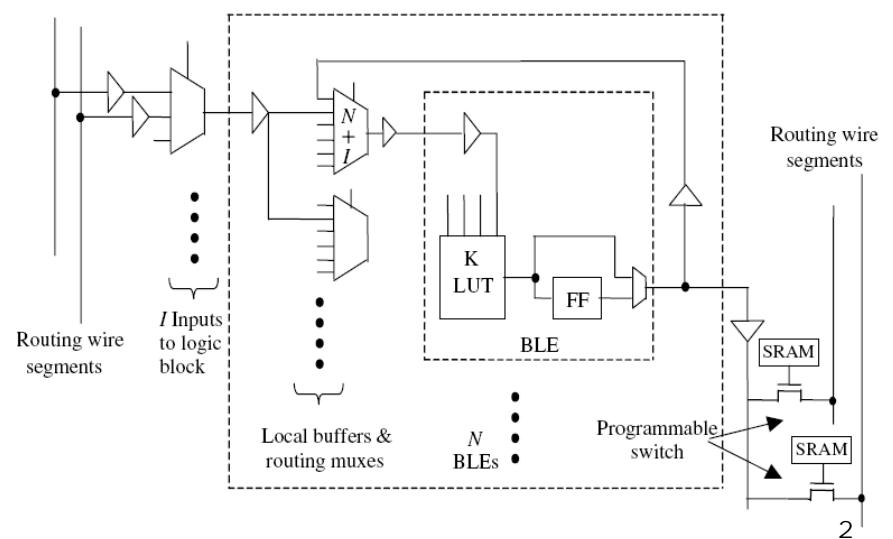


Circuit Design and Architecture Exploration of FPGAs

1

Topics

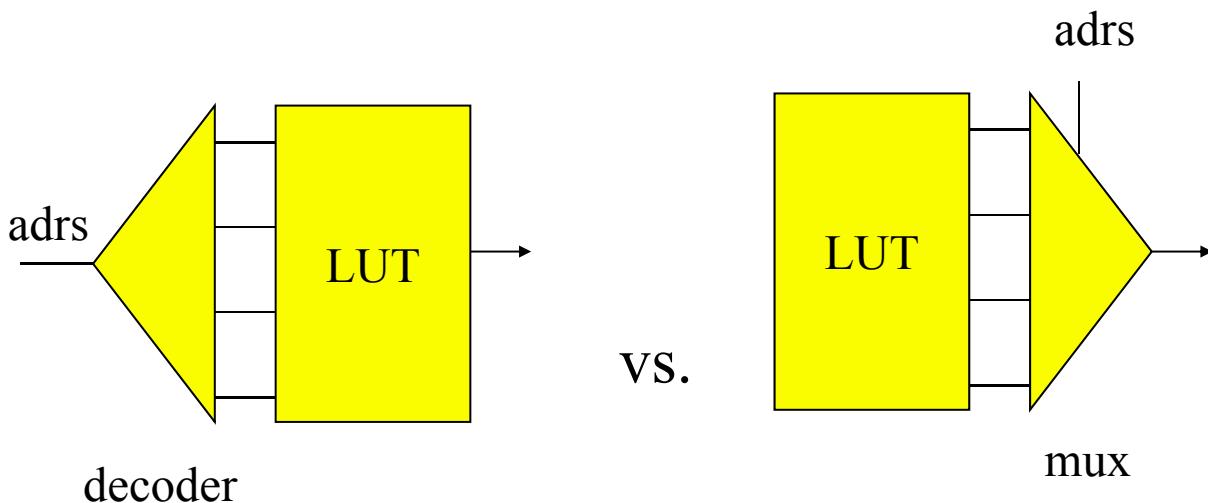
- Circuit Design for FPGAs
- Logic Block Architecture Study
- Routing Architecture Study



2

Lookup Table Circuitry

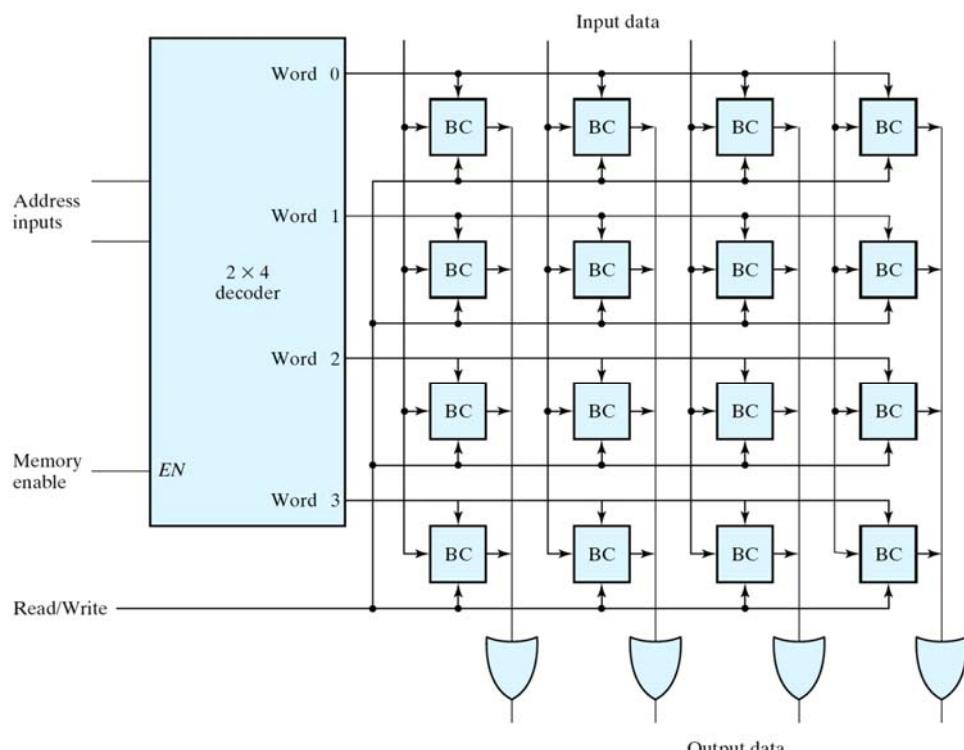
- 2 options for bit-selection
 - decoder or multiplexer?



3

Bit-Selection in Conventional RAM/ROM

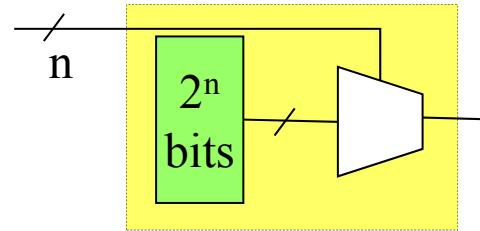
- RAM/ROM typically uses decoder for bit-selection



4

Bit-Selection for LUT

- FPGA's LUT uses a multiplexer for bit-selection.

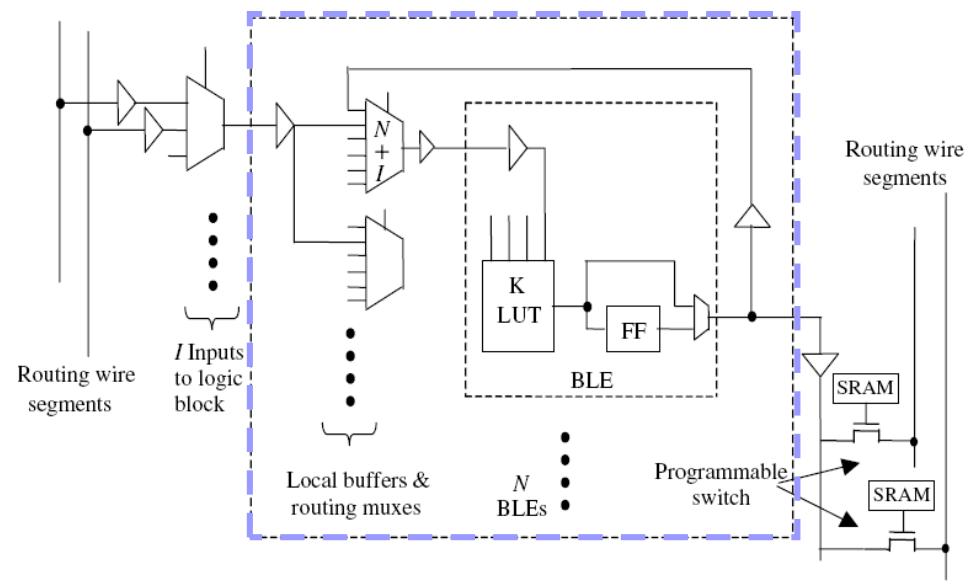


- Multiplexer presents smaller load to memory cells.
 - Allows smaller memory cells

