

# CS 33

## Data Representation (Part 3)

# Floating-Point Operations: Basic Idea

- $x +_f y = \text{Round}(x + y)$
- $x \times_f y = \text{Round}(x \times y)$
- Basic idea
  - first **compute exact result**
  - make it fit into desired precision
    - » possibly overflow if exponent too large
    - » possibly **round to fit into** frac

# Rounding

- Rounding modes (illustrated with \$ rounding)

	<b>\$1.40</b>	<b>\$1.60</b>	<b>\$1.50</b>	<b>\$2.50</b>	<b>-\$1.50</b>
<b>towards zero</b>	<b>\$1</b>	<b>\$1</b>	<b>\$1</b>	<b>\$2</b>	<b>-\$1</b>
<b>round down (<math>-\infty</math>)</b>	<b>\$1</b>	<b>\$1</b>	<b>\$1</b>	<b>\$2</b>	<b>-\$2</b>
<b>round up (<math>+\infty</math>)</b>	<b>\$2</b>	<b>\$2</b>	<b>\$2</b>	<b>\$3</b>	<b>-\$1</b>
<b>nearest integer</b>	<b>\$1</b>	<b>\$2</b>	<b>?</b>	<b>?</b>	<b>?</b>
<b>nearest even (default)</b>	<b>\$1</b>	<b>\$2</b>	<b>\$2</b>	<b>\$2</b>	<b>-\$2</b>

# Floating-Point Multiplication

- $(-1)^{s1} M1 2^{E1} \times (-1)^{s2} M2 2^{E2}$
- **Exact result:**  $(-1)^s M 2^E$ 
  - sign  $s$ :  $s1 \wedge s2$
  - significand  $M$ :  $M1 \times M2$
  - exponent  $E$ :  $E1 + E2$
- **Fixing**
  - if  $M \geq 2$ , shift  $M$  right, increment  $E$
  - if  $E$  out of range, overflow (or underflow)
  - round  $M$  to fit `frac` precision
- **Implementation**
  - biggest chore is multiplying significands

# Floating-Point Addition

- $(-1)^{s1} M1 2^{E1} + (-1)^{s2} M2 2^{E2}$

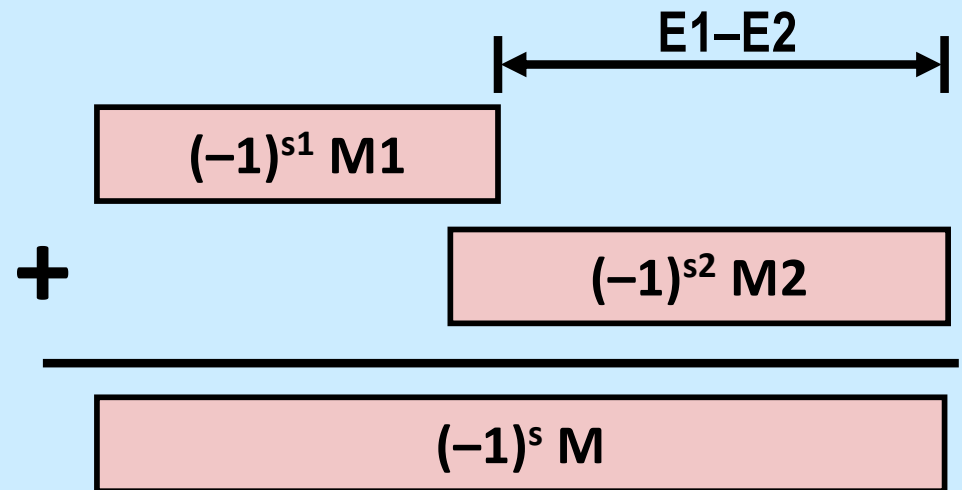
–assume  $E1 > E2$

- **Exact result:**  $(-1)^s M 2^E$

–sign  $s$ , significand  $M$ :

» result of signed align & add

–exponent  $E$ :  $E1$



- **Fixing**

–if  $M \geq 2$ , shift  $M$  right, increment  $E$

–if  $M < 1$ , shift  $M$  left  $k$  positions, decrement  $E$  by  $k$

–overflow if  $E$  out of range

–round  $M$  to fit **frac** precision

# Floating Point

- **Single precision (float)**



– range:  $\pm 1.8 \times 10^{-38}$  –  $\pm 3.4 \times 10^{38}$ , ~7 decimal digits

