ELEC 424/553 Mobile & Embedded Systems

Lecture 7 - System Calls



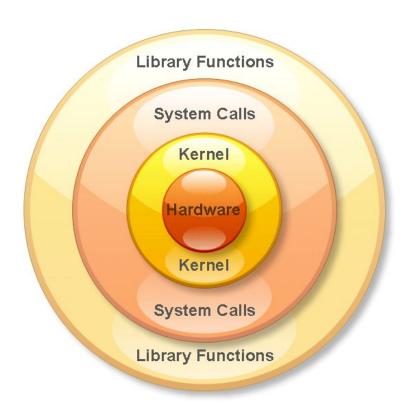
Housekeeping

ARM has hardware support for OS privilege levels similar to x86

Note from prior exercise on factorial code:

- In C, ++n is pre-increment and n++ is post-increment
- For the condition of a while loop, ++n will happen before the condition is checked and n++ will happen after the condition is checked

Soldering session, extra office hours



Balakumaran Kannan, "Anatomy of Linux system call in ARM64". URL: https://eastrivervillage.com/Anatomy-of-Linux-system-call-in-ARM64/

Questions In The Past

- Circular Queue vs. O(1) scheduler
 - Scheduling classes
- Scheduler overhead
 - O(log n)
 - Maybe a few µs for a context switch
 - Sources:
 - https://unix.stackexchange.com/questions/553120/linux-scheduling-overhead-in-orders-of-magnitude
 - https://eli.thegreenplace.net/2018/measuring-context-switching-and-memory-overheads-for-linux-threads/
 - https://events.static.linuxfound.org/sites/events/files/slides/lemoal-nvme-polling-vault-2017-final 0.pdf
- Processes going from to sleep to awake
 - One response indicates min_vruntime normalization does occur
 - https://stackoverflow.com/questions/11297285/how-can-the-linux-cfs-scheduler-prevent-a-task-with-a-very-small-vrunt ime-from-s

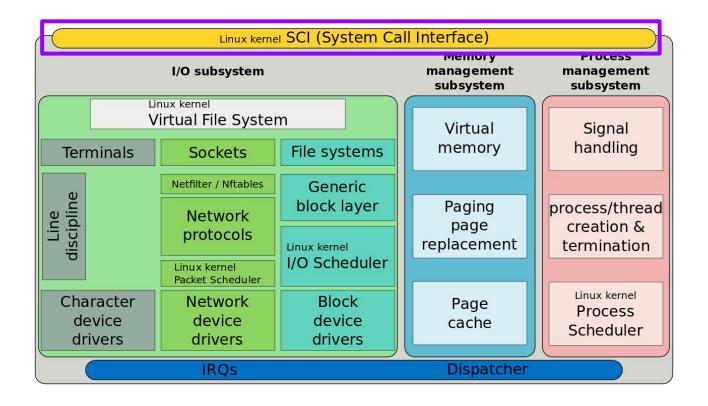
Goals of Today's Lecture

- Enable you to become a kernel hacker
- Learn the basics of system calls
 - How they're used
 - How they work
 - How to make your own



"Anonymous - Cool Hacker Guy". URL: https://i.kym-cdn.com/photos/images/original/001/109/423/cb4.jpg

Where Are We In The Linux Kernel?

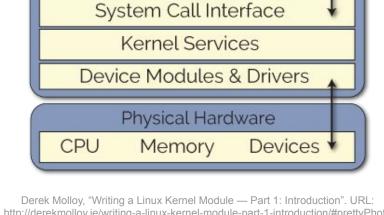


Kernel Space Must Be Protected At All Costs

- User applications have **no direct** access to kernel
- Indirect access via system calls
- System calls provide applications with access to kernel services

Kernel acts on behalf of application





User Space

User-Level Programs

/sbin/init User Code Linux Terminal

Kernel Space

GNU C Library (glibc)

Rachelle Hampton, "There's Nothing Wrong With Server-Speak, In Fact, It's Polite!". Image source: SeventyFour / Thinkstock https://slate.com/human-interest/2018/02/food-service-grammar-is-perfectly-polite.h

pchdtvr - A Case Study of System Calls In a Program

Home / Personal Media Networks

frequal.com

pchdtvr -- record digital television on your PC

pchdtvr is a great program for recording digital television on your PC. It lets you:

- 1. Set up schedules
- Record programs instantly
- 3. Check station signal strength
- 4. And much more!

http://frequal.com/pmn/pchdtvr.html

pchdtvr - A Case Study in Terrible Naming

pchdtvr - a good reason why techs should never be allowed to name their products :-)

- user10776 Dec 18 '16 at 8:04

"Can GPL licensed code be close sourced later by the author? [duplicate]". URL: https://softwareengineering.stackexchange.com/questions/98774/can-gpl-licensed-code-be-close-sourced-later-by-the-author/98778

pchdtvr - A Case Study of License Legal Issues Too

Home / Personal Media Networks

frequal.com

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pchdtvr is a great program for recording digital television on your PC. It lets you:

- 1. Set up schedules
- 2. Record programs instantly
- 3. Check station signal strength
- 4. And much more!

