# 1 Bianchi

## 1.1 probabilities

 $\pi$ , probability of transmission, p, probability of collision,  $b_{i,k}$  stationary probability of state i,k:

$$\begin{split} p &= 1 - (1 - \pi)^{N - 1} \\ \pi &= \frac{2}{1 + W_{\min} + pW_{\min} \sum_{k=0}^{m-1} (2p)^k} \\ &= \frac{2(1 - p)}{(1 - 2p)(W_{\min} + 1) + pW_{\min}(1 - (2p)^m)} \\ b_{i,k} &= \frac{CW_i - k}{CW_i} \cdot \begin{cases} (1 - p) \sum_{j=0}^m b_{j,0} & i = 0 \\ p \cdot b_{i-1,0} & 0 < i < m \\ p \cdot (b_{m-1,0} + b_{m,0}) & i = m \end{cases} \end{split}$$

# 1.2 Saturation throughput

$$\begin{split} \tau &= \frac{E[\text{Payload Transmitted by user i in a slot time}]}{E[\text{Duration of slot time}]} \\ &= \frac{P_{\text{s}}P_{\text{tr}}L}{P_{\text{s}}P_{\text{tr}}T_{\text{s}} + P_{\text{tr}}(1-P_{\text{s}})T_{\text{c}} + (1-P_{\text{tr}})T_{\text{id}}}, \\ P_{\text{s}} &= \frac{N\pi(1-\pi)^{N-1}}{1-(1-\pi)^{N}}, \\ P_{\text{tr}} &= 1-(1-\pi)^{N}, \\ T_{\text{s}} &= t_{\text{header}} + t_{\text{payload}} + \text{SIFS} + t_{\text{ACK}} + \text{DIFS} + \sigma, \\ T_{\text{c}} &= t_{\text{header}} + t_{\text{payload}} + \text{SIFS} + \sigma \end{split}$$

# 2 Trunk dimensioning

For a trunk of N channels, an offered load  $A = \lambda E[X]$ X the call duration, Y the call arrival per sec  $\sim$ Poisson( $\lambda$ ) and  $\rho$  the traffic carried by each channel:

$$\begin{split} P_{\rm Blocking} &= P(\text{Drop a call because busy line}) \\ &= \frac{A^N}{N! \sum_{i=0}^N (\frac{A^i}{i!})} \\ \rho &= \frac{(1-P_{\rm blocking})A}{N} \end{split}$$

Cellular efficiency  $E = \frac{Conversations}{cells \times MHz}$ 

# 3 Cellular Geometry: Hexagons

**Area**:  $A = 1.5R^2\sqrt{3}$ 

Distance btw. adjacent cells:  $d = \sqrt{3}R$ 

### 3.1 Co-channel interference

Co-channel reuse ratio :  $Q = \frac{D}{R} = \sqrt{3N}$  with D the distance to the nearest co-channel cell, R the radius of a cell and N the cluster size.

Signal to Interference ratio (SIR): 
$$SIR = \frac{S}{T} = \frac{S}{\sum_{i=1}^{i_0} I_i}$$
. With  $S$  the desired signal power,  $I_i$  the interference power from the  $i$ th interfering co-channel base-station,  $i_0$  the number of co-channel interfering cells.

Signal to Interference plus Noise ratio (SINR) : SINR = $\frac{S}{I+N_0}$ 

Average received power  $P_r: P_r = P_0(\frac{d}{do})^{-\alpha}$  or

 $P_r(dBm) = P_0(dBm) - 10\alpha \log(\frac{d}{d\alpha})$  with  $P_0$  the power received from a small distance  $d_0$  from the transmitter and  $\alpha$  the path loss exponent.