- **Double Precision (double)**



– range:  $\pm 2.23 \times 10^{-308}$  –  $\pm 1.8 \times 10^{308}$ , ~16 decimal digits

# Floating Point in C

- **Conversions/casting**

- casting between `int`, `float`, and `double` changes bit representation
- `double/float`  $\rightarrow$  `int`
  - » truncates fractional part
  - » like rounding toward zero
  - » not defined when out of range or NaN: generally sets to TMin
- `int`  $\rightarrow$  `double`
  - » exact conversion, as long as `int` has  $\leq 53$ -bit word size
- `int`  $\rightarrow$  `float`
  - » will round according to rounding mode

# Quiz 1

Suppose  $f$ , declared to be a `float`, is assigned the largest possible floating-point positive value (other than  $+\infty$ ). What is the value of  $g = f + 1.0$ ?

- a)  $f$
- b)  $+\infty$
- c) NaN
- d) 0



# Float is not Rational ...

- **Floating addition**
  - **commutative:  $a +_f b = b +_f a$** 
    - » **yes!**
  - **associative:  $a +_f (b +_f c) = (a +_f b) +_f c$** 
    - » **no!**
      - **$2 +_f (1e38 +_f -1e38) = 2$**
      - **$(2 +_f 1e38) +_f -1e38 = 0$**

# Float is not Rational ...

- **Multiplication**

- **commutative:**  $a *_f b = b *_f a$

- » **yes!**

- **associative:**  $a *_f (b *_f c) = (a *_f b) *_f c$

- » **no!**

- $1e37 *_f (1e37 *_f 1e-37) = 1e37$

- $(1e37 *_f 1e37) *_f 1e-37 = +\infty$

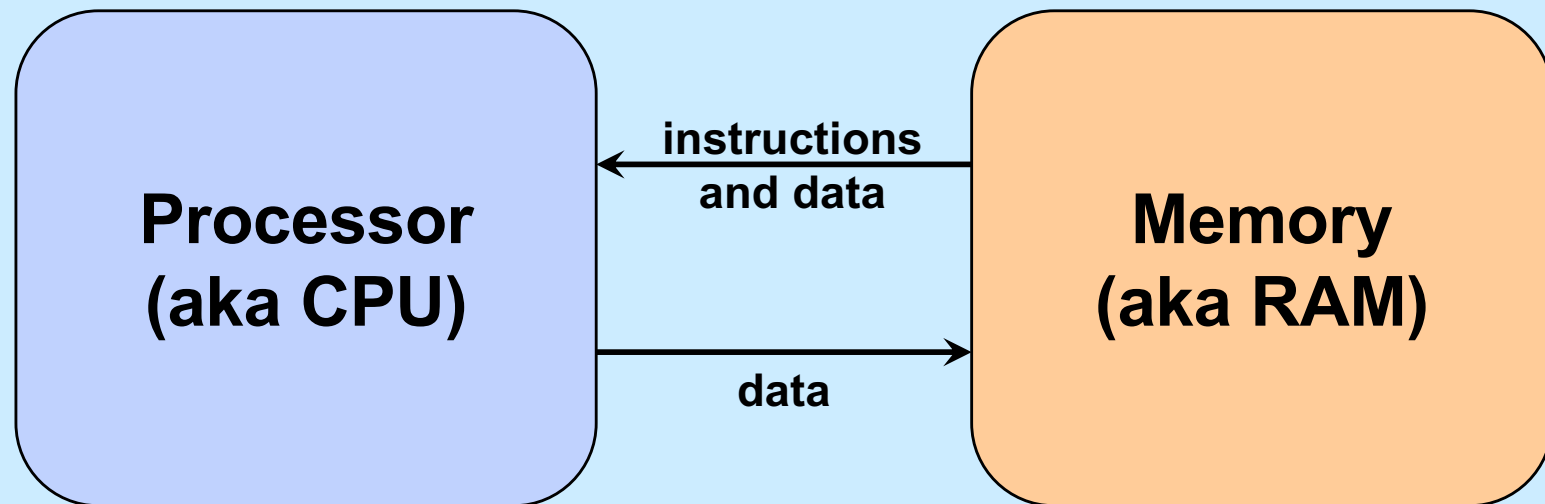
# Float is not Rational ...

