Lecture 5 Stack and Queue

Our Roadmap

Stack

Queue

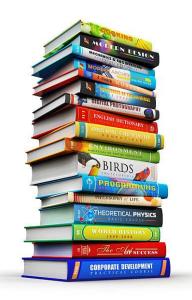
Stack vs. Queue

Stack

A stack is a sequence in which:



- Items can be added and removed only at one end (the top)
- You can only access the item that is currently at the top
- Stack Analogy

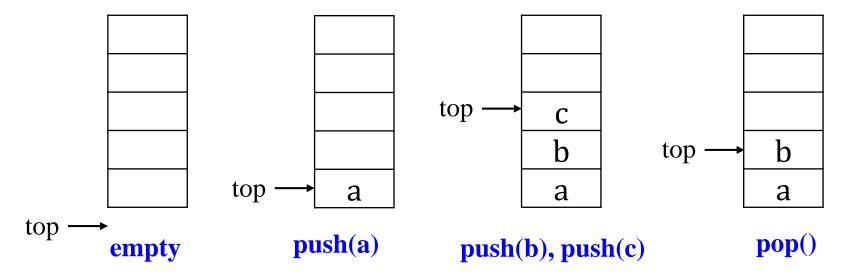






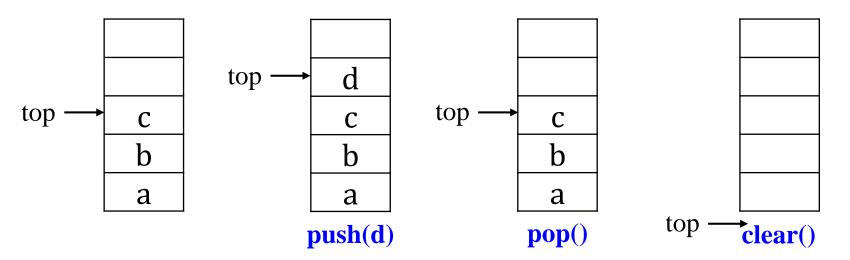
Stack

- First In Last Out (FILO)
 - Constrained item access
- Major Operations
 - push: add an item to the top of the stack
 - pop: remove the item at the top of the stack
- Illustration



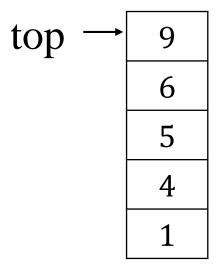
Stack Operation

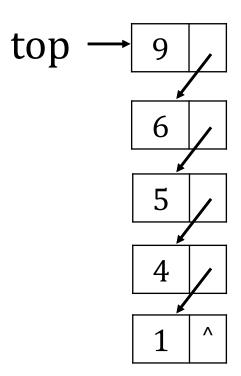
- push: add an item to the top of the stack
- pop: remove the item at the top of the stack
- top/peek: get the item at the top of the stack, but do not remove it
- isEmpty: test if the stack is empty
- isFull: test if the stack is full
- clear: clear the stack, set it as empty stack
- size: return the current size of the stack



Implementation of Stack

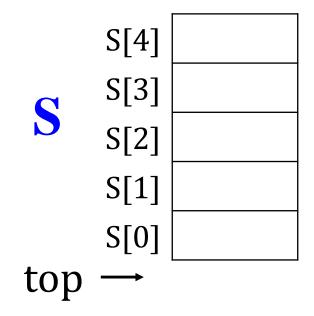
- Array-based Stack
- Linked Stack



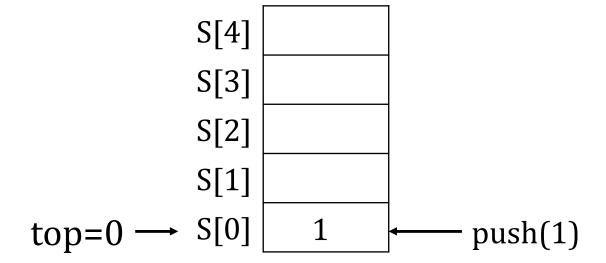


Implementation of Stack

- Array based Stack
 - MAX_SIZE = n // the max size of stack
 - \bullet top = -1 // the current top position
 - Array S with n elements
- Example
 - ♦ MAX_SIZE = 5
 - top = -1
 - Array S



- push(item):
 - 1. top++;
 - 2. S[top] = item
- push(1)



push(item):
 1. top++;
 2. S[top] = item

push(1)
push(4)
 S[4]
 S[3]

S[2]

4

 $top=1 \longrightarrow S[1]$

 $top=0 \longrightarrow S[0]$

push(4)

push(1)

push(item): 1. top++; 2. S[top] = itempush(1) push(4) S[4] push(5) S[3] $top=2 \longrightarrow S[2]$ push(5) $top=1 \longrightarrow S[1]$ push(4)4 $top=0 \longrightarrow S[0]$ push(1)

push(item): 1. top++; 2. S[top] = itempush(1) push(4)S[4] push(5) $top=3 \rightarrow S[3]$ 6 push(6)push(6) $top=2 \longrightarrow S[2]$ push(5)5 $top=1 \longrightarrow S[1]$ push(4)4 $top=0 \longrightarrow S[0]$

push(1)

top=5

push(9)

top=4
$$\rightarrow$$
 S[4] 9 \leftarrow push(9)
S[3] 6
S[2] 5
S[1] 4
S[0] 1

- push(10)
 - OVERFLOW
 - How to avoid that?

