

CS 33

More OS; Shells and Files

A Random Program ...

```
int main(int argc, char *argv[]) {  
    if (argc != 2) {  
        fprintf(stderr, "Usage: random count\n");  
        exit(1);  
    }  
    int stop = atoi(argv[1]);  
    for (int i = 0; i < stop; i++)  
        printf("%d\n", rand());  
    return 0;  
}
```

Passing It Arguments

- **From the shell**

```
$ random 12
```

- **From a C program**

```
if (fork() == 0) {  
    char *argv[] = {"random", "12", (void *)0};  
    execv("./random", argv);  
}
```

Quiz 1

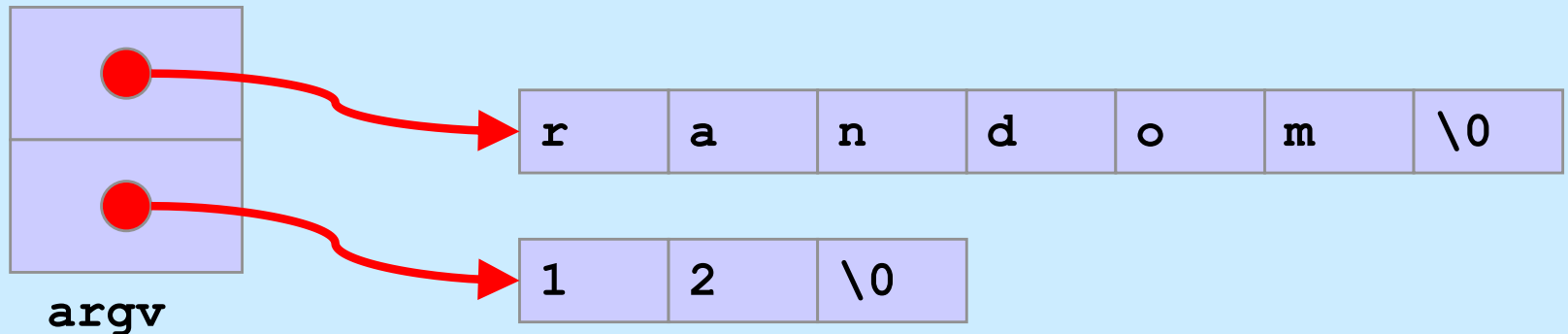
```
if (fork() == 0) {  
    char *argv[] = {"random", "12", (void *)0};  
    execl(".", "random", argv);  
    printf("random done\n");  
}
```

The *printf* statement will be executed

- a) only if execl fails
- b) only if execl succeeds
- c) always

Receiving Arguments

```
int main(int argc, char *argv[]) {  
    if (argc != 2) {  
        fprintf(stderr, "Usage: random count\n");  
        exit(1);  
    }  
    int stop = atoi(argv[1]);  
    for (int i = 0; i < stop; i++)  
        printf("%d\n", rand());  
  
    return 0;  
}
```



Not So Fast ...

- How does the shell invoke your program?

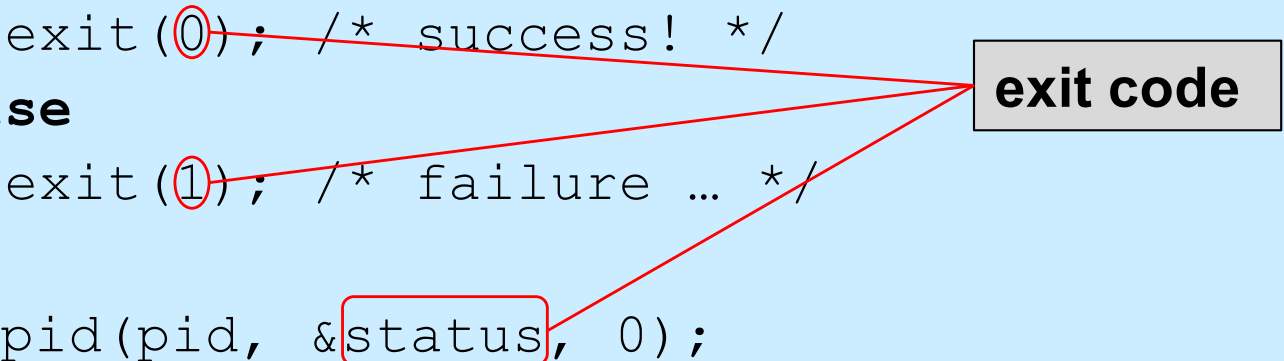
```
if (fork() == 0) {  
    char *argv = {"random", "12", (void *)0};  
    execv("./random", argv);  
}  
/* what does the shell do here??? */
```

Wait

```
#include <unistd.h>
#include <sys/wait.h>
...
pid_t pid;
int status;
...
if ((pid = fork()) == 0) {
    char *argv[] = {"random", "12", (void *)0};
    execv("./random", argv);
}
waitpid(pid, &status, 0);
```

Exit

```
#include <unistd.h>
#include <stdlib.h>
#include <sys/wait.h>
int main( ) {
    pid_t pid;
    int status;
    if ((pid = fork()) == 0) {
        if (do_work() == 1)
            exit(0); /* success! */
        else
            exit(1); /* failure ... */
    }
    waitpid(pid, &status, 0);
    /* low-order byte of status contains exit code.
       WEXITSTATUS(status) extracts it */
}
```



The diagram illustrates the flow of exit codes from the code to a box labeled "exit code". Red lines connect the exit codes in the code to the box: one from the 0 in `exit(0);`, one from the 1 in `exit(1);`, and one from the `status` variable in `&status`.

Shell: To Wait or Not To Wait ...

```
$ who
```

```
    if ((pid = fork()) == 0) {  
        char *argv[] = {"who", 0};  
        execv("who", argv);  
    }  
    waitpid(pid, &status, 0);  
    ...
```

```
$ who &
```