SIR in the corner of a cell :  $\frac{S}{I} = \frac{R^{-\alpha}}{\sum_{i=0}^{i_0} \frac{R^{-\alpha}}{R^{-\alpha}}}$ 

First interfering layer approximation :  $\frac{S}{I} = \frac{(\frac{D}{R})^{\alpha}}{\frac{1}{l_{\Omega}}} =$  $\frac{(\sqrt{3N})^\alpha}{i_0}$ eg. =  $(\frac{D}{R})^2\frac{1}{2}$  for two first layer interferers (cell divided into 3 sectors with directional

# 3.2 Capacity of a cellular network

For  $B_t$  the total allocated spectrum and  $B_c$  the channel bandwidth:

$$m = \frac{B_t}{B_c \frac{Q^2}{3}} = \frac{B_t}{B_c \left(\frac{6}{3 \frac{\alpha}{2}} \left(\frac{S}{I}\right)_{\min}\right)^{\frac{2}{\alpha}}}$$

For a cluster size N,  $N = (i+j)^2 - ij$  for i, j = 0, 1, 2, ...and number of channels C we have  $m = \lfloor \frac{C}{N} \rfloor$ 

## 3.2.1 CDMA Capacity: single cell case

For the bitrate R, available bandwidth W, noise spectral density  $N_0$ , thermal noise  $\eta$ , received user signal (at base station) S, we have a possible number N of users:

$$N = 1 + \frac{W/R}{E_b/N_0} - (\frac{\eta}{S})$$

With a duty cycle  $\delta$  (Discontinuous transmission mode: takes advantage of intermittent nature of speech):

$$N = 1 + \frac{1}{\delta} \frac{W/R}{E_b/N_0} - (\frac{\eta}{S})$$

And if we have m sectors, the effective capacity becomes m.N

### 3.2.2 CDMA multiple cells

Frequency reuse factor on the uplink  $f = \frac{N_0}{N_0 + \sum_i U_i N_{ai}}$ where  $N_0 = \text{total interference power received from } N-1$ in-cell users,  $U_i$  = number of users in thei<sup>th</sup> adjacent cell and  $N_{ai}$  = average interference power from a user located in the  $i^{\rm th}$  adjacent cell

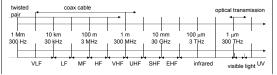
Average received power from users in adjacent cell  $N_{ai} =$  $\sum_{i} N_{ij}/U_{i}$  where  $N_{ij}$  = power received at the base station of interest from the  $i^{th}$  user in the  $i^{th}$  cell

### 4 Noise

Categories: Thermal Noise, Intermodulation Noise, Cross-talk, Impulse Noise.

Thermal Noise  $N_0 = kT$  (W/Hz)For signal power S, bitrate  $R, k = 1.3806 \cdot 10^{-23} JK^{-1}$ the Boltzmann constant and T the temperature:  $\frac{E_b}{M^2}$  $\frac{S/R}{N_0} = \frac{S}{kTR}$ 

# 5 Wireless Misc Stuff



Mobile IP Requirements : Transparency, Compatibility, 5.4.1 Pure ALOHA Security, Efficiency, Scalability.

Mobile IP Issues : Security(Authentication to FA is problematic), Firewalls, QoS

Network Layers Top-down: Application, Transport, (HIP layer), Network, Data-link, Physical.

### 5.1 Ad-hoc Netowrks

**Upper Bound for the Throughput** If we have n identical randomly located nodes each capable of transmitting W bits/s. Then the throughput  $\lambda(n)$  obtainable by each node for a randomly chosen destination is

Routing proactive: DSDV, OLSR. reactive: AODV,

### 5.2 Antennas & Propagation

Free space propagation, received power:  $P_R$  =  $P_{\rm T}\frac{A_{\rm R}}{4\pi d^2}\eta_{\rm R}$  with  $\eta_{\rm R}$  an efficiency parameter,  $A_{\rm R}$  the receiving antenna area.