S[4]	9
S[3]	6
S[2]	5
S[1]	4
S[0]	1

Push / Pop Operator

push(item):
 1. if(top == MAXSIZE-1)
 2. Stack is FULL! No push!
 3. else
 4. top++;
 5. S[top] = item

```
pop(): // should avoid underflow
1. if(top == -1)
2.    Stack is EMPTY! No pop!
3. else
4.    top--;
```

Application of Stacks

- Making sure the delimiters (parens, brackets, etc.) are balanced:
 - Push open (i.e., left) delimiters onto a stack
 - When you encounter a close (i.e., right) delimiter, pop an item off the stack and see if it matches
- Evaluating arithmetic expressions
 - Parsing arithmetic expressions written using infix notation
- The runtime stack in memory
 - Converting a recursive algorithm to an iterative one by using a stack to emulate the runtime stack

Brackets Balance Problem

- \bullet a+{2-[b+c]*(8* [8+g]/[m-e]-7)-p}
- * {[]([][])}
- Skip operators and notations
- Is the bracket expression balanced or not?
 - () Yes
 - ♦ No
 - ♦ { [()} No

Brackets Balance Problem

- Given a bracket expression, determine whether it is balanced or not?
- * {[]([][])}
 - Mow to solve it by using stack?
 - Bracket pairs: (), [], { }
 - Any ideas?
- Methodology
 - Employ stack store checked left bracket
 - Pop out left bracket if it is matched

Arithmetic Expression Evaluation

- Arithmetic expression
 - operands (a, b, c), operator (+, *)
 - a + b * c
- Prefix expression
 - ♦ + a * b c
- Infix expression
 - \diamond a + b * c
- Postfix expression
 - a b c * +

Postfix Expression

- Infix expression
 - 5 * ((9 + 3) * (4*2) + 7)
- Postfix expression
 - § 5 9 3 + 4 2 * * 7 + *
- Parse postfix expression is somewhat easier problem than directly parsing infix (why)
- Postfix has a nice property that parentheses are unnecessary
- Postfix Expression Evaluation
 - Convert from infix to postfix
 - Evaluate a postfix expression

Postfix Expression

- Postfix expression
 - ⋄ 5 9 3 + 4 2 * * 7 + *
- Methodology
 - Read the tokens in one at a time
 - If it is an operand, push it on the stack
 - If it is a binary operator:
 - pop top two elements from the stack,
 - apply the operator,
 - and push the result back on the stack

Postfix Expression Evaluation

- § 593+42**7+*
- Postfix Expression Evaluation

Stack operations

- push(5)
- push(9)
- push(3)
- push(pop() + pop())
- push(4)
- push(2)
- push(pop() * pop())
- push(pop() * pop())
- push(7)
- push(pop() + pop())
- push(pop() * pop())

Stack elements

5

59

593

5 12

5 12 4

5 12 4 2

5 12 8

5 96

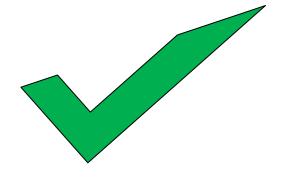
5 96 7

5 103

515

Our Roadmap

Stack



Queue

Stack vs. Queue

Queue

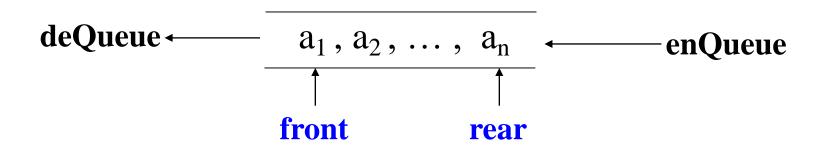
- A queue is a sequence in which:
 - items are added at the rear and removed from the front
 - You can only access the item that is currently at the front
- Queue Analogy





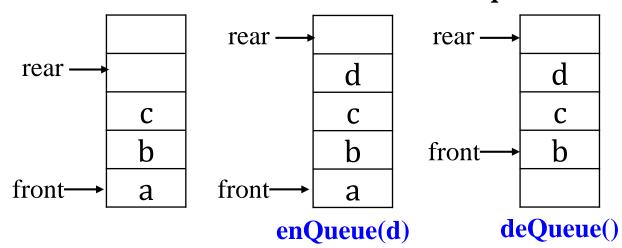
Queue

- First In First Out (FIFO)
 - Items access constrained
- Major elements
 - front: the first element in the queue (remove)
 - rear: the last element in the queue (add)
- Illustration