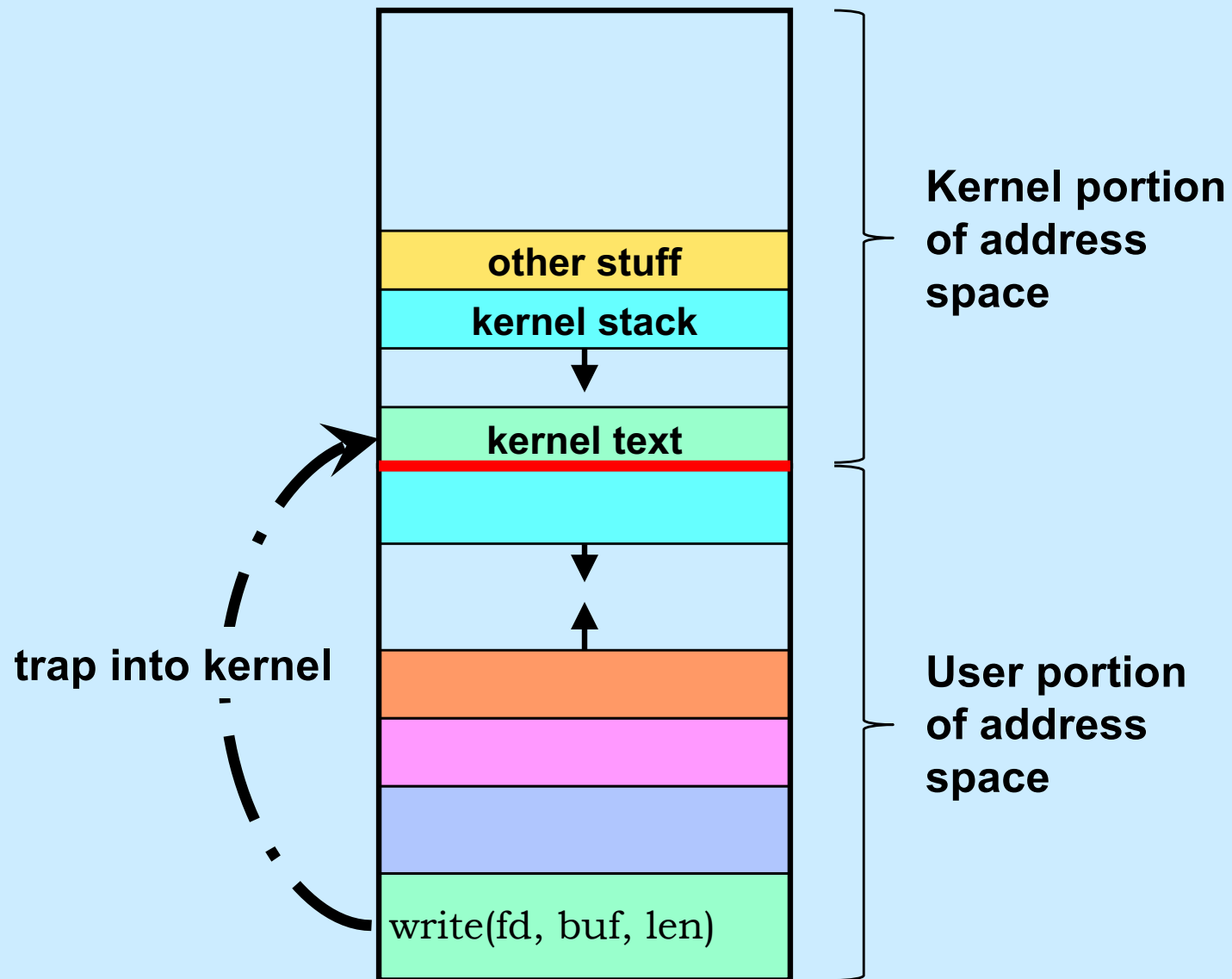
```
    if ((pid = fork()) == 0) {  
        char *argv[] = {"who", 0};  
        execv("who", argv);  
    }  
    ...
```

System Calls

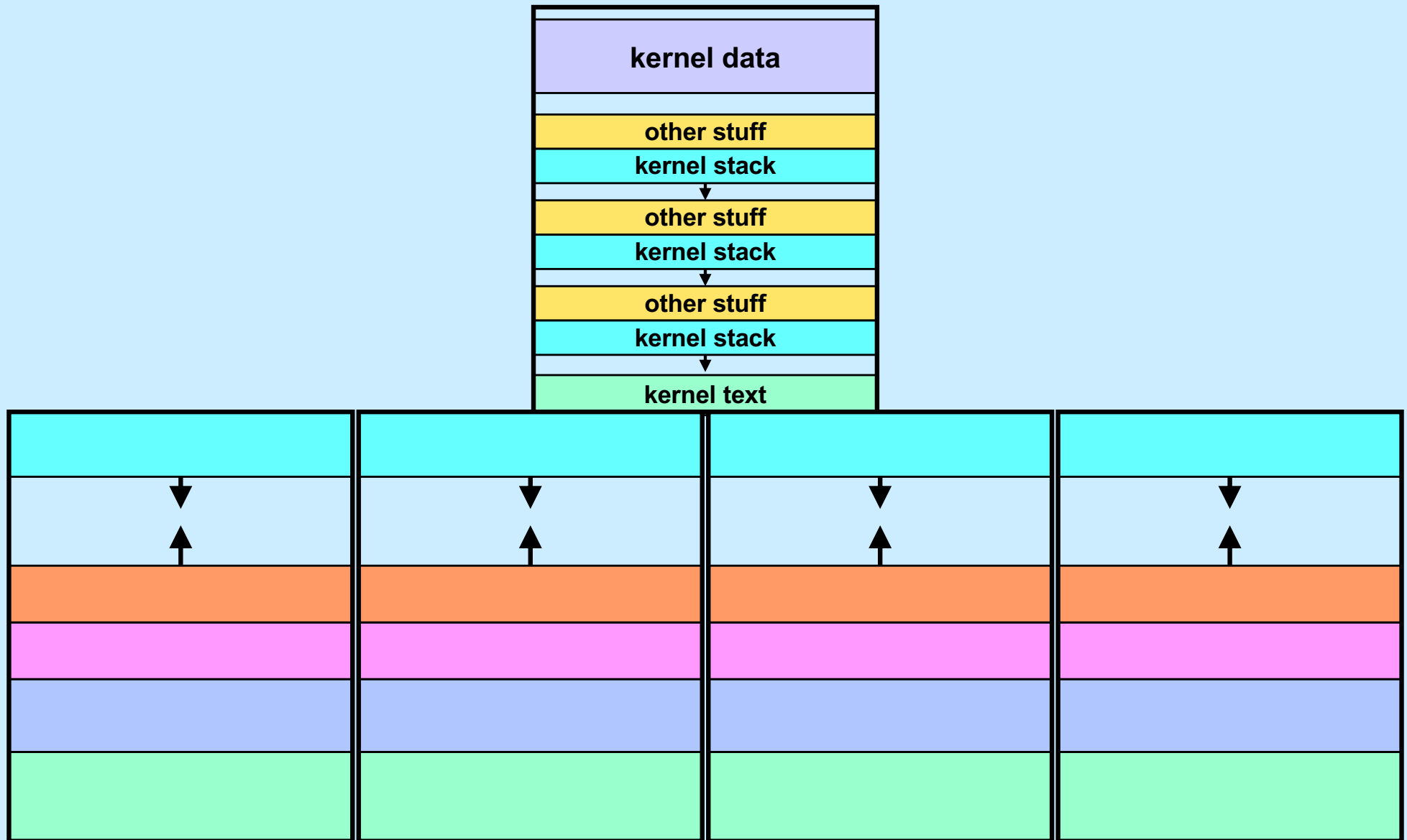
- Sole direct interface between user and kernel
- Implemented as library function that execute *trap* instructions to enter kernel
- Errors indicated by returns of -1 ; error code is in global variable *errno*

```
if (write(fd, buffer, bufsz) == -1) {  
    // error!  
    printf("error %d\n", errno);  
    // see perror  
}
```