- More ...
  - multiplication distributes over addition:  
$$a *_f (b +_f c) = (a *_f b) +_f (a *_f c)$$
    - » no!
    - »  $1e38 *_f (1e38 +_f -1e38) = 0$
    - »  $(1e38 *_f 1e38) +_f (1e38 *_f -1e38) = \text{NaN}$
  - insignificance:  
$$x = y +_f 1$$
$$z = 2 /_f (x -_f y)$$
$$z == 2?$$
    - » not necessarily!
      - consider  $y = 1e38$

# CS 33

## Intro to Machine Programming

# Machine Model



# Memory



**Instructions**

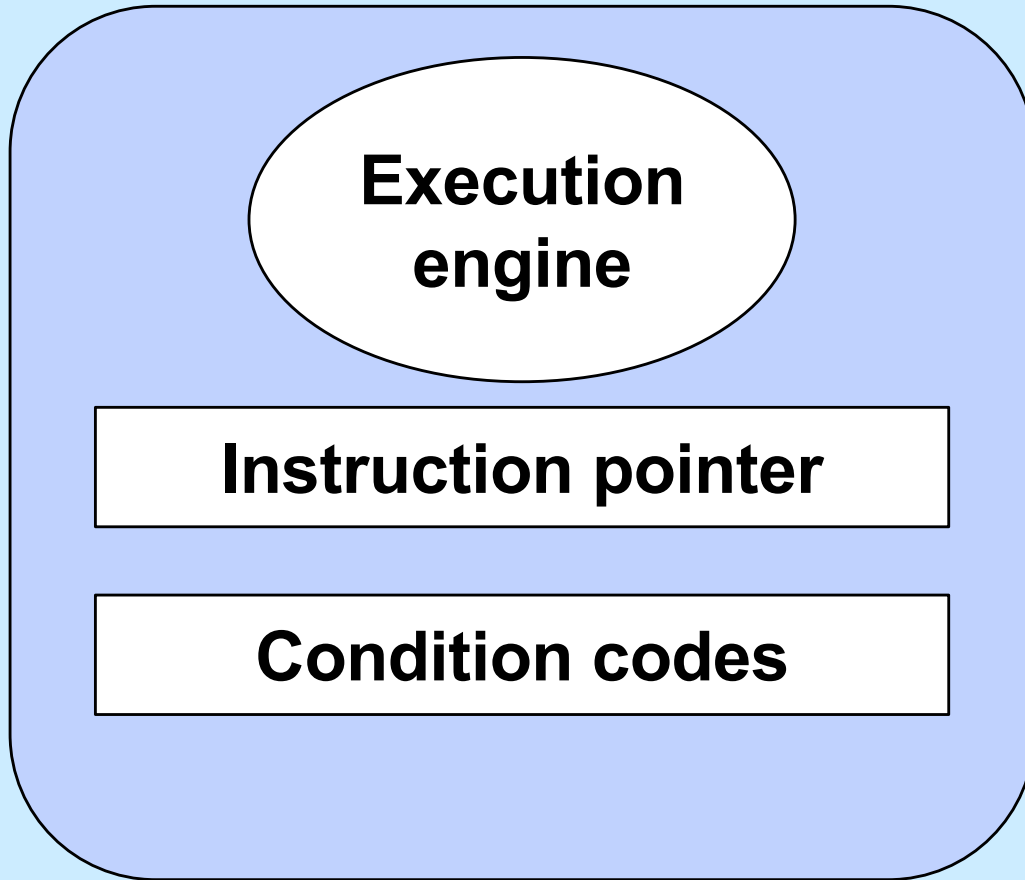
**Data**

**or**



**Instructions  
are Data**

# Processor: Some Details



# Processor: Basic Operation

```
while (forever) {  
  fetch instruction IP points at  
  decode instruction  
  fetch operands  
  execute  
  store results  
  update IP and condition code  
}
```



# Instructions ...

<b>Op code</b>	<b>Operand1</b>	<b>Operand2</b>	<b>...</b>
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# Operands

- **Form**
  - immediate vs. reference
    - » value vs. address
- **How many?**
  - 3
    - » add a,b,c
      - $c = a + b$
  - 2
    - » add a,b
      - $b += a$

# Operands (continued)

- **Accumulator**
  - special memory in the processor
    - » known as a *register*
    - » fast access
  - allows single-operand instructions
    - » add a
      - `acc += a`
    - » add b
      - `acc += b`

