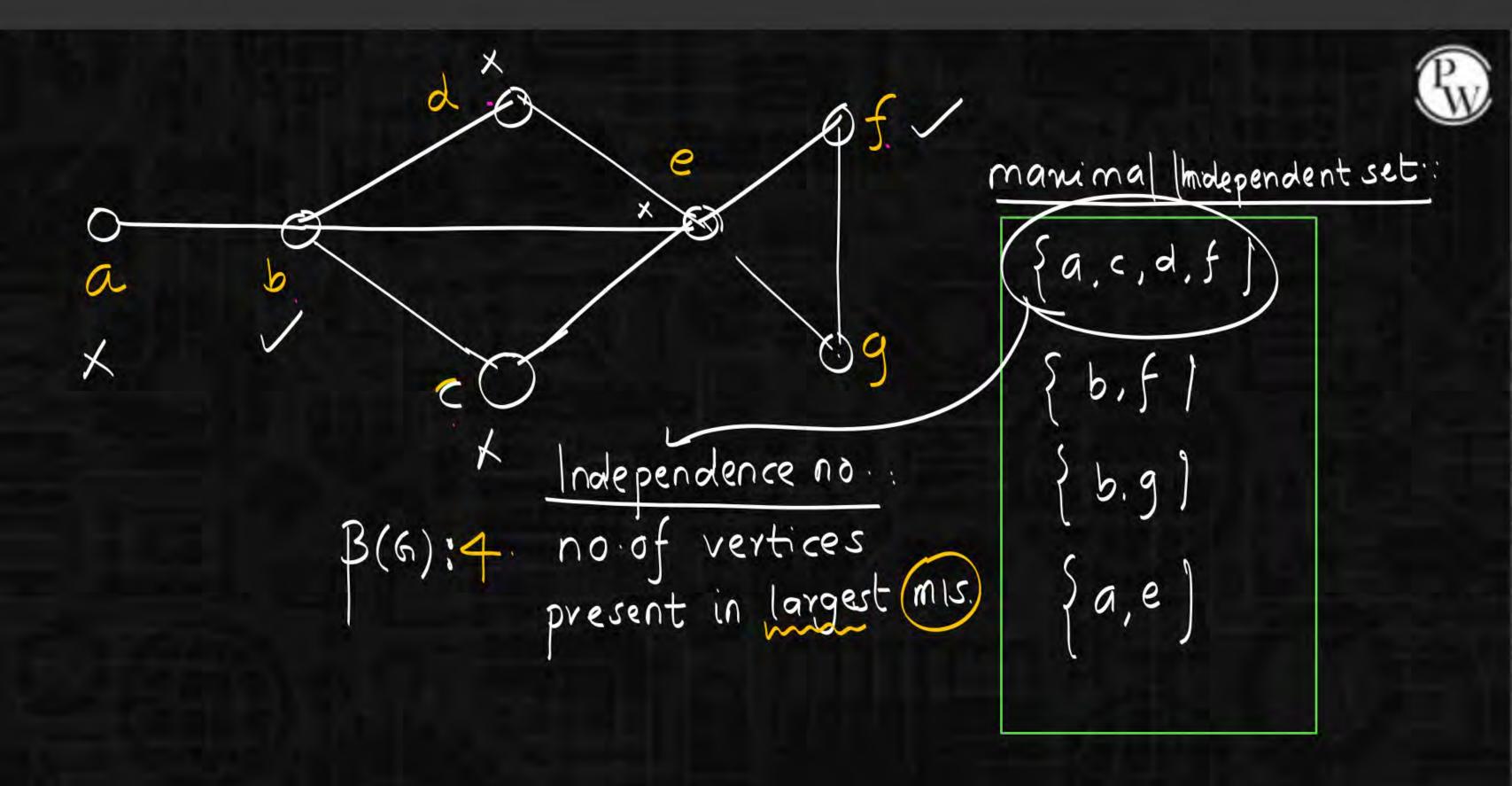




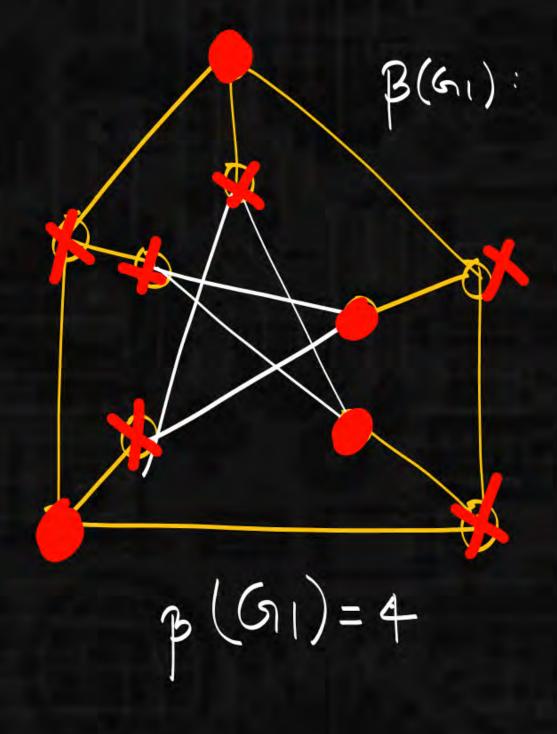


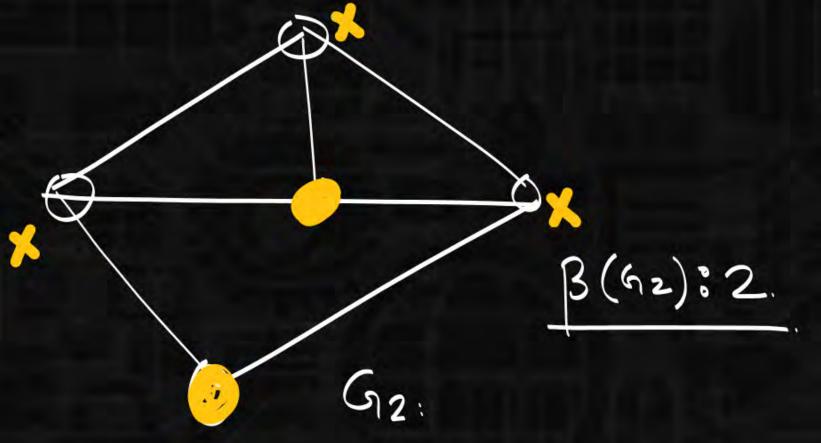




