Essentials of Computer Architecture

Susan B. Davidson and Zachary G. Ives

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CIS 545 – Big Data Analytics





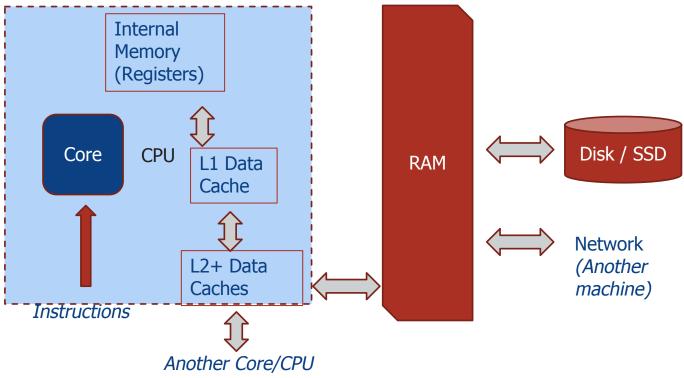
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Understanding a Bit of Computer Architecture Is Key

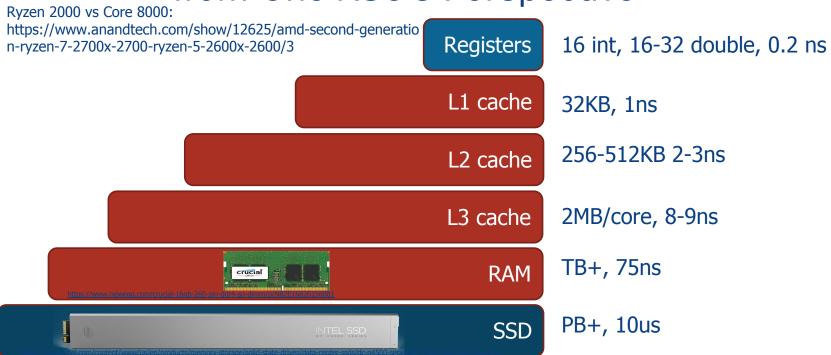
A few ideas we'll need from Computer Architecture:

- 1) Different program instructions take different amounts of work and time
 - Lines of code (in Python / C / etc) turn into multiple machine instructions
 - Machine instructions don't all have equal performance!
- 2) The computer's *memory hierarchy* means accessing data is also not uniform!
- 3) Modern hardware includes special optimizations to:
 - Do the same operation on multiple data items simultaneously
 - https://tinyun.multiple.independentopieces of code at the same time

Key Aspects of a Computer: Simplified View of a Microprocessor



The Memory Hierarchy from One X86's Perspective



How Are Caches Used? (We Are Simplifying to 1 Level)

```
# Create an array 1,000,000 by 1
my_list = np.empty(1000000)

for i in range(0, len(my_list)):
    x = x + my_list[i] + \
    random.random()
```

```
Core
1ns 1ns 1ns
                 ~4KB
my list[0:999]
    75ns
    my_list[0:1000000]
           ~4MB
```

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```

```
Core
1ns Ins 1ns 1ns Ins
                 ~4KB
           75ns
   75ns
   my_list[0:1000000]
           ~4MB
```

Interactive Visualizer (Colin Scott):

https://people.eecs.berkeley.edu/~rcs/research/interactive_latency.html

Original list from Jeff Dean and Peter Norvig (Google)

Latency Numbers Every Programmer Should Know





2019 Main memory reference: Send 2,000 bytes over Read 1.000.000 bytes 1ns commodity network: 62ns sequentially from SSD: 100ns 62,000ns ≈ 62µs L1 cache reference: 1ns SSD random read: 16,000ns 1,000ns ≈ 1µs ≈ 16us Disk seek: 3,000,000ns ≈ Branch mispredict: 3ns Compress 1KB wth Zippy: Read 1,000,000 bytes 2,000ns ≈ 2µs sequentially from memory: Read 1,000,000 bytes 12 cache reference: 4ns 4,000ns ≈ 4µs sequentially from disk: 10,000ns ≈ 10us = ■ 947,000ns ≈ 947us Round trip in same Mutex lock/unlock: 17ns datacenter: 500,000ns ≈ Packet roundtrip CA to 500µs Netherlands: 150,000,000ns ≈ 150ms 100ns = | 1,000,000ns = 1ms = =

How Do We Take Advantage of the Memory Hierarchy?

