

1. $D_1 = 2 - 6 = -4$

$D_2 = 6 - 3 = 3$

$g_{12} = D_1 + D_2 - 2C_{12}$
 $= -1 - 2C_{12}$

$\because C_{12}$ is non-negative

$\therefore g_{12} < 0$

交換 vertex 1, vertex 2 會使結果變好!



2. $D_1, D_n, D_{n+1}, D_{2n}$ 皆為 0

$D_2 \sim D_{n-1}, D_{n+2} \sim D_{2n-1}$ 皆為 -1 ($E - I = 1 - 2$) + 4

$C_{ij} = 1$ if $|i - j| = n$

$C_{ij} = 0$ if $|i - j| \neq n$

$g_{ij} = D_i + D_j - 2C_{ij}$

$\because D_i + D_j \leq 0$ 且 $C_{ij} \geq 0$

$\therefore g_{ij}$ 必定 ≤ 0

因此經過一個 pass 之後

得到之最大 partial sum 必 ≤ 0

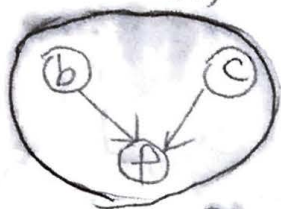
根據 partitioning 講義第 14 頁之虛擬臨
 , 結果不會做任何的 swap!



3. $l(b)=1, l(c)=1, l(e)=3, l(g)=3$

f:

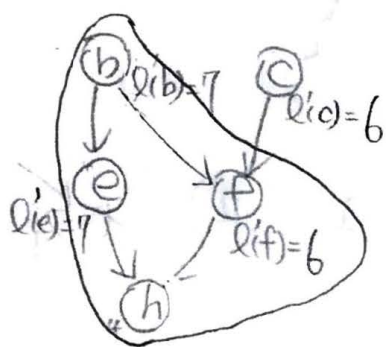
cluster(f)



$$\begin{aligned}
 l(f) &= \max\{l_1(f), l_2(f)\} \\
 &= \max\{\max\{2, 2\}, 0\} \\
 &= 2
 \end{aligned}$$

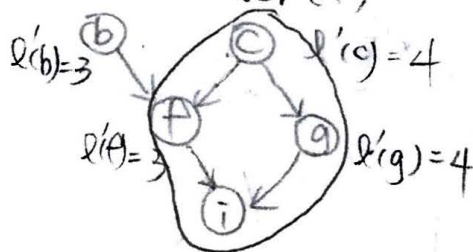
h:

cluster(h)



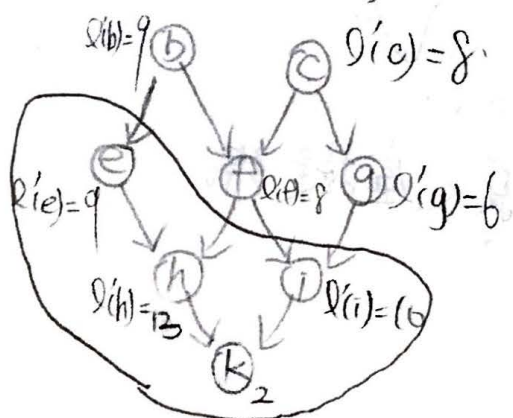
$$\begin{aligned}
 l(h) &= \max\{l_1(h), l_2(h)\} \\
 &= \max\{\max\{7\}, \max\{6+5\}\} \\
 &= \max\{7, 11\} \\
 &= 11
 \end{aligned}$$

i: cluster(i)



$$\begin{aligned}
 l(i) &= \max\{l_1(i), l_2(i)\} \\
 &= \max\{\max\{4\}, \max\{3+5\}\} \\
 &= \max\{4, 8\} \\
 &= 8
 \end{aligned}$$