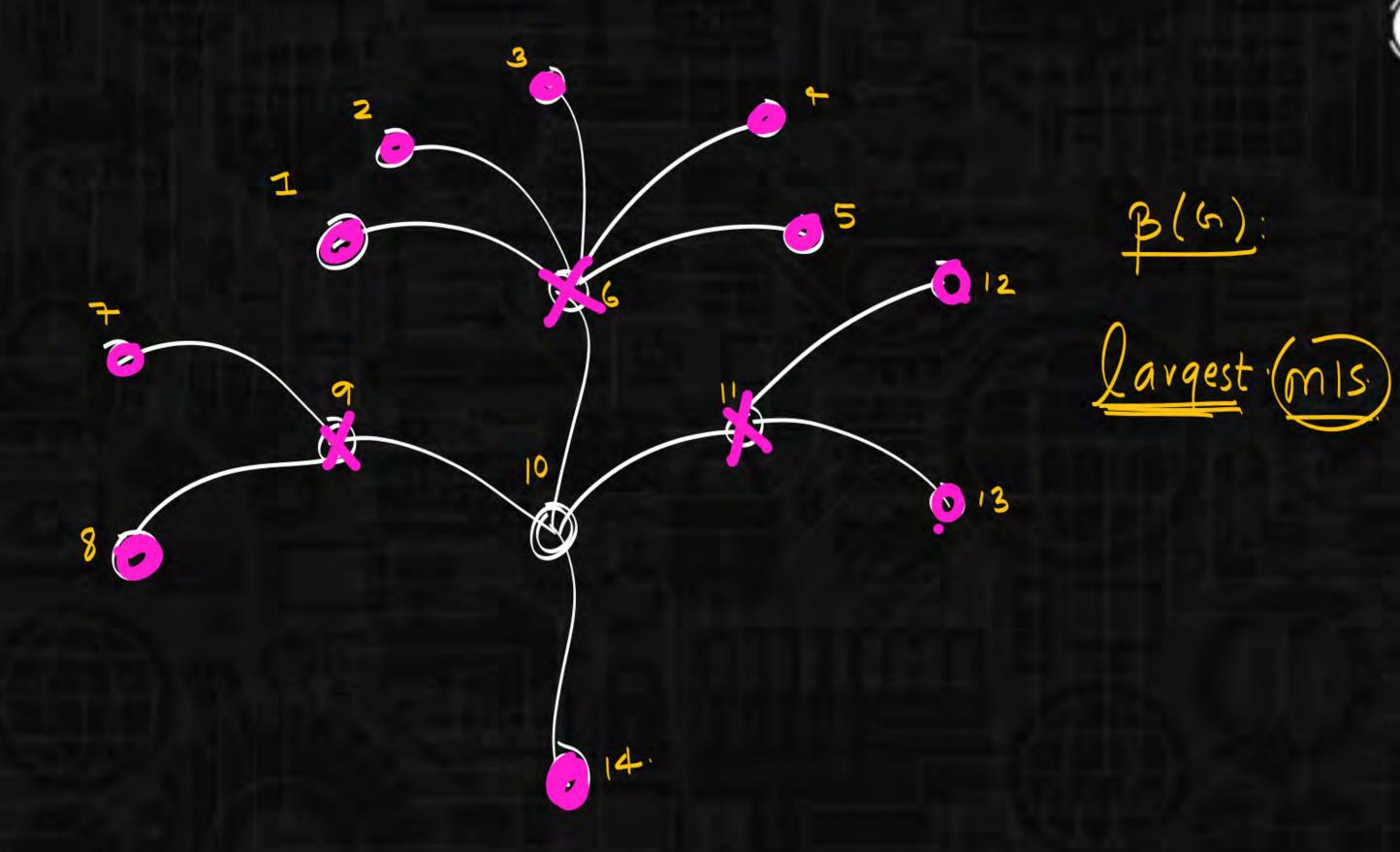




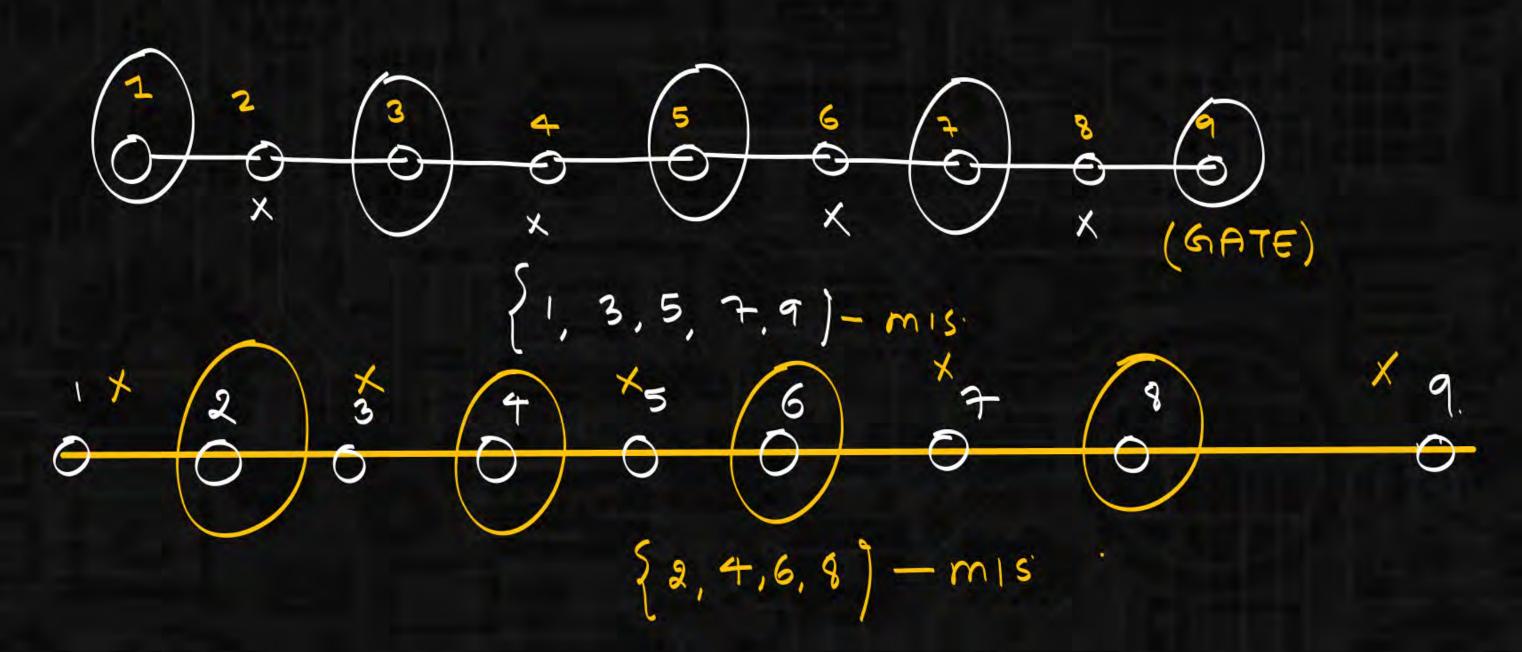




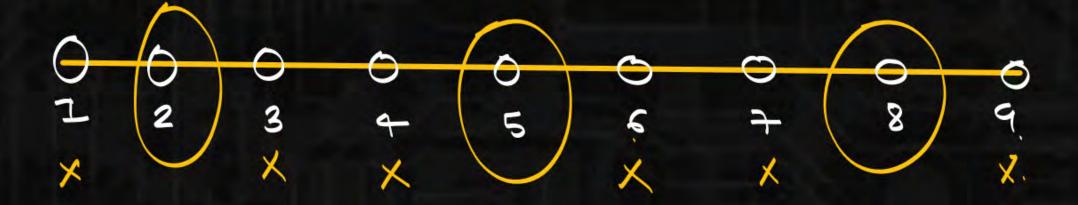