# From C to Assembler ...

```
a = (b + c) * d;
```

```
mov    b,%acc  
add    c,%acc  
mul    d,%acc  
mov    %acc,a
```

```
if (a<b)  
    c = 1;  
else  
    d = 1;
```

```
cmp    a,b  
jge    .L1  
mov    $1,c  
jmp    .L2  
.L1  
mov    $1,d  
.L2
```

immediate operand

immediate operand

# Condition Codes

- **Set of flags giving status of most recent operation:**
  - **zero flag**
    - » result was zero
  - **sign flag**
    - » for signed arithmetic interpretation: sign bit is set
  - **overflow flag**
    - » for signed arithmetic interpretation
  - **carry flag (generated by carry or borrow out of most-significant bit)**
    - » for unsigned arithmetic interpretation
- **Set implicitly by arithmetic instructions**
- **Set explicitly by compare instruction**
  - **cmp a,b**
    - » sets flags based on result of  $b-a$

# Examples (1)

- **Assume 32-bit arithmetic**
- **x is 0x80000000**
  - **TMIN** if interpreted as two's-complement
  - **$2^{31}$**  if interpreted as unsigned
- **x-1 (0x7fffffff)**
  - **TMAX** if interpreted as two's-complement
  - **$2^{31}-1$**  if interpreted as unsigned
  - **zero flag** is not set
  - **sign flag** is not set
  - **overflow flag** is set
  - **carry flag** is not set

# Examples (2)

- **x is 0xffffffff**
  - -1 if interpreted as two's-complement
  - UMAX ( $2^{32}-1$ ) if interpreted as unsigned
- **x+1 (0x00000000)**
  - zero under either interpretation
  - zero flag is set
  - sign flag is not set
  - overflow flag is not set
  - carry flag is set

# Examples (3)

- **x is 0xffffffff**
  - -1 if interpreted as two's-complement
  - UMAX ( $2^{32}-1$ ) if interpreted as unsigned
- **x+2 (0x00000001)**
  - (+)1 under either interpretation
  - zero flag is not set
  - sign flag is not set
  - overflow flag is not set
  - carry flag is set



# Quiz 2

- **Set of flags giving status of most recent operation:**
  - zero flag
    - » result was zero
  - sign flag
    - » for signed arithmetic interpretation: sign bit is set
  - overflow flag
    - » for signed arithmetic interpretation
  - carry flag (generated by carry or borrow out of most-significant bit)
    - » for unsigned arithmetic interpretation
- **Set explicitly by compare instruction**
  - `cmp a,b`
    - » sets flags based on result of `b-a`

**Which flags are set to one by “`cmp 2,1`”?**

- a) overflow flag only**
- b) carry flag only**
- c) sign and carry flags only**
- d) sign and overflow flags only**
- e) sign, overflow, and carry flags**

# Jump Instructions

- **Unconditional jump**
  - just do it
- **Conditional jump**
  - to jump or not to jump determined by condition-code flags
  - field in the op code indicates how this is computed
  - in assembler language, simply say
    - » **je**
      - jump on equal
    - » **jne**
      - jump on not equal
    - » **jg**
      - jump on greater than (signed)
    - » **etc.**

# Addresses

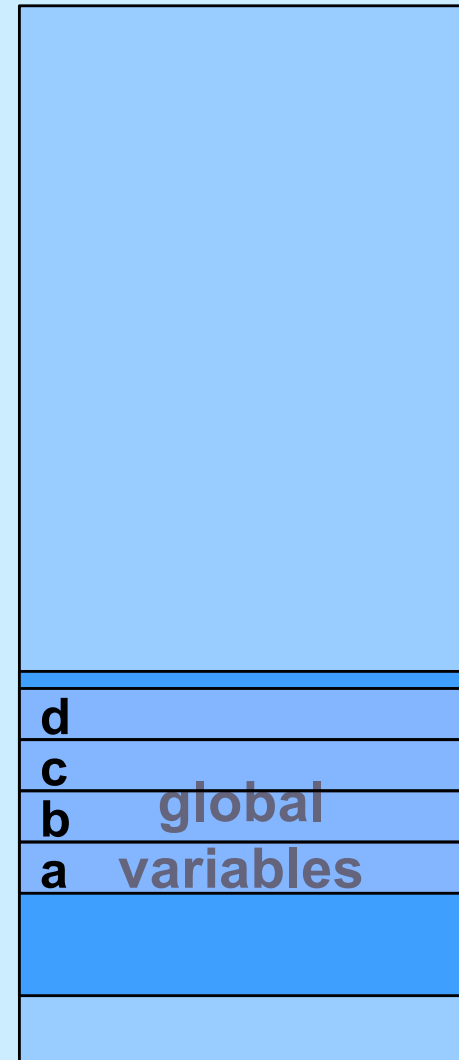
```
int a, b, c, d;
```

```
int main() {  
    a = (b + c) * d;  
    ...  
}
```

```
mov    b, %acc  
add    c, %acc  
mul    d, %acc  
mov    %acc, a
```

```
mov    1004, %acc  
add    1008, %acc  
mul    1012, %acc  
mov    %acc, 1000
```

