transm

, X_2 , ... X_N and Y_1 , ..., Y_N : these are not itted.

 $r=r_1+r_2$ for k=1

generated by h $c_k \in \set{-1,0}$

as random seed1...N

nsional c_k is nashing r, giving 0, 1

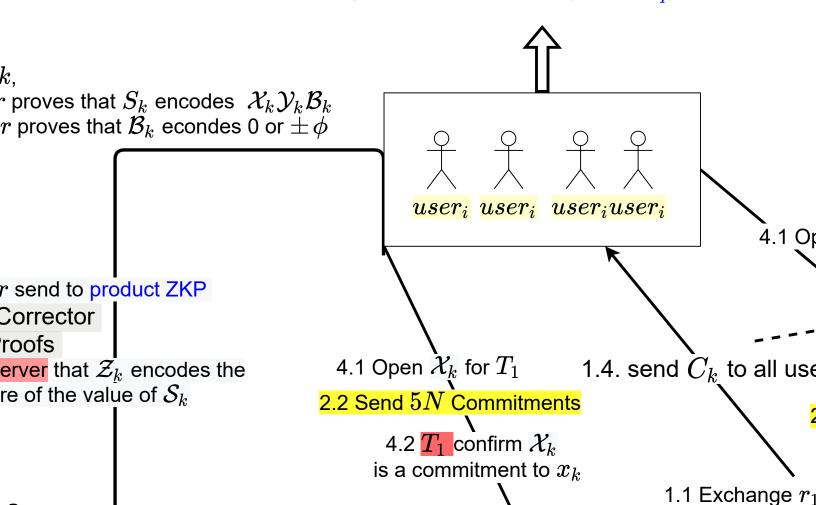
for each 5.1 *User* 5.2 *Use*

 $5.3 \ Use$

scP to s

squa

2.1 Users compute Commitments X_k , Y_k , S_K , B_K Z_K (computed over the large field \mathbb{Z}_q)



For Challenge Vectors

Each m-dimensional c_k is generated by hashing r, ightharpoonup giving $c_k \in \Set{-1,0,1}^m$

ben \mathcal{Y}_k for T_2

ers

and r_2

2.2 Send 5N Commitments

4.2 T_2 confirm \mathcal{Y}_k is a commitment to y_k

 c_k : challenge vector

N: num of challenges

talliers: 2-way setting carried out betwee

Privacy Peer[Talliers]



mdCorrector: The modular reduction corrector (the B's in the paper). They sthe commitment to 0 or +/-F.

nd



- should be



 T_1 Server

6.1 Server computes the product

$$Z = \sum_{k=1}^{N} Z_k$$

6.2 Use ZKP to ${\cal Z}$ encodes a value in the range

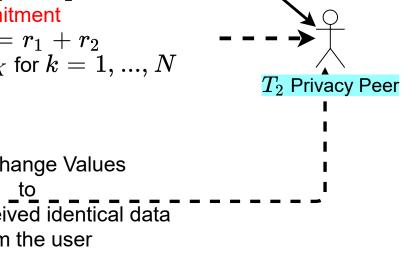
$$\left[0,N\frac{L^2}{2}\right.$$

1.2 Verify Comm

1.3 Compute r=1.4 generate C_R

check if rece fror

3 . Exc



<u>Vector Addition</u> uses <u>small-field</u> operat <u>logarithmic</u> number of crypto operations user data.

 $G,\,F$ are in general non-linear d_i are computed locally by each user





_Commitment

- ~ g: NativeBigInteger
- ~ h: NativeBigInteger
- ~ val BigInteger
- ~ r BigInteger // randomness used in the commitment
- + sanityCheck(): void
- + Commitment(g, h)
- ~ ‰computeCommitment(val, r) BigInteger
- + commit(long val) BigInteger
- + commit(BigInteger val)BigInteger
- // commit to Z_q using given randomness
- + commit(BigInteger val, BigInteger r)
- + getRandomness() BigInteger
- + getValue() BigInteger
- + verify(BigInteger c, BigInteger val, BigInteger r)// check a