Calibre D2S R&D

C++ Internship Project

Rectangles Intersection Problem

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Submitted To:
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1) Platform used:

- Operating System ➡ **Windows 10**
- (Integrated Development Environment) IDE ➡ **Visual Studio 2017**

2) Assumptions and Important Points:

- Touching Rectangles are not considered as Overlapping Rectangles
  I'm not sure about this point but I've done some reasearch and found that there is a difference betwwen Touching Rectangles and Overlapping Rectangles.

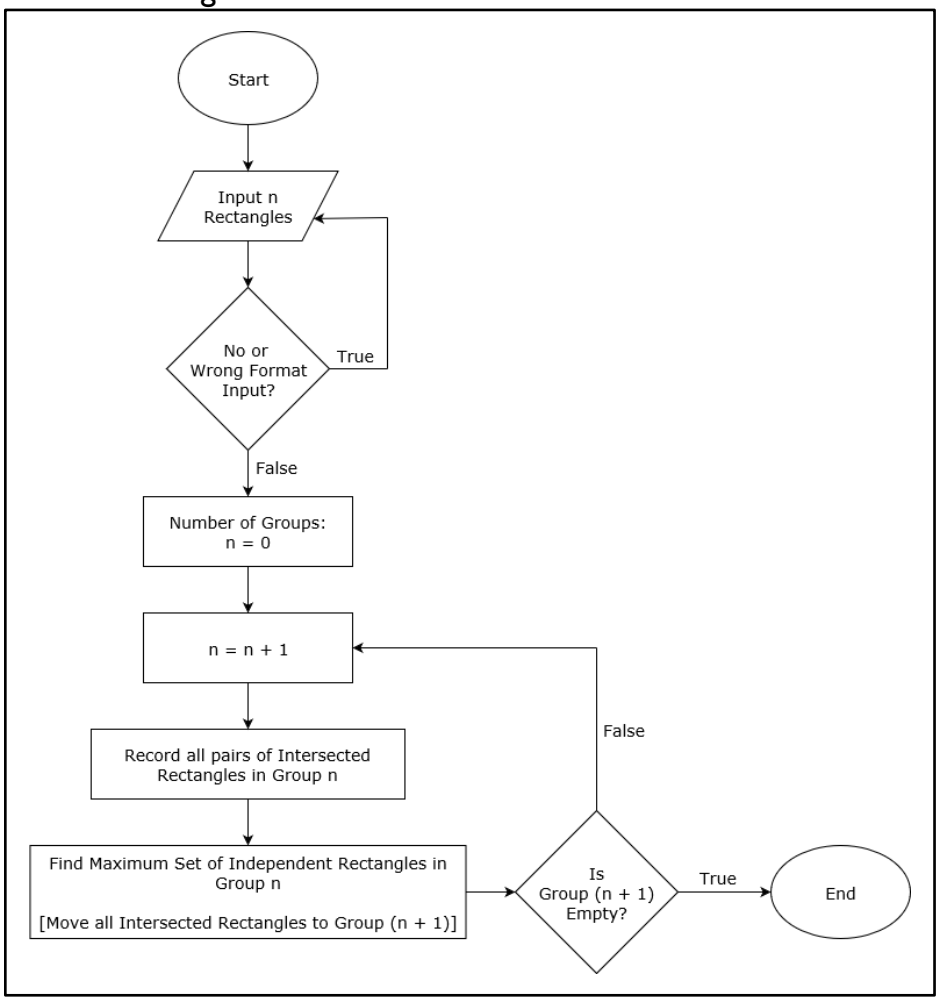
- I tried my best to run Dataset 16 but everytime I try to run I get nothing except runtime error. I was thinking about splitting the file and merging it later on but I think that this type of solution is not allowed.

- I used 2 different algorithms to solve this problem, and I compared the results.
  I used Brute Force as my first algorithm but it was not generating the minimum number of sets (except for Dataset 9 which was 12 Files). However, this algorithm is naive and time consuming. So, I thought about the Divide and Conquer algorithm. This algorithm was much better in all datasets except for dataset 9, where the output was more files (14 exactly) and it was a little bit time consuming. I think that because this problem is NP-hard so it may be because of the approximations involved in the algorithm.

