



Official Course and Teaching Evaluation

- ▶ OCTE system: <https://octe.cuhk.edu.cn/a/login>



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香港中文大學 (深圳)
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CSC3100 Data Structures

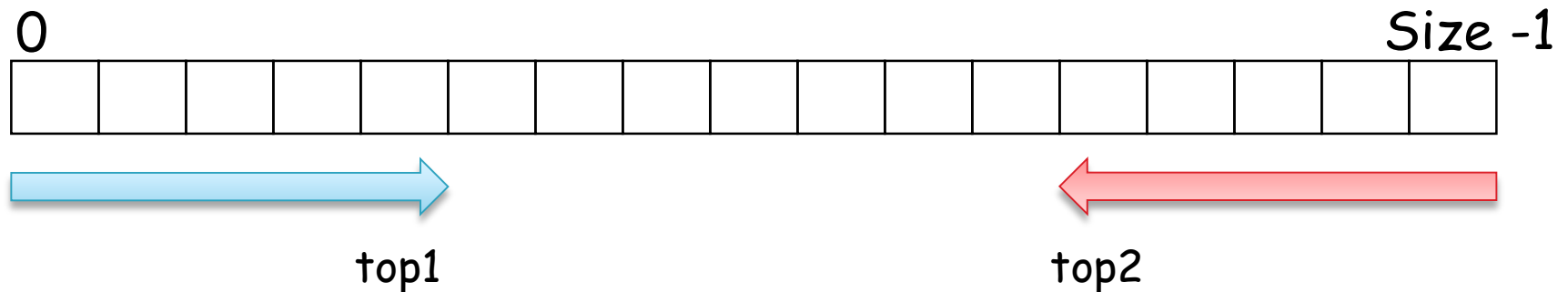
Lecture 26: Exercises

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Exercise 1: Two stacks in one array

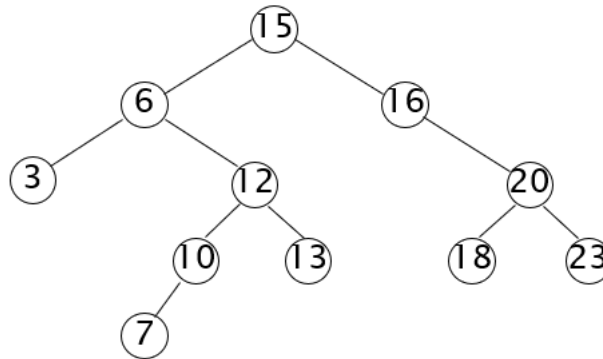
- ▶ Design two stacks in one array $A[1...n]$ ($n > 1$) in such that neither stack overflows unless the total number of elements in both stacks is n . The PUSH and POP operations should run in $O(1)$ time.
 - Please briefly explain how to implement the two stacks, and then use pseudocodes to explain the steps of PUSH and POP for them.





Exercise 2: Binary search tree

- ▶ Insertion and deletion on binary search tree:
 - Given an initial binary search tree as shown below, perform two consecutive operations: **first insert a new key 14 into the tree, and then delete key 15** from the updated tree. Please draw these two updated trees after these operations respectively.

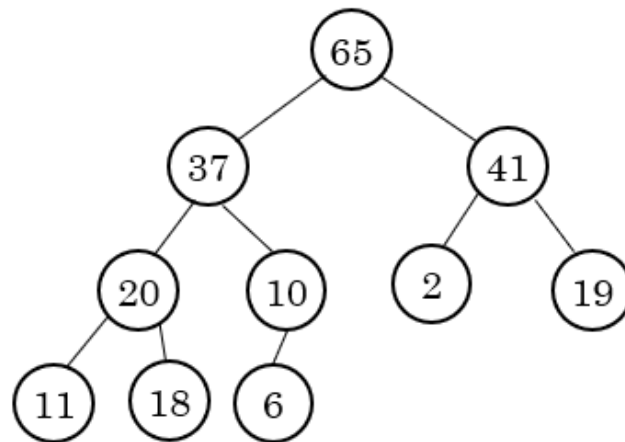


- Analyze the time complexities of key insertion and key deletion on binary search tree with n nodes, in the best case and worst case respectively.