- Focus on data reduction filter rows and columns
- Put the most frequently used items together
- •Find ways to pass over the data once, instead of many times

•Some of this is done by our big data engines – but there

htare to the factors in optimization we'll see shortly

Recap: Consider the Memory Hierarchy When Thinking about Scale

- Accessing data in predictable ways is better
 - CPU predicts and "pre-fetches" the data into caches
 - Repeated requests to related data keep memory in the caches
- Smaller "memory footprints" are better (at tension with the previous)

Packing multiple data values into the same memory, and http://lying.a/repeated.computation, is better

Brief Review

https://canvas.upenn.edu/courses/1636888/quizzes/2771492/ (07B)

How many instructions can a typical machine run in one second, at peak?

- a. billions (10⁹)
- b. trillions (10^{12})
- c. thousands (10^3)
- d. millions (10⁶)

How does a cache help with performance?

- e. it stores all of system memory in a faster place
- f. it stores the contents of the registers
- g. it allows us to amortize expensive memory fetches across multiple requests
- h. it stores all of the disk state

Optimizing Relational Expressions to Improve Performance

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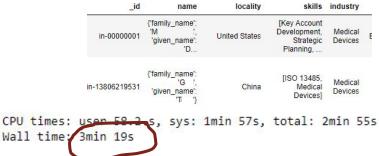
Roadmap for Data Processing

https://tinyurl.com/cis545-notebook-03

- How computer architecture and memory affect performance
- Optimizing relational algebra to minimize memory costs
- Optimizing join orders
- Algorithmic techniques

Big Data Takes A Long Time to Process

Even simple operations are expensive at scale...



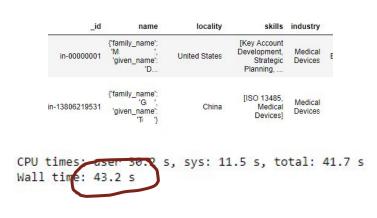
What makes this so slow?

- Scanning through memory
- Parsing strings
- Updating the list data structure
- Converting to Pandas DF

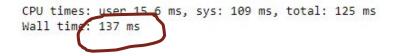
Optimizing for Costs

- Reduce data as soon as possible... "Push down" operations!
 Select (filter) at the earliest point possible
 - Less fetching from disk (if relevant), updating of data structures, processing of rows!

Project once columns aren't needed: Data better fits into the CPU cache



SQL Does this Automatically (and Doesn't Parse Every Time)



%%time	
<pre>pd.read_sql_query('select * from people where industry="Medical Devices"', conn)</pre>	

locality	industry	summary
United States	Medical Devices	SALES MANAGEMENT / BUSINESS DEVELOPMENT / PROJ
China	Medical Devices	
	United States	United Medical States Devices

Even Better:

Indexing Speeds Selection

Can we speed up fetches for specific data, e.g., people by industry?

An **index** is a map from an **index key** to a set of values

- It allows us to directly find matches to the key without scanning the data
- Can be in-memory or on disk

Two types of indices:

- Tree indices (B+ Trees) allow us to find all values <= key, >= key, = key
- Hash indices allow us to look up all values equal to the key

Q: Which is a **dict** in Python? https://tinyurl.com/cis545-lecture-02-09-22

```
conn = sqlite3.connect('linkedin.db')
conn.execute('begin transaction')
conn.execute("create index
people industry on people(industry)")
conn.execute('commit')
%%time
pd.read sql query('select * from
people where industry="Medical
Devices"', conn)
 CPU times: user 0 ns, sys: 0 ns, total: 0 ns
```

```
Wall time, 26.9 ms
```

To Consider

Can indexing also speed projection?

The Summary So Far: Selections and Projections

- A good rule of thumb is to "push down" selection and projection operations
- SQL DBMSs do this automatically

- Index data structures allow us to evaluate predicates on a key and directly return the matches
 - Hash indices help with exact-matches (we'll revisit this soon)
 - Tree indices (B+ Trees in relational DBMSs) help with equality and inequality

Brief Review

https://canvas.upenn.edu/courses/1636888/quizzes/2771528 (07C)

How does "pushing down" filter conditions typically help performance? Choose the best answer.