5

Multiplexers in FPGA

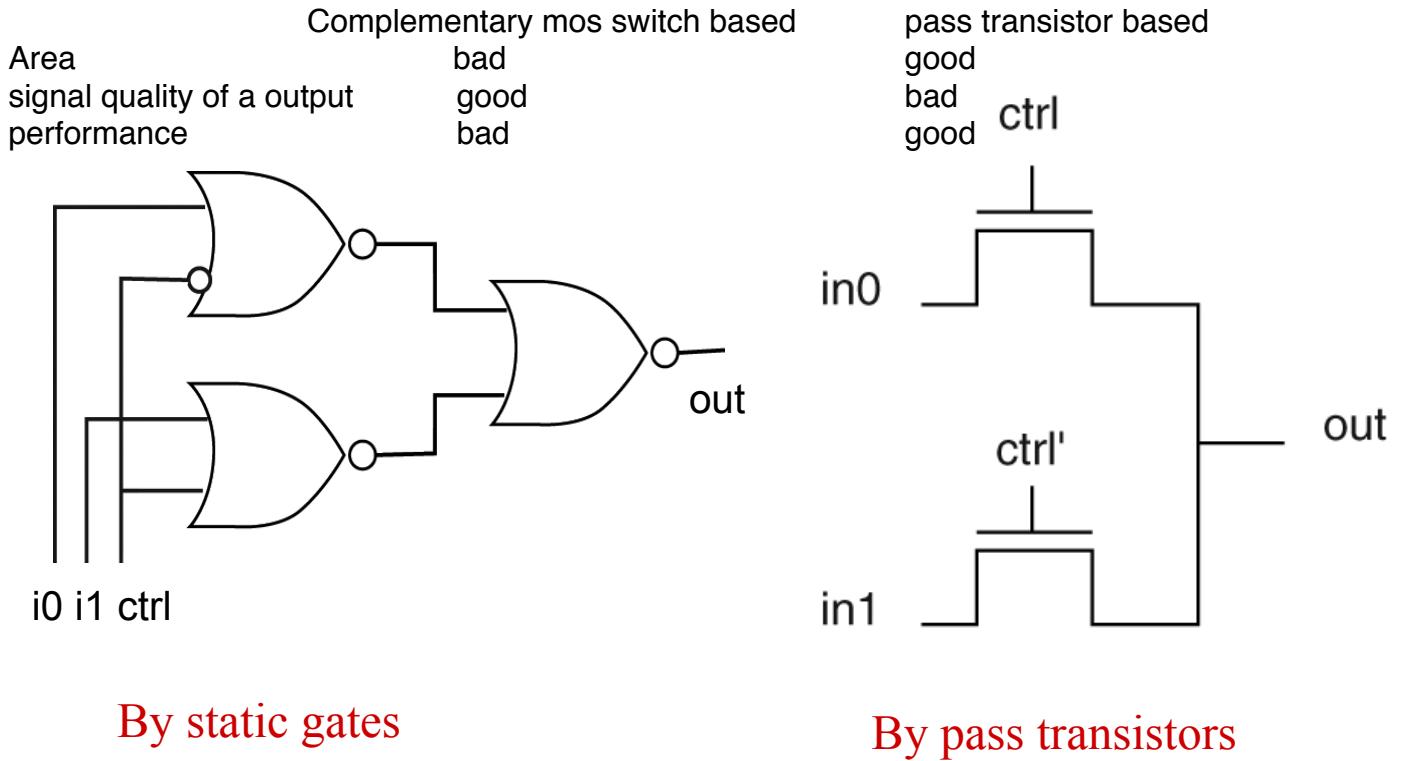
- Muxes used inside LUTs and for routing (intra-block and inter-block).



A logic block and its periphery

6

Multiplexer Design



7

Multiplexer Design

- Pass transistor multiplexer uses fewer transistors than fully complementary gates.
- Pass transistor is faster than complementary switch:
 - Equal-strength p-type is 2.5X n-type width.
 - Total resistance is 0.5X, total capacitance is 3.5X.
 - RC delay is $0.5 \times 3.5 = 1.75$ times n-type switch.

nmos
Good at pass 0
Poor at pass 1

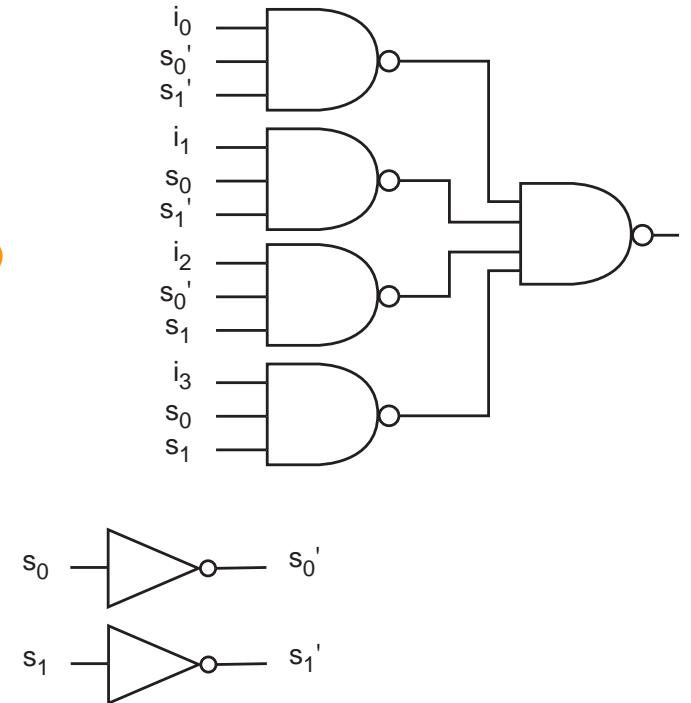
pmos
Good at pass 1
Poor at pass 0

cmos
結合兩個的優點

Performance of Static Gate MUX

- Delay through n -input NAND is $(n+2)/3$ using logical effort computation.
- For b -to-1 MUX
 - $\lg b + 1$ inputs at first level, so delay is $(\lg b + 3)/3$.
 - Delay at second level is $(b+2)/3$.
- Delay grows as b .

$O(b)$

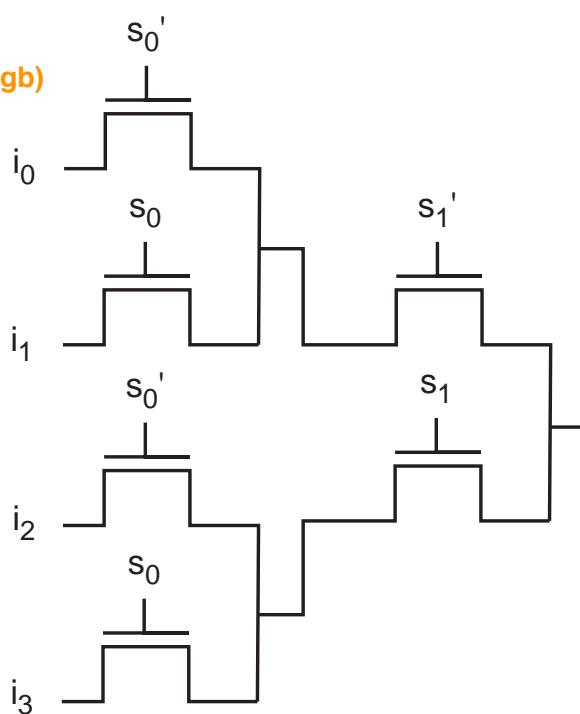


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Performance of Tree-based Pass Transistor MUX

- Delay proportional to square of path length.
- Delay grows as $(\lg b)^2$.

$O(\lg b)$

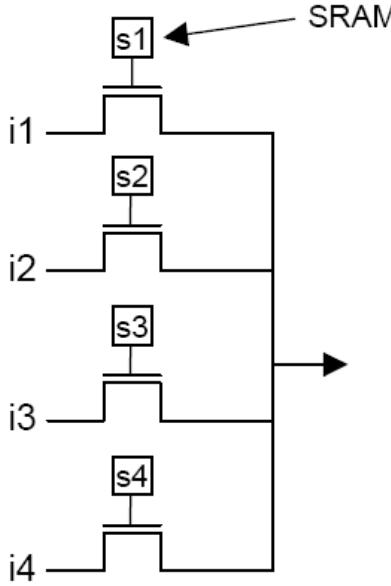


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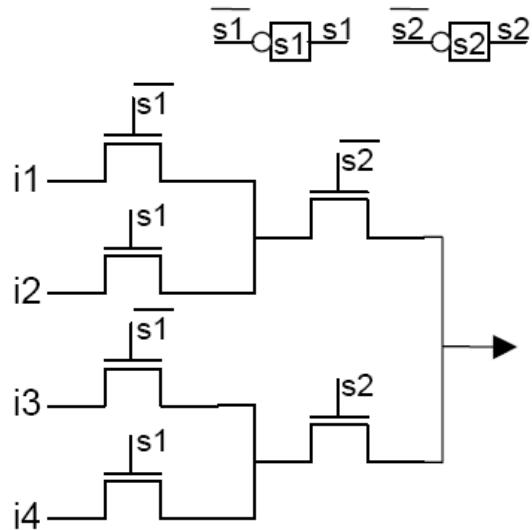
Encoded MUX vs Decoded MUX

4 to 1 mux

- Tradeoff between transistor count and delay



a) decoded multiplexer



b) encoded multiplexer

一般的SRAM 5 transistors

能夠吐出Q和!Q的SRAM 6 transistors

$4 \text{ SRAM} + 4 \text{ transistors} = 4 \times 5 + 4 = 24 \text{ transistors}$