Best of all, since pchdtvr is licensed under the GNU GPL, nobody can ever take away your right to use it, change it, or give it away!

Download

There are several versions available:

- 1. pchdtvr-1.0-frequal, the latest and greatest version available.
 - 1. New % file suffix to attach date to recordings
 - 2. Append instead of rename for restarted recording
- 2. pchdtvr-1.0
- 3. pchdtvr-0.9s

Last modified on 26 Jan 2008 by AO

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I have revoked the licensing under the GNU General Public License (herein after referred to as "the GPL") for the atscap version 1.1 codebase, all prior versions of the atscap codebase and all the various release candidates of the atscap codebase.

[...]

As sole author of both the atscap and the pchdtvr codebases, the licenses under the GPL were granted at my sole discretion and the licenses under the GPL are now hereby revoked at my sole discretion.

[...]

If you are currently using the atscap or pchdtvr packages, or any part thereof, it is in your best interest to remove the software from your system(s) and destroy all copies in your possession.

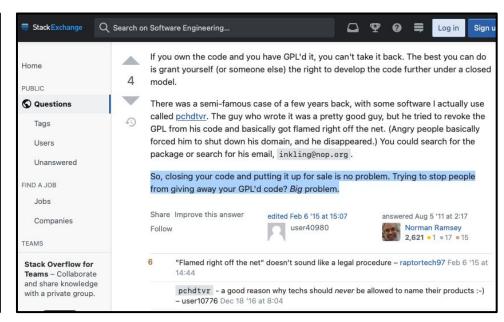
[...]

If you have incorporated the atscap or pchdtvr codebase, or any part thereof, into any of your projects, it is in your best interest to remove any and all of my code from your project(s).

pchdtvr - Not a Totally Satisfying Conclusion in Regards to Legal Issues

I've never used this package (apparently its used for HDTV scheduling/recording on Linux), but this link on Slashdot caught my eye: http://sourceforge.net/developer/diary.php? diary_id=26407&diary_user=147583. Apparently, the developer of this software package is seeking to revoke the GPL license not just for his current code, but his past code/package as well. I have a difficult time believing that this is possible, but I'm sure we will soon find out. My guess is that this guy is productizing his software and has a good idea who is currently using, selling and distributing his source so there will likely be some kind of legal challenge to the GPL as well. It's always interesting to see how these kind of things play out in the U.S. courts which can sometimes be a little schizophrenic, though I'd have a difficult time believing that this type of retroactive license change is actually possible.

User: reeset, "atscap and pchdtvr GPL revoked or can it be". URL: https://blog.reeset.net/archives/493



"Can GPL licensed code be close sourced later by the author? [duplicate]". URL: https://softwareengineering.stackexchange.com/questions/98774/can-gpl-licensed-code-be-close-sourced-later-by-the-author/98778

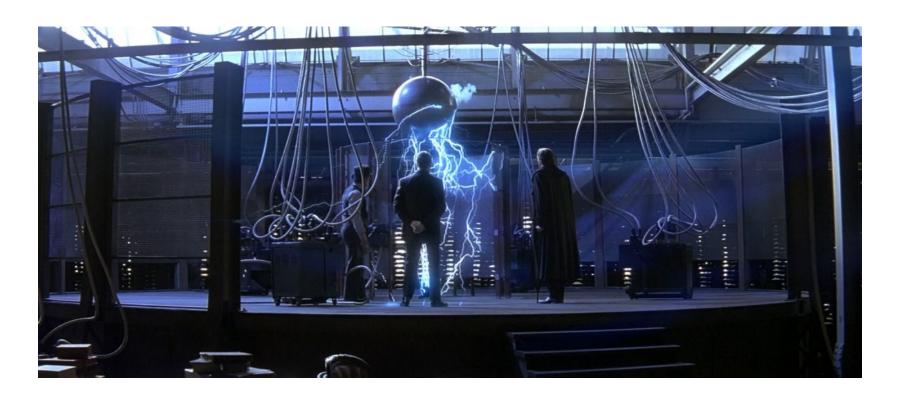
pchdtvr - GPL Means We Still Have the Source

```
pchdtvr.c
```