- a. Reduces memory "footprint" and number of operations
- b. Requires an index
- Reduces number of operations only
- d. Reduces memory "footprint" only

A typical index:

- e. requires time to access data that's linear in the number of rows in the table
- f. requires roughly constant time for every access
- g. requires time to access data that's linear in the number of columns in the table
- h. doesn't work unless all data fits in system memory

Join Ordering and Optimization

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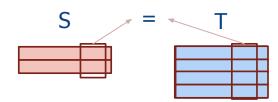




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Recall Join Compares all Pairs of Tuples for a Match

S join T on s_on = t_on:



How we implement it:

```
for every tuple s in S
  for every tuple t in T
   if s[s_on] == t[t_on] then combine & return
```

Joins: Expensive Operations

How many operations does this require? How big is the output?

Using Our Intuitions in Order of Evaluation (Abstract Example)

Suppose we have three simple Dataframes describing people at Penn (roughly 10.1K undergrad, 10K grad/professional, 0.1K both, 10K faculty/staff):

```
people(id, name) category(id, grad_undergrad_facstaff) grade(id, sem, score)

30K people 30.1K entries (G, U, F/S) 450K entries

1. Roughly how much work is it to do: ~10K 300M comp, ~10K res

people.merge(category[category['id']=='G']):merge(grade)

2. Roughly how much work is it to do: ~10K

people.merge(grade).merge(category[category['id']=='G']))

13.5B comp, ~450K res

4.5B comp, ~150K res
```

Looking at Order of Evaluation on Our Real, LinkedIn Dataset

```
people_df = pd.read_sql_query('select * from people limit 500', conn)
experience_df = pd.read_sql_query('select * from experience limit 5000',
conn)
skills_df = pd.read_sql_query('select * from skills limit 8000', conn)

temp = merge(people_df, experience_df, '_id', 'person') # 2228 rows
mktg_df = skills_df[skills_df['value'] == 'Marketing'][['person']] # 23 rows
```

```
%%time
merge(merge(people_df, exp

Merge compared 2500000 tup
Merge compared 51244 tuple
CPU times: user 32.4 s, sy
Wall time: 32.4 s
Wall time: 1.23 s

%%time
merge(merge(people_df, mktg_df, '_id', 'person'), experience_df, '_id', 'person')

Merge compared 11500 tuples
Merge compared 85000 tuples
CPU times: user 1.23 s, sys: 31.2 ms, total: 1.27 s
Wall time: 1.23 s)
```

Do We Have to Choose Orders Manually?

Order of evaluation matters because intermediate result sizes matter SQL (and query optimization) decide for us!

Do We Have to Do this Manually?

Order of evaluation matters because intermediate result sizes matter

SQL (and query optimization) decide for us!

Recap

 Joins are expensive, requiring a quadratic (in the number of inputs) number of comparisons

 Cleverly ordering our joins to reduce intermediate result sizes makes a huge performance difference

•SQL databases can choose this automatically for us https://tinyurl.com/cis545-lecture-02-09-22

Brief Review

https://canvas.upenn.edu/courses/1636888/quizzes/2771532 (07D)

If we join R (1,000 tuples) and S (1,000) tuples, our join result will be:

- a. anywhere from 0 to 1 million tuples
- b. 1,000 tuples
- c. 10,000 tuples
- d. 25,000 tuples

If we are joining three tables, our best strategy is to

- e. choose a join order that minimizes the intermediate result
- f. choose a join order by random sampling
- g. join the first table with the second, then the third
- h. choose a join order that maximizes the intermediate result

Improving Algorithmic Efficiency

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Two Key Ideas

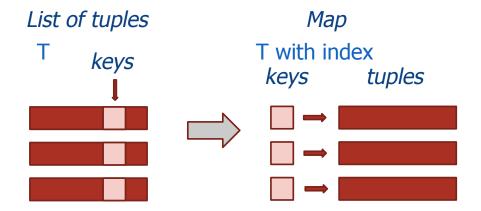
To this point, we've looked at reducing data

- Now: two ideas for improving performance by changing how we access data
 - Dictionaries / in-memory indices / maps
 - Buffering / blocks