{2,5,8}

Smallest mis Bizeis 3.

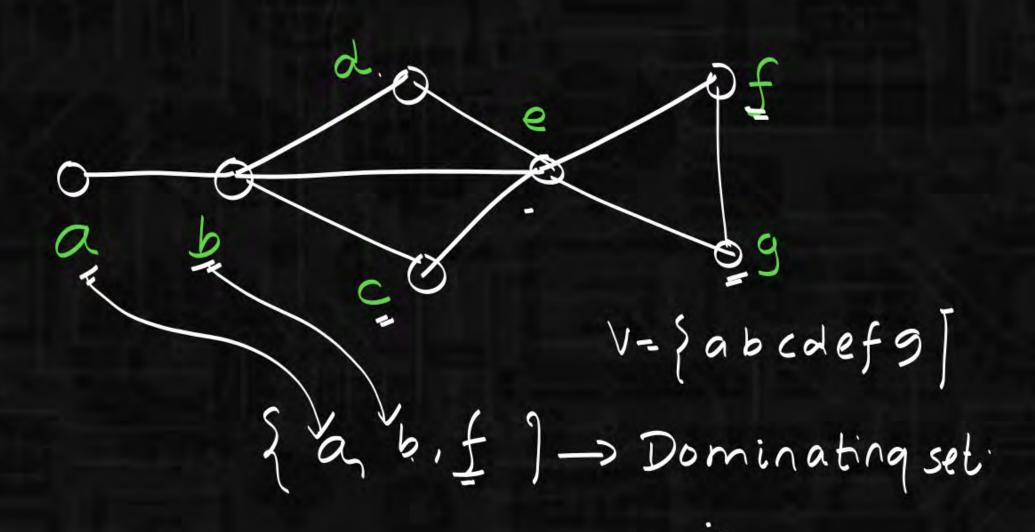


## Dominating set:

Graph G=(V,E) Dominating set D