1012:  
1008:  
1004:  
1000:



**Memory**

# Addresses

```
int b;
```

```
int func(int c, int d) {  
    int a;  
    a = (b + c) * d;  
    ...  
}
```

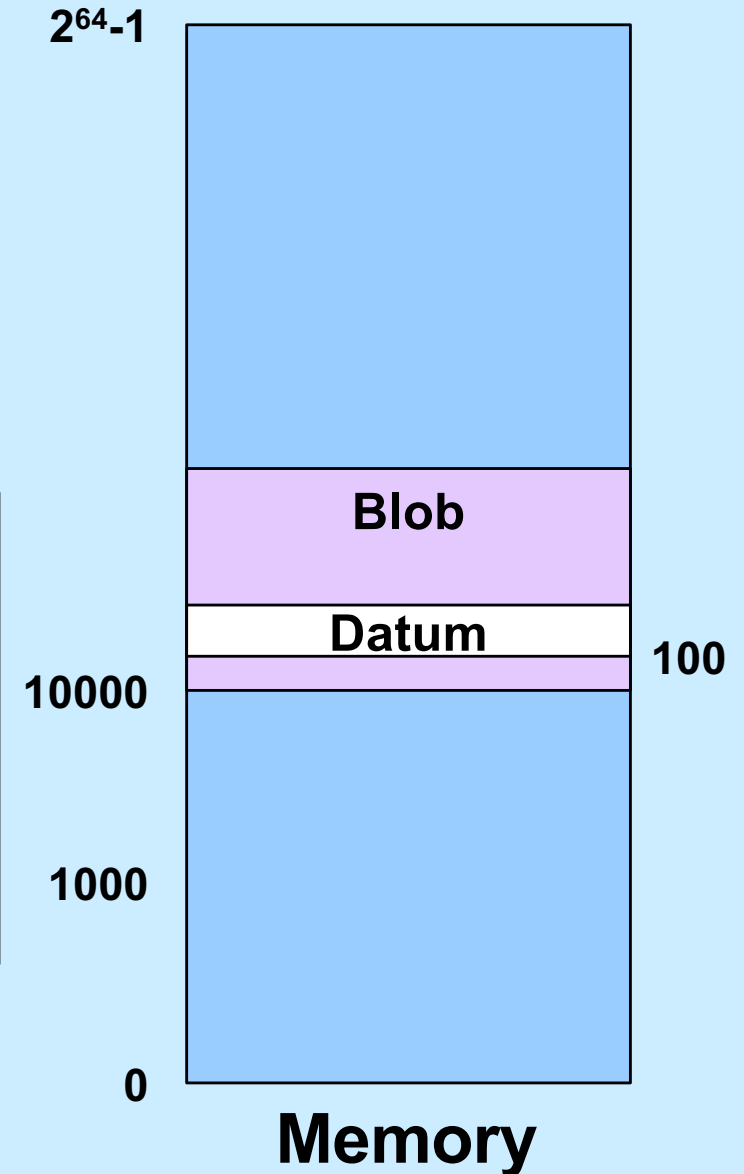
```
mov    ?, %acc  
add    ?, %acc  
mul    ?, %acc  
mov    %acc, ?
```

- One copy of *b* for duration of program's execution
  - *b*'s address is the same for each call to *func*
- Different copies of *a*, *c*, and *d* for each call to *func*
  - addresses are different in each call