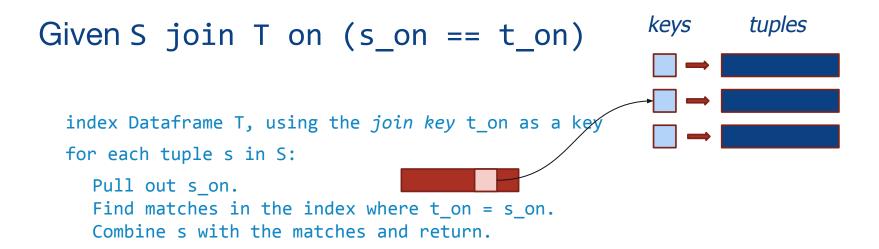
Pandas Merge Is Based on Exact-Matches

Can we find exact-matches between the values in one tuple (from some table S) with another tuple (from some table T)?

We had a way of doing fast lookups: maps from keys to values (i.e., dicts, indices)



Rethinking Join (For Equality)



A More Efficient Equality Join

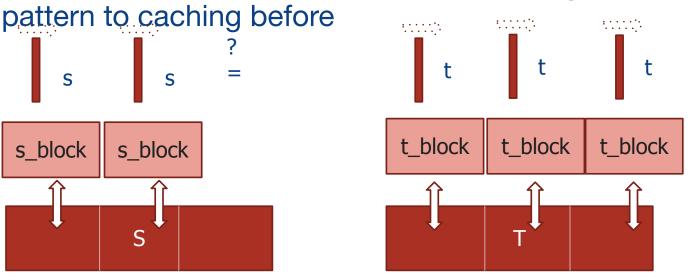
```
def merge_map(S,T,s_on,t_on):
    ret = pd.DataFrame()
    T map =
             %%time
    # Take
    # make
             # Here's a test join, with people and their experiences. We can see how many # comparisons are made
     T_map[
             merge man(experience df, people df, 'person', 'id')
    # Now f Merge compared 5500 tuples
    for s i CPU times: user 7.12 s, sys: 0 ns, total: (7.12 s
        cou Wall time: 7.14 s
         if S.loc[s index, s on] in T map:
                 ret = ret.append(S.loc[s index].\
              merge(people df, experience df, 'id', 'person')
        igno
              Merge compared 2500000 tuples
              CPU times: user 30.8 s, sys: 0 ns, total 30.8 s
return ret
              Wall time: 30.9 s
```

For Further Thought

How does this work if you are joining on more than one field? (first + last name)?

Suppose We Have Big Data, so S and T Don't Fit in Memory!

We need to read a *block* at a time, following a similar



Joins with Big Data Need to Read a Table in Blocks

Now we read blocks(S) * blocks(T) pages, and compare each s and t from these 100us per block read, 75ns per comparison!!!

Beyond the Relational Algebra

- For big data, sometimes we'll need to supply our own operations
 - Functions to be called via apply (or applymap)
 - Functions that take collections of data

•Again, we'll want to rely on using maps and indices to reduce data usage, and intelligent buffering

Brief Review

https://canvas.upenn.edu/courses/1636888/quizzes/2771493 (07E)

If we use a dictionary / map / hash table to join two tables ${\tt S}$ and ${\tt T}$, our cost is proportional to

- a. the product of the cardinalities (numbers of rows) of S and T
- b. the cardinality (number of rows) of the bigger of S and T
- c. the sum of the cardinalities (numbers of rows) of S and T
- d. the square root of the sums of the sizes of the tables, squared

How does buffering or blocked access help speed performance when tables are bigger than memory?

- e. we reduce the size of the index
- f. we reduce the overall number of disk fetches
- g. we reduce the size of the join output
- h. we reduce the overall number of join comparisons

Summary:

Making Dataframes / Queries Efficient

General rule of thumb for efficiency: consider order of evaluation

- Minimize intermediate results
- •For Pandas, "ballpark estimate" which order is better
- •SQL database optimizes using statistical information it collects on tables!

For some joins, can use faster algorithm:

- Joins, by default, iterate over both tables
- •An *index* makes the lookups more efficient IF the join is on the *index* https://tinyurl.com/cis545-lecture-02-09-22

Processing Data at Scale

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How Do We Handle Large Volumes of Data?

