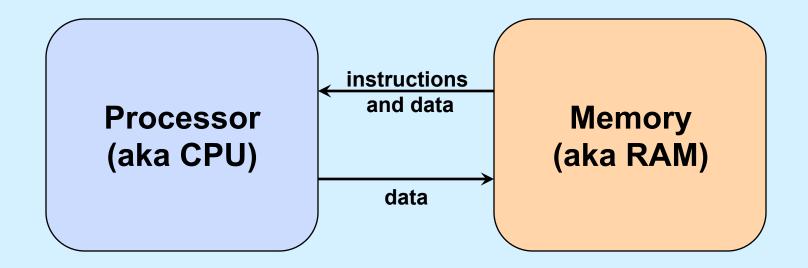
CS 33

Intro to Machine Programming

Machine Model



Memory

Instructions Instructions or are Data **Data**

Processor: Some Details

Execution engine

Instruction pointer

Condition codes

Processor: Basic Operation

while (forever) {
 fetch instruction IP points at
 decode instruction
 fetch operands
 execute
 store results
 update IP and condition code
}

Instructions ...

Op code Operand1 Operand2 ...

Operands

- Form
 - immediate vs. reference
 - » value vs. address
- How many?
 - **3**
- » add a,b,c
 - \cdot c = a + b
- **2**
- » add a,b
 - b += a

Operands (continued)

- Accumulator
 - special memory in the processor
 - » known as a register
 - » fast access
 - allows single-operand instructions
 - » add a
 - acc += a
 - » add b
 - acc += b

From C to Assembler ...

Condition Codes

- Set of flags giving status of most recent operation:
 - zero flag
 - » result was or was not zero
 - sign flag
 - » for signed arithmetic interpretation: sign bit is or is not set
 - overflow flag
 - » for signed arithmetic interpretation
 - carry flag (generated by carry or borrow out of mostsignificant bit)
 - » for unsigned arithmetic interpretation
- Set implicitly by arithmetic instructions
- Set explicitly by compare instruction
 - cmp a,b
 - » sets flags based on result of b-a

Quiz 1

- Set of flags giving status of most recent operation:
 - zero flag
 - » result was or was not zero
 - sign flag
 - » for signed arithmetic interpretation: sign bit is or is not set
 - overflow flag
 - » for signed arithmetic interpretation
 - carry flag (generated by carry or borrow out of most-significant bit)
 - » for unsigned arithmetic interpretation
- Set explicitly by compare instruction
 - cmp a,b
 - » sets flags based on result of b-a

Which flags are set by "cmp 2,1"?

- a) overflow flag only
- b) carry flag only
- c) sign and carry flags only
- d) sign and overflow flags only
- e) sign, overflow, and carry flags

Jump Instructions

- Unconditional jump
 - just do it
- Conditional jump
 - to jump or not to jump determined by conditioncode flags
 - field in the op code indicates how this is computed
 - in assembler language, simply say
 - » je
 - jump on equal
 - » jne
 - jump on not equal
 - » jgt
 - jump on greater than
 - » etc.

Addresses

```
int a, b, c, d;
int main() {
   a = (b + c) * d;
   ...
}
```

mov b,%acc add c,%acc mul d,%acc mov %acc,a

 mov
 1004,%acc

 add
 1008,%acc

 mul
 1012,%acc

 mov
 %acc,1000

1012: d 1008: c 1004: b global 1000: a variables

Memory

Addresses

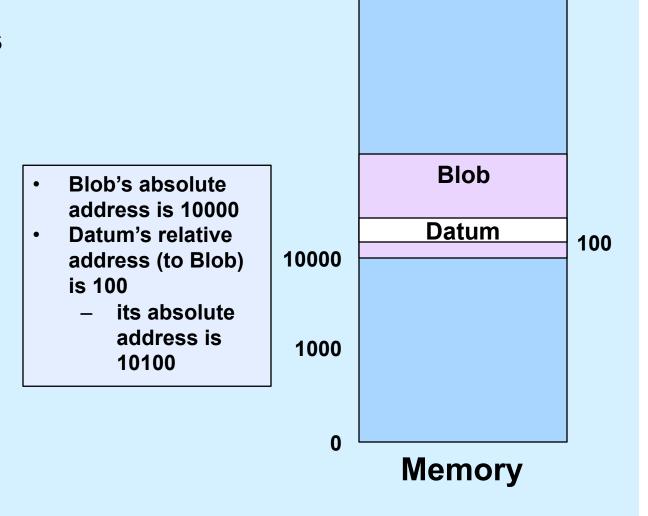
```
int b;
int func(int c, int d) {
   int a;
  a = (b + c) * d;
  mov ?, %acc
  add ?, %acc
  mul ?,%acc
      %acc,?
```

- One copy of b for duration of program's execution
 - *b*'s address is the same for each call to func
- Different copies of a, c, and d for each call to func
 - addresses are different in each call

mov

Relative Addresses

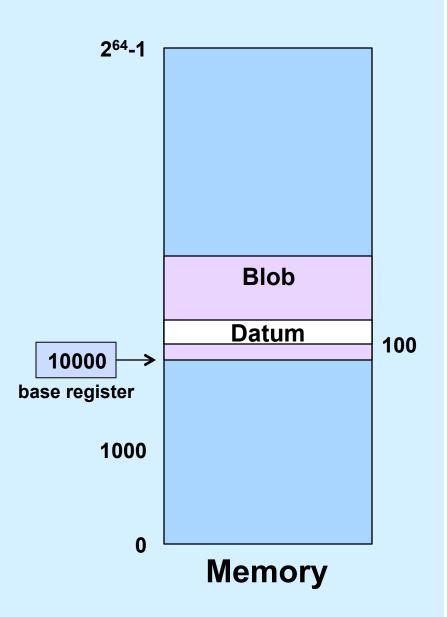
- Absolute address
 - actual location in memory
- Relative address
 - offset from some other location



264-1

Base Registers

mov \$10000, %base mov \$10, 100(%base)



Addresses

```
frame
int b;
                                      previous stack
                                          frame
                               base →
                                        func stack
int func(int c, int d) {
                                          frame
   int a;
   a = (b + c) * d;
                                 1000:
                                         globai
                                        variables
   mov 1000, %acc
         c rel(%base),%acc
   add
         d rel(%base),%acc
   mul
         %acc, a rel(%base)
   mov
                                        Memory
```

earlier stack

Quiz 2

Suppose the value in *base* is 10,000 and *c_rel* is -8. What is the address of *c*?

- a) 9992
- b) 9996
- c) 10,004
- d) 10,008

mov 1000,%acc
add c_rel(%base),%acc
mul d_rel(%base),%acc
mov %acc,a_rel(%base)

earlier stack frame previous stack frame base → func stack frame 1000: global variables

Registers