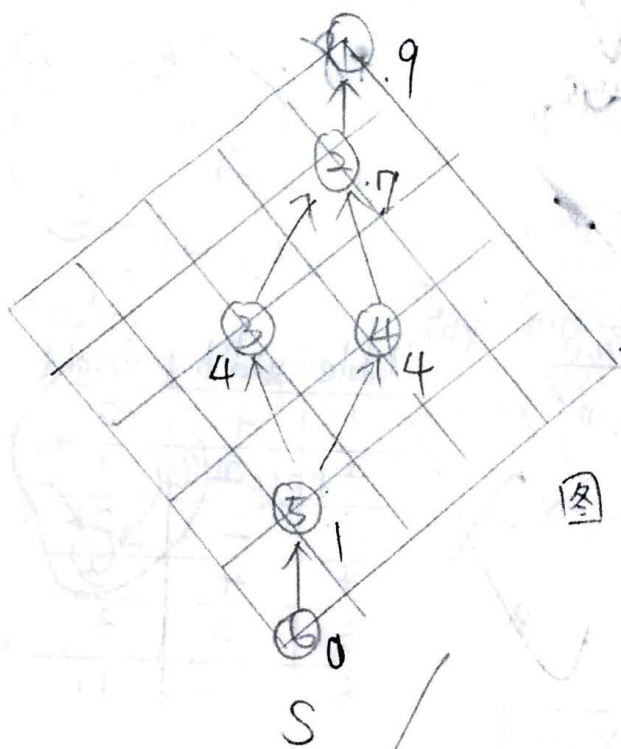
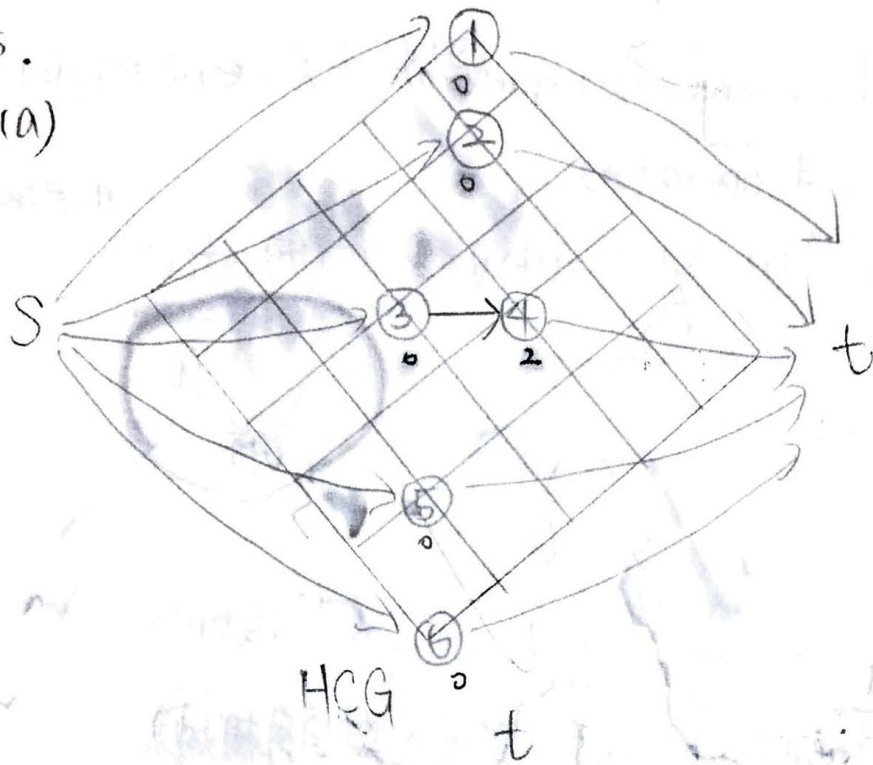
k: cluster(k)