System Calls



Multiple Processes

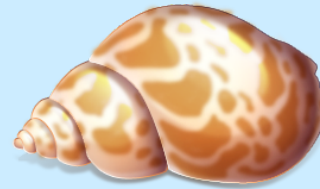
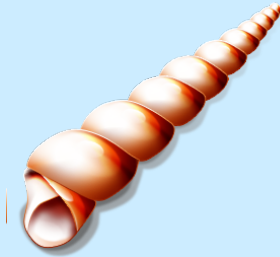


Shells




- **Command and scripting languages for Unix**
- **First shell: Thompson shell**
 - sh, developed by Ken Thompson
 - released in 1971
- **Bourne shell**
 - also sh, developed by Steve Bourne
 - released in 1977
- **C shell**
 - csh, developed by Bill Joy
 - released in 1978
 - tcsh, improved version by Ken Greer

More Shells



- **Bourne-Again Shell**
 - bash, developed by Brian Fox
 - released in 1989
 - found to have a serious security-related bug in 2014
 - » shellshock
- **Almquist Shell**
 - ash, developed by Kenneth Almquist
 - released in 1989
 - similar to bash
 - dash (debian ash) used for scripts in Debian Linux
 - » faster than bash
 - » less susceptible to shellshock vulnerability

Roadmap

- **We explore the file abstraction**
 - what are files
 - how do you use them
 - how does the OS represent them
 - **We explore the shell**
 - how does it launch programs
 - how does it connect programs with files
 - how does it control running programs
- 
- shell 1
- shell 2

The File Abstraction

- A file is a simple array of bytes
- A file is made larger by writing beyond its current end
- Files are named by paths in a naming tree
- System calls on files are synchronous

Naming

- **(almost) everything has a path name**
 - files
 - directories
 - devices (known as *special files*)
 - » keyboards
 - » displays
 - » disks
 - » etc.

I/O System Calls

- **int** file_descriptor = open(pathname, mode [, permissions])
- **int** close(file_descriptor)
- **ssize_t** count = read(file_descriptor, buffer_address, buffer_size)
- **ssize_t** count = write(file_descriptor, buffer_address, buffer_size)
- **off_t** position lseek(file_descriptor, offset, whence)

Standard File Descriptors

```
int main( ) {  
    char buf[BUFSIZE];  
    int n;  
    const char *note = "Write failed\n";  
  
    while ((n = read(0, buf, sizeof(buf))) > 0)  
        if (write(1, buf, n) != n) {  
            write(2, note, strlen(note));  
            exit(1);  
        }  
    return (0);  
}
```

Standard I/O Library

Formatting

printf ... **scanf**

Buffering

stdin **stdout** **stderr** ...

Syscalls

fd 0 **fd 1** **fd 2** ...

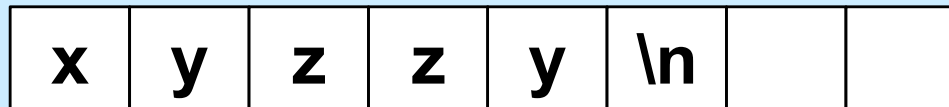
Standard I/O

```
FILE *stdin;           // declared in stdio.h
FILE *stdout;          // declared in stdio.h
FILE *stderr;          // declared in stdio.h
```

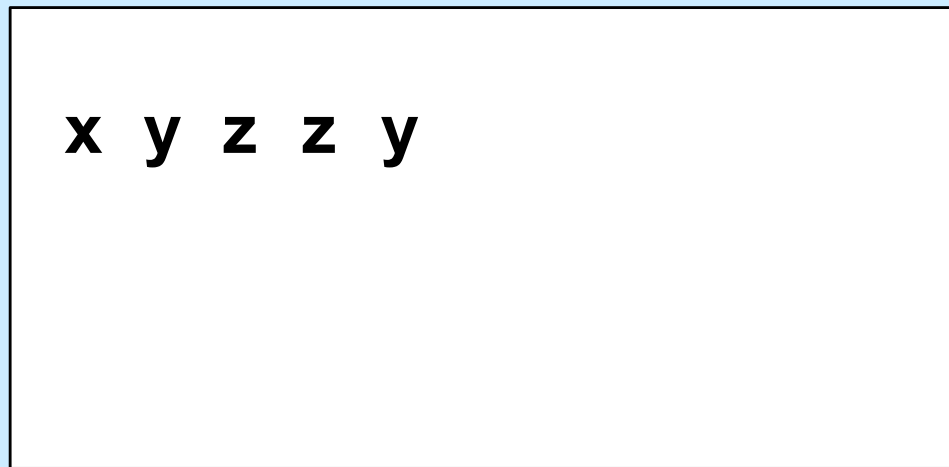
```
scanf("%d", &in);      // read via f.d. 0
printf("%d\n", in);    // write via f.d. 1
fprintf(stderr, "there was an error\n");
    // write via f.d. 2
```

Buffered Output

```
printf("xy");  
printf("zz");  
printf("y\n");
```



buffer



display

Unbuffered Output

```
fprintf(stderr, "xy");  
fprintf(stderr, "zz");  
fprintf(stderr, "y\n");
```

x y z z y

display

A Program

```
int main(int argc, char *argv[]) {  
    if (argc != 2) {  
        fprintf(stderr, "Usage: echon reps\n");  
        exit(1);  
    }  
    int reps = atoi(argv[1]);  
    if (reps > 2) {  
        fprintf(stderr, "reps too large, reduced to 2\n");  
        reps = 2;  
    }  
    char buf[256];  
    while (fgets(buf, 256, stdin) != NULL)  
        for (int i=0; i<reps; i++)  
            fputs(buf, stdout);  
    return (0);  
}
```


From the Shell ...

```
$ echon 1
```

- ***stdout*** (fd 1) and ***stderr*** (fd 2) go to the display
- ***stdin*** (fd 0) comes from the keyboard

```
$ echon 1 > Output
```

- ***stdout*** goes to the file “Output” in the current directory
- ***stderr*** goes to the display
- ***stdin*** comes from the keyboard

