C\$205
C / C++
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As we have seen last time, although most layers can be (and usually are) written in C, we'll just consider here network exchanges at the application level.

Where does C network programming hit?

Application
Host-to-Host
Internet
Network Access

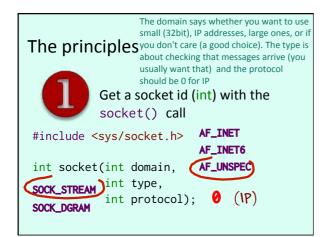
Based on a SOCKET

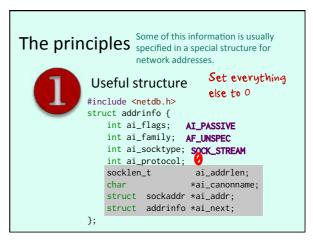
To send a message you need a "socket", which is a stream, like a file. When you ask the system for a "socket", the system returns to you an int that identifies the socket. You need to pass some information to specify your target address.

Lots of weird structures

If you are really interested in the details, there is a very good guide on network programming on the web at http://beej.us/guide/bgnet/







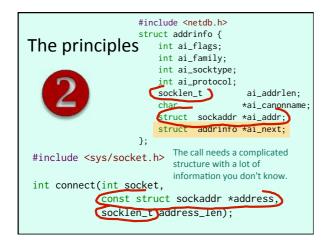
You can also optionally specify if your calls will block until something comes on the network (the default), or return an error if there is no message. Optional - specify socket behavior It can be done in different ways. #include <sys/socket.h> int socket(int domain, int type, sock_stream | sock_nonblock int protocol);

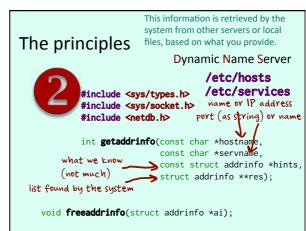
```
The principles

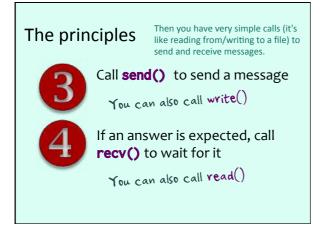
Establish a connection using the connect() call that takes the socket id and a structure with host/port

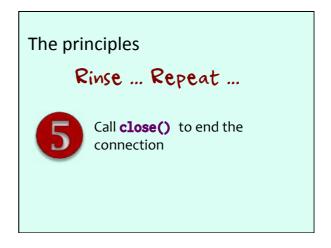
#include <sys/socket.h>

int connect(int socket, const struct sockaddr *address, socklen_t address_len);
```









What about writing a server?

What we have just seen is about writing a client program, a program that talks to a listening program, in other words a server. Writing a server program is hardly more complicated. The only difference is that if you want client programs to be able to talk to you, they must not only know on which computer your program is running, but on which port it's listening. We must decide on a port number.

Must use a fixed port

Read from configuration file!



First rule: port numbers below 1024 are a definite and resounding NO, they are reserved for well defined, often highly technical services.

Beware for other ports!

Your port number must be a value between 1025 and 65532, but check with system engineers because some ports in this range may already be in use (Database Management Systems all use port numbers over 1024). Read the port number from a configuration file, it will be easy to change.

bind()

listen()

connect()

accept()

When your server starts, it must therefore read its port number, then call the bind() function that states that it's using that port number (if another program is already using it, the call will fail). Then it must call in a loop the listen() function that does nothing but wait for an incoming connection. When a client program issues a connect() call, listen() returns out of its slumber, and the program must acknowledge the connection request by calling accept(). Then it can call recv() to read the incoming message, and reply to it with send().

Needless to say, both sides must agree on a high level protocol: what should be the format of the answer to every request?

Turning it to C++

http://vichargrave.com/network-programming-design-patterns-in-c/

TCPStream

TCPConnector

TCPAcceptor (for writing a server)

This page explains in detail and provides interesting C++ classes to ease programming.