# Relative Addresses

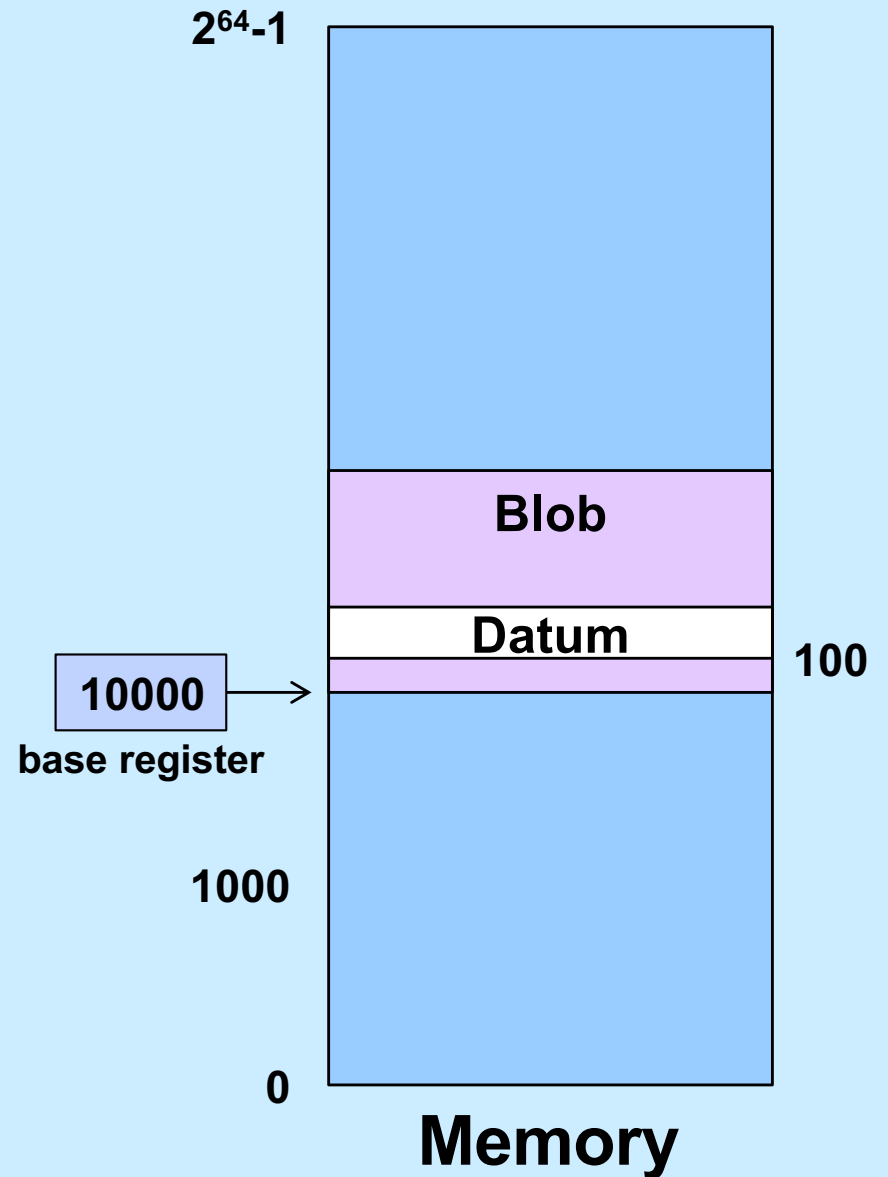
- **Absolute address**
  - actual location in memory
- **Relative address**
  - offset from some other location

- Blob's absolute address is 10000
- Datum's relative address (to Blob) is 100
  - its absolute address is 10100