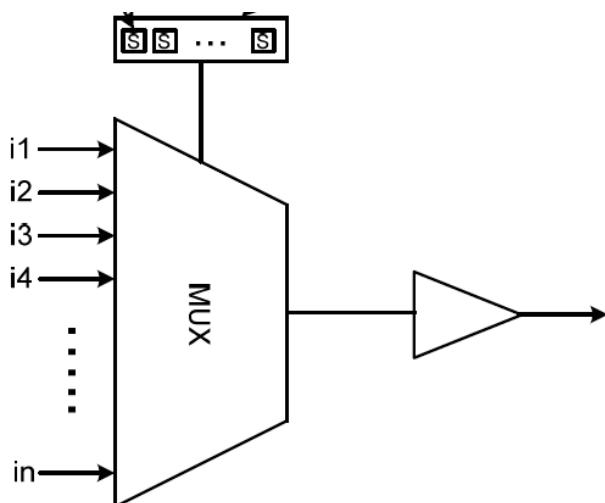
$2 \text{ SRAM} + 6 \text{ transistors} = 2 \times 6 + 6 = 18 \text{ transistors}$

11

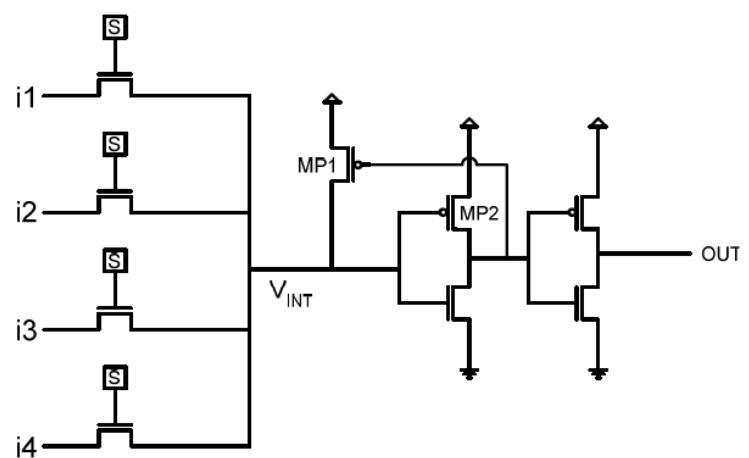
Leakage in MUX-based Routing Switch

- Level restoring buffer to avoid leakage at MP2 due to a weak V_{INT}

是用nmos傳的如果傳1 V_{INT} 會有個很爛的1 可能會出問題
加MP1 來解決這個問題



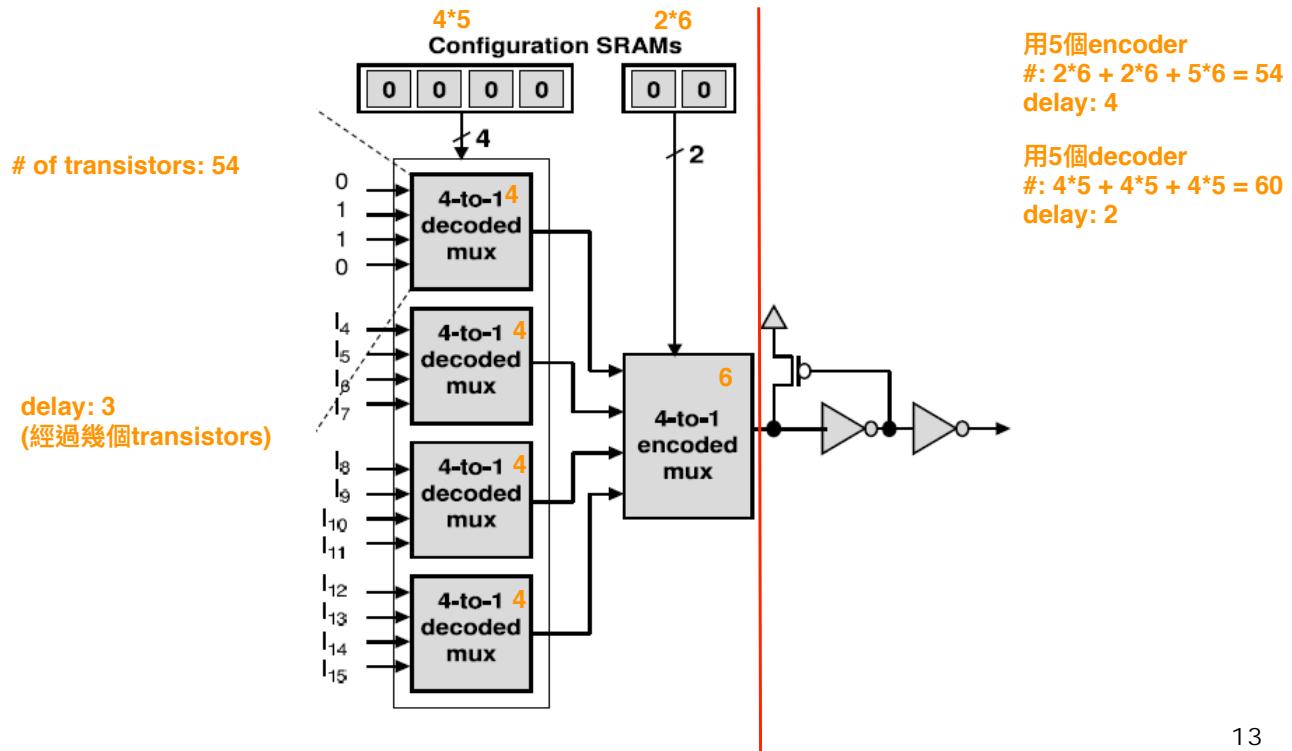
a) Routing switch (abstract)



b) 4-input routing switch (transistor-level view)

MUX-based Routing Switch Design

- Optimize: transistor count, delay, leakage



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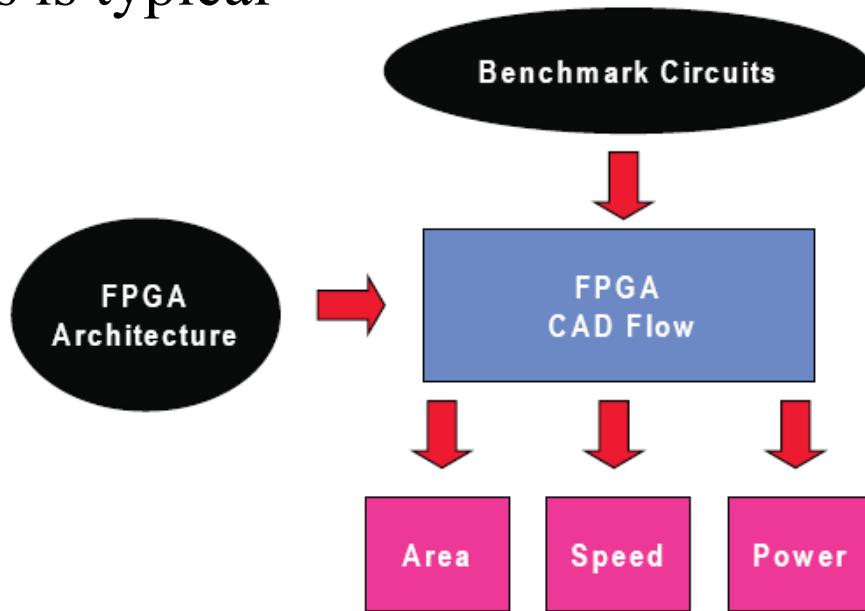
Architectural Issues

- Granularity of logic elements in the FPGA?
- LE structure:
 - What functions?
 - How many inputs?
 - Dedicated logic?
- What types of interconnect?
 - How much of each type?
- How long should interconnect segments be?
- How should we vary interconnect?
 - Uniform or non-uniform over chip?

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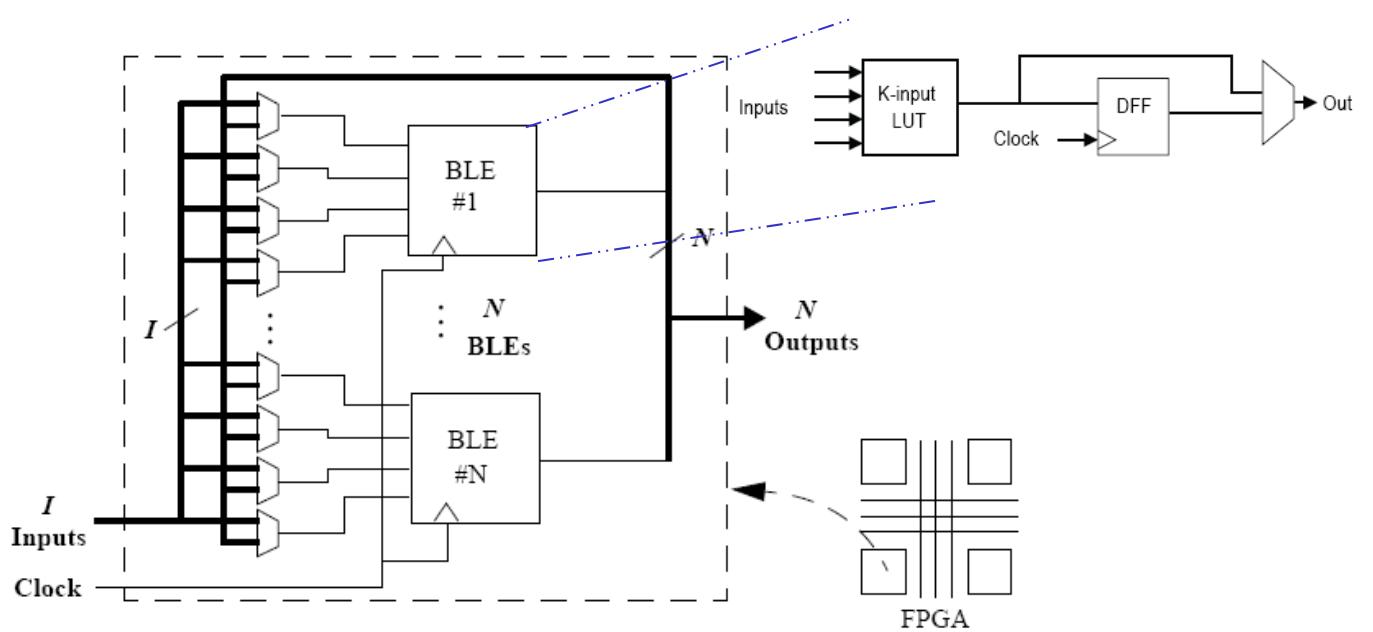
FPGA Architecture Evaluation Methodology