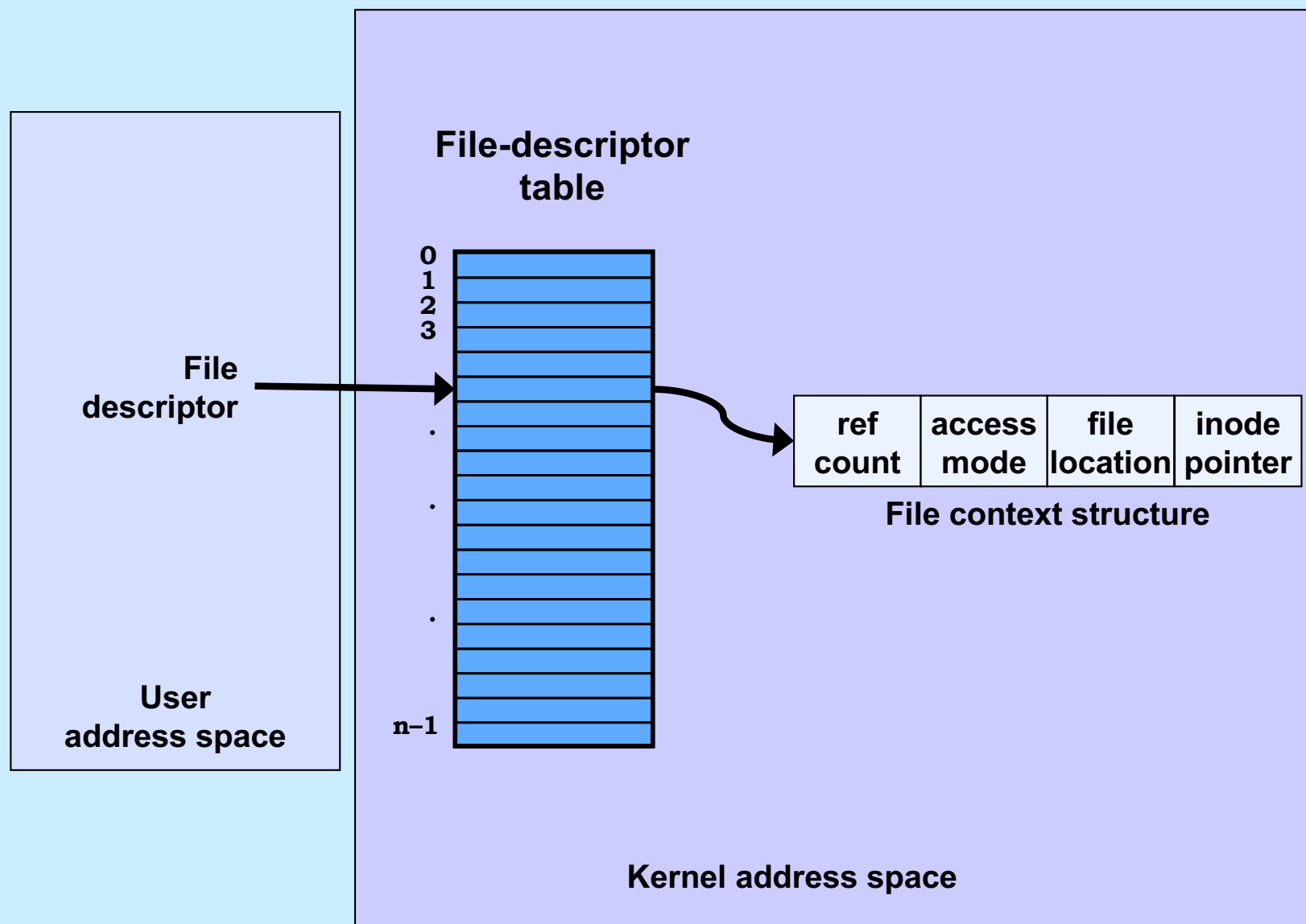
```
$ echon 1 < Input
```

- ***stdin*** comes from the file “Input” in the current directory

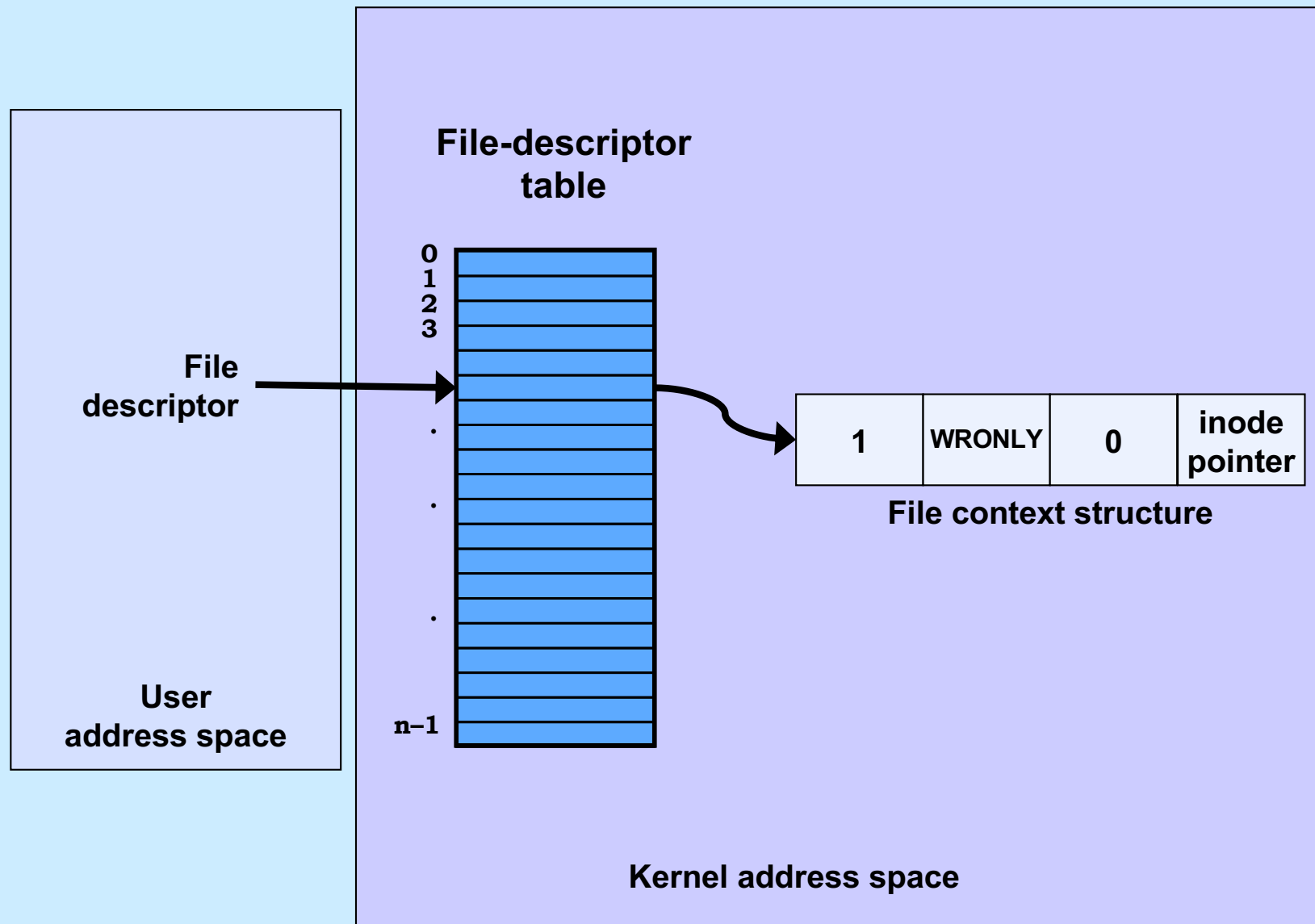
Running It

```
if (fork() == 0) {  
    /* set up file descriptor 1 in the child process */  
    close(1);  
    if (open("/home/twd/Output", O_WRONLY) == -1) {  
        perror("/home/twd/Output");  
        exit(1);  
    }  
    char *argv[] = {"echon", "2", NULL};  
    execv("/home/twd/bin/echon", argv);  
    exit(1);  
}  
  
/* parent continues here */  
  
while(pid != wait(0))      /* ignore the return code */  
    ;
```