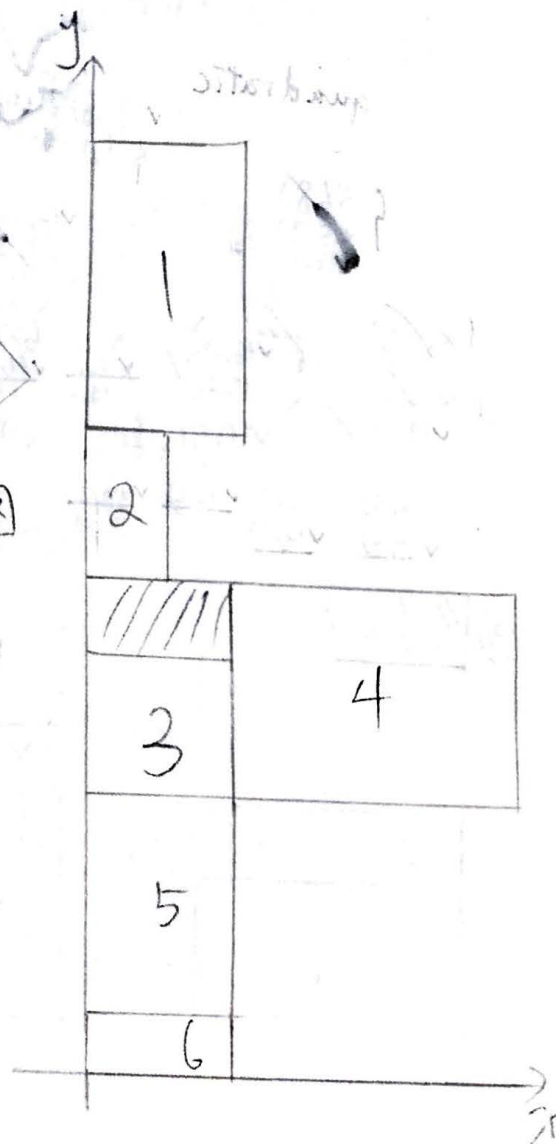
$$\begin{aligned}
 l(k) &= \max\{l_1(k), l_2(k)\} \\
 &= \max\{0, \max\{9+5, 8+5, 8+5, 6+5\}\} \\
 &= \max\{0, 14\} \\
 &= 14
 \end{aligned}$$

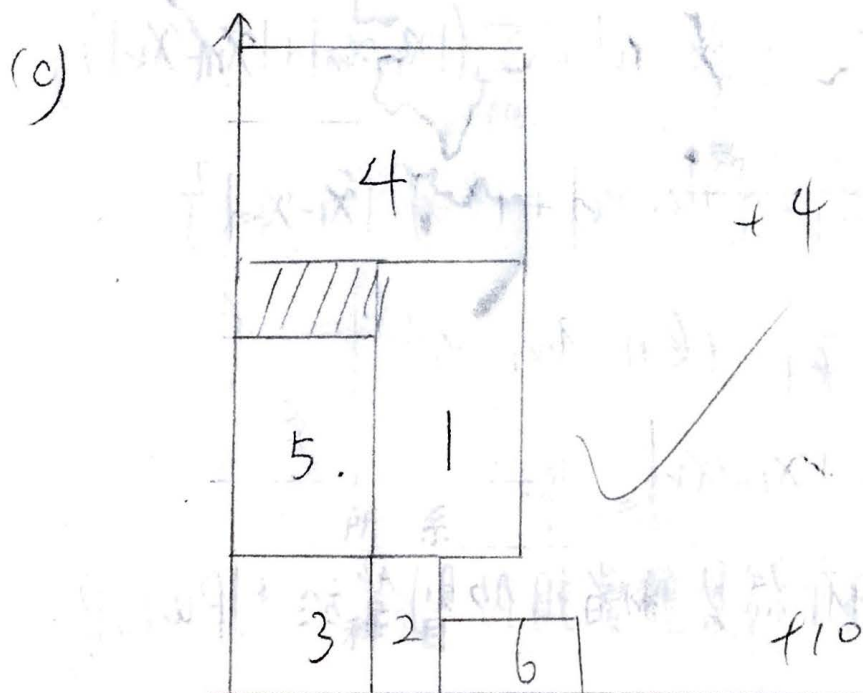
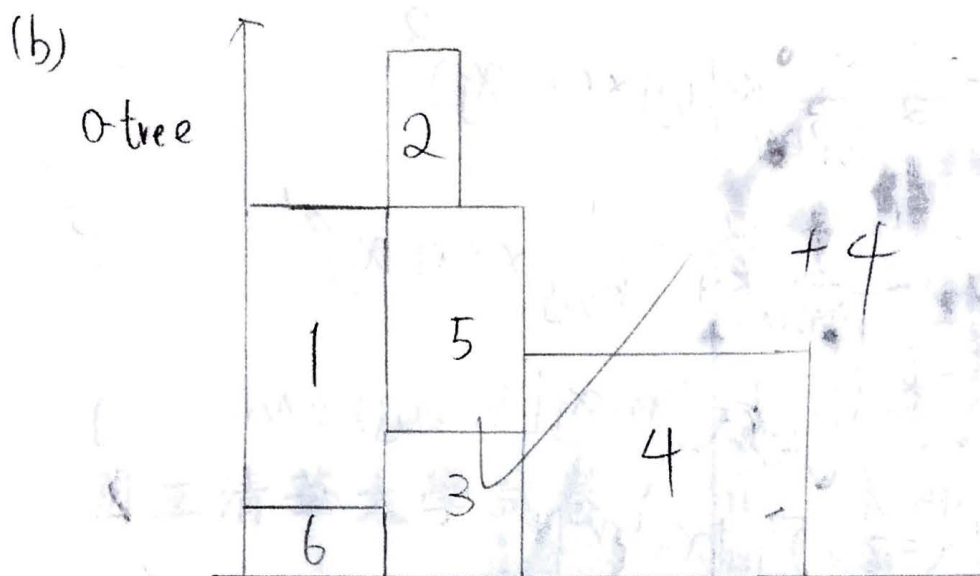
5.