- Empirical approach to explore different architectures is typical



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Logic Block Structure



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Logic Block Granularity Study

- How large should the LUT size (K) be?

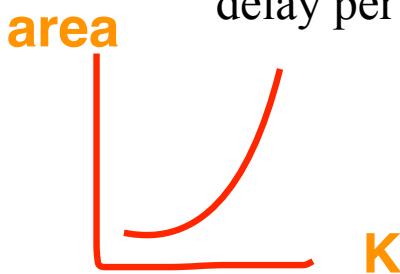
- Effects on area & speed

- Area

- As K increases, fewer logic blocks are needed for a design but area per block increases (LUT's SRAM bits is 2^K)

- Speed

- As K increases, each critical path contains fewer blocks but delay per block increases

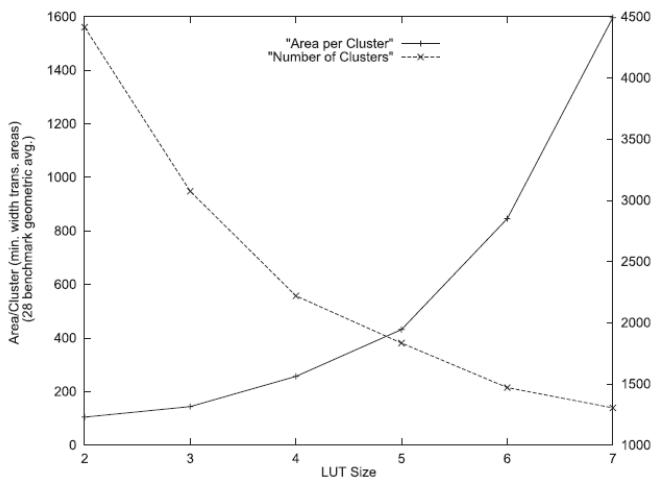


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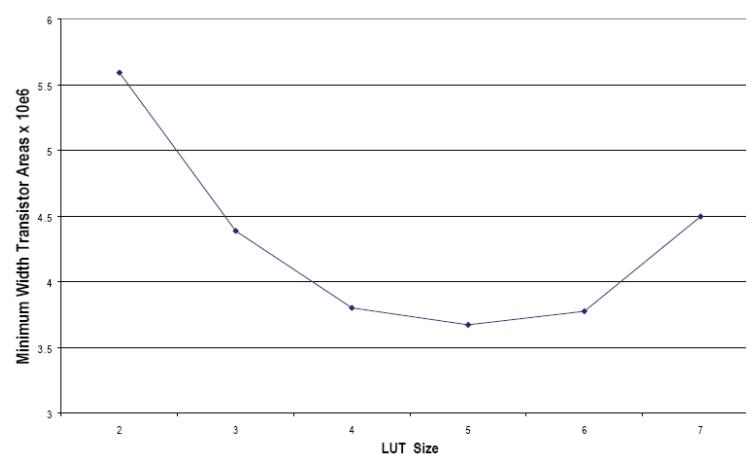
Effect of LUT Size on Area

- As LUT size (K) increases

- Total FPGA area first decreases and then increases



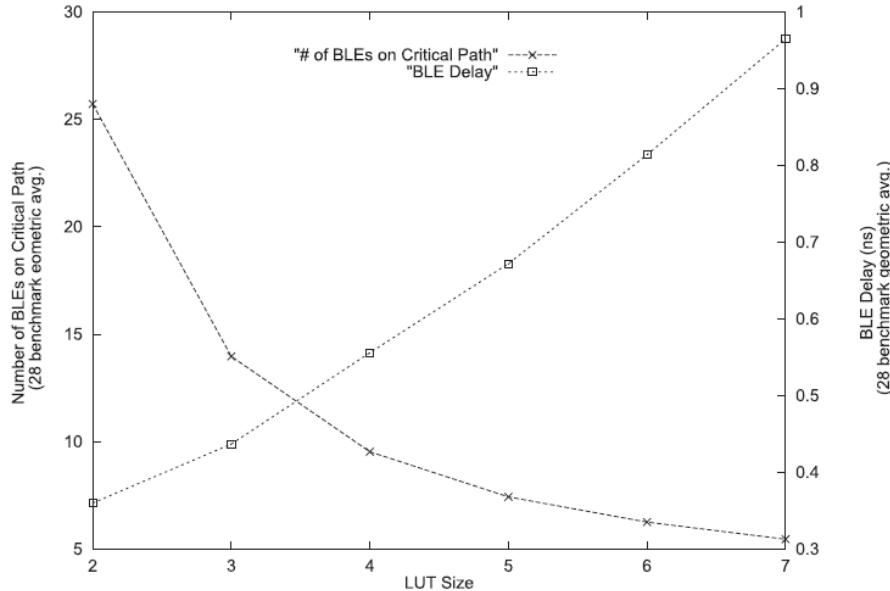
blocks required & area per block for different LUT sizes



Total FPGA area for different LUT sizes

Effect of LUT Size on Speed

- As K increases, each critical path contains fewer blocks but delay per block increases

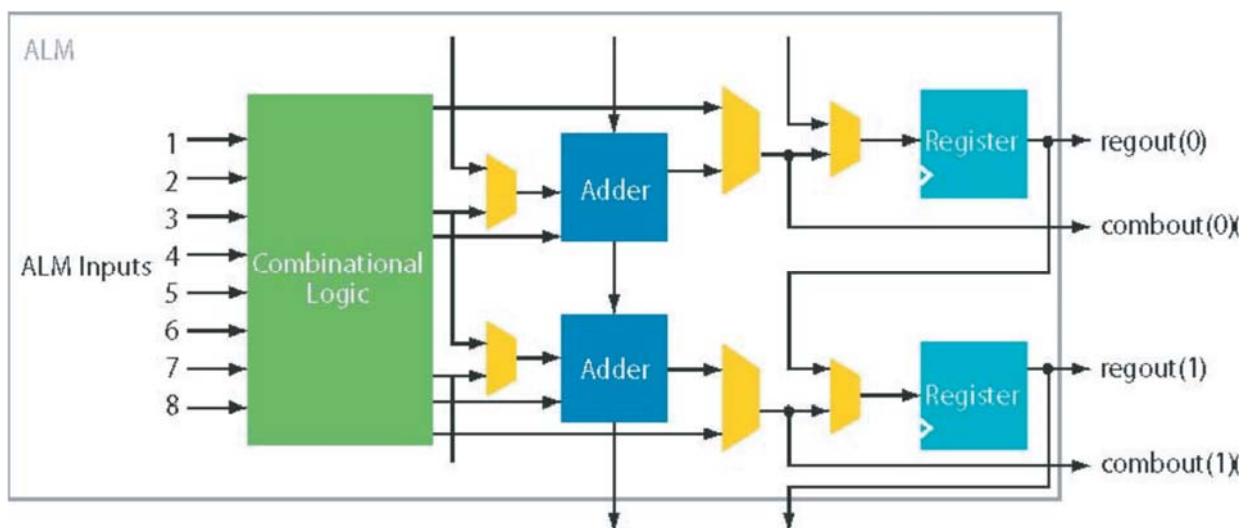


#LUTs on a critical path & delay per LUT for different LUT sizes

19

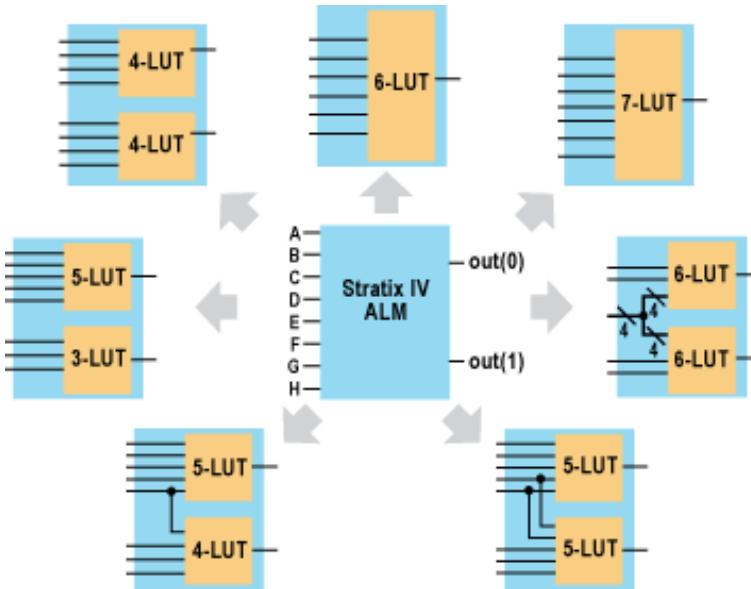
Innovative Idea – Adaptive Logic Module

- Altera Stratix ALM (adaptive logic module)