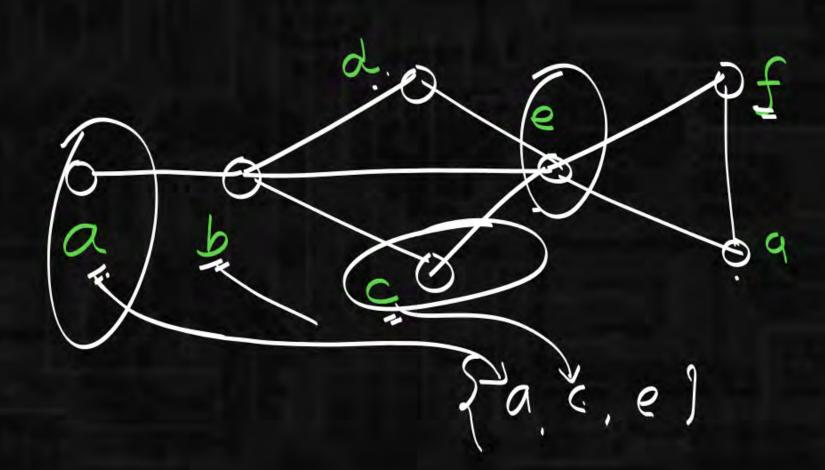
if we take any verten from G either that verten belongs to D or its adjacent will belongs to D.