Queue Operations

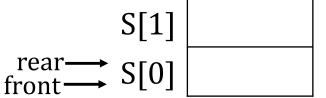
- enQueue: add an item at the rear of the queue
- deQueue: remove the item at the front of the queue
- front: get the item at the front of the queue, but do not remove it
- isEmpty: test if the queue is empty
- isFull: test the queue is full
- clear: clear the queue, set it as empty queue
- size: return the current size of the queue



Implementation of Queue

- Array based Queue
 - MAX_SIZE = n // the max size of stack
 - front = 0 // the current front
 - rear = 0 // the current rear
 - Array S with n elements
- Example
 - ϕ MAX_SIZE = 5
 - front = 0
 - $\Rightarrow rear = 0$
 - Array S

S[4]	
S[3]	
S[2]	



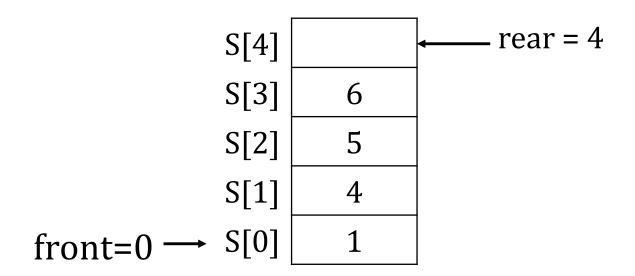
S

enQueue Operator

enQueue(item):

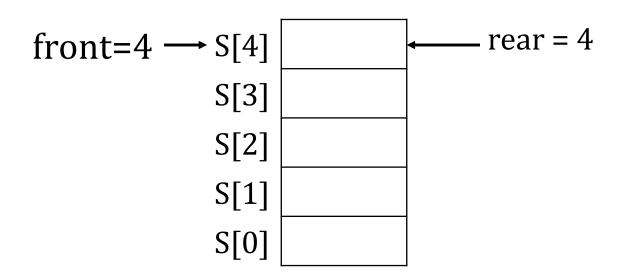
```
    if(rear < MAXSIZE)</li>
    S[rear] = item
    rear++
```

- 3. else
- 4. Queue is FULL, no enQueue
- enQueue (1), enQueue(4), enQueue(5), enQueue(6)



deQueue Operator

- deQueue():
 - 1. if(front < rear)</pre>
 - 2. front++
 - 3. else
 - 4. Queue is empty, no deQueue
- deQueue(), deQueue(), deQueue()



enQueue and deQueue

enQueue(9)

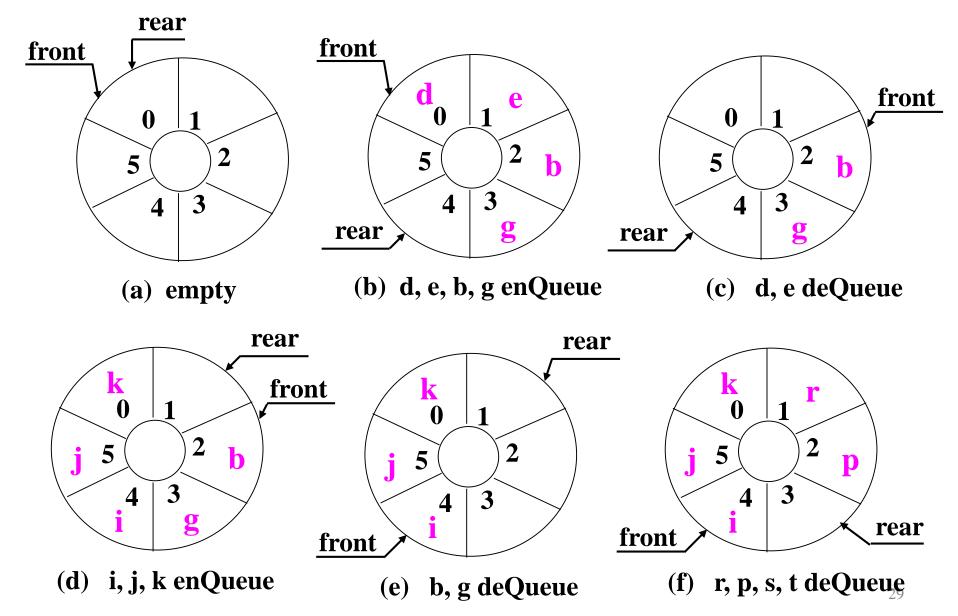
- enQueue(10)
 - ⋄ rear >= MAXSIZE
 - Queue is FULL!!!
 - Wrong OVERFLOW
 - \bullet S[0] to S[3] is empty?
 - How to address it?

rear = 5

- front= $4 \longrightarrow S[4]$ 9
 - S[3]
 - S[2]
 - S[1]
 - S[0]

_ rear = 5

Ring Queue