Flexibility of Adaptive Logic Module

- Fracturable into two

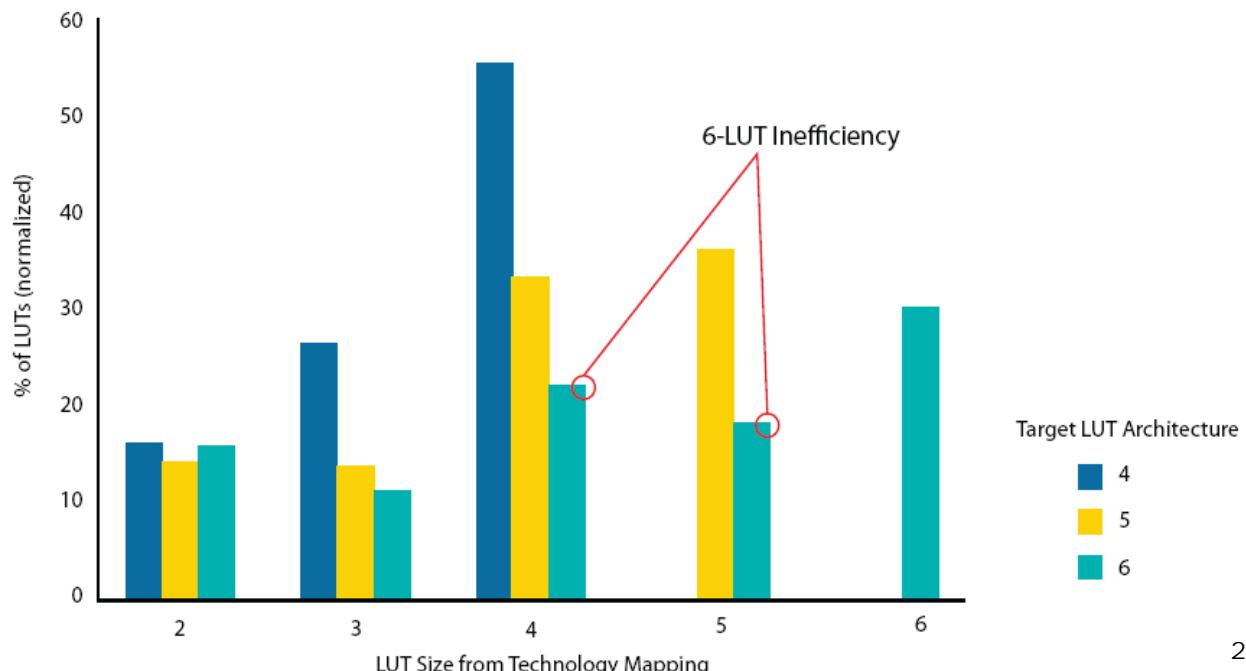


Output 1	Output 2	Shared inputs (min)
6-LUT	-	-
5-LUT	5-LUT	2
5-LUT	4-LUT	1
5-LUT	3-LUT	0
4-LUT	4-LUT	0
4-LUT	3-LUT	0
3-LUT	3-LUT	0

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Adaptive Logic Module

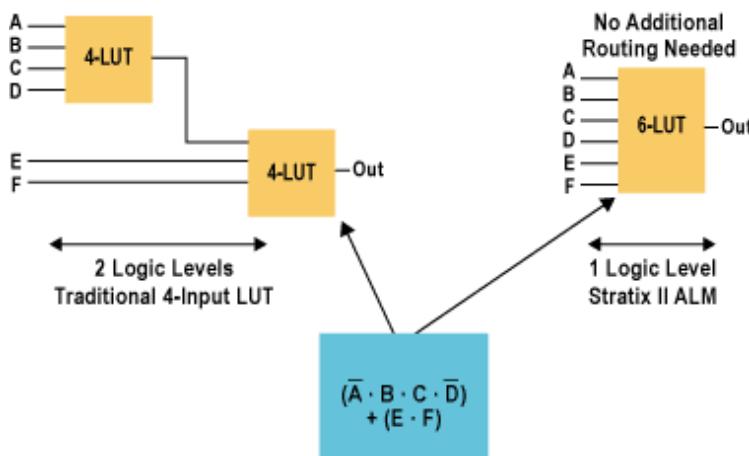
- *Observation:* Functions generated by synthesis have different input sizes.



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Advantages of Adaptive Logic Module

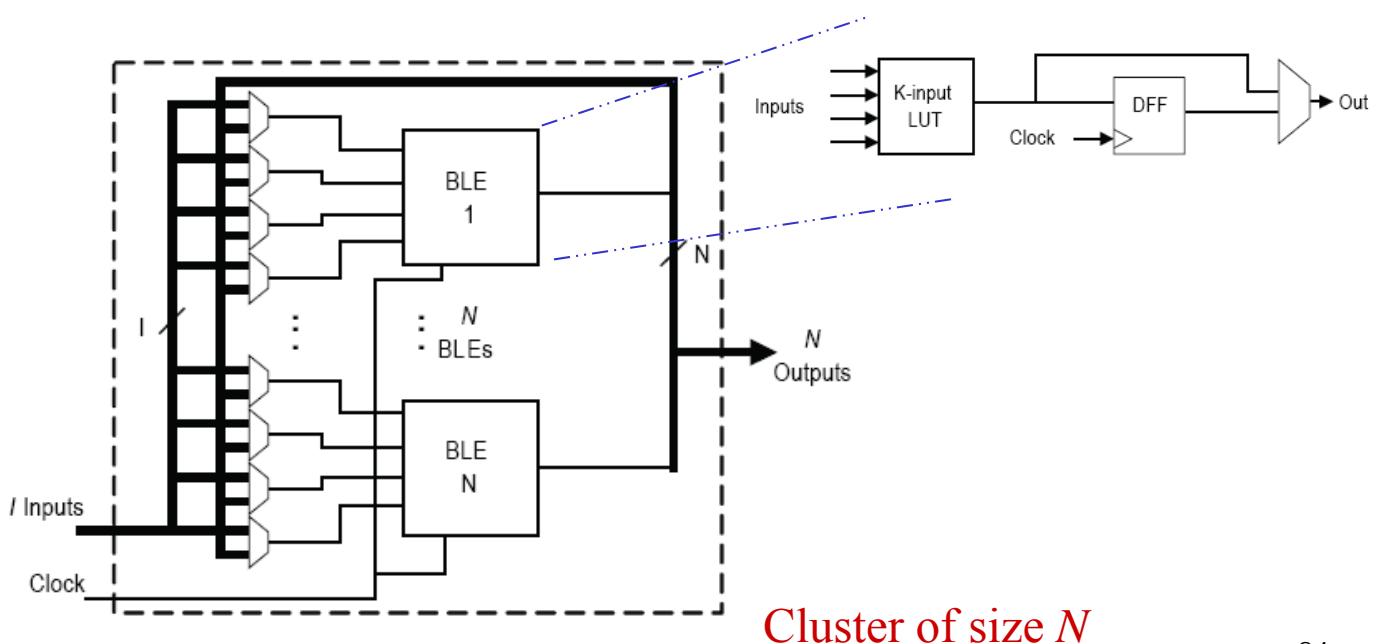
- ALM-based architecture vs traditional 4-LUT-based architecture
 - Improved area efficiency
 - Improved timing performance



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Logic Block Clustering

- Logic block made up of a cluster of LUTs and FFs



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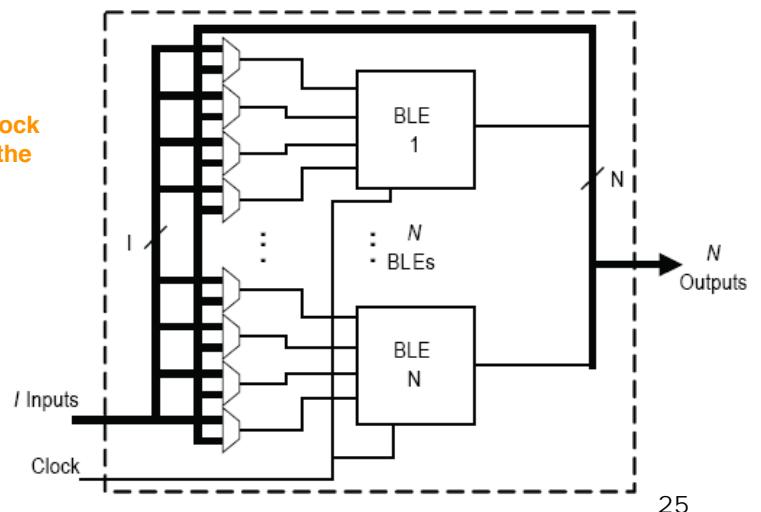
Logic Cluster Study

■ How many cluster input pins (I) are needed?