3) Flowcharts of the main program and algorithms involved:

- Main Program:

❖ Flowchart Diagram:



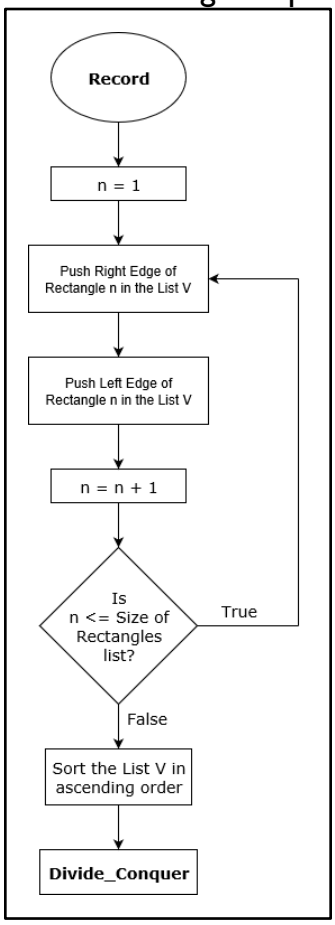
❖ Pseudocode:

```
procedure main()
    Read a file to Input n Rectangles
    Check that the file is not empty and that the format is correct
    Do {
            Record all pairs of intersected rectangles in Group n
            Find the Maximum Set of Independent rectangles in Group n
    } while (Group (n + 1) is not empty)
```

- **Algorithm** to Record all pairs of Intersected Rectangles in Group n:
  Since this problem involves massive input data, one of the best cases to solve such a problem can be achieved by using Divide and Conquer algorithm. This may to help us to finish this process in average time complexity of O(n logn) rather than $O(n^2)$ by using Brute Force.
  In this algorithm, we usually use the following 3 steps:

  a) Divide: Break the given problem into subproblems of same type.
  b) Conquer: **Recursively** solve these subproblems.
  c) Combine: Appropriately combine the answers.

❖ Flowchart Diagram: procedure **Record**(S, n):