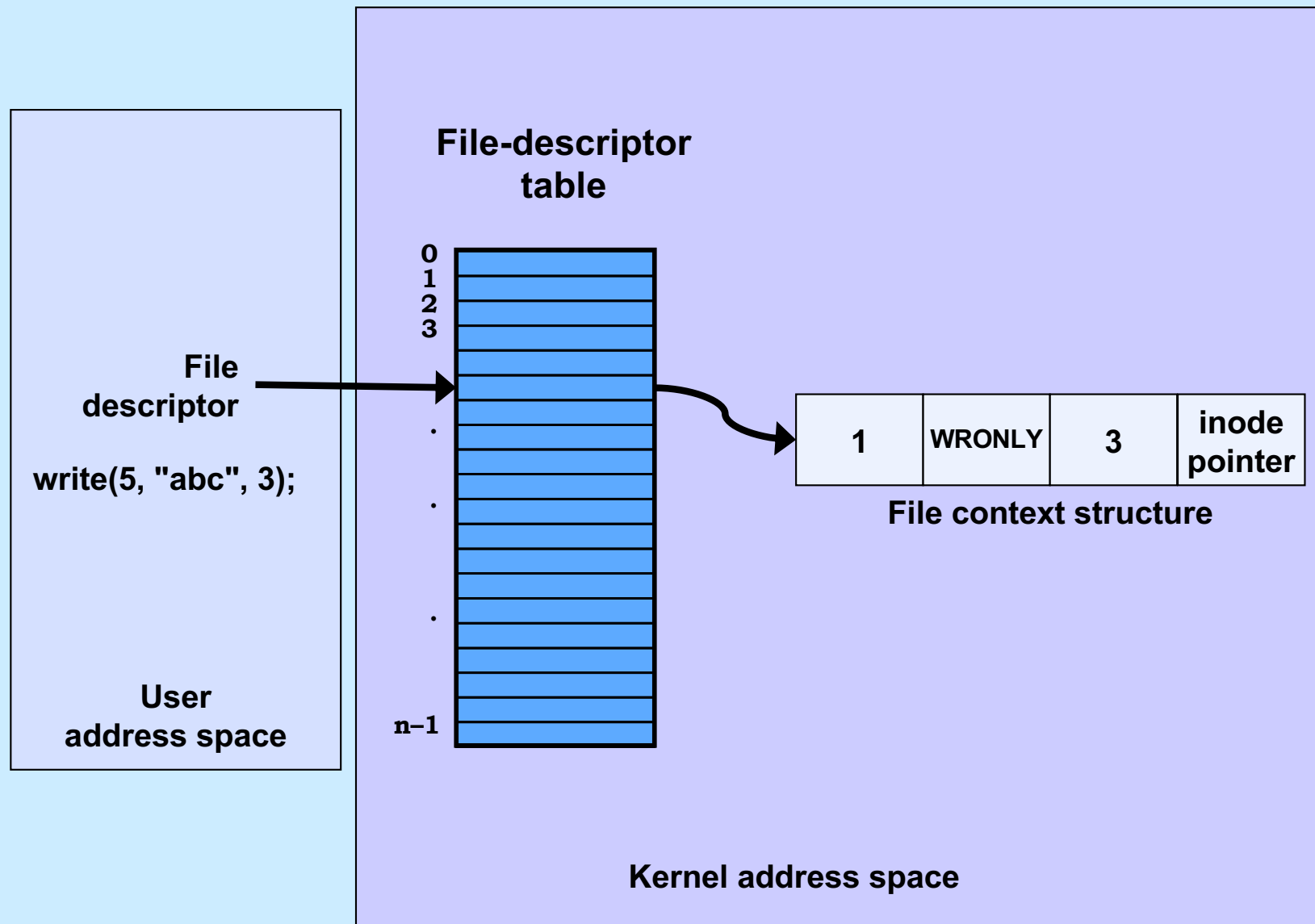
File-Descriptor Table



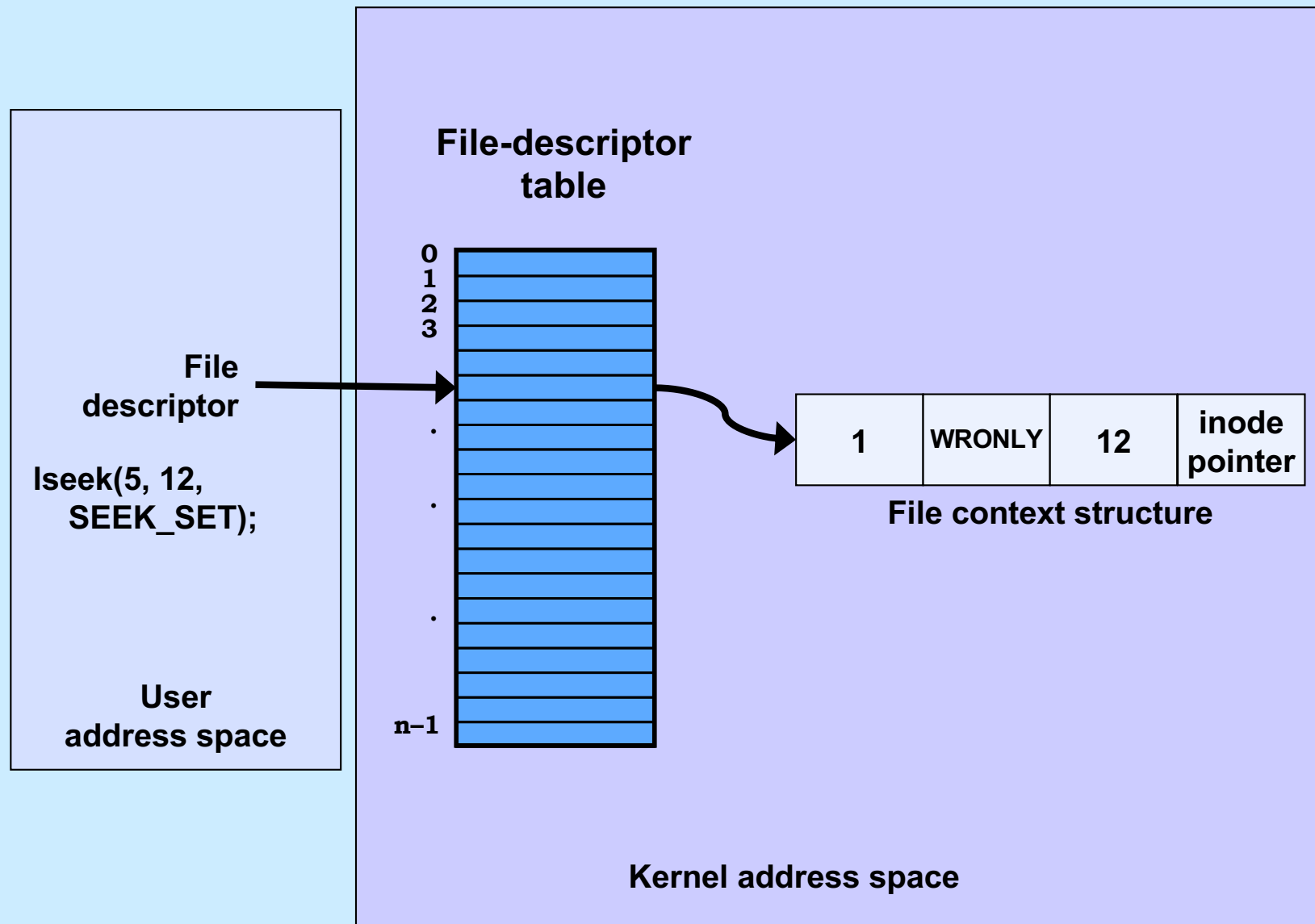
File Location



File Location



File Location



Allocation of File Descriptors

- Whenever a process requests a new file descriptor, the lowest-numbered file descriptor not already associated with an open file is selected; thus

```
#include <fcntl.h>
#include <unistd.h>
```

```
close(0);
fd = open("file", O_RDONLY);
```

- will always associate *file* with file descriptor 0 (assuming that *open* succeeds)