pchdtvr - We Can Find Our Old Friend fork ()

```
pchdtvr.c
                                                                Remember: pid t is an
                                                                opaque type (not known by
  if (0 != arg_detach)
                                                                application)
pid_t fp;
fp = fork();
if (0 == fp)
                                                                fork() called
    pid_m = getpid();
    asyslog( LOG_INFO,
      "daemon forked, main() is now pid %d", pid_m );
} else {
    /* old forker must die right here */
    console_exit(0);
```

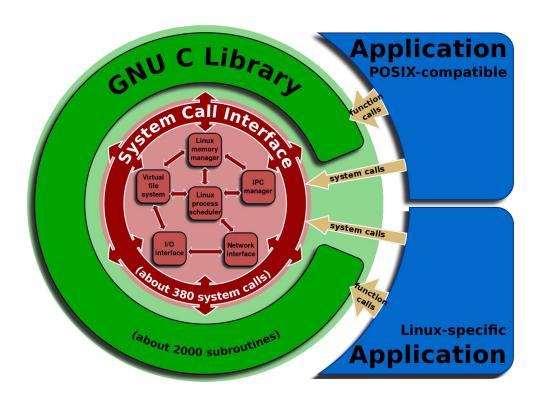
Christopher Nolan's Explanation of fork() Return Value



pchdtvr - We Can Find Our Old Friend fork ()

```
pchdtvr.c
                                                                         Remember: pid t is an
                                                                         opaque type
           if (0 != arg_detach)
                                                                         fork() called
         pid_t fp;
         fp = fork();
         if (0 == fp)
                                                                         Am I the child process?
                                                                         What is my PID?
             pid_m = getpid();
             asyslog( LOG_INFO,
               "daemon forked, main() is now pid %d", pid_m );
                                                                         "say my name"
         } else {
22262
             console_exit(0);
                                                                         I'm the parent process
                                                                         Death imminent
```

Using System Calls Directly & Indirectly



Indirectly Using System Call Via glibc

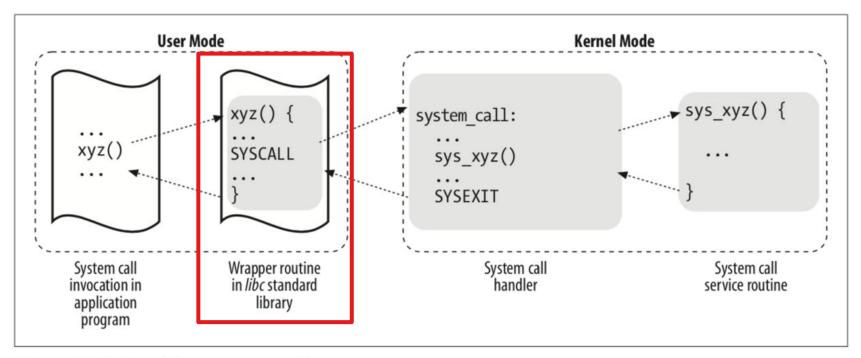
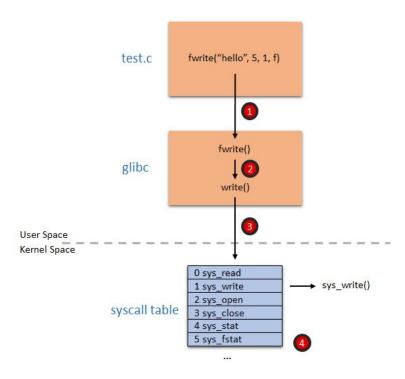


Figure 10-1. Invoking a system call

Example of write () System Call Being Used



fork() Does Not Call fork()



C library/kernel differences

Since version 2.3.3, rather than invoking the kernel's **fork**() system call. the glibc **fork**() wrapper that is provided as part of the NPTL threading implementation invokes clone(2) with flags that provide the same effect as the traditional system call. (A call to **fork**() is equivalent to a call to clone(2) specifying flags as just **SIGCHLD**.) The glibc wrapper invokes any fork handlers that have been established using pthread_atfork(3).

Remember when?...