# Base Registers

```
mov $10000, %base  
mov $10, 100(%base)
```

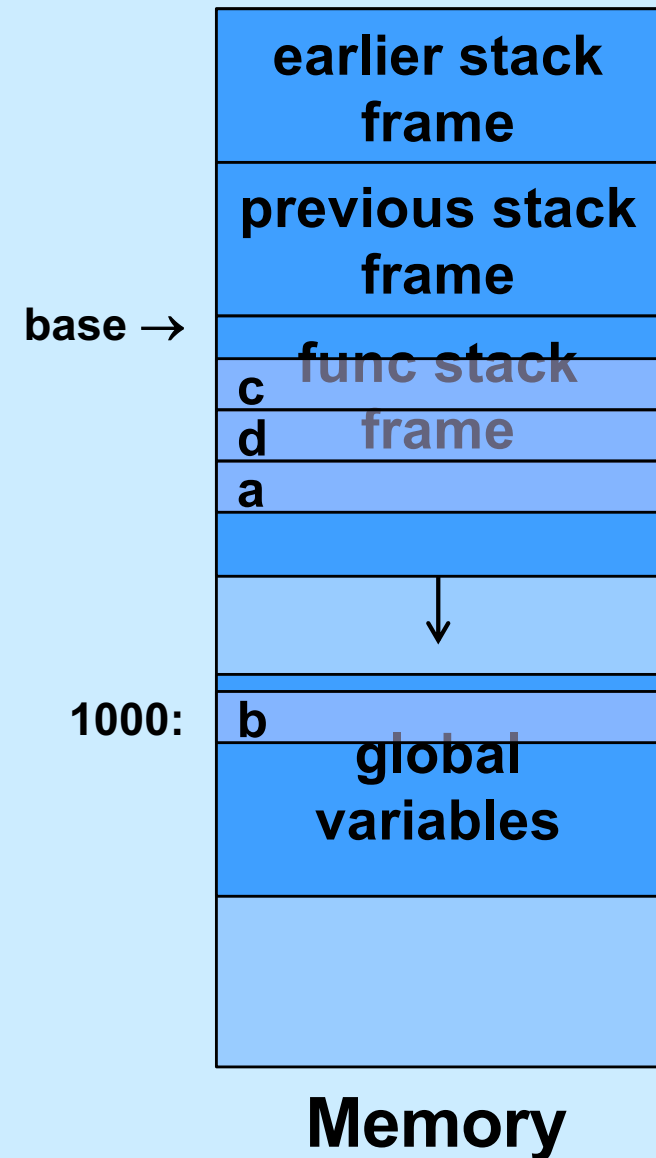


# Addresses

```
long b;
```

```
int func(long c, long d) {  
    long a;  
    a = (b + c) * d;  
    ...  
}
```

```
mov    1000,%acc  
add    -8(%base),%acc  
mul    -12(%base),%acc  
mov    %acc,-16(%base)
```

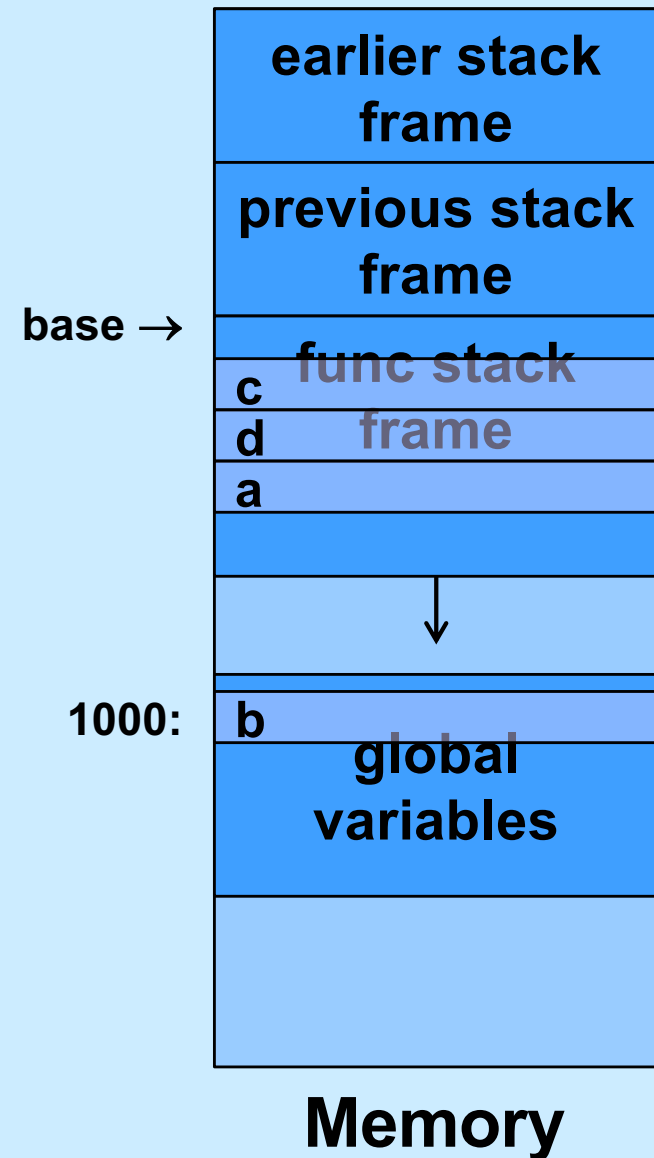


# Quiz 3

Suppose the value in *base* is 10,000. What is the address of *c*?