WARNING

The code is ALMOST complete (small gaps to fill)

You should use send() and recv() rather than write() and read() as suggested in the post

There are some errors in the test application

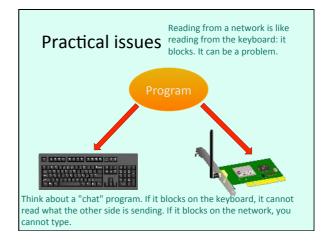
(confuses NULL and \0)

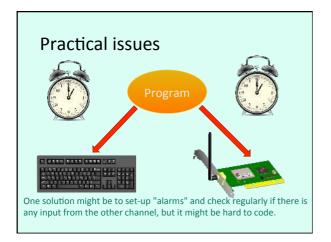
```
Many interesting things in this class:
class TCPStream {
              m_sd;
                              you have friend classes. Like friend
  std::string m_peerIP;
                              functions, all their methods can access
  int
              m_peerPort;
                              what is private in this class.
                              Constructors, in particular, are private,
  friend class TCPAcceptor; which is rather unusal. You'll never
  friend class TCPConnector; create a TCPStream directly, but used
                             one created by a TCPAcceptor or
  ~TCPStream();
                              TCPConnector object.
  ssize_t send(const char* buffer, size_t len);
  ssize_t receive(char* buffer, size_t len, int timeout);
  std::string getPeerIP();
              getPeerPort();
  int
  TCPStream(int sd, struct sockaddr_in* address);
  TCPStream();
  TCPStream(const TCPStream& stream);
```

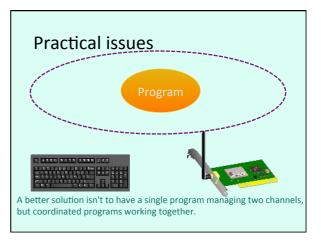
```
class TCPConnector {
    public:
        TCPStream* connect(const char* server, int port);
    private:
        int resolveHostName(const char* host, struct in_addr* addr);
};

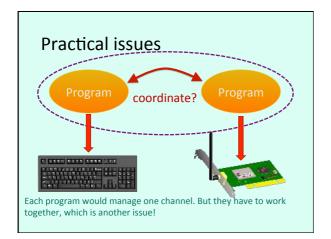
It's the connect() method of a TCPConnector object (note: no coinstructor, it uses the default C++ constructor) that instantiates (creates) a TCPStream object and returns to you a pointer to this object.

TCPAcceptor not shown here
```





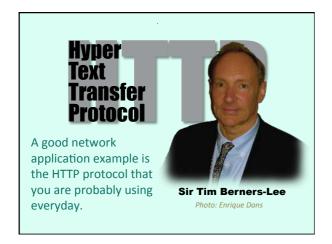


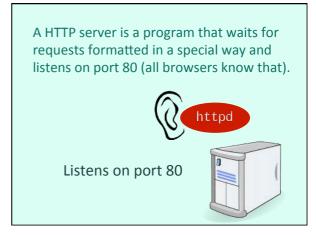


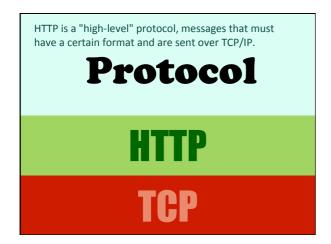
Practical issues

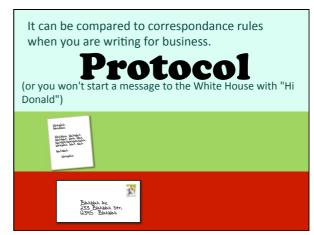
Multithreading

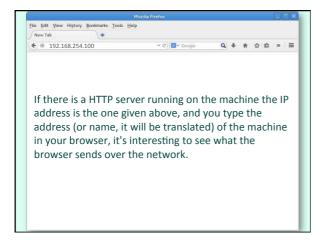
This is an excellent case where having what is called "multiple threads" (parts of the program running at the same time and independently) would be an interesting solution. We'll soon see related topic, and perhaps we'll have time to talk a little about miltithreading proper.

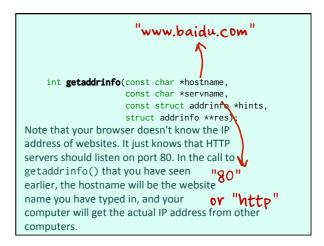












GET / HTTP/1.1

Host: 192.168.254.100

User-Agent: Mozilla/5.0 (Windows NT 5.1; rv:2.0.1) [...]

Accept: text/html,application/xhtml+xml,application/xml;[...]