1) Reduce the amount of data

2) Use more efficient algorithms or data structures

3) Split work among multiple processors / workers!

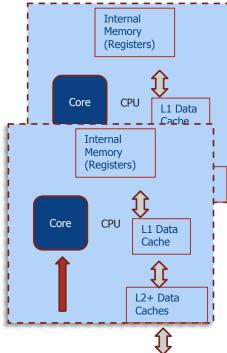
Multiple Processors Are Important!

For decades, CPU designers made cores dramatically faster each generation:

Roughly 2x performance every 18 months – "Moore's Law" (and "Dennard Scaling")

No longer true – but now we can put often more *cores* into each CPU generation

But: how to coordinate among the cores?



Bulk Operations and Scale

Relational and other **bulk** operations are key:

•They express how to do the same thing over many values

•Selections, projects, joins, apply, etc. can often be done on different rows in parallel!

•(Also true for certain other operations, as we'll see later)
https://tinyurl.com/cis545-lecture-02-09-22

Road Map: Variations on "Data Parallelism" "Data parallelism" = same instructions, multiple data items

- Dask: multicore processing of relational algebra on a single machine
- Cluster computing and the cloud
- Data processing in a cluster
- Apache Spark and cluster-based processing of dataframes

Multicore Processing of Dataframes







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Multicore and Pandas

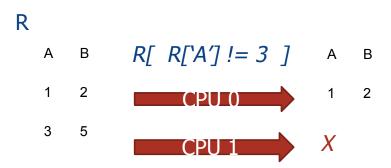
 Relational operators inherently apply the same value-based computation to each row

 There's a natural way of assigning different data to different CPU cores, and processing them in parallel

Not done by Pandas, but there are alternatives

Parallel Processing of Relational Algebra

So far: we did our work one step at a time!
Relational operators can be done *in parallel* across records



Dask: Parallel Execution



Dask provides its own replacement dataframes for Pandas (and other libraries) that automatically supports parallelism on a single machine

Almost identical code to Pandas and Apache Spark,

which we'll talk about instead (and next)...

Let's Revisit our LinkedIn Data Example

https://tinyurl.com/cis545-005

```
# 100K records from linkedin
linked in = open('/content/linkedin small.json.txt')
people = []
for line in linked in:
    person = json.loads(line)
    people.append(person)
people df = pd.DataFrame(people)
print ("%d records"%len(people df))
people df
```

We should split out the nested lists for skills, education, experience, and honors:

```
skills_df = people_df[['_id','skills']].\
   explode('skills')
education_df = people_df[['_id','education']].\
   explode('education')
experience_df = people_df[['_id','experience']]
honors_df = people_df[['_id', 'honors']]

linkedin_df = people_df.copy().drop(\
   columns=['skills','education','experience',\
   'honors'])
```

Rejoining the Tables Could Be Expensive!

%%time

```
linkedin_df.merge(experience_df, on='_id').\
  merge(skills_df, on='_id').\
  merge(honors_df, on='_id').\
  merge(education_df, on='_id')
Wall time: 3.71 s
```

Full LinkedIn data is 90x bigger!

Dask Has Its Own Dataframes

```
import dask
import dask.dataframe as dd
linkedin ddf = dd.from pandas(linkedin df, npartitions=100)
skills ddf = dd.from pandas(skills df,npartitions=100)
experience ddf = dd.from pandas(experience df, npartitions=100)
education ddf = dd.from pandas(education df,npartitions=100)
honors ddf = dd.from pandas (honors df, npartitions=10)
88time
linkedin ddf.merge(experience ddf, on=' id').
 merge(skills ddf, on=' id').merge(honors ddf, on=' id').\
  merge(education ddf, on=' id')
Wall time: 160 ms
```

We Saw Dask Does Parallel Dataframes – What about Parallel SQL?

 Many DBMSs like PostgreSQL can support parallel processing, but on a single machine disk I/O is typically the bottleneck

•So they largely benefit with clusters as we'll see soon!

Summary of Multicore Processing

Through parallel libraries like Dask, we can process data items independently on different cores

Not true for every operation, but for some relational operations we can conceptually get speedups up to min(# cores, # rows)

But for big data, we often have too much data to handle on a single (even multi-core) computer!