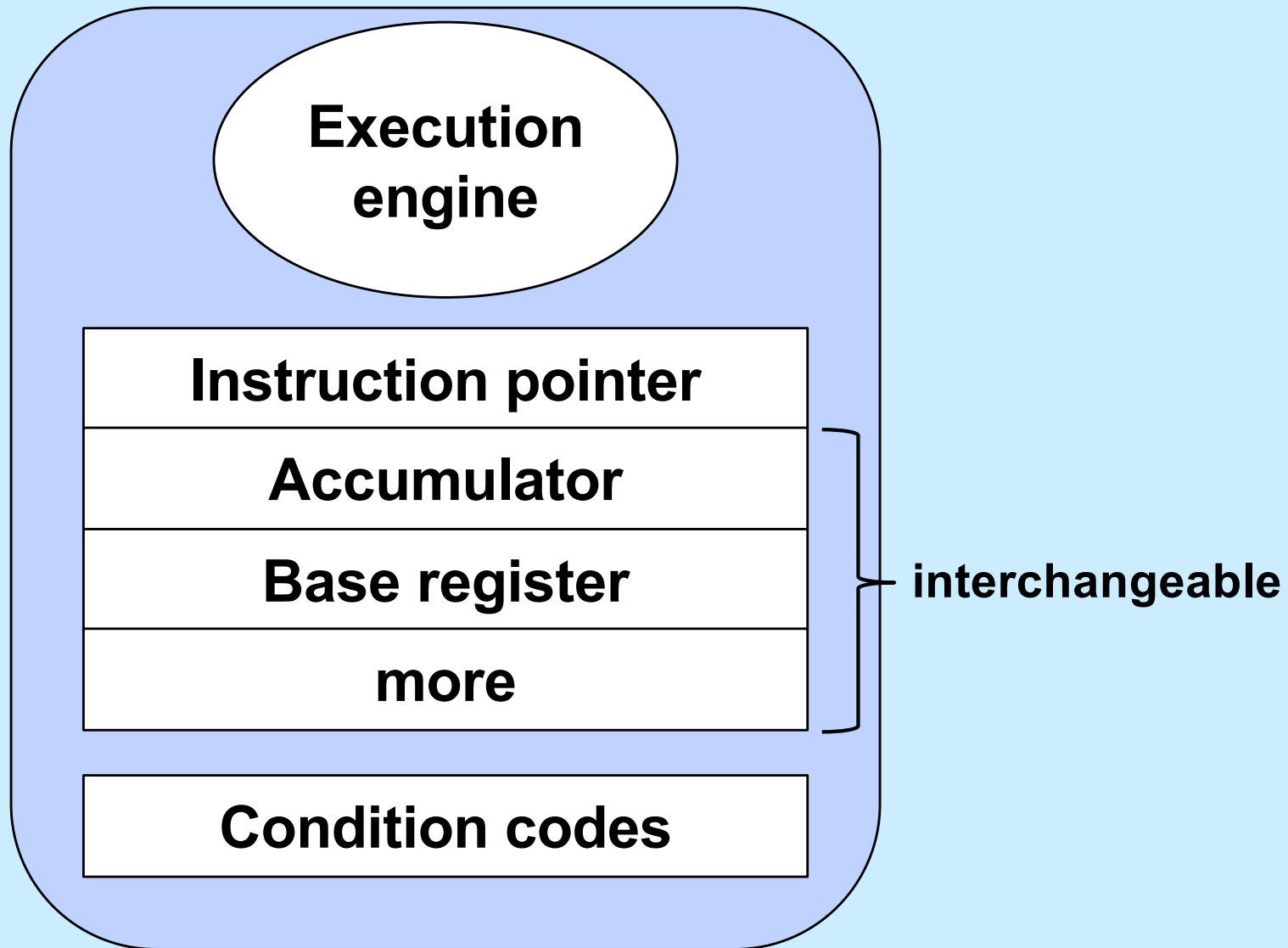
- a) 9992
- b) 9996
- c) 10,004
- d) 10,008

```
mov    1000, %acc
add    -8(%base), %acc
mul    -12(%base), %acc
mov    %acc, -16(%base)
```





# Registers



# Registers vs. Memory