Accept-Language: en-gb,en;q=0.5

Accept-Encoding: gzip, deflate

Accept-Charset: ISO-8859-1,utf-8;q=0.7

Keep-Alive: 115

Connection: keep-alive
Cache-Control: max-age=0

[empty line]

Cache-Control: max-age=0

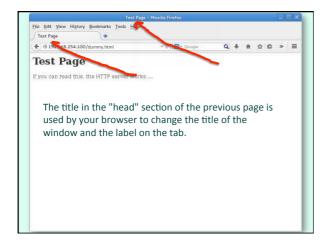
[empty line]

A message starts with a header

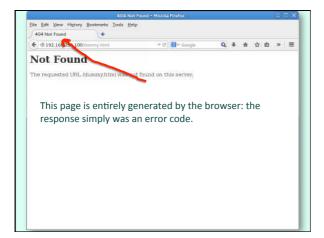
(sometimes it's just a header) always terminated by an empty line. The first and second line are mandatory. The first one says which page should be returned (/ = homepage) and which is the HTTP version in use.

The second one repeats the name of the target (here, on a local network).

HTTP/1.1 200 OK A reply from a webserver Server: Apache/2/2/15 (Linux/SUSE) also follows a protocol. The Last-Modified: [date] first line specifies the ETag: "2d7d-b7-4a3c52ef29046" protocol and provides a Accept-Ranges: bytes numerical code that gives the Content-Length: 175
Keep-Alive: timeout=15, max=100 outcome of the previous Connection: Keep-Alive request, followed by the Content-Type: text/html meaning in clear, then the <html> header contains various <head: <title>Test Page</title> information, and the actual </head>
<body> page displayed follows the empty line. <h1>Test Page</h1> If you can read this, the HTTP server works ... </body> </html>







Minimum HTTP message

"GET <page> HTTP/1.1\nHOST: <host>\n\n"

Many websites will drop the connection if there is no "User-Agent:" line, though (they also expect good manners). If you are interested in testing a program that sends HTTP requests and gets the answer, though, you can try this site, that was created for testing:

"GET /html HTTP/1.1\nHOST: httpbin.org\n\n"

Networking classes in action

I've written a small (obviously not complete) program that uses Vic Hargrave's classes to get HTML pages using the HTTP protocol, just to show how easy it becomes when the complicated code is wrapped in easy-to-use objects and methods.

I have created a HTTPCnx class that uses a TCPConnector and the TCPStream returned by the connect() method of the TCPConnector. The constructor only needs teh host name, as the prort is known.

class HTTPCnx {
 std::string host;
 TCPConnector* connector;
 TCPStream* stream;

HTTPCnx(const char *host);

std::string get(const char *page);

~HTTPCnx();

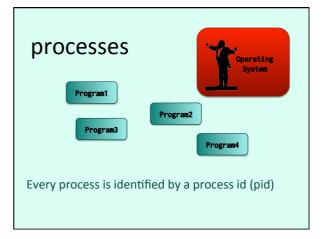
};

```
int main(int argc, char **argv) {
    if ((argc < 2) || (argc > 3)) {
        cerr << "Usage : " << argv[0] << " website [<page>]\n";
        return 1;
    }

HTTPCnx *http = new HTTPCnx(argv[1]);
    if (http) {
        char line[BUFFER_SZ];
        int len;
        int code;
        string response = "";
        response = http->get((argc == 3 ? argv[2] : "/"));
        if (response.length()) {
            cout << response << endl;
        } else {
            cerr << "No response" << endl;
        } else {
            perror("Connection");
            cerr << "(httptest) Failed to connect\n";
        }
        return 0;
}</pre>
```

Back to System Calls and processes

In today's complex software architectures, programs are permanently controlled, sometimes automatically restarted, or request flows diverted elsewhere to ensure service continuity. We are going to see a few useful functions for monitoring programs.



\$ ps
PID TTY TIME CMD
1800 ttys000 0:00.09 -bash
826 ttys001 0:00.34 -bash
\$
If you run the "ps" command on a Unix-like operating

If you run the "ps" command on a Unix-like operating system, it will list you the running processes (TTY means "terminal"; I had two open console windows, "bash" is the program that accepts commands in them)

A process has easy access to two pids:

- Its own

- Its parent's

#include <unistd.h>

pid_t getpid(void); pid of the current process

pid_t getppid(void); pid of the parent process

A "pid_t" is a typedef'd integer

Important

Every process except process 1 is created by another process!