Focusing capability, depends on size in wavelength  $\lambda$ :  $G_T = 4\pi \eta_T A_T / \lambda^2$ 

Directional emitter, received power:  $P_{\rm R}$ 

Free space received power:  $P_{\rm R} = P_{\rm T} G_{\rm T} G_{\rm R} (\frac{\lambda}{4\pi d})^2$ 

Loss: 
$$L = \frac{P_T}{P_R} = \frac{(4\pi d)^2}{G_R G_T \lambda^2}$$
  
 $c = 3 \cdot 10^8$   
Parabola:  $G = \frac{7A}{2}$ 

Mobnet Decibels :  $B = 10 \log(\frac{P}{P_0})$ 

**Propagation modes** Ground Wave:  $f \leq 2 \text{ Mhz}$ , Sky Wave, Line of Sight: f > 30 Mhz

### 5.2.1 Line of sight equations

Horizon distance d[km] in **kilometers**, antenna height h[m] and refraction adjustment factor K = 4/3:

Optical LOS:  $d = 3.57\sqrt{h}$ 

Effective LOS:  $d = 3.57\sqrt{Kh}$ 

Max LOS distance for two antennas :

$$3.57(\sqrt{Kh_1} + \sqrt{Kh_2})$$

## 5.3 Free Space Loss

Free space loss, ideal isotropic antenna:

$$\frac{P_t}{P_r} = \frac{(4\pi d)^2}{\lambda^2} = \frac{(4\pi f d)^2}{c^2}$$

Free space loss equation can be recast:

$$L_{DB} = 10 \log \frac{P_t}{P_r} = 20 \log(f) + 20 \log(d) - 147.56 dB$$

Free space loss accounting for gain of other antennas:

$$\frac{P_t}{P_r} = \frac{(4\pi d)^2}{G_r G_t \lambda^2} = \frac{(cd)^2}{f^2 A_r A_t}$$

 $G_t = gain of transmitting antenna$  $A_r$  = effective area of receiving antenna

## 5.4 Protocol performances

G: Total load, S arrival rate of new packets.

$$P(k \text{ transm. in } 2Xs) = \frac{(2G)^k}{k!}e^{-2G}$$
 
$$S = G \cdot P(0) = Ge^{-2G}$$

### 5.4.2 Slotted ALOHA

Probability of k packets generated during a slot:

$$P(k) = \frac{G^k e^{-G}}{k!}$$

Throughput:

$$P(1) = Ge^{-G}$$

## 5.4.3 CSMA

Non-persistent If channel is busy, directly run back off algorithm.

**p-persistent** If it is busy, they persist with sensing until the channel becomes idle. If it is idle:

- With probability p, the station transmits its packet
- With probability 1-p, the station waits for a random time and senses again

Performance of Unslotted nonpersistent CSMA : For a = $t_{\text{prop}}/X$ , the normalized one-way propagation delay.

$$S = \frac{G^{-aG}}{G(1+2a) + e^{-aG}}$$

Performance of Slotted nonpersistent CSMA

$$S = \frac{aG^{-aG}}{1 - e^{-aG} + a}$$

# 5.5 DOMINO Cheating detection

Cheating Method	Detection Test	
Frame scrambling	Number of retransmissions	
Oversized NAV1	Comparison of the declared	
	and actual NAV values	
Transmission be-	Comparison of the idle time af-	
fore DIFS	ter the last ACK with DIFS	
Backoff manipula-	Actual Backoff/ Consecutive	
tion	Backoff	
Frame scrambling	Periodic dummy frame injec-	
with MAC forging	tion	

# 5.6 Forward Error Correction (FEC)

Redundancy in packets to allow limited error correction at the receiver: used in 802.11a (Convolutional), HS-DPA (Turbo Codes) and 802.11n (LDPC).

# 6 TCP

## 6.1 Standard

**Tahoe** Basic TCP. Three duplicate ACK's provoke fast retransmit (resend 1st missing packet), set ssthresh to cwnd/2, cwnd to 1 and provoke slow start.

**Reno** Three duplicate ACK's provoke fast retransmit. ssthresh to cwnd/2, cwnd to ssthresh + 3 and enter fast recovery.