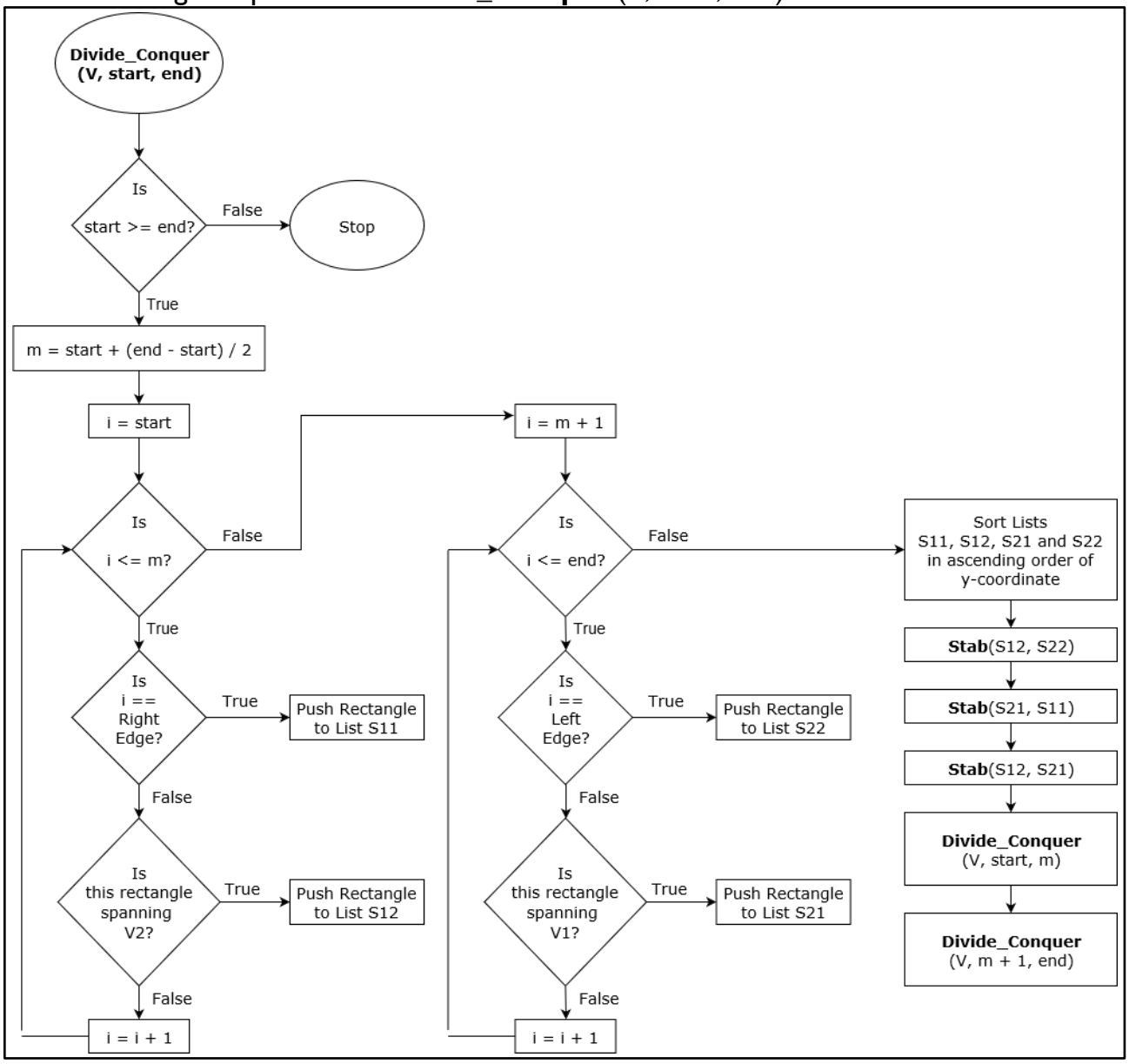


❖ Pseudocode: procedure **Record**(S, n):

Let S be the set of n iso-oriented rectangles from a specific dataset.

procedure **Record**(S, n)
  Let V be the list of x-coordinates of the 2n vertical edges in S.
  Sort V in ascending order of the x-coordinate values.
  Call procedure **Divide_Conquer**(V, 0, 2n).

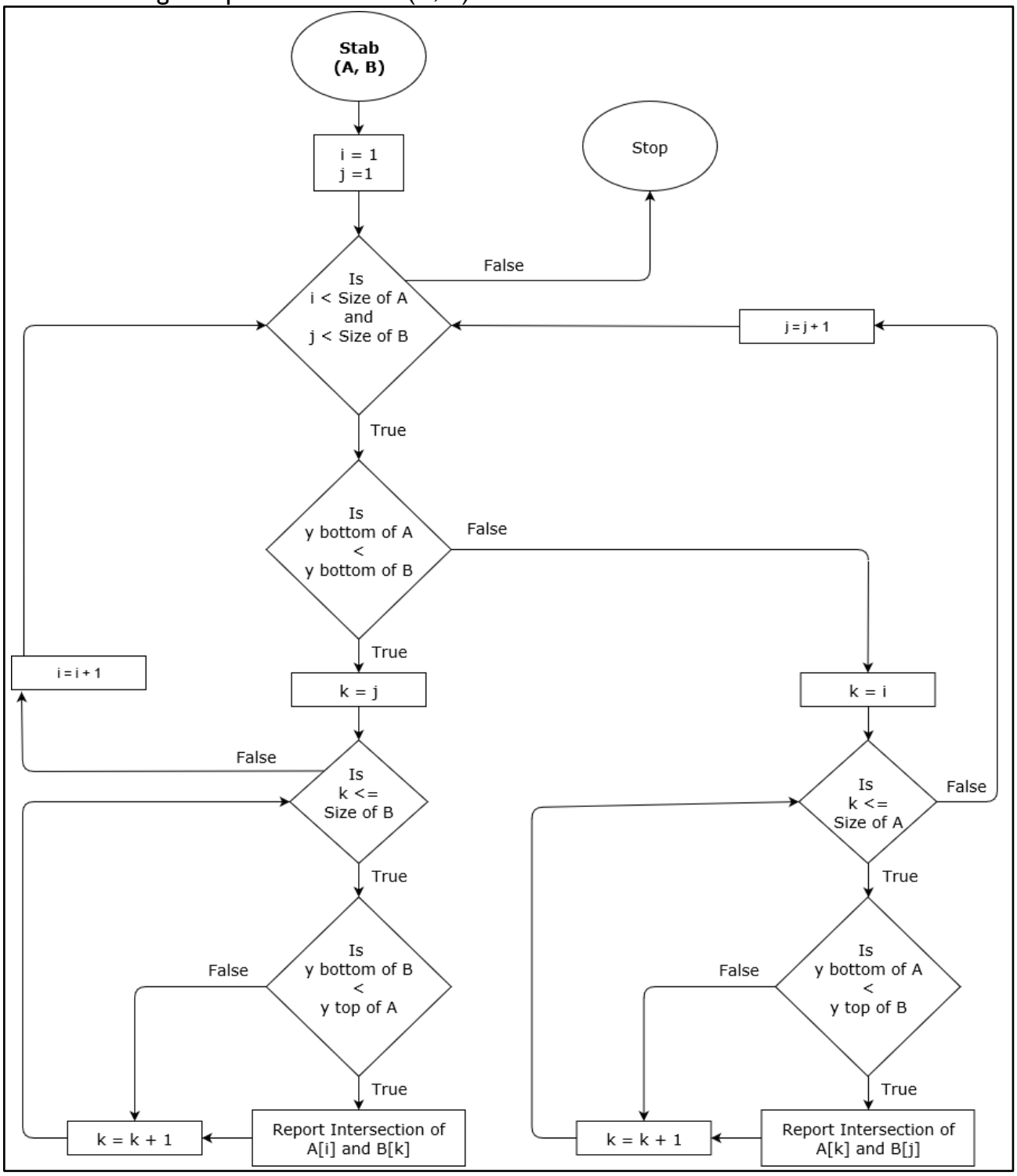
❖ Flowchart Diagram: procedure **Divide_Conquer**(V, start, end):



❖ Pseudocode: procedure **Divide_Conquer**(V, start, end):

```
procedure Divide_Conquer(V, start, end)
    if start ≥ end:
            return
    else
            m = start + (end - start) / 2;

            Let V1 be the first m elements of the set V.
            Let V2 be the rest of the elements remaining in the set V.

            Scan the list V1:
                    if the corresponding element is a Right Edge:
                            Push the corresponding rectangle to the list S11.
                    else if the corresponding element is to the right of V2:
                            Push the corresponding rectangle to the list S12.
            End Scanning V1.

            Scan the list V2:
                    if the corresponding element is a Left Edge:
                            Push the corresponding rectangle to the list S22.
                    else if the corresponding element is to the left of V1:
                            Push the corresponding rectangle to the list S21.
            End Scanning V2.

            Sort the list S11, S12, S22, and  S21 in ascending order of the bottom y-values

            Call procedure Stab(S12, S22).
            Call procedure Stab(S21, S11).
            Call procedure Stab(S12, S21).

            Call procedure Divide_Conquer(V1, start, m).
            Call procedure Divide_Conquer(V2, m + 1, end).
```

❖ Flowchart Diagram: procedure **Stab**(A, B):

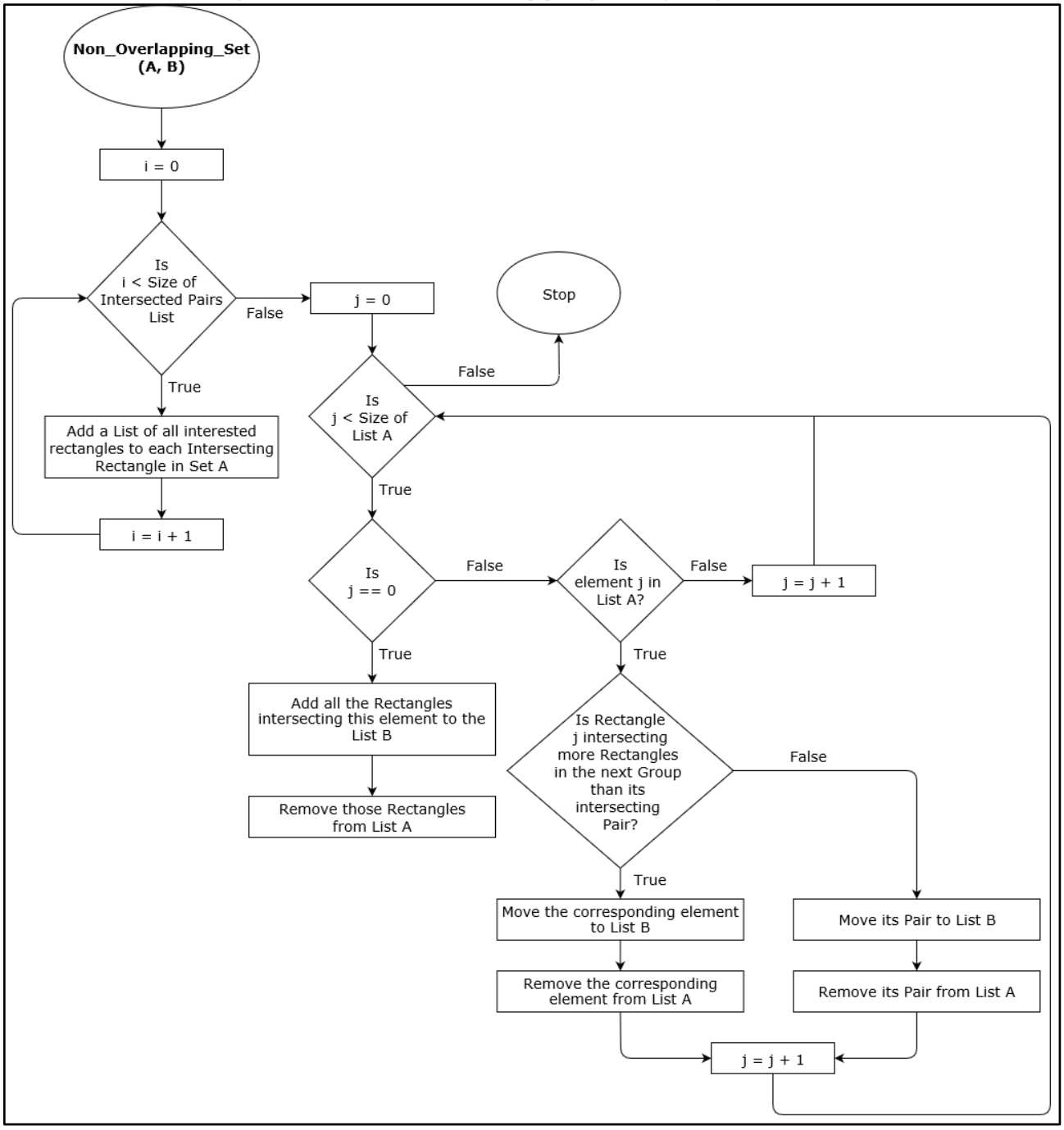


❖ Pseudocode: procedure **Stab**(A, B):

```
procedure Stab(A, B)
        i := 1 and j := 1
        while (Size of A ≥ i) and (Size of B ≥ j):
            if Bottom y-value of A[i] < Bottom y-value of B[j]:
                k := j
                while (Size of B ≥ k):
                    if Bottom y-value of B[k] < Top y-value of A[i]:
                        Report_Intersection(A[i], B[k])
                    k := k + 1
                    End if
                i := i + 1
                End while
            else
                k := i
                while (Size of A ≥ k):
                    if Bottom y-value of A[k] < Top y-value of B[j]:
                        Report_Intersection(A[k], B[j])
                    k := k + 1
                    End if
                j := j + 1
                End while

        End while
```

- **Algorithm** to Find the Maximum Set of Independent Rectangles in Group n:
  Since this is a NP-hard Problem, the approach of detecting the maximum set of non-overlapping rectangles in a single set can be base on a kind of approximation. My approximation was mainly based on picking TWO intersected rectangles in Group n and compare it to the rectangles added to the Group (n + 1) and then the rectangle with the fewer number of intersection with rectangles in Group (n + 1) should be removed from Group n and added to the Group (n + 1).

❖ Flowchart Diagram: procedure **Non_Overlapping_Set**(A, B):



❖ Pseudocode: procedure **Non_Overlapping_Set**(A, B):

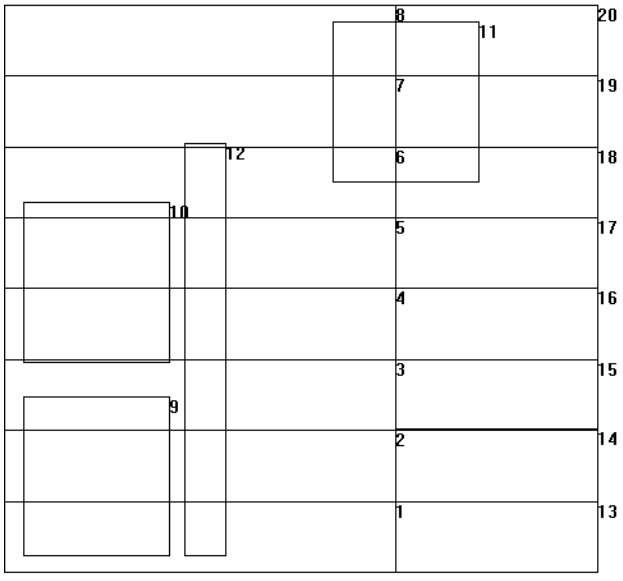
```
Let A be the set of n iso-oriented rectangles in Group n.
Let A be the set of n iso-oriented rectangles in Group (n + 1).

procedure Non_Overlapping_Set(A, B)
    Scan the list of Intersected Pairs:
        For each INTERSECTING rectangle in Set A add a list of all the Rectangles IDs
        intersecting A.
    End Scanning the list of Intersected Pairs.

    Scan the list A:
        if the corresponding element is the First Rectangle in the List:
            Add all the Rectangles intersecting this element to the List B
            Remove those Rectangles from List A
        End if.

        else if the corresponding element is in List A:
            Scan the list of all Rectangles Intersecting this element:
                if the Rectangle n in the corresponding List intersects more
                Rectangles in List B than the corresponding element:
                    Move the corresponding element to List B
                    Remove the corresponding element from List A
                End if

                else
                    Move the Rectangle n to List B
                    Remove the Rectangle n from List A
                End else
            End Scanning.
        End else.
    End Scanning A.
```

4) Table of Information:

Dataset	Number of Inputs	Number of Output Groups	Runtime (sec)	Memory usage (MB)
data_set_1.txt	5	1	0.005	0.9
data_set_2.txt	7	2	0.009	0.9
data_set_3.txt	20	2	0.017	0.9
data_set_4.txt	39	3	0.03	0.9
data_set_5.txt	77	3	0.048	0.9
data_set_6.txt	136	4	0.234	1
data_set_7.txt	216	5	0.443	1
data_set_8.txt	460	5	1.134	2
data_set_9.txt	741	14	7.523	4
data_set_10.txt	981	5	2.82	2
data_set_11.txt	5793	6	44.65	8
data_set_12.txt	6775	6	56.88	9
data_set_13.txt	7538	6	61.194	9
data_set_14.txt	8774	6	110.234	10
data_set_15.txt	9188	6	125.341	11
data_set_16.txt	25,000,000	-	-	-

5) Testing methodology:
In order to test my algorithm efficiently. I prepared a C++ Project (Including the Graphics Library ccc_win.h) in order to be able to plot the input file as well as the output files of the dataset. This will make it more clear to identify the minimum number of sets and run my program to compare it with my expectations. I provided the project folder and the source file for this program in the deliverable files. I also provided some screenshots of the program output below:

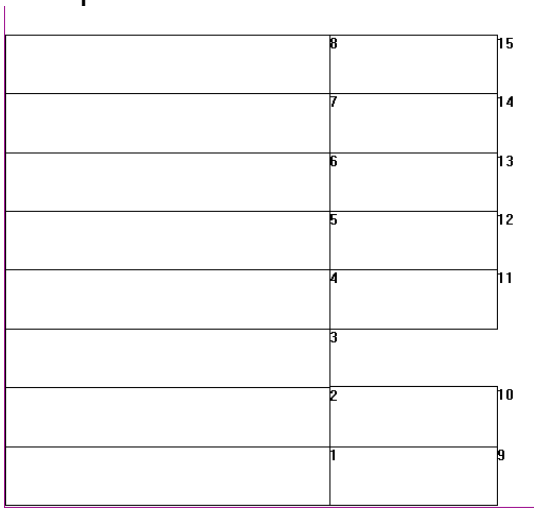
- Input plot of Dataset 3:



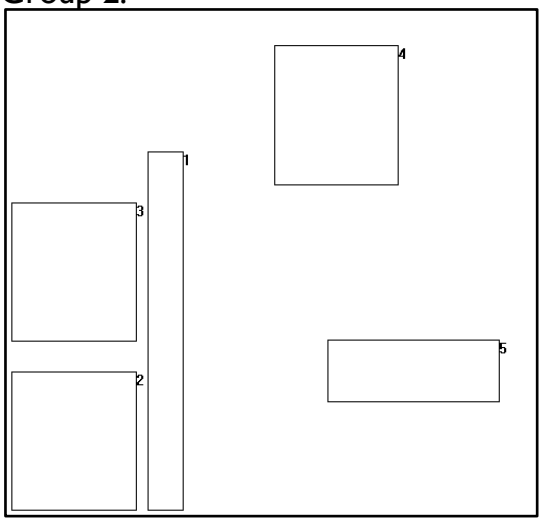
By Observing this figure, we can see the intersected rectangles (9, 10, 11, 12 and 15) and hence none of these rectangles intersect with each other. We can guess that the minimum number of non-overlapping groups will be 2. So, let's check it and plot the output files.

- Output Plots of Dataset 3:

Group 1:

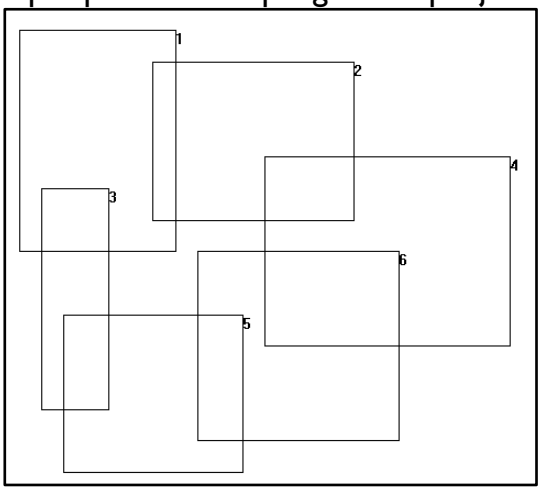


Group 2:

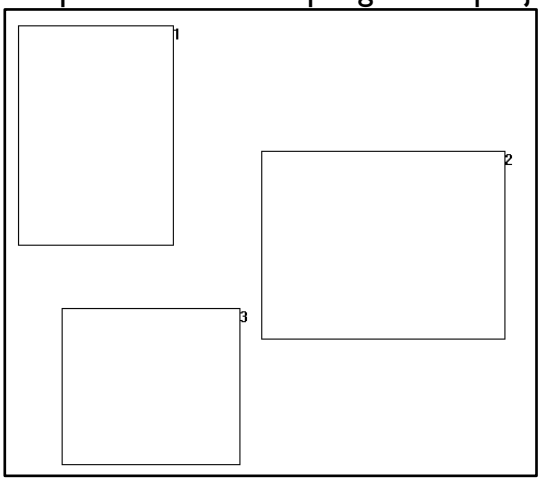


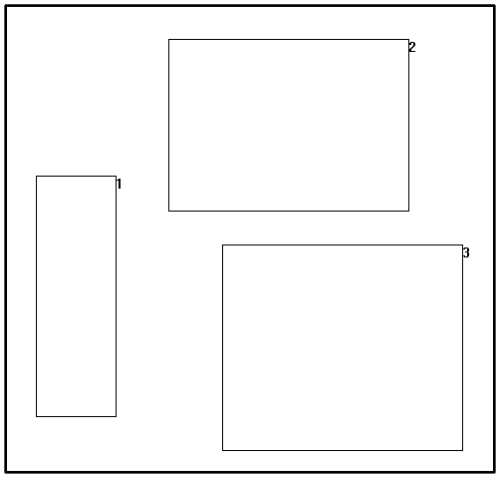
As shown, the output is as expected and thus. We tested our algorithm efficiently. Of course this will be more efficiently using 'small' dataset so that we can ensure our algorithm and implement it our larger datasets.

- Input plot of example given in project description:



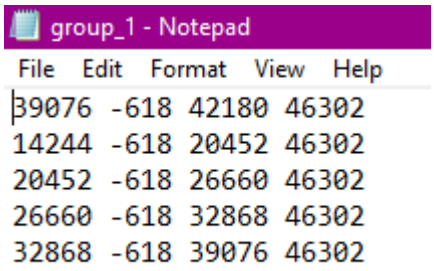
Output Plots of example given in project description:





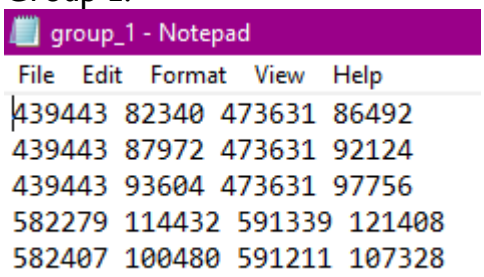
6) Screen shots for the output of your program:

- Dataset 1:

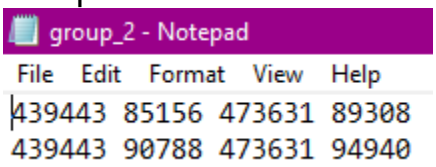


- Dataset 2:

Group 1:

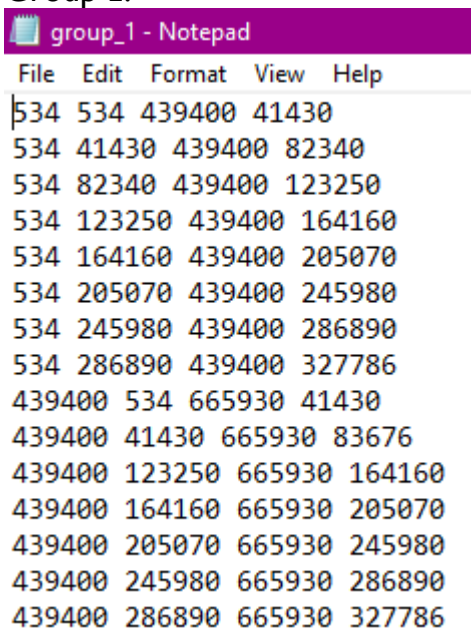


Group 2:

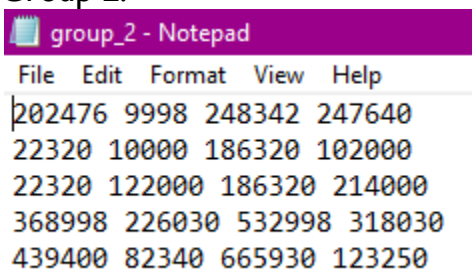


- Dataset 3:

Group 1:



Group 2:



7) Final Notes:

- I hope that you liked both of the programs and the Report.
- I read some papers in order to achieve the best algorithm for Rectangles Intersection problem.
- If you didn't like anything about my program or the report, I would love to hear your feedback so that I can improve in the long run.
- Finally, Thank you for taking the time to check my work and rate it. It was a great pleasure for me to work on this project.