Exercise 3: Heap operations

- ▶ A priority queue is implemented as a max-heap. Use a list of tree graphs to show how the heap Q shown below would look like after a series of operations:
 - (1) $Q.Enqueue(12)$;
 - (2) $Q.Enqueue(102)$; (use the heap from step (1))
 - (3) $Q.Dequeue()$. (use the heap from step (2))





Exercise 4: Hashing

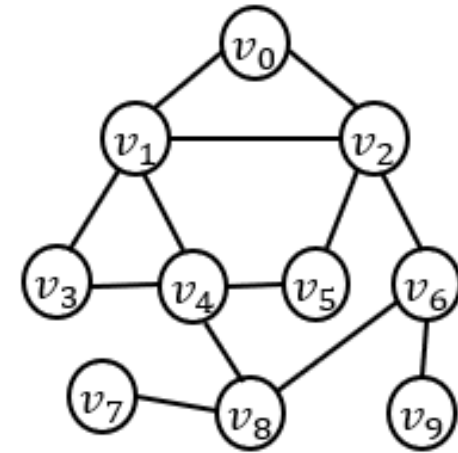
- ▶ Assume that we have a sequence of 9 keys: 6, 12, 34, 29, 18, 17, 1, 5, 4, and the size of the hash table is $m=11$. Show the filled hash tables after inserting these keys using the following collision resolution strategies and hashing functions respectively.
 - Use chaining with $h(k) = k \% 11$;
 - Use linear probing with $h(k) = k \% 11$;

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Exercise 5: BFS & DFS

- ▶ Consider an undirected graph G_1 :
 - Draw the BSF tree of G_1 by assuming that the starting node is set to v_1
 - Draw the DSF tree of G_1 by assuming that the starting node is set to v_2
 - Compare BFS with DFS in terms of time cost and extra space cost, if we use the adjacent list to represent the graph. Besides, for each of them, show a specific application that favors it and explain the reason briefly.

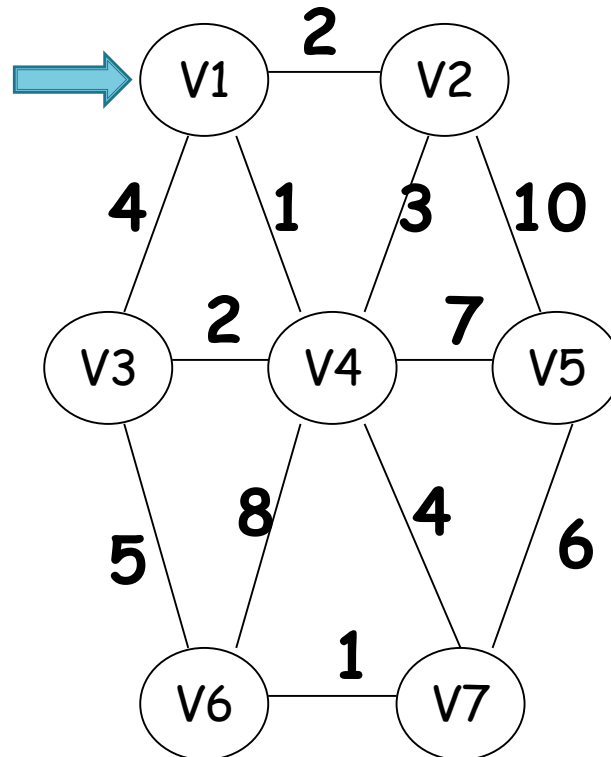


Graph G_1



Exercise 6: Find an MST

- ▶ Use Prim's algorithm and show its steps
- ▶ Use Kruskal's algorithm and show its steps





Exercise 7: “maximum” spanning tree

- ▶ Find an algorithm for the “maximum” spanning tree. That is, given an undirected weighted graph G , find a spanning tree of G of maximum cost, and prove the correctness of your algorithm
 - Consider choosing the “heaviest” edge (i.e., the edge associated with the largest weight) in a cut
 - The generic proof can be modified easily to show that this approach will work
 - Alternatively, multiply the weights by -1 and apply either Prim’s or Kruskal’s algorithms without any modification at all!



Exercise 8: shortest path with k edges

- ▶ Given a directed, weighted graph that might have negative weight edges, and we know that all the shortest paths use at most k edges. How to compute shortest paths from a source vertex s in $O(k(V+E))$ time?
 - Run the Bellman-Ford algorithm, but stop after k iterations of the outer loop.
 - After the i -th iteration of the outer loop, Bellman-Ford has found paths that are at least as good as the shortest paths from s that use at most i edges.
 - Thus, after k iterations, it has found paths that are at least as good as the shortest paths from s that use at most k edges, which under the given assumptions are the shortest paths overall.



Recommended reading

- ▶ Next lecture
 - Review of all course materials