(a)



图





6. 在一個包含所有 p_i 的 Bounding Box 中, 若欲建一個 RSMT 連接所有 p_i , 最理想之情形為任一 x 或 y 座標, 只會被此 RSMT 中的一線段經過, 不可有第二個線段經過之, 在此情形下將水平及垂直方向之所有線段長相加, 可分別得到 $(\text{Max of } x - \text{Min of } x)$ 及 $(\text{Max of } y - \text{Min of } y)$, 兩者相加為 HPWL. 故 HPWL 為 RSMT 線長之 lower bound!

$$\begin{aligned}
 1. \quad L^{\sim BB} &= \frac{1}{2} \sum_{\{i,j\} \in N} w_{\{i,j\}} \times (x_i - x_j)^2 \\
 &= \frac{1}{2} \sum_{\{i,j\} \in N} \frac{2}{k-1} \times \frac{1}{l_{\{i,j\}}} \times (x_i - x_j)^2 \\
 &\quad \text{if } l_{\{i,j\}} = |x_i - x_j| \quad \forall \{i,j\} \in N
 \end{aligned}$$

$$\begin{aligned}
 \Rightarrow L^{\sim BB} &= \frac{1}{k-1} \sum_{\{i,j\} \in N} |x_i - x_j| \\
 &= \frac{1}{k-1} \left\{ |x_1 - x_k| + \sum_{2 \leq m \leq k-1} (|x_1 - x_m| + |x_m - x_k|) \right\} \\
 &= \frac{1}{k-1} \cdot \left\{ |x_1 - x_k| + (k-2) \cdot |x_1 - x_k| \right\} \\
 &= \frac{1}{k-1} \cdot (k-1) \cdot |x_1 - x_k| \\
 &= |x_1 - x_k|
 \end{aligned}$$

y 方向亦然! 兩者相加則等於 HPWL!

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8.

7	6	7					
6	5	6	7				
5	4	5	6	7			
4	3						
3	2						6
2	1		1	2	3	4	5
3	2	1	2	3	4	5	6
4	3	2	3	4	5	6	7

+3

9.

將 lee-algorithm 之水波擴散方向改成 8 個
，往斜方行進之格子標記為現有步數加上 $\sqrt{2}$

，若欲標記之格子已被標記則不標記之

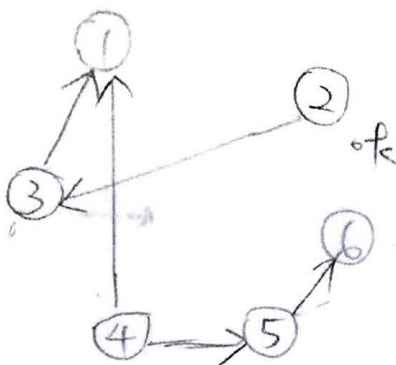
每次從標記最小步數的點向外擴 (同 lee alg)

，backtrace 之方向改為從終點開始找，
每次皆從八個方向找到小於自身標記
步數之格子中的最小值做為 backtrace
之方向。

10.

VCG

(a)



前边反了

+2

(b)

Track 1: $I_4[4,7]$

Track 2: $I_5[3,7]$

Track 3: $I_6[1,3], I_2[5,6]$

Track 4: $I_3[2,6]$

Track 5: $I_1[2,4]$

+7