Fast Recovery Increase cwnd by 1 segment for every received duplicate ACK. (Warning, unlogical: When new ACK is received, cwnd = ssthresh and enter congestion avoidance). If a timeout occurs, set cwnd to 1 and enter slow start.

New Reno Fast Recovery More intelligent fast recovery where you remember the last received ACK.

Indirect TCP (I-TCP) Connection split at FA. Standard TCP on the wire line, wireless optimized TCP on the wifi side: shorter timeout, faster retransmission. Loss of end-to-end semantics, security issues.

Mobile TCP (M-TCP) Split connection at FA. Monitor packets, if a disconnect is detected, report receiver window = 0: sender will go into persist mode and doesn't timeout or modify his congestion window. Preserves end-to-end semantics. Disadv.: wifi losses propagate to the wire network, link-errors pkt loss must be resent by sender, security issues. Summary: only handles mobility errors, no transmission errors.

Snooping-TCP TCP-aware link layer: Split connections, FA buffers and retransmits segments, does not ACK buffered packets (preserves end-to-end semantics).

Transaction oriented TCP (T-TCP) TCP phases: connection setup, data transmission, connection release. T-TCP combines does steps and only 2-3 packets are needed for short messages. Efficient for single packet transactions, but requires TCP modifications on all

# 7 Security

Security Requirements : Confidentiality, Authenticity, Replay Detection, Integrity, Access Control, Jamming Protection.

GSM Shared secret and challenge responses, one-way authentication.

**3GPP** (Improvements from GSM) Two-way authentication, avoid fake base station, cipher keys and auth data is now encrypted, integrity. Privacy/Anonymity not completely protected however.

# 8 Privacy

Privacy Related Notions Anonymity, untraceability, unlinkeability, unobservability, pseudonymity

Best to worst against information leakage: GPS: no thirdparty, determined 'alone'. Cell-ID: requires the operator database that is relatively protected (they won't easily mine you). Wireless: requires one or several third-party owned databases that can track you, and it is relatively precise due to short radio range.

## 8.1 Privacy Metrics

**Entropy-Based Anonymity** A the anonymity set,  $p_x$  the probability for an external observer that the action was performed by x:

$$\sum_{\forall x \in A} p_x \log(p_x)$$

**Entropy-Based Unlinkability**  $I_1, I_2$ , sets of elements to be related,  $p_r$ , the probability two elements are related for an external observator:

$$\sum_{\forall R \subseteq I_1 \times I_2} p_r \log(p_r)$$

### **8.2 RFID**

Standard tags possibilities : Kill, Sleep, Rename, Block, (Legislation).

Crypto enable tags possibilities : Tree-approach, synchronization approach, hash chain based approach.

Singulation (determining which tags are present around the reader) Binary tree walking: reader first asks the tags to emit the first bit of their ID. If every answer is 0 (or 1) the reader knows on which side the ID's are. This is done recursively until all ID's are determined.

Privacy zone A tag ID can be changed so that it lies in the private zone of the tree. A special device simulates collisions for every query in this area, so an exhaustive search would be required to find a tag.

Pseudonyms Tags can be set to use different ID's that an authorized reader would know how to correlate. To avoid having too complex tags, the reader will generally be responsible for refilling the pseudonyms. This will be done in cleartext and assumes an attacker does not always listen.

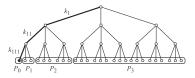
KISS Keep It Simple and Stupid

LDPC Low Density Parity Check

LEAP Light EAP

## 8.2.1 Key Tree

Tags are the leaves of a tree with branching factor b and depth d, and each edge to arrive to a tag has an associated key: hence, a tag has d associated keys. Maximize branching factor at the first level for strong anonymity.



Anonymity set has minimum size of 1, maximum size of all the tags. Compromising a tag yields all the keys leading to it and permit to partition the other tags (neighbors in the tree share common keys):  $P_0$  contains the compromised tag,  $P_1$  contains the compromised tag's brothers not being in  $P_0$ , etc. Tags that belong to larger partitions have better privacy (e.g. tags in  $P_3$  are not distinguishable, attacker only knows they don't use  $k_1$ .)