All processes have a parent process. We'll see this more in detail later. Nothing is born of nothing.

```
$ ps -o pid,ppid,time,comm
PID PPID TIME COMM
1800 1799 0:00.09 -bash
826 825 0:00.35 -bash
```

The "ps" command takes a lot of options. A particularly useful one is -0, followed by a comma separated list of things you want to see, here process id, parent process id, the time it has been running and the name of the command.

An easy way to create a subprocess:

You may already know system(). This function creates a subprocess that runs the command passed as a parameter.

#include <stdlib.h>

int system(const char *command);

Waits for command completion

Returns the return code of the command

```
#include <stdio.h>
#include <stdlib.h>

int main() {
    system("ps -o pid,ppid,time,comm");
    return 0;
}
A simple example that runs the previous command. Such a
"wrapper" program can sometimes be useful; for instance,
you can take advantage of conditional compiling for catering
for subtle syntax variations between systems. Sometimes it's
easier to do it that way than try to write a shell script that
runs everywhere.
```

```
$ ./run_ps
  PID
        PPID
                       TIME COMM
         1799
 1800
                   0:00.09 -bash
  (826)
           825
                   0:00.44 -bash
 2223
           (826)
                   0:00.00 ./run_ps
$
Here is the output of the previous command, which
shows the run_ps program as a child process of the
"bash" that runs in one of the consoles.
```

But how can two programs talk to each other? We can contemplate various options. How can processes communicate with each other? OK if sequential operations FILE Files are easy, Program1 Program2 synchronization isn't: the reader must make sure that the writer is done with writing. Socket for different hosts Sockets are a bit of an overkill on a single computer, "listening" and assume that someone is listening.

If you want to grab the attention of a program on the same computer that isn't listening, the best is to send it a signal.

Signals

#include <signal.h>

int kill(pid_t pid, int sig);
The "kill" name only refers to ONE usage of signals.

> 0 specific process

0 processes in the same group

-1 other processes of the same user

```
There may be some slight
$ kill -1 differences between systems
           2) SIGINT
1) SIGHUP
                          3) SIGQUIT 4) SIGILL
5) SIGTRAP 6) SIGABRT
                          SIGEMT
                                    8) SIGFPE
9) SIGKILL 10) SIGBUS
                         11) SIGSEGV 12) SIGSYS
13) SIGPIPE 14) SIGALRM
                         15) SIGTERM 16) SIGURG
17) SIGSTOP 18) SIGTSTP
                         19) SIGCONT 20) SIGCHLD
21) SIGTTIN 22) SIGTTOU
                         23) SIGIO 24) SIGXCPU
25) SIGXFSZ 26) SIGVTALRM 27) SIGPROF 28) SIGWINCH
                         31) SIGUSR2
29) SIGINFO 30) SIGUSR1
The "kill -l" command lists all available signals.
The default behavior when a program receives a signal
depends on the signal.
```

```
$ kill -1
 1) SIGHUP 2) SIGINT
                                      3) SIGQUIT 4) SIGILL
 5) SIGTRAP 6) SIGABRT
                                 7) SIGEMT 8) SIGFPE
 9) SIGKILL 10) SIGBUS
                                    11) SIGSEGV 12) SIGSYS
                                  15) SIGTERM 16) SIGURG
13) SIGPIPE 14) SIGALRM
                                  19) SIGCONT 20) SIGCHLD
17) SIGSTOP 18) SIGTSTP

      21)
      SIGTTIN
      22)
      SIGTTOU
      23)
      SIGIO
      24)
      SIGXCPU

      25)
      SIGXFSZ
      26)
      SIGVTALRM
      27)
      SIGPROF
      28)
      SIGWINCH

29) SIGINFO 30) SIGUSR1
                                    31) SIGUSR2
$
              Ignored by default
```