- BLEs in a cluster often share many input signals
- Empirically, # input pins I needed to fully utilize a cluster of N K -LUT is

$$\blacksquare I = K(N+1)/2$$

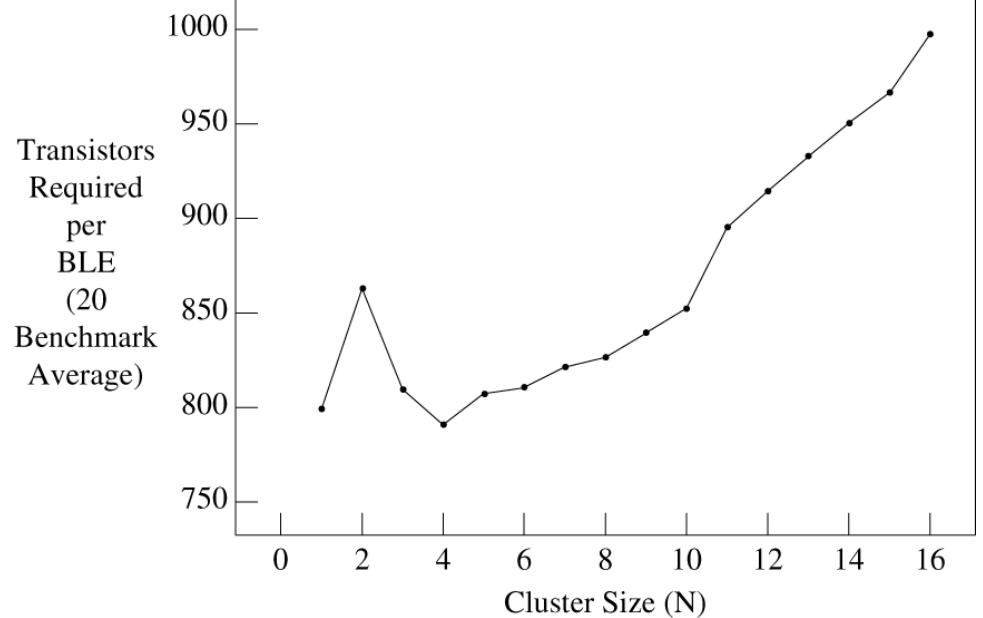
larger I will increase the overhead inside a logic block and also requires more access point(switches) to the routing channel



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Area Efficiency of Different Cluster Sizes

■ Clusters in size 1-8 are area-efficient.

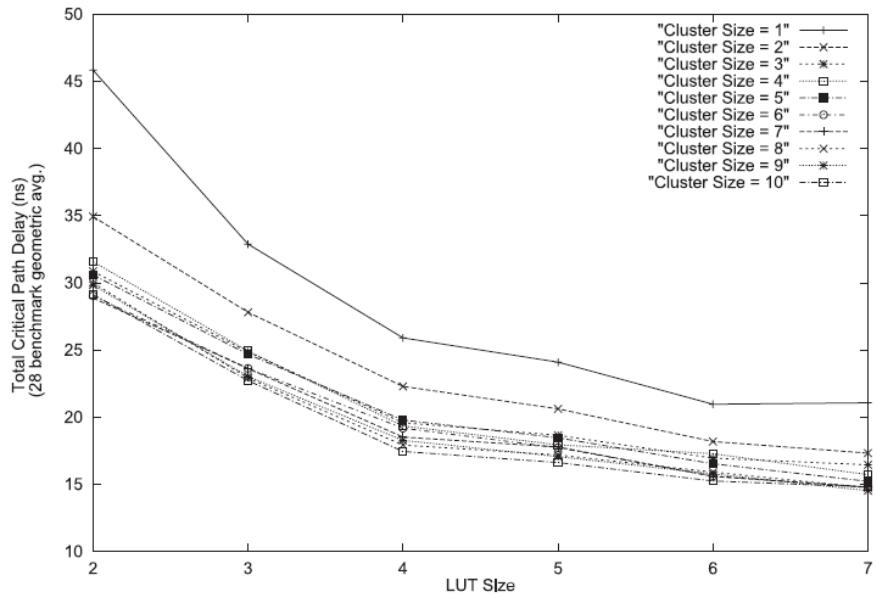


Transistors per BLE vs. cluster size (includes overhead circuits)

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Effect of Cluster Size and LUT Size on Speed

- As LUT and cluster size increase, critical path delay monotonically decreases with diminishing returns
- Significant returns to increase LUT size up to 6 and cluster size up to 3 or 4

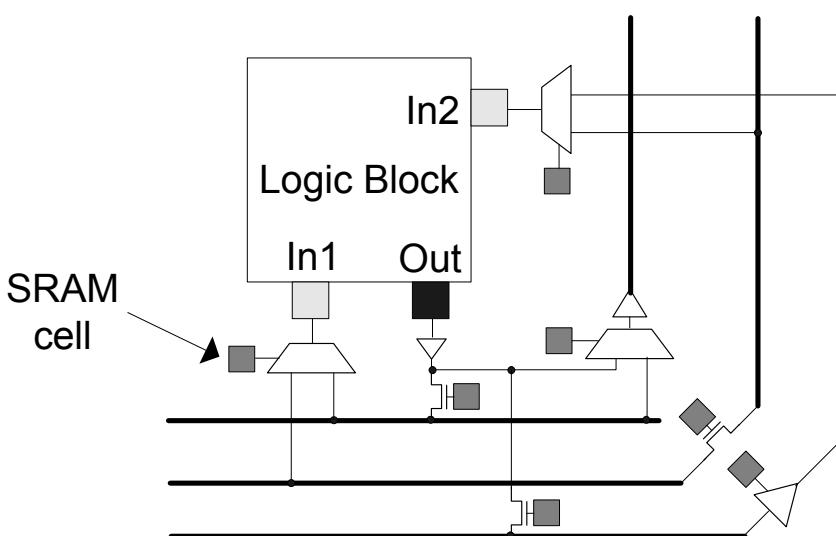


Critical path delay for different LUT and cluster sizes

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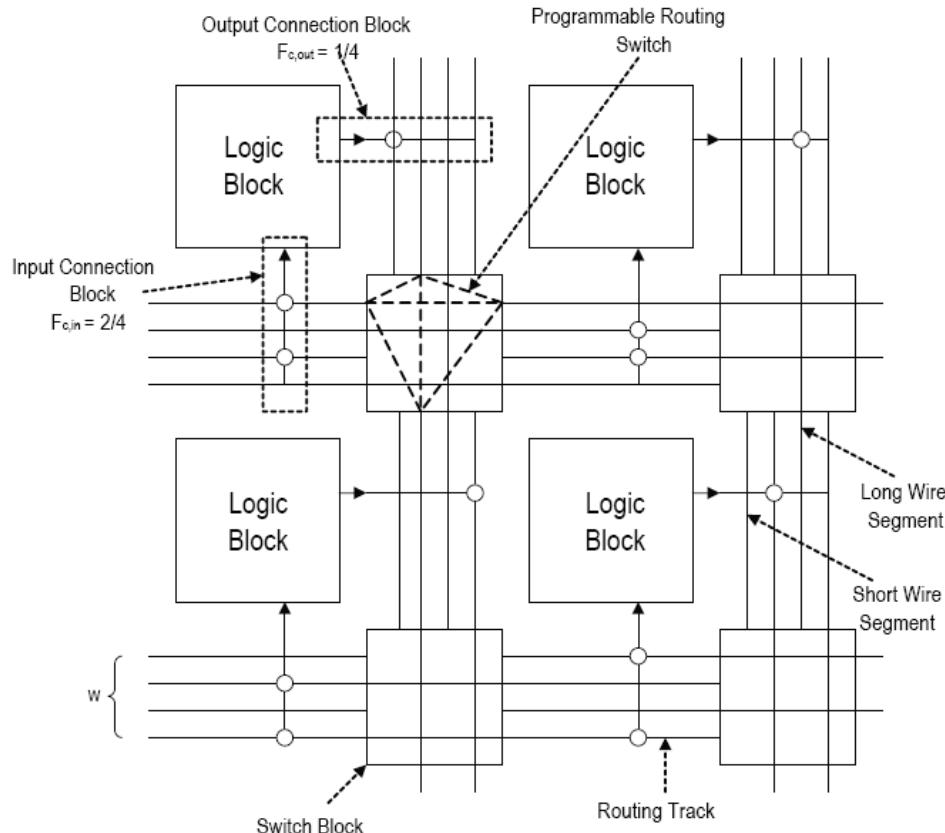
Programmable Routing