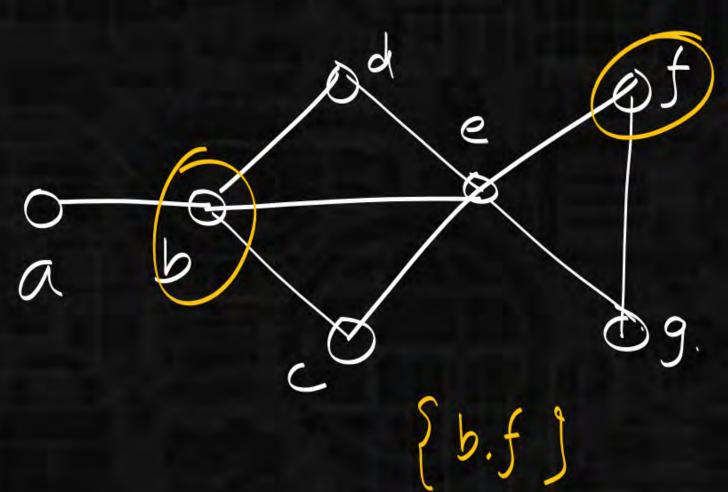




Directly
or belongs
adjacent

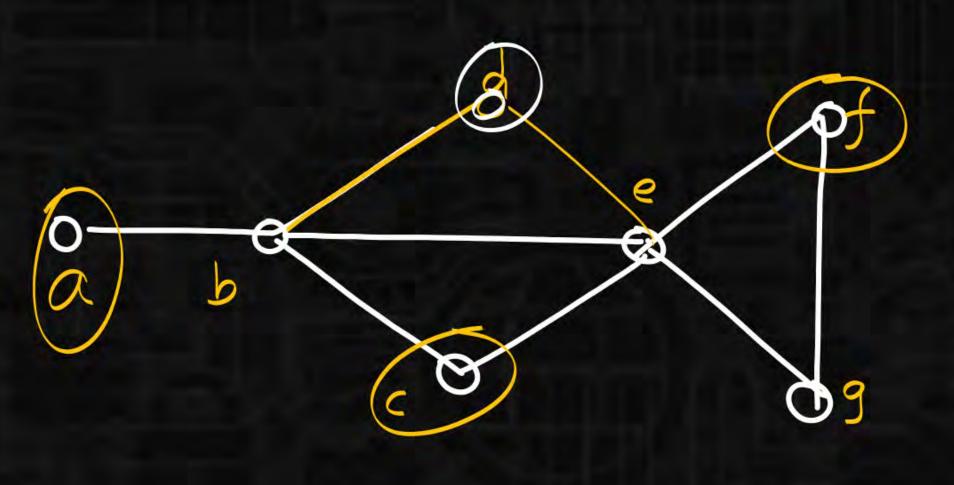






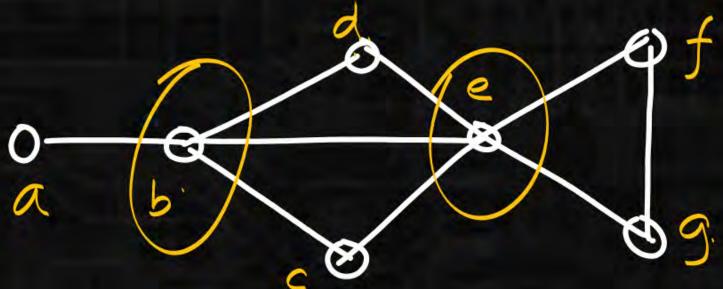


Sabcdefg) -> Ds. {abcdef | -> D.S. fabcaf | -> D.s.





{abcdefg] - Ds {abcdef | -Ds Sábcde | - D.S { b/de | - 0.5 } be ] -> (DS) minimally Dominating set



\* manima | - cannot add

\* minimal -> can not Remove.



## minimal Dominating set:

Dominating set such that we cannot remove new vertex from it.

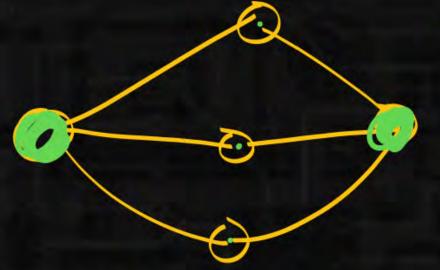
## Domination no.

no of vertices present in Sale! Sale! Sale!