Expected size of the anonymity set for a random tag : for n the total number of tags and  $|P_i|/n$  the probability of selecting a tag from partition  $P_i$ 

$$\bar{S} = \sum_{i=0}^{d} \frac{|P_i|}{n} |P_i| = \sum_{i=0}^{d} \frac{|P_i|^2}{n}$$

Normalized expected anonymity : Using  $n = b^d$  and  $|P_0| = 1, |P_1| = b - 1, |P_2| = (b - 1)b, \dots, |P_l| =$ 

$$R = \frac{\bar{S}}{n} = \sum_{i=0}^{d} \frac{|P_i|^2}{n^2} = \frac{b-1}{b+1} + \frac{2}{(b+1)n^2}$$

For **one** tag in  $P_i$ , the linkability probability is  $1/|P_i| \to \text{global linkability in } P_i$  is  $|P_i| \frac{1}{|P_i|} = 1$ . For l partitions, the probability that two transactions from a randomly chosen tag are linkable is (with  $n = b^d$ ):

$$\frac{1}{n}\sum_{i=1}^{l}(|P_i|\frac{1}{|P_i|}) = \frac{l}{n}$$

# 9 Comparisons

Approach	SDMA	TDMA	FDMA	CDMA
Idea	segment space into cells/sectors	segment sending time into disjoint time-slots, demand driven or fixed patterns	segment the frequency band into disjoint sub-bands	spread the spectrum using orthogonal codes
Terminals	only one terminal can be active in one cell/one sector	all terminals are active for short periods of time on the same frequency	every terminal has its own frequency, uninterrupted	all terminals can be active at the same place at the same moment, uninterrupted
Signal separation	cell structure, directed antennas	synchronization in the time domain	filtering in the frequency domain	code plus special receivers
Advantages	very simple, increases capacity per km²	established, fully digital, flexible	simple, established, robust	flexible, less frequency planning needed, soft handover
Dis- advantages	inflexible, antennas typically fixed	guard space needed (multipath propagation), synchronization difficult	inflexible, frequencies are a scarce resource	complex receivers, needs more complicated power control for senders
Comment	used in all cellular systems	standard in fixed networks, together with FDMA/SDMA used in many mobile networks	typically combined with TDMA (frequency hopping patterns) and SDMA (frequency reuse)	higher complexity

This amazing cheat-sheet was brought to you by Julien Perrochet. Christopher Chiche and Tobias Schlatter. Follow us on GitHub:

https://github.com/Shastick/mobnet2012!

	-	
ACO Authenticate	d Cipher Offs	et
AIFS Arbitrary Int	er-Frame Spa	ice
AMF Authentication	on and Key m	anageme
AODV Ad Hoc Vector	On-demand	Distanc
AP Access Point		
AP Access Point		
ATIM Ad-hoc Traf	fic Indication	Map
AUTN Authenticat		-
AV Authentication		
BO BackOff		
BSSID Basic Service	e Set Identifi	ier
BSS Basic Service		
CARMA Collision tion Multip	Avoidance ar	nd Resolu
CA Collision Avoi	dance	
CCA Clear Channe	el Assessment	
CDMA Code Divisi	on Multiple	Access

DA Destination Address

DBPSK Differential Binary Phase Shift

Shift Keving DSDV Destination Sequenced Distance Vector DSRC Dedicated Short Range Communica-DSR Dynamic Source Routing DSSS Direct Sequence Spread Spectrum DS Differentiated Service DS Distribution System DTIM Delivery Traffic Indication Map DoS Denial of Service CH Correspondant Host EAP-TLS TLS over EAP
EAPOL EAP Over LAN CN Correspondant Node COA Care-Of Address EAP Extensible Authentication Protocol CRC packet received CoRreCtly EDCA Enhanced Distributed Channel Ac-CSMA/CD CSMA with Collision Detection EHF Extra High Frequency CSMA Carrier Sense Multiple Access EPC Electronic Product Code CTS Clear To Send ESP Encapsulating Security Payload Contention Window ESS Extended Service Set DAMA Demand-Assigned Multiple Access FAMA Floor Acquisition Multiple Access