- Programmable switches connect fixed metal wires



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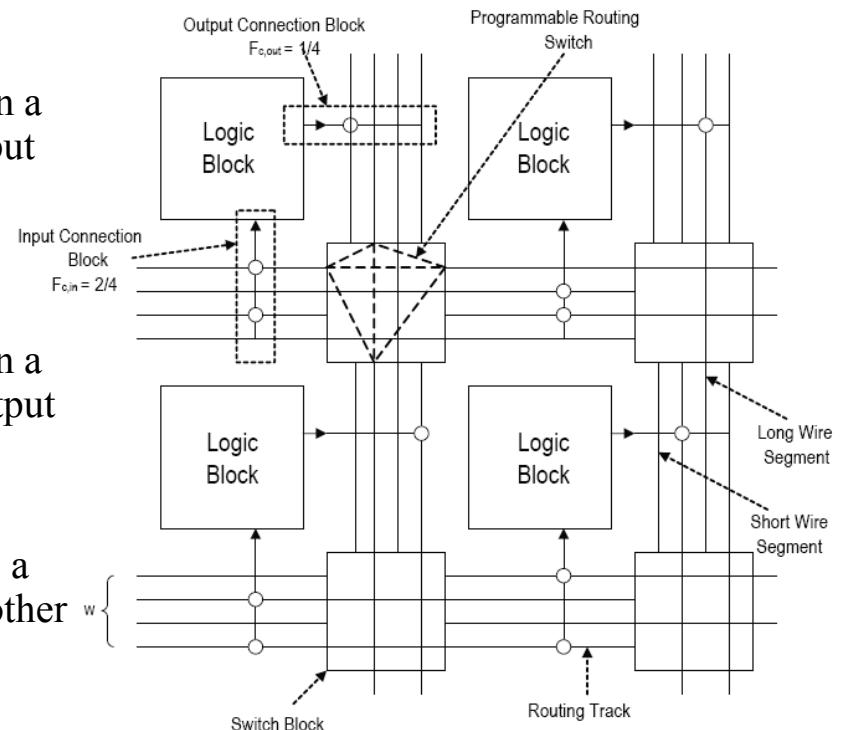
Routing Architecture



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Some Parameters of Routing Architecture

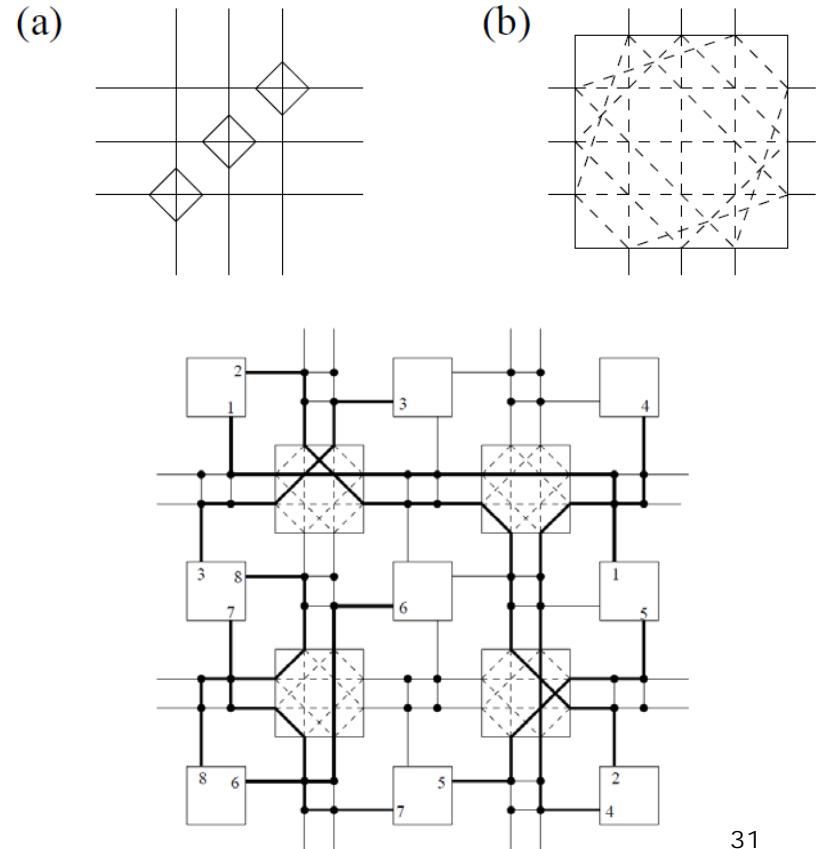
- Input connection block flexibility $F_{c,in}$
 - Fraction of wire segments in a channel connected to an input pin of a block
- Output connection block flexibility $F_{c,out}$
 - Fraction of wire segments in a channel connected to an output pin of a block
- Switch block flexibility F_s
 - No. of possible connections a wire segment can make to other wire segments



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Switch Block Structure

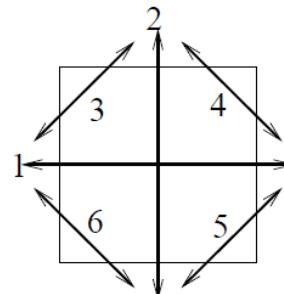
- E.g. Xilinx XC4000 switch block and its abstract representation



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Switch Block Structure

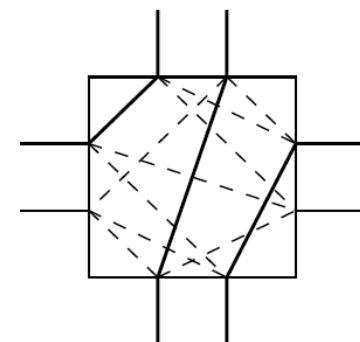
- 6 types of connections



- Routing requirement vector $(n_1, n_2, n_3, n_4, n_5, n_6)$ where n_i denotes #type- i connections

□ e.g. $(0,1,1,0,1,0)$ is routable but $(2,0,0,0,0,0)$ is not with the switch block below

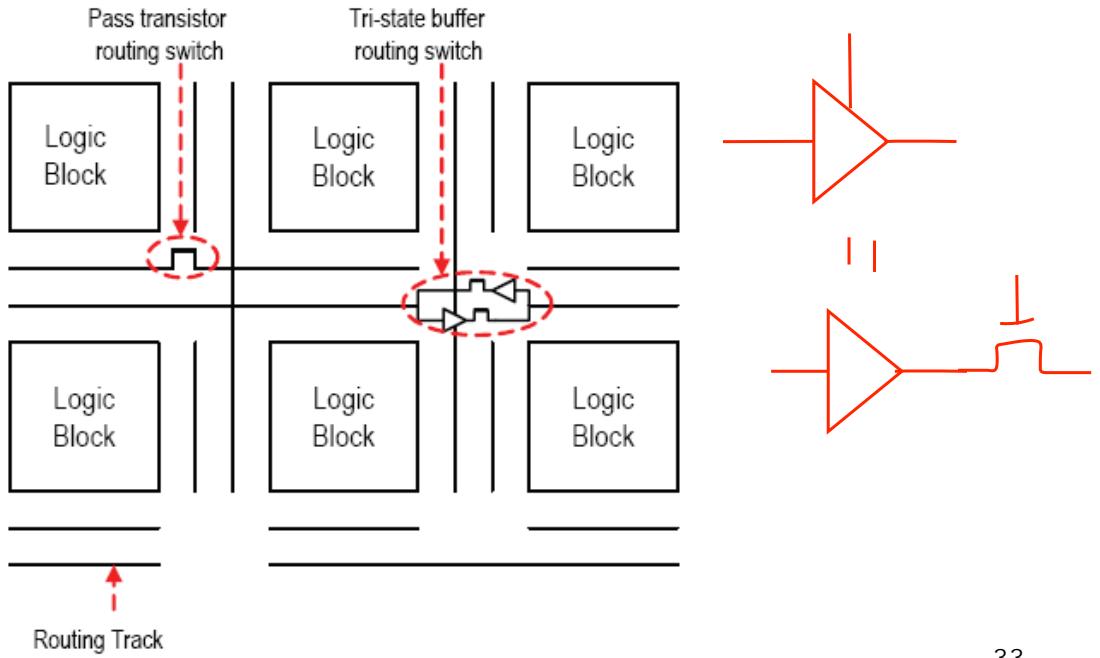
- Want a switch block structure with max no. of routable RRVs



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2 Types of Routing Switches

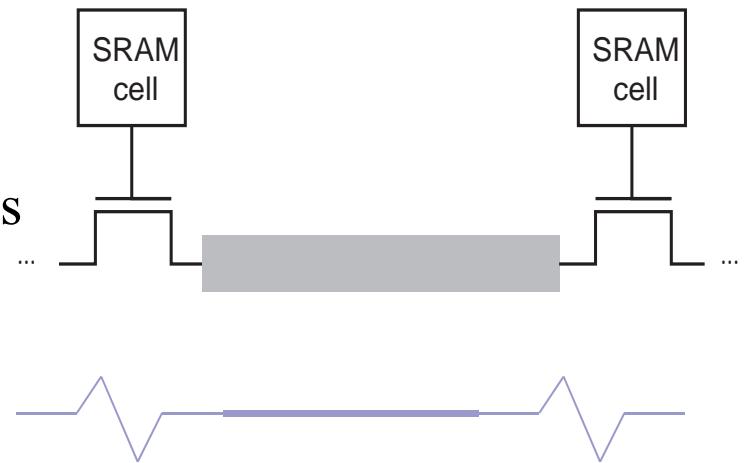
- Typically, mix pass transistor switches & tri-state buffer switches (*why?*)