Brief Review

https://canvas.upenn.edu/courses/1636888/quizzes/2771523 (08B)

How does Dask speed up common dataframe operations?

- a. it partitions dataframe by row so they can be processed by different CPU cores
- b. it replaces Pandas code with native Python
- c. it partitions dataframes by column so they can be more efficiently executed
- d. it optimizes the cache

Why don't DBMSs generally show big benefits from multicore processors?

- e. typically most of their processing is limited by I/O
- f. they are just fundamentally slower
- g. SQL can't be executed in parallel
- h. they need to incorporate Apache Spark

Beyond a Single Machine: Clusters and Data Centers

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The First Approach: Bigger Machines and Dask / Swifter



2019 laptop: 16GB RAM 6 4GHz cores 2TB SSD



2019 Server: 1TB RAM GPU 28 4GHz cores 1PB SSD



Cluster: 30 servers, fast network among them

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Non-uniform Costs in Clusters and Parallel Systems

Recall we needed to pay a lot of attention to algorithms when we had a disk

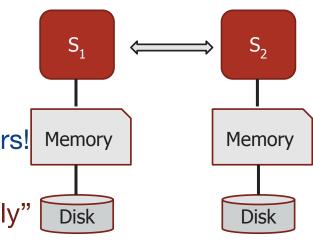
Internal CPU state (registers, cache)

Memory (latency 1000s of times slower)

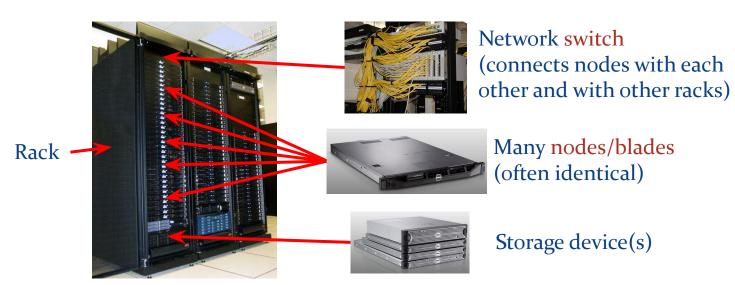
SSD / disk (latency 1000+ times slower)

•Same is true if we have multiple computers!

Each processor should work "independently" and "locally" as much as possible https://tinyurl.com/cis545-lecture-02-09-22



Building Clusters of Computers



- Many similar machines, close interconnection (same room?)
- Often special, standardized hardware (racks, blades)
- Usually owned and used by a single organization

https://tinyurl.com/cis545-lecture-02-09-22

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Scaling Out



What if your cluster is too big (hot, power hungry) for an office building?

- Build a separate building for the cluster a data center!
- Building can have lots of cooling and power

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"The Cloud"

Commercial cloud: data centers run by commercial providers

Amazon Web Services, Microsoft Azure, Google Cloud, IBM Cloud, Rackspace, Linode, ...

- Machines you can access "elastically", maintained by the cloud provider's staff
- MANY services, ranging from "configurable Linux machines" to "Storage" or "Jupyter"
- Often less expensive than maintaining local resources due to economies of scale

Hybrid cloud: data centers run by commercial providers, combined with some local clusters administered the same way

Google data centers



https://tinyurl.com/cis545-lecture-02-09-22

Brief Review

https://canvas.upenn.edu/courses/1636888/quizzes/2771570 (08C)

On average, how much system memory (RAM) is typical in a laptop today?

- a. single-digit to 3-digit GB
- b. 1 PB
- c. 1TB 4TB
- d. 64 128 TB

When might we need a data center?

- e. we need a server with more memory
- f. we want to set up a compute cluster
- q. we don't have enough power and cooling for a machine room
- h. we don't have enough Internet bandwidth and would need to upgrade to fiber

Summary: Clusters and Data Centers Provide the Hardware for Big Data

- Today we build parallel computation by linking together many compute nodes
- Each has its own local CPU cache, memory, and disk
- Communication costs vary dramatically at different levels

- •The challenge: we can't program these just using Python and Pandas!
 - •How do we get 30 multicore servers to work together?
- •At scale, machines crash all the time how do we handle failures? https://tinyurl.com/cis545-lecture-02-09-22

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