{acdf} 
$$\{b,f\}$$
  $(5)=2$ .  
{b,9}  
{b,e}



Every mis Will always be MDS.



manimal Independent set:

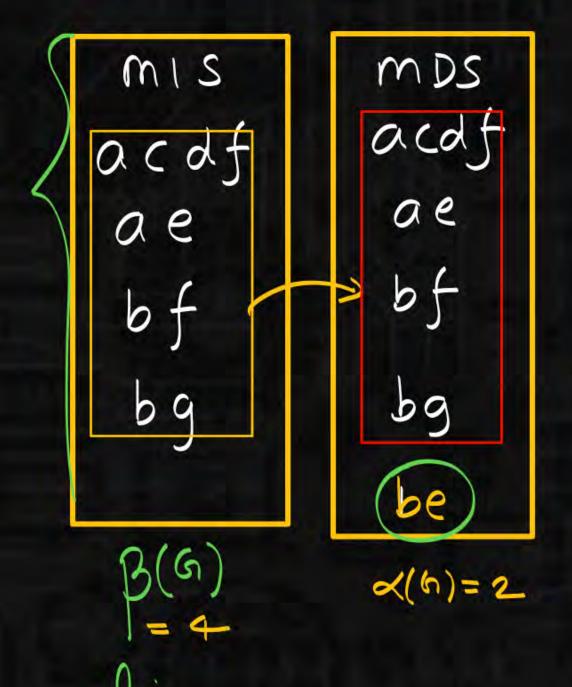
tannot non adjacent

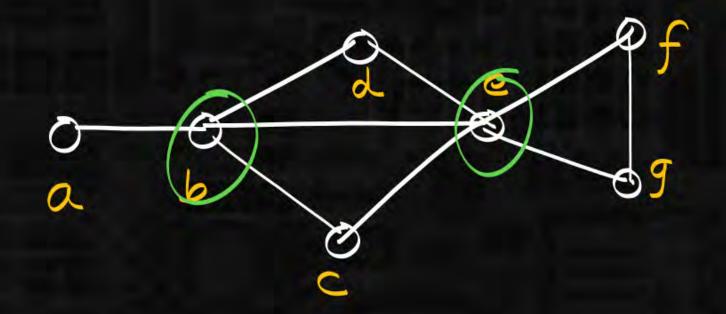
cannot ever

addiever

new

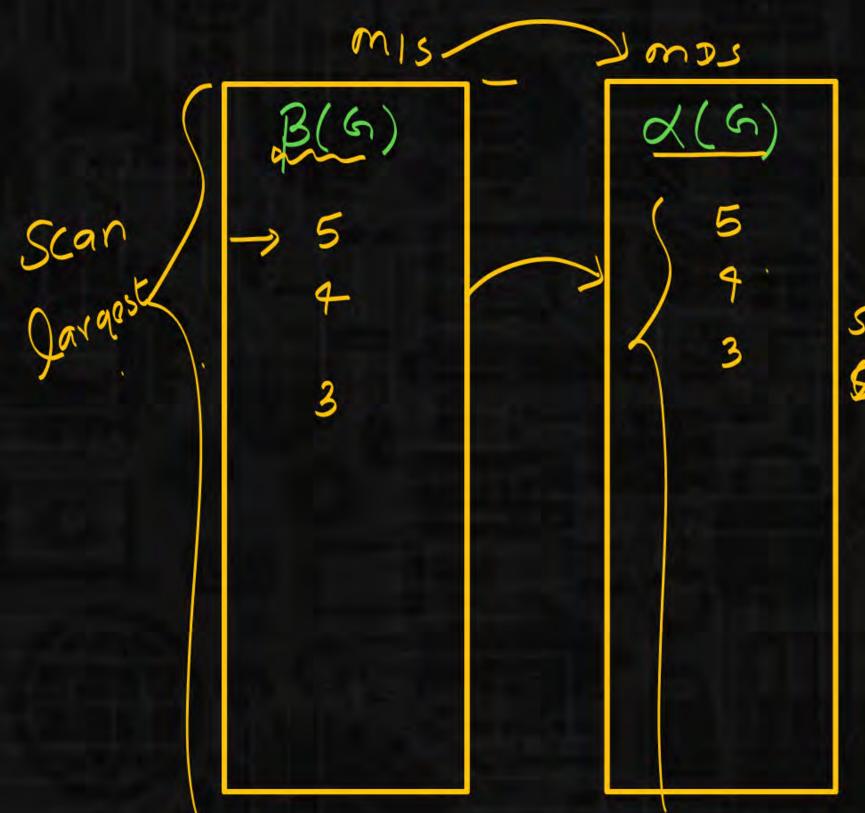






every mis will always be mDs.
but Every mDs need not be mis.
eq: {be}





scan grayest