Of those, SIGCHLD, which means that the status of a child process has changed, can be useful for monitoring.

```
$ kill -1

1) SIGHUP 2) SIGINT 3) SIGQUIT 4) SIGILL
5) SIGTRAP 6) SIGABRT 7) SIGEMT 8) SIGFPE
9) SIGKILL 10) SIGBUS 11) SIGSEGV 12) SIGSYS
13) SIGPIPE 14) SIGALRM 15) SIGTERM 16) SIGURG
17) SIGSTOP 18) SIGTSTP 19) SIGCONT 20) SIGCHLD
21) SIGTTIN 22) SIGTTOU 23) SIGIO 24) SIGXCPU
25) SIGXFSZ 26) SIGVTALRM 27) SIGPROF 28) SIGWINCH
29) SIGINFO 30) SIGUSR1 31) SIGUSR2

$

Stop (suspend) process
You can restart the process by sending SIGCONT to it.
```

```
$ kill -1
1) SIGHUP 2) SIGINT
                          3) SIGQUIT 4) SIGILL
5) SIGTRAP 6) SIGABRT
                          7) SIGEMT 8) SIGFPE
9) SIGKILL 10) SIGBUS
                          11) SIGSEGV 12) SIGSYS
13) SIGPIPE 14) SIGALRM
                        15) SIGTERM 16) SIGURG
                          19) SIGCONT 20) SIGCHLD
17) SIGSTOP 18) SIGTSTP
21) SIGTTIN 22) SIGTTOU
                          23) SIGIO 24) SIGXCPU
25) SIGXFSZ 26) SIGVTALRM 27) SIGPROF 28) SIGWINCH
29) SIGINFO 30) SIGUSR1
                          31) SIGUSR2
          Terminate process
SIGINT is what Ctrl-C delivers, SIGSEGV is a segmentation
violation (dangling pointer), a zero-divide causes SIGFPE
```

(floating-point exception), and so forth.

```
$ kill -l
 1) SIGHUP 2) SIGINT
                             3) SIGQUIT 4) SIGILL
5) SIGTRAP 6) SIGABRT
                           7) SIGEMT 8) SIGFPE
9) SIGKILL 10) SIGBUS
                            11) SIGSEGV 12) SIGSYS
13) SIGPIPE 14) SIGALRM 15) SIGTERM 16) SIGURG
17) SIGSTOP 18) SIGTSTP
                            19) SIGCONT 20) SIGCHLE
21) SIGTTIN 22) SIGTTOU
                           23) SIGIO 24) SIGXCPU
25) SIGXFSZ 26) SIGVTALRM 27) SIGPROF 28) SIGWINCH 29) SIGINFO 30) SIGUSR1 31) SIGUSR2
         ... with a core dump
A core dump is a file that contains an image of memory
when the program crashes, to help with debugging. This files
are very big and their generation is often disabled by system
administrators.
```

Special signal:

0 = test if process is alive

(kill() returns 0 if no such process)

Finally, signal 0 is a dummy signal which isn't delivered to the "target", but allows the sender to know whether the process corresponding to the pid is up and running. This is very much used by monitoring system that check the good health of a system, and generate alerts when a critical process is down.

Almost all signals can be caught and handled!

"signal handler"

void handler(int sig)

typedef'd as sig_t

Beware:

In some systems, handlers are automatically deactivated after being called.

Must reset themselves.

```
SIGKILL and SIGSTOP cannot be caught or ignored

The easy way to trap a signal is by associating a handler to a signal with function signal(), which isn't considered a "system call"

#include <signal.h>

Sig_t signal(int sig, sig_t func);

SIG_IGN lanore

SIG_DFL behavior void handler(int sig)

Previous settings
```

The hard way to trap a signal is calling the sigaction() function which, also declared in the same signal.h as signal(), is documented in section 2 (system calls) of the manual.

sigaction() allows for finer handling of signals than signal() can. For instance, nothing prevents with signal() a handler from being itself interrupted while processing a signal. Function sigaction() allows to mask interrupts and work uninterrupted when needed, and so forth

```
#include <stdio.h>
#include <unistd.h>
                                       beeper.c
#include <time.h>
                                 We can easily illustrate
#define SLEEP_TIME 10
                                 signals with a simple
                                 program that runs an
int main() {
   int nap = SLEEP_TIME; infinite loop, and displays the time_t now; infinite loop, and displays
    struct tm *t;
                                 10 seconds.
                                 We can stop this program
   while (1) {
   now = time(NULL);
                                 with Ctrl-C.
        t = localtime(&now);
       sleep(nap);
    return 0;
```

```
If we add this function to the program:

void handler(int sig) {
    printf("Tada!\n");
}

then define it as the SIGINT handler before the start of the loop:

(void)signal(SIGINT, handler);

Ctrl-C will no longer stop the program, but display Tada!

If we want to stop the program, we must open another window, run "ps" to find out the process id of the beeper, and issue something such as

$ kill -SIGKILL <pid> ("kill pid" sends SIGTERM)
```