- 1983: UNIX System V released
 - Richard Stallman also started GNU Project
 - Started developing free programs functionally equivalent to UNIX programs
- The UNIX wars (late 80s to mid 90s)
 - What is UNIX? Many versions were floating around
- Efforts for peace were made
 - 1985: AT&T introduced the System V Interface Definition (SVID)
 - 1988: IEEE POSIX (Portable Operating System Interface)
 - AT&T worked with others to merge other UNIX versions into System V
 - The Open Software Foundation countered with OSF/1
 - All the while, more commercial distributions of UNIX appeared that expanded beyond AT&T's standard



Richard Stallman

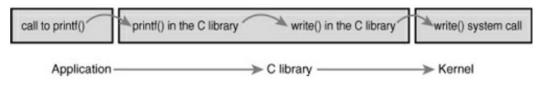
From the cover of: Free as in Freedom, Sam Williams. https://www.amazon.com/Free-Freedom-Richard-Stallmans-Software-dp-059600287 74/dp/0596002874/ref=mt other? encoding=UTF8&me=&qid=1630343742

POSIX - Portable Operating System Interface

- POSIX encourages compatibility among OS's for applications via set of standards
- UNIX Wars Guarantee application worked across UNIX flavors
- Standardizes:
 - Shell
 - API (application programming interface) to kernel
- Many things are not standardized
 - GUIs
- This does not talk about binary compatibility

POSIX - API Specifics

- Common API enables applications to access kernel services across UNIX and UNIX-like systems
- POSIX talks about System Interfaces:
 - "the functions, macros, and external variables to support applications portability at the C-language source level" [IEEE & The Open Group: <u>link</u>]
- Implementation unspecified can include:
 - System calls via the C library
 - Staying entirely within the C library (user space)
- glibc The GNU C Library
 - C standard library realization by GNU Project

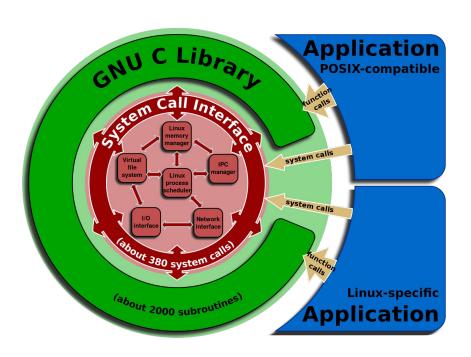


POSIX - System Interface Examples

- List of POSIX system interfaces available on The Open Group's site <u>here</u>
 - Again, not the same, but we have a list of Linux system calls <u>here</u>
- System interface: <u>fork()</u>
 - System call: <u>fork()</u>
 - But we know glibc will not use the fork() system call, it will use the clone() system call
 - Note: no clone () system interface
- System interface: <u>fopen()</u>
 - System call: open ()
- System interface: <u>sqrt()</u>
 - System call: Assumably none

Using System Call Directly or Indirectly

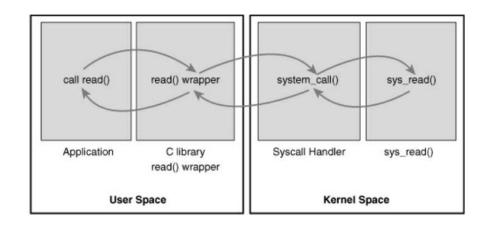
```
#include <syscall.h>
#include <unistd.h>
#include <stdio.h>
#include <sys/types.h>
int main(void) {
     long ID1, ID2;
     /* direct system call */
     /* SYS getpid (func no. is 20) */
     ID1 = syscall(SYS getpid);
     printf ("syscall(SYS getpid)=%ld\n", ID1);
     /* "libc" wrapped system call */
     /* SYS getpid (Func No. is 20) */
     ID2 = qetpid();
     printf ("getpid()=%ld\n", ID2);
     return(0);
```



ScotXW. CC Attribution-Share Alike 3.0 Unported license (<u>link</u>). URL: https://commons.wikimedia.org/wiki/File:Linux_kernel_System_Call_Interface_and_glibc.svg

How Does This Really Work? (x86 Edition)

- Cause a software interrupt: exception/trap/fault
- Kernel mode activated
- Exception handler (system call handler) runs
- Which system call?
 - Use eax register to give system call # to kernel
 - eax read by handler