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Pass Transistor Routing Switch

- Small area
- Resistive switch
- Faster for short paths
- Delay grows as the square of no. of switches



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Tri-state Buffer Routing Switch

- Larger area
- Regenerative driver
- Faster for long paths passing through many switches
- Delay grows linearly as no. of switches



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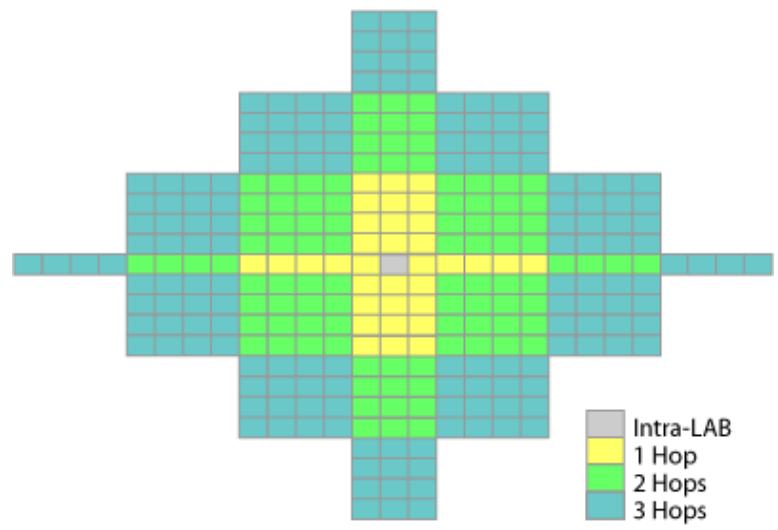
Other Routing Architecture Factors and Parameters

- Speed, Area and Power also depend on
 - Channel segmentation
 - Transistor size
 - Buffer size
 - Ratio of pass transistor switches & tri-state buffer switches
 - Metal width
 - Wire spacing
 - ...

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Connectivity

- What is the # hops required to get from one logic block to another?
- Fewer hops → better performance
- More predictable pattern → easier CAD tool optimization

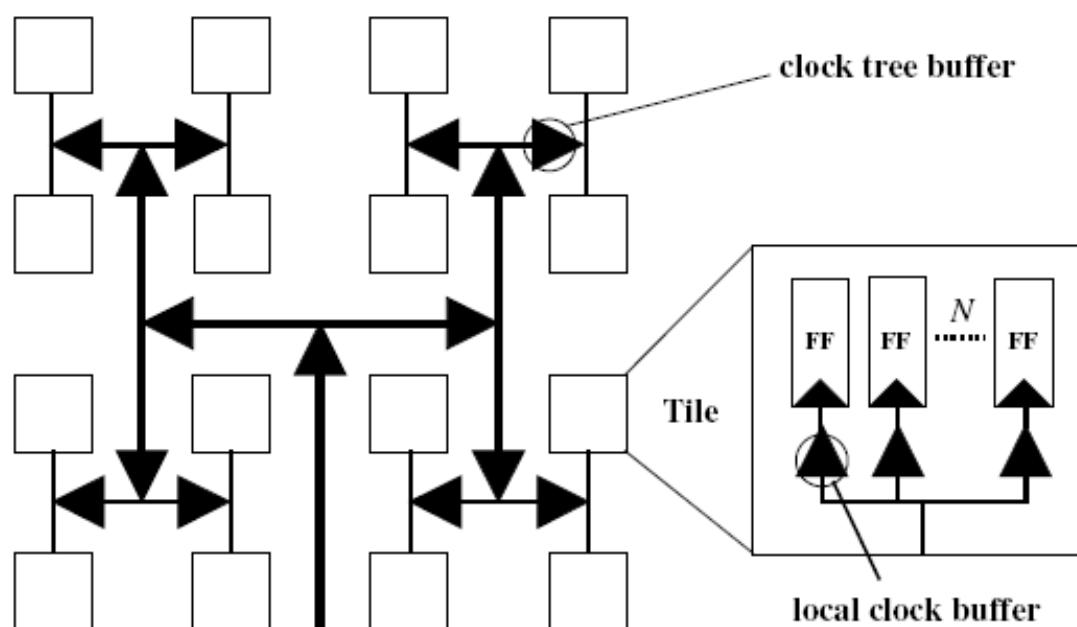


Stratix FPGA series connectivity

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Clock Nets

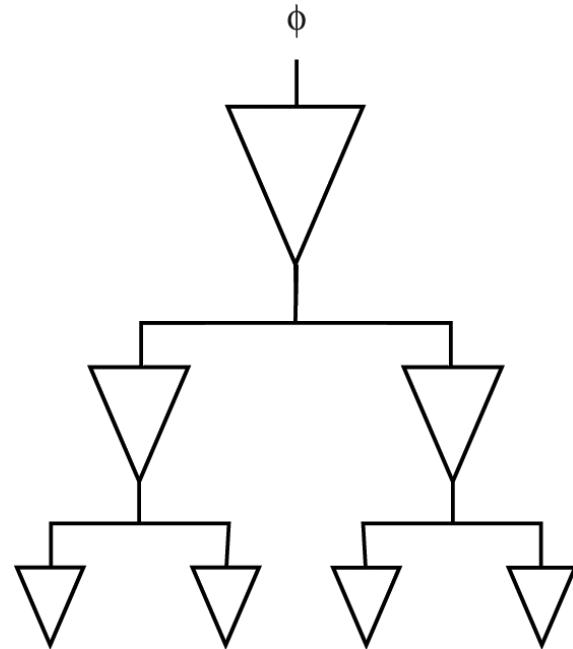
- Must drive all LEs.



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Clock Drivers

- Clock driver tree.
- Determine optimal buffer sizes.



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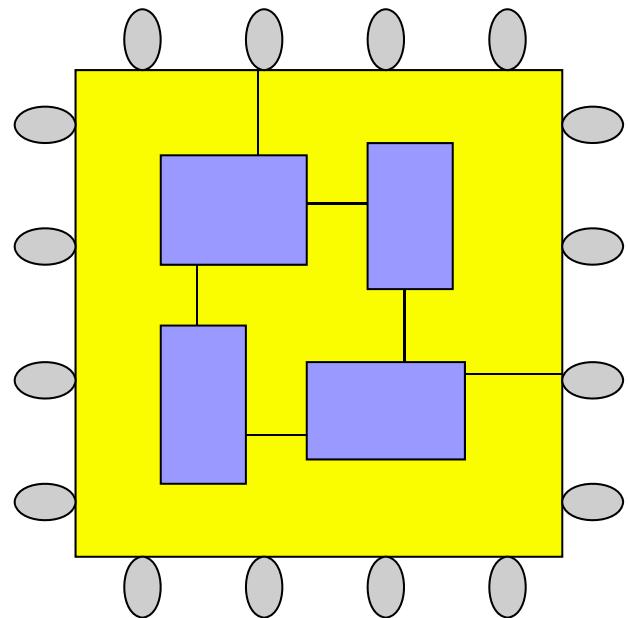
Track Distribution

- Is wiring concentrated near the center of the FPGA?
 - No.
- Is wiring directional (horizontal/vertical)?
 - No.
- Make channels to I/O pins about 25% larger improves routability.

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Pinout

- How many pins?
 - Limited by technology.
 - Too much logic, not enough pins means we can't get signals off-chip.
 - Too many pins means logic won't be fully utilized.



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Rent's Rule

- Developed by E. F. Rent (IBM) in 1960.
 - Experimentally derived from sample designs.
- Number of pins vs. number of components is a line on a log-log plot:
 - $N_p = K_p N_s^\beta$
- Parameters may vary based on technology:
 - Rent measured $\beta = 0.6$, $K_p = 2.5$.
 - Modern microprocessor has $\beta = 0.455$, $K_p = 0.82$.

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FPGAs and Pins

- Chip capacity is growing faster than package pinout.

chip 進步速度很快 PIN的數量進步速度很慢

- Harder to use logic in a multi-FPGA design
 - must try to fit a large function with a small interface into the FPGA
 - may use time-division multiplexing for I/Os

輪流用IO pin

43

References

- “Leakage Control in FPGA Routing Fabric”, in *ASP-DAC’05*.
- “The Effect of LUT and Cluster Size on Deep-Submicron FPGA Performance and Density”, in *FPGA ’00*.
- “Improving FPGA Performance and Area Using an Adaptive Logic Module”, in *FPL’04*.
- “Flexibility of interconnection structures for field programmable gate arrays”, *IEEE J. Solid-State Circuits*, vol. 26(3), 1991.
- “Mixing Buffers and Pass Transistors in FPGA Routing Architectures”, in *FPGA ’01*.
- “FPGA Architecture: Survey and Challenges”, *Foundations and Trends in Electronic Design Automation*, vol.2(2), 2007.
- “VPR 5.0: FPGA CAD and Architecture Exploration Tools with Single-Driver Routing, Heterogeneity and Process Scaling”, in *FPGA ’09*.

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