

CS & IT ENGINEERING

COMPUTER NETWORKS

Medium Access Control

Lecture No-06

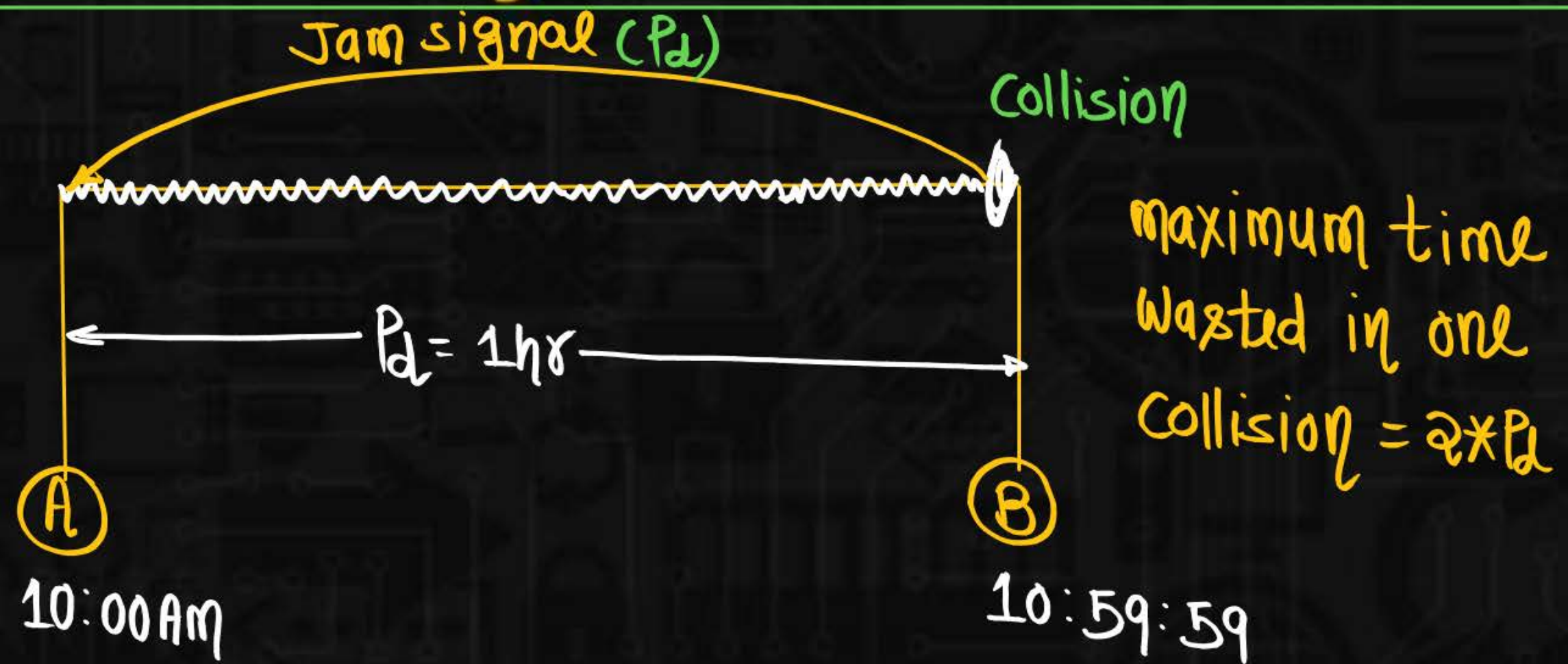


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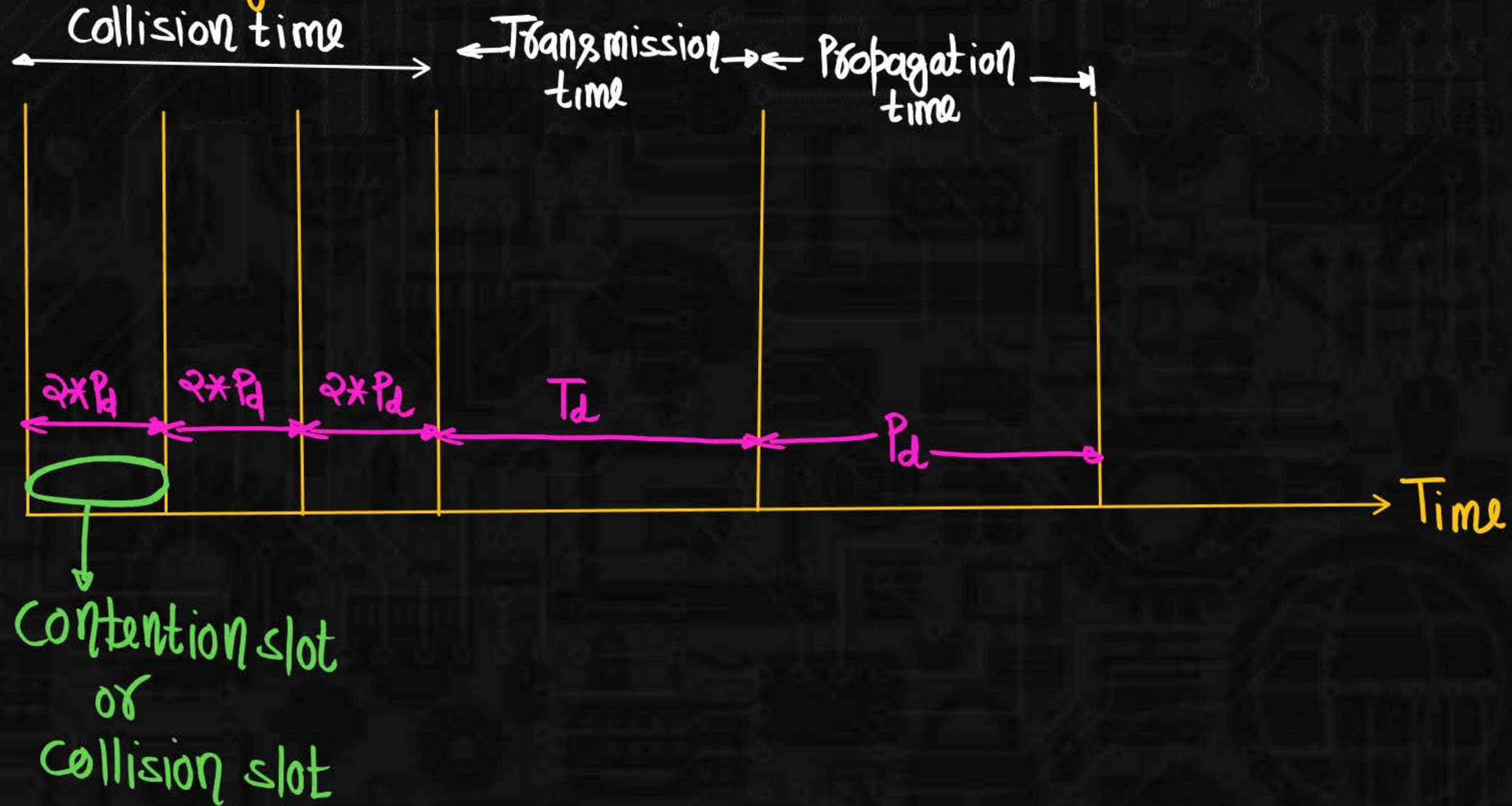
TOPICS TO
BE
COVERED

**Multiple Access
Protocols-6**

Efficiency of Ethernet (CSMA/CD)



Efficiency Calculation of Ethernet (CSMA/CD)



$$\text{efficiency} = \frac{\text{useful time}}{\text{Total time}}$$

$$\eta = \frac{\text{Transmission delay}}{\text{collision time} + \text{Transmission time} + \text{Propagation time}}$$

$$\eta = \frac{T_d}{C * \alpha * P_d + T_d + P_d}$$

$$\eta = \frac{T_d}{e * \alpha * P_d + T_d + P_d}$$

$$\eta = \frac{\cancel{T_d}}{\cancel{T_d} \left[e * \alpha * \frac{P_d}{\cancel{T_d}} + 1 + \frac{P_d}{\cancel{T_d}} \right]}$$