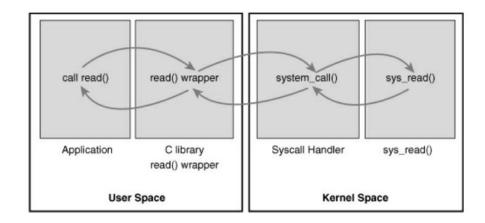


System Call Numbers (x86-64)

```
/ arch / x86 / entry / syscalls / syscall 64.tbl
                                                            Search Identifier
                                                  All sy V
     #
     # 64-bit system call numbers and entry vectors
     #
       The format is:
 5
       <number> <abi> <name> <entry point>
 6
       The x64 sys *() stubs are created on-the-fly for sys *() system calls
 8
     #
 9
       The abi is "common", "64" or "x32" for this file.
10
     #
11
                      read
                                               sys read
              common
12
              common write
                                               sys write
13
                      open
                                               sys open
              common
14
              common close
                                               sys close
15
              common stat
                                               sys newstat
16
              common fstat
                                               sys newfstat
17
              common 1stat
                                               sys newlstat
18
                      poll
                                               sys poll
              common
```

Parameter Passing & Return Value

- Use mechanism similar to syscall # transmission
- Place parameters in registers
- Shown registers are for 5 arguments
- If more needed, put user-space pointer in 1 register
 - No other registers used
- How do we get the return value?
 - o eax





4 Simple Steps to Making Your Very Own System Call

- 1. What is the point of its existence?
- Write the code & add to file
- 3. Add call to the system call table
- 4. Compile

1. What Is The Point of The System Call?

- Linux philosophy: syscall does one thing
- Determine:
 - Input arguments
 - Return value
 - Error codes
- Unix motto: "Provide mechanism, not policy"
- Portability, robustness
 - Fundamental syscalls have been helpful for 30 years!

2. Write The System Call Code: Parameter Verification

- We're in kernel space, tread carefully
- File I/O: Is file descriptor legitimate?
- Processes: Is PID legitimate?
- Permissions: Is the process allowed to see this?
- Pointers ••
 - o Is this really pointing to user space?
 - o Is this really pointing to process's space?
 - Is the memory readable/writable/executable?

2. Example of System Call Code

```
SYSCALL_DEFINE1(stephen, char *, msg)
{
  char buf[256];
  long copied = strncpy_from_user(buf, msg, sizeof(buf));
  if (copied < 0 || copied == sizeof(buf))
    return -EFAULT;
  printk(KERN_INFO "stephen syscall called with \"%s\"\n", buf);
  return 0;
}</pre>
```

Stephen Brennan, "Tutorial - Write a System Call". URL: https://brennan.io/2016/11/14/kernel-dev-ep3/

Include syscall code in kernel file

• E.g. sys.c

3. Let's Make It Official: Add to Syscall Table

- Append call to end of syscall table
 - Number defined by ranking in table
 - Architecture-specific

Add to syscall_64.tbl

| / ard | ch / x86 / e | entry / syscalls | s / syscall_64.tbl | All syml∨ Search |
|-------|--------------|------------------|---|---|
| 361 | 437 | common | openat2 | sys openat2 |
| 362 | 438 | common | pidfd getfd | sys pidfd getfd |
| 363 | 439 | common | faccessat2 | sys faccessat2 |
| 364 | 440 | common | process madvise | sys process madvise |
| 365 | 441 | common | epoll pwait2 | sys epoll pwait2 |
| 366 | 442 | common | mount setattr | sys mount setattr |
| 367 | 443 | common | quotactl fd | sys quotactl fd |
| 368 | 444 | common | landlock create ruleset | sys landlock create ruleset |
| 369 | 445 | common | landlock add rule | sys landlock add rule |
| 370 | 446 | common | landlock restrict self | sys landlock restrict self |
| 371 | 447 | common | memfd secret | sys memfd secret |
| 372 | | | Countries, destroyed species —— and species, consistent products of 1900. | TO A STATE OF THE |

4. Final Step

Compile and cross your fingers

Test Ride Your Beautiful New Syscall

- glibc contains the accepted system calls
- Yours is brand new!
- Use syscall() function

```
* Test the stephen syscall (#329)
#define GNU SOURCE
#include <unistd.h>
#include <sys/syscall.h>
#include <stdio.h>
* Put your syscall number here.
#define SYS stephen 329
int main(int argc, char **argv)
 if (argc <= 1) {
   printf("Must provide a string to give to system call.\n");
   return -1:
 char *arg = argv[1];
 printf("Making system call with \"%s\".\n", arg);
 long res = syscall(SYS stephen, arg);
  printf("System call returned %ld.\n", res);
  return res;
```