Redirecting Output ... Twice

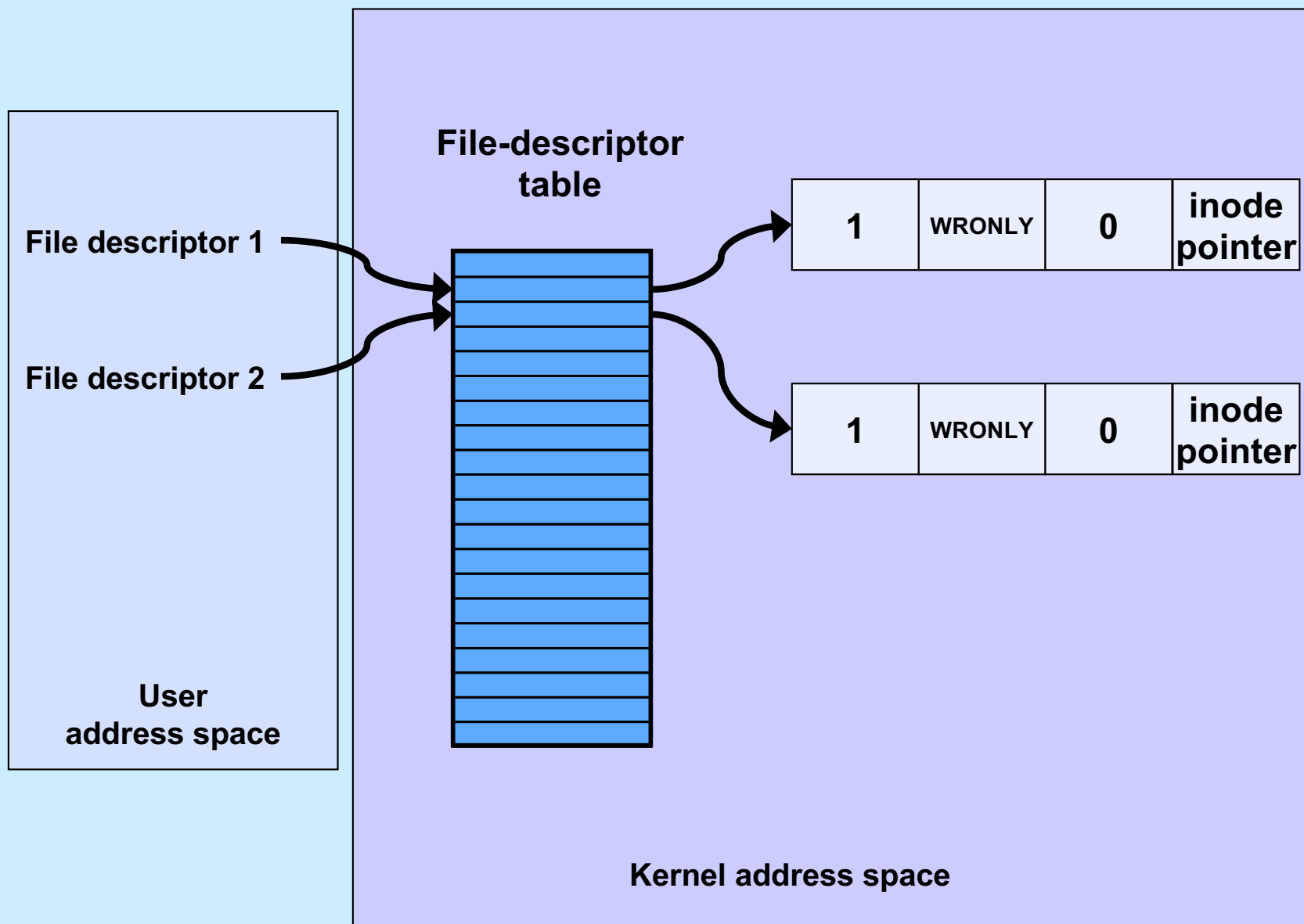
```
if (fork() == 0) {  
    /* set up file descriptors 1 and 2 in the child process */  
    close(1);  
    close(2);  
    if (open("/home/twd/Output", O_WRONLY) == -1) {  
        exit(1);  
    }  
    if (open("/home/twd/Output", O_WRONLY) == -1) {  
        exit(1);  
    }  
    char *argv[] = {"echon", 2, NULL};  
    execv("/home/twd/bin/echon", argv);  
    exit(1);  
}  
/* parent continues here */
```


From the Shell ...