$$\left(e = 2.72 \right)$$

$$\left(\frac{P_d}{T_d} = a \right)$$

C = No. of contention slots
or
No. of collision slot



$$\eta = \frac{1}{c \times 2 \times a + 1 + a}$$

$$\eta = \frac{1}{2 \cdot 72 \times 2 \times a + 1 + a}$$

$$\eta = \frac{1}{5.44a + 1 + a}$$

$$\eta = \frac{1}{6.44a + 1}$$

$$\eta = \frac{1}{1 + 6.44a}$$

$$\eta = \frac{1}{1 + 6.44a}$$

$$= \frac{1}{1 + 6.44 \frac{P_d}{T_d}}$$

$$= \frac{1}{1 + 6.44 \times \frac{d}{U} \times \frac{B}{L}}$$

$B \rightarrow \text{Fixed}, U \rightarrow \text{Fixed}$

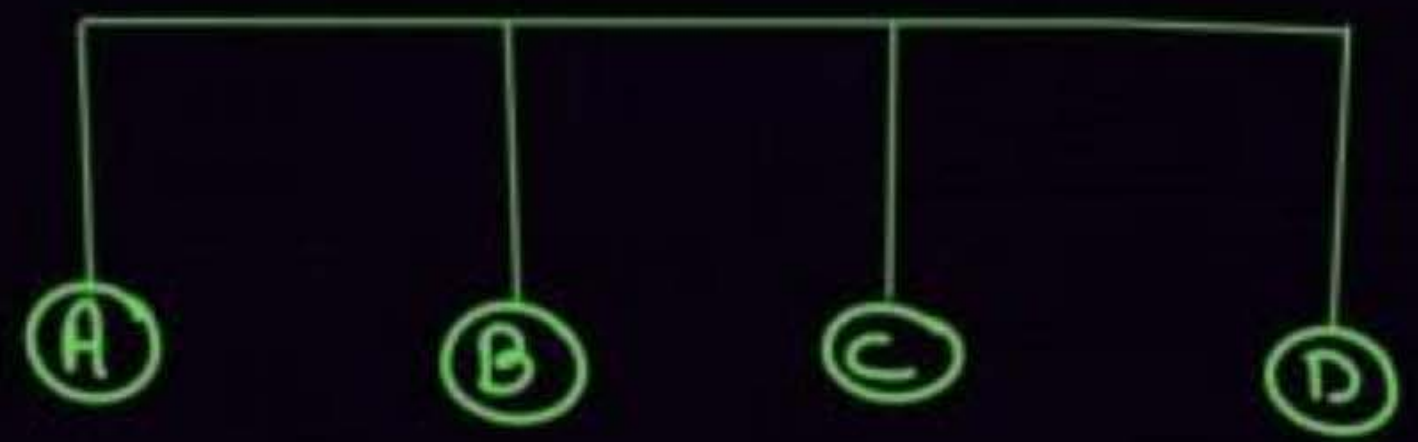
$d \uparrow$

$\eta \downarrow \rightarrow \text{Good for LAN not for WAN}$

$L \uparrow$

$\eta \uparrow \rightarrow \text{Good for Large Packet size}$

1. In Ethernet efficiency is High when Transmission delay is High and Propagation delay is Low
2. In Ethernet efficiency is Low when Propagation delay is High and Transmission delay is Low



N = Total Number of station in Ethernet

p = Probability of station to transfer the data packet

$(1-p)$ = Probability of station Not to transfer the data packet

"For the successful transmission of one station then Remaining $(N-1)$ station should Not transfer the data PKT"

$(1-p)^{N-1}$ = Probability of (N-1) station Not to transfer the data packet

$p(1-p)^{N-1}$ = Probability of success for a single station

$N \cdot p(1-p)^{N-1}$ = Probability of success For any station among 'N' stations

$$P_{\text{succ}} = N \cdot p(1-p)^{N-1}$$

For Max (P_{succ})

$$\frac{d}{dp}(P_{\text{succ}}) = 0$$

$$\frac{d}{dp} \left[N \cdot p(1-p)^{N-1} \right]$$

$$N \left[p \cdot \frac{d}{dp} (1-p)^{N-1} + (1-p) \frac{d}{dp} p \right]$$

$$p = \frac{1}{N}$$

$$\begin{aligned} P_{\text{suc}} &= N \cdot p(1-p)^{N-1} \\ &= \cancel{N} * \frac{1}{\cancel{N}} \left(1 - \frac{1}{N} \right)^{N-1} \\ &= \left(1 - \frac{1}{N} \right)^{N-1} \end{aligned}$$

IF there are sufficiently large
Number of stations i.e. $n \rightarrow \infty$ then
we have -

$$\begin{aligned} \lim_{N \rightarrow \infty} (P_{\text{suc}})_{\text{max}} &= \lim_{N \rightarrow \infty} \left(1 - \frac{1}{N} \right)^{N-1} \\ &= \frac{1}{e} \end{aligned}$$

Number of times we need to try
 Before getting the 1st success = 'e'
 From Here, we conclude -

Average Number of collision that might occur Before a
Successful transmission = e



Q.1

The efficiency of Ethernet

- ☒ A. Increases when propagation delay and transmission delay are low
- ☒ B. Increase When propagation delay and transmission delay are high
- ☒ C. Increase when Propagation delay is low and transmission delay is high
- ☒ D. Increases when propagation delay is high and transmission delay is low

Q.2

Which of the following statements is TRUE about CSMA/CD

[GATE – 2005]



CSMA/CD

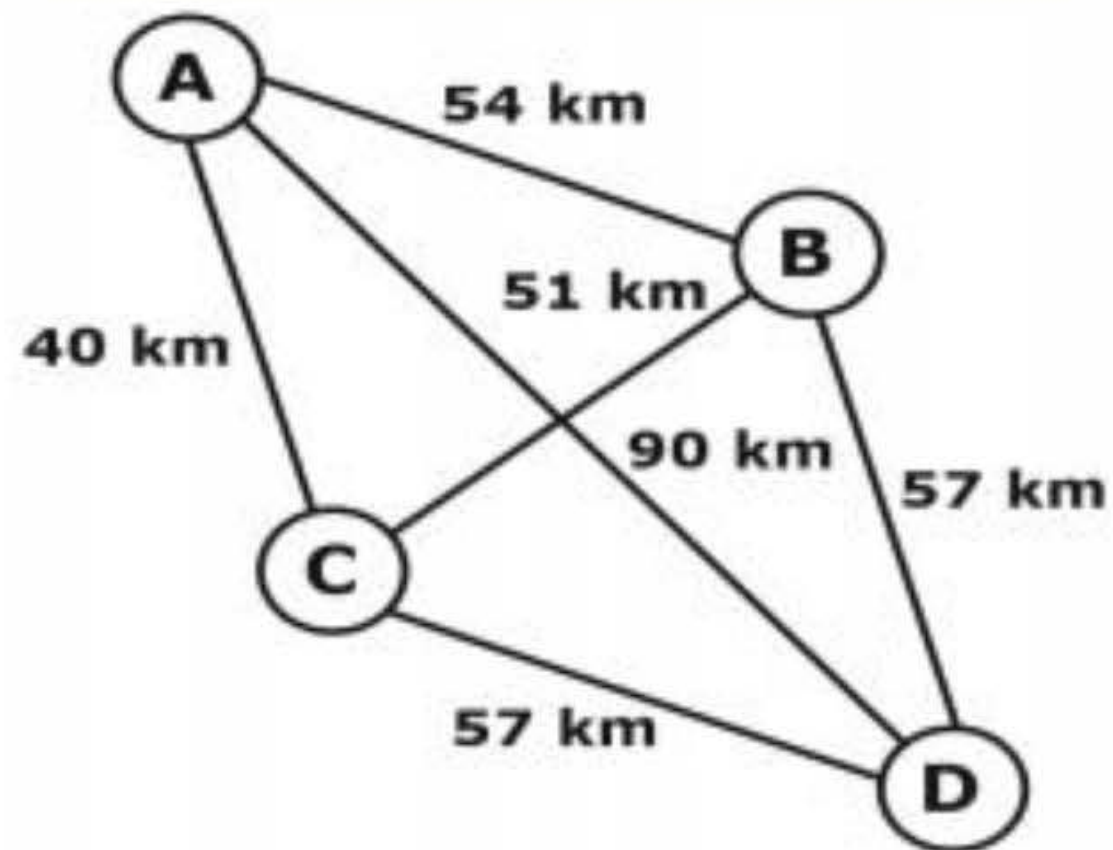
- ☒ A. IEEE 802.11 wireless LAN runs CSMA/CD protocol
- ☒ B. Ethernet is not based on CSMA/CD protocol
- ☒ C. CSMA/CD is not suitable for a high propagation delay network like satellite network
- ☒ D. There is no contention in a CSMA/CD network

Q.3 The network consists of 4 hosts distributed as shown below.

Assume this network uses CSMA/CD and signal travel at 3×10^5 km/sec.

If sender send at 1 Mbps, what could be minimum size of the packet?

- A. 600 bits
- B. 400 bits
- ☒ C. 6000 bits
- D. 1500 bits



$$U = 3 \times 10^5 \text{ km/sec}$$

$$B = 1 \text{ Mbps} = 10^6 \text{ bits/sec}$$

$$L = ?$$

$$T_d(F) \geq 2 * P_d + T_d(JS)$$

$$\frac{L}{B} \geq 2 * P_d$$

$$L \geq 2 * P_d * B$$

$$L \geq 2 * 30 * 10^{-5} \text{ sec} * 10^6 \text{ bits/sec}$$

$$L \geq 2 * 30 * 10 \text{ bits}$$

$$L \geq 600 \text{ bits}$$

$$P_d = \frac{d}{U}$$

$$P_d = \frac{90 \text{ km}}{3 \times 10^5 \text{ km/sec}}$$

$$P_d = 30 \times 10^{-5} \text{ sec}$$

TDMA(Time division Multiple Access)

In time division multiple access-

- Time of the link is divided into fixed size intervals called as time slots.
- Time slots are allocated to the stations in round robin manner.
- Each station transmit its data during the time slot allocated to it.
- In case, station does not have any data to send, its time slot goes waste.

Disadvantage

- If any station does not have data to send during its time slot, then its time slot goes waste.
- This reduce the efficiency.
- This time slot could have been allocated to some other station willing to send data.

Note:

- Effective bandwidth/bandwidth utilization/Throughput
= Efficiency * Bandwidth
- Maximum available effective bandwidth
= Total no of stations* Bandwidth requirement of 1 station.

Problem:

If transmission delay and propagation delay of a packet in Time division multiplexing is 1 msec each at 8Mbps bandwidth, then-

- Find the efficiency?
- Find the effective bandwidth/throughput?
- How many maximum stations can be connected to the network if each station requires 4kbps bandwidth ?