munications

DNS Domain Name System

DH Diffie-Hellman

FA Foreign Agent

FDD Frequency Division Duplex

DOPSK Differential

DCF Distributed Coordination Function FDMA Frequency Division Multiple Access DECT Digital Enhanced Cordless Telecom FEC Forward Error Correction FHSS Frequency Hopping Spread Spec-DHCP Dynamic Host Configuration Proto-FQDN Fully Qualified Domain Name GFSK Gaussian Frequency Shift Keying GMK Group Master Key Quadrature Phase

GPRS General Packet Radio Service GSM Global System for Mobile Communi-MAP Mobility Anchor Point cation MD Mobile Device **HA** Home Agent MF Medium Frequency HCCA HCF Controlled Channel Access MH Mobile Host HCF Hybrid Coordination Function MIB Management Information Base **HF** High Frequency MIC Message Integrity Code HIP Host Identity Protocol MN Mobile Node HIT Host Identity Tag MSC Mobile service Switching Center HI Host Identifier HMIP Hierarchical Mobile IP HSPDA High Speed Downlink Packet Ac-ICMP Internet Control Message Protocol Set Size NAT Network Address Translation IFS Inter Frame Spacing IHI Internet Header Length IKE Internet Key Exchange Multiple Access IMSI International Mobile Subscriber OLSR Optimized Link- State Bouting OTP One-Time Password ISI InterSymbol Interference PCF Point Coordination Function

LF Low Frequency LTE Long Term Evolution MACA-BI MACA By Invitation MACA Multiple Access with Collision Avoidance MAC Message Authentication Code MAHO Mobile Assisted Handover MTSO Mobile Telecommunications Switching Office NAASS Normalized Average Anonymity NAV Net Allocation Vector
OFDMA Orthogonal Frequency-Division

PEAP Protected EAP

PEP Performances Enhancing Proxies

PIN Personal Identification Number

LFSR Linear Feedback Shift Register

RADIUS Remote Authentication Dial-In User Service RA Receiver Address RERR Route ERRor

RFID Radio Frequency Identification RREP Route REPly RREQ Route REQuests RSN Robust Security Network RTCP Real Time Control Protocol RTM Retransmission TimeOut RTP Real Time Protocol RTS Request To Send RVS Rendez-Vous Server RWND Receiver Window SACK Selective ACKnowledgment

SA Security Association SA Source Address SDMA Space Division Multiple Access SHF Super High Frequency SIFS Short Inter Frame Spacing SIM Subscriber Identity Module

SIP Session Initiation Protocol

PLCP Physical Layer Convergence Proto-

PSTN Public Switched Telephone Network

col
PMD Physical Medium Dependent

PMK Pairwise Master Key

PN Pseudo-random Noise

QoS Quality of Service

PTK Pairwise Transient Key

SPI Security Parameter Index SSTresh Slow Start Threshold STA STAtion

STA Station TA Transmitter Address TCP Transmission Control Protocol TDD Time Division Duplex TDMA Time Division Multiple Access

TIM Traffic Indication Map TKIP Temporal Key Integrity Protocol TLS Transport Layer Security

TMSI Temorary Mobile Subscriber Iden-

TOS Type Of Service TSF Timing Synchronisation Function

TTL Time To Live

UHF Ultra High Frequency

UMTS Universal Mobile Telecommunications System

UV Ultraviolet Light

VANET Vehicular Ad-hoc NETwork VHF Very High Frequency VLF Very Low Frequency

WAP Wireless Access Point WEP Wired Equivalent Privacy

WLAN Wireless Local Area Network WMN Wireless Mesh Network WPAN Wireless Personal Area Network

WPA WiFi Protected Access