Our 1) Would linear Search pseudorade to Search an element in a Sosted away with minimum Compositions?

```
Jos ( und l=0 ton)

"y (acci [2] = Value)

"I clement feran d

vecture Value;

eretur -1;
```

Soul is called online Souting why? what about other Sorting algorithm

```
9 desative:
```

Recussive: -

```
Void unsertion_Soxt (unt aux), int n)

if (n<=1) vection;

unsertion_Soxt (aux, n-1);

unt last = aux(n-1);

unt j=n-2;

while (j>=0 18 coun[j]>last)

aux(j+1) = aux(j);
```

an [j+1] = lost;

Hoofst

Inscrition sout is called online sout because it does not need to know anything about values it will soot and information its requested while algorithy is Funning

Other Soxting Alganithm: -

·) Butthe Sort

·) Selection Sost ·) Bruck soxt

·) Heap Soxt. ·) Mesge Sout

Quy 3) comploxity of all sorting algorithm that has been discussed un

Soliding Algo Solicion Solid	Best	woest	Aug.
Bubble Sout	0(n <sup>2</sup> )	0(n2)	$O(n^2)$
9 nsection Sost Heat Sost	O(n) O(nlogn)	0(12)	oln²)
Buick sost	Olnlagn)	$O(n\log n)$	OlnJagn)
Merge Sout	o(nlagn)	olnJogn)	OlnJagn)

Dur 4) Divide all Societing algorithm unto implace stable Online

-	
INPLACE SORTING	STABLE SORTING   ONLINE SORTING
Bubble Soxt	Mersge Sost ghoseution Sort.
Selection Sout	Bubble Sost
gnoeution sost	greedin sout
Dwick Sost	Count Sost
Heap soist	I was to the same of the same

Our. 5) what execusive literative Pseudo Code for beinany secuch what is the Time and Space Complexity of Lineau & Rivery Starch,

Iterative:

and behavy Search (int and I, int &, int Key) while (d<=8)

ind m= l+(8-1)/2;

Time Complexity: -Lineau Search-O(n) Binary Search · O (logn)

Bus-6) white securions to be binary Search (securios) T(n) = T(n|2) + 1 - 1 T(n|2) = T(n|4) + 1 - 2 T(n|4) = T(n|8) + 1 - 3 T(n) = T(n|8) + 1 - 3

$$T(n) = T(n|2) + 1$$
  
=  $T(n|4) + 1 + 1$   
=  $T(n|8) + 1 + 1 + 1$   
 $T(n|2K) + 1 (K dimes)$ 

Let gk=n k=logn T(n)=T(n|n)+logn T(n)=T(1)+logn T(n)=0(logn)

Should

Quy 7) Find two under Such that A[1]+ A[j]=K in runinum time Complaid jox(i=0;i<n; 1++) 1 200 ( int j=0; j <=n; j++) ig (a(1)+ a(j) == K)

psuidy("%d%d", i, j); Dues-8) which sosteng is best jox practical uses! Durck Soxt is the Jastest general purpose soxt in most practical Situations, quet sont is the method of choice. As istability is impostan and in space is available, merge sost might be best, Ourse) what do you mean by inversion in an away? Court the number of inversion in Assay awr [] [7,21,31,8,10,1,20,4,54 using marge Sorot. A pain (A[1], A[1]) is Said to ancested by

A[1], > A[1] Total no. of inversion in given away are is 31 using marge soxt

Our lo) In which Case, quick Sout will give deast & woust lase T.C.

Worst Case O(n2); The worst case occases when the privat cloment is an extreme (smalles/larges) clonered. This hoppens when input away is souted or severes Souted is cither juint or last cloner is selected as pivot

Best Case Oldogn): The best Case occurs when we Select pivot element as a mean element.

Bus II) Woulde recuerrence relation of merge quick Sout in best is worst Can

Mesgr Sout! Best Case - Tln)=2T(n/2)+0(n) } 6(nlogn)

Buick 50xxx)

Best Case - T(n)=2T(n/2)+0(n) - 0(n)ogn)

Worst Case - T(n)=T(n-1)+0(n) - 0(n)2

```
aug 12) selection sost is not stable by defauet but can you write a visua
        of Istable Selection South Code?
            dos(1°=0; 1 < n-1; 1++)
               unt min = " '
                100 (int i = it1; icn ; jt+)
                   ij (almin) sa[i])
               unt Key =a [min];
                 while (minzi)
                  a [min] = a [min - i];
                  a[i]= Key,
Ours 13) Bubble Sout Scans away even when away is sould. Can you
       modify the bubble Sost So that it closs not Scan the whole away
       Once it is sorted.
```

```
Void bubble Sost (und auxC], und n)

Jos (und i=0; i< n-1; i++)

und Swaps =0;

Jos (und i=0; i<n-1-i; j++)

y (aux[i] > aux[i+1])

und d = aux[i];

aux[i] = aux[i+1];

aux[i] = d;

Swap ++;

y (swap = =0) evolueum

i
```

Shershit