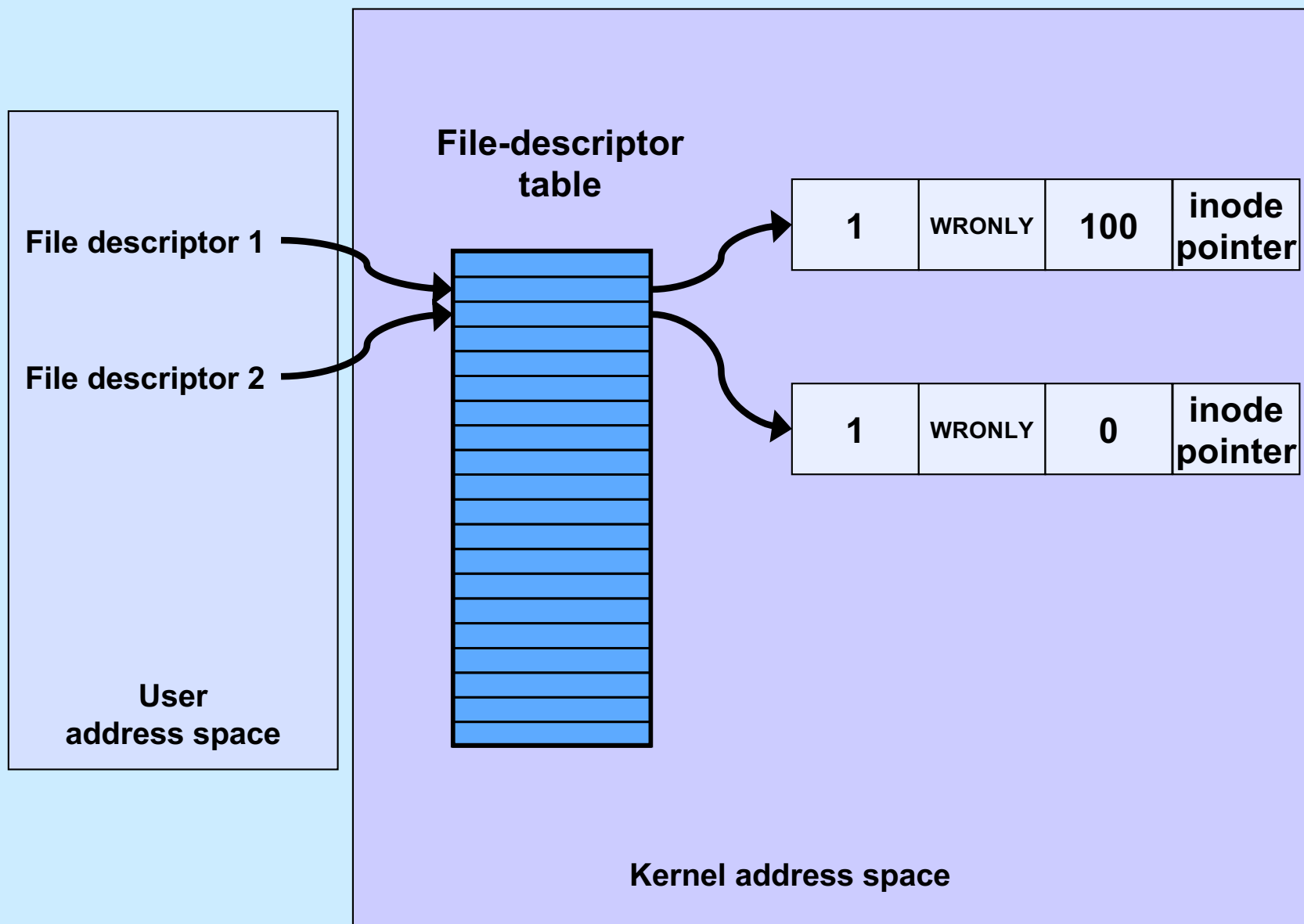
```
$ echon 1 >Output 2>Output
```

– **both stdout and stderr go to Output file**

Redirected Output



Redirected Output After Write



Not a Quiz

- **Suppose we run**

```
$ echo 3 >Output 2>Output
```

- **The input line is**

X

- **What is the final content of Output?**

a) reps too large, reduced to 2\nX\nX\n

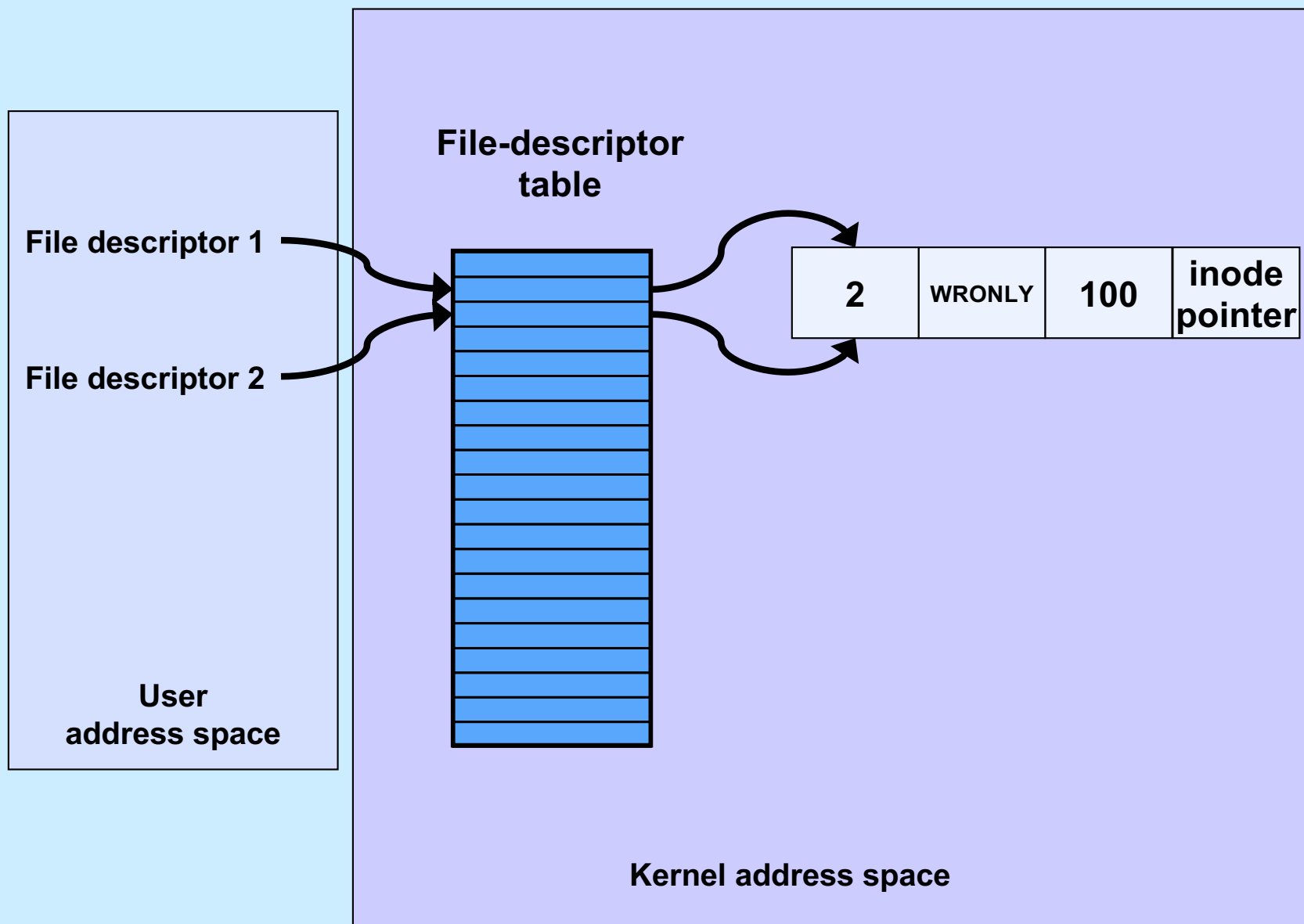
b) X\nX\nreps too large, reduced to 2\n

c) X\nX\n too large, reduced to 2\n

Sharing Context Information

```
if (fork() == 0) {  
    /* set up file descriptors 1 and 2 in the child process */  
    close(1);  
    close(2);  
    if (open("/home/twd/Output", O_WRONLY) == -1) {  
        exit(1);  
    }  
    dup(1); /* set up file descriptor 2 as a duplicate of 1 */  
    char *argv[] = {"echon", 2};  
    execv("/home/twd/bin/echon", argv);  
    exit(1);  
}  
/* parent continues here */
```

Redirected Output After Dup



From the Shell ...

```
$ echon 3 >Output 2>&1
```

- **stdout goes to Output file, stderr is the dup of fd 1**

- **with input “X\n” it now produces in Output:**

```
reps too large, reduced to 2\nX\nX\n
```

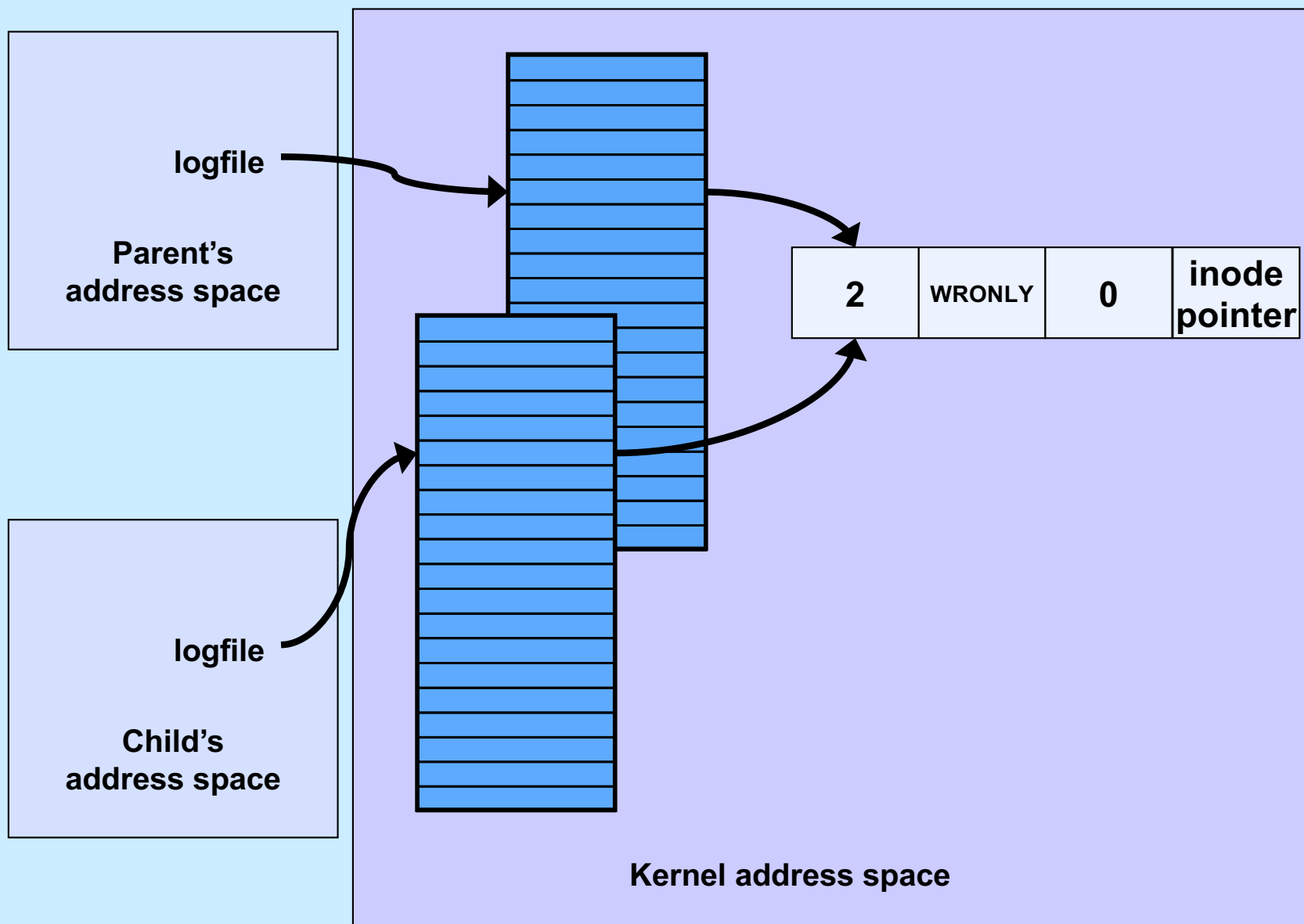
Fork and File Descriptors

```
int logfile = open("log", O_WRONLY);
if (fork() == 0) {
    /* child process computes something, then does: */
    write(logfile, LogEntry, strlen(LogEntry));
    ...
    exit(0);
}

/* parent process computes something, then does: */

write(logfile, LogEntry, strlen(LogEntry));
...
```


File Descriptors After Fork



Quiz 2

```
int main() {  
    if (fork() == 0) {  
        fprintf(stderr, "Child");  
        exit(0);  
    }  
    fprintf(stderr, "Parent");  
}
```

Suppose the program is run as:

```
$ prog >file 2>&1
```

What is the final content of file? (Assume writes are “atomic”.)

- a) either “ChildParent” or “ParentChild”
- b) either “Childt” or “Parent”
- c) either “Child” or “Parent”