Devres - Managed Device Resource

设备资源管理

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1. Intro

devres came up while trying to convert libata to use iomap. Each iomapped address should be kept and unmapped on driver detach. For example, a plain SFF ATA controller (that is, good old PCI IDE) in native mode makes use of 5 PCI BARs and all of them should be maintained.

As with many other device drivers, libata low level drivers have sufficient bugs in ->remove and ->probe failure path. Well, yes, that's probably because libata low level driver developers are lazy bunch, but aren't all low level driver developers? After spending a day fiddling with braindamaged hardware with no document or braindamaged document, if it's finally working, well, it's working.

For one reason or another, low level drivers don't receive as much attention or testing as core code, and bugs on driver detach or initialization failure don't happen often enough to be noticeable. Init failure path is worse because it's much less travelled while needs to handle multiple entry points.

So, many low level drivers end up leaking resources on driver detach and having half broken failure path implementation in ->probe() which would leak resources or even cause oops when failure occurs. iomap adds more to this mix. So do msi and msix.

2. Devres

devres is basically linked list of arbitrarily sized memory areas associated with a struct device. Each devres entry is associated with a release function. A devres can be released in several ways. No matter what, all devres entries are released on driver detach. On release, the associated release function is invoked and then the devres entry is freed.

Managed interface is created for resources commonly used by device drivers using devres. For example, coherent DMA memory is acquired using dma_alloc_coherent(). The managed version is called dmam_alloc_coherent(). It is identical to dma_alloc_coherent() except for the DMA memory allocated using it is managed and will be automatically released on driver detach. Implementation looks like the following.

```
struct dma devres
      size t
                      size:
      void
                      *vaddr;
                      dma handle;
      dma addr t
};
static void dmam coherent release (struct device *dev, void *res)
      struct dma devres *this = res;
      dma free coherent (dev, this->size, this->vaddr, this->dma handle);
dmam alloc coherent (dev, size, dma handle, gfp)
      struct dma devres *dr;
      void *vaddr:
      dr = devres alloc (dmam coherent release, sizeof (*dr), gfp);
      /* alloc DMA memory as usual */
      vaddr = dma alloc coherent(...);
      /* record size, vaddr, dma_handle in dr */
      dr->vaddr = vaddr:
      devres add(dev, dr);
      return vaddr;
```

使用devres机制,资源会在驱动detach时 自动释放 }

If a driver uses dmam_alloc_coherent(), the area is guaranteed to be freed whether initialization fails half-way or the device gets detached. If most resources are acquired using managed interface, a driver can have much simpler init and exit code. Init path basically looks like the following.

使用devm机制device manage resource 无论初始化中途失败还是驱动卸载都会正常释放资源

```
my_init_one()
       struct mydev *d:
       d = devm kzalloc(dev, sizeof(*d), GFP KERNEL);
       if (!d)
              return -ENOMEM;
                                使用devm机制,初始化失败时可直接返回错误,不用使用goto来回滚。
       d->ring = dmam alloc coherent(...);
       if (!d->ring)
              return -ENOMEM:
       if (check something)
              return -EINVAL:
       return register to upper layer(d);
And exit path,
  my remove one()
       unregister from upper layer(d);
       shutdown my hardware():
```

As shown above, low level drivers can be simplified a lot by using devres. Complexity is shifted from less maintained low level drivers to better maintained higher layer. Also, as init failure path is shared with exit path, both can get more testing.

3. Devres group

Devres entries can be grouped using devres group. When a group is released, all contained normal devres entries and properly nested groups are released. One usage is to rollback series of acquired

```
resources on failure. For example,
  if (!devres open group(dev, NULL, GFP KERNEL))
        return -ENOMEM:
  acquire A;
  if (failed)
        goto err;
  acquire B;
  if (failed)
        goto err;
  . . .
  devres remove group (dev, NULL);
  return 0:
 err:
 devres release_group(dev, NULL);
  return err code;
As resource acquisition failure usually means probe failure, constructs
like above are usually useful in midlayer driver (e.g. libata core
layer) where interface function shouldn't have side effect on failure.
For LLDs, just returning error code suffices in most cases.
Each group is identified by void *id. It can either be explicitly
specified by @id argument to devres open group() or automatically
created by passing NULL as @id as in the above example. In both
cases, devres open group() returns the group's id. The returned id
can be passed to other devres functions to select the target group.
If NULL is given to those functions, the latest open group is
selected.
For example, you can do something like the following.
  int my midlayer create something()
        if (!devres open group(dev, my midlayer create something, GFP KERNEL))
                return -ENOMEM:
        devres close group (dev, my midlayer create something);
        return 0:
```

```
void my_midlayer_destroy_something()
{
    devres_release_group(dev, my_midlayer_create_something);
}
```

4. Details

Lifetime of a devres entry begins on devres allocation and finishes when it is released or destroyed (removed and freed) - no reference counting.

devres core guarantees atomicity to all basic devres operations and has support for single-instance devres types (atomic lookup-and-add-if-not-found). Other than that, synchronizing concurrent accesses to allocated devres data is caller's responsibility. This is usually non-issue because bus ops and resource allocations already do the job.

For an example of single-instance devres type, read pcim_iomap_table() in lib/devres.c.

All devres interface functions can be called without context if the right gfp mask is given.

5. Overhead

Each devres bookkeeping info is allocated together with requested data area. With debug option turned off, bookkeeping info occupies 16 bytes on 32bit machines and 24 bytes on 64bit (three pointers rounded up to ull alignment). If singly linked list is used, it can be reduced to two pointers (8 bytes on 32bit, 16 bytes on 64bit).

Each devres group occupies 8 pointers. It can be reduced to 6 if singly linked list is used.

Memory space overhead on ahci controller with two ports is between 300 and 400 bytes on 32bit machine after naive conversion (we can certainly invest a bit more effort into libata core layer).

6. List of managed interfaces

```
CLOCK
  devm clk get()
  devm clk put()
  devm clk hw register()
  devm of clk add hw provider()
DMA
  dmaenginem async device register()
  dmam alloc coherent()
  dmam alloc attrs()
  dmam declare coherent memory()
  dmam free coherent()
  dmam pool create()
  dmam pool destroy()
GP10
  devm gpiod get()
  devm gpiod get index()
  devm gpiod get index optional()
  devm gpiod get optional()
  devm gpiod put()
  devm gpiochip add data()
  devm gpiochip remove()
  devm gpio request()
  devm gpio request one()
  devm gpio free()
110
  devm iio device alloc()
  devm iio device free()
  devm iio device register()
  devm iio device unregister()
  devm iio kfifo allocate()
  devm iio kfifo free()
  devm iio triggered buffer setup()
  devm iio triggered buffer cleanup()
  devm iio trigger alloc()
  devm iio trigger free()
  devm iio trigger register()
  devm iio trigger unregister()
  devm iio channel get()
  devm iio channel release()
  devm iio channel get all()
  devm iio channel release all()
INPUT
  devm input allocate device()
```

```
IO region
  devm release mem region()
  devm release region()
  devm release resource()
  devm request mem region()
  devm request region()
  devm request resource()
IOMAP
  devm ioport map()
  devm ioport unmap()
  devm ioremap()
  devm ioremap nocache()
  devm ioremap wc()
  devm ioremap resource(): checks resource, requests memory region, ioremaps
  devm iounmap()
  pcim iomap()
  pcim iomap regions() : do request region() and iomap() on multiple BARs
  pcim iomap table()
                        : array of mapped addresses indexed by BAR
  pcim iounmap()
IRQ
  devm free irq()
  devm_request_any_context irq()
  devm request irq()
  devm request threaded irq()
  devm irq alloc descs()
  devm_irq_alloc_desc()
  devm irq alloc desc at()
  devm irg alloc desc from()
  devm irq alloc descs from()
  devm irg alloc generic chip()
  devm irq setup generic chip()
  devm irq sim init()
LED
  devm led classdev register()
  devm led classdev unregister()
MDIO
  devm mdiobus alloc()
  devm mdiobus alloc size()
  devm mdiobus free()
MEM
  devm free pages()
```

```
devm get free pages()
  devm kasprintf()
  devm kcalloc()
  devm kfree()
  devm kmalloc()
  devm kmalloc array()
  devm kmemdup()
  devm kstrdup()
  devm kvasprintf()
  devm kzalloc()
MFD
  devm mfd add devices()
MUX
  devm mux chip alloc()
  devm mux chip register()
  devm mux control get()
PER-CPU MEM
  devm alloc percpu()
  devm free percpu()
PCI
  devm pci alloc host bridge() : managed PCI host bridge allocation
  devm pci remap cfgspace()
                                : ioremap PCI configuration space
  devm pci remap cfg resource(): ioremap PCI configuration space resource
  pcim enable device()
                                : after success, all PCI ops become managed
  pcim pin device()
                                : keep PCI device enabled after release
PHY
  devm usb get phy()
  devm usb put phy()
PINCTRL
  devm pinctrl get()
  devm pinctrl put()
  devm pinctrl register()
  devm pinctrl unregister()
POWER
  devm reboot mode register()
  devm reboot mode unregister()
PWM
  devm pwm get()
  devm pwm put()
```

```
REGULATOR
  devm_regulator_bulk_get()
  devm_regulator_get()
  devm_regulator_put()
  devm_regulator_register()
RESET
  devm_reset_control_get()
  devm_reset_controller_register()
SERDEV
  devm_serdev_device_open()
SLAVE DMA ENGINE
  devm_acpi_dma_controller_register()
SPI
  devm_spi_register_master()
WATCHDOG
  devm_watchdog_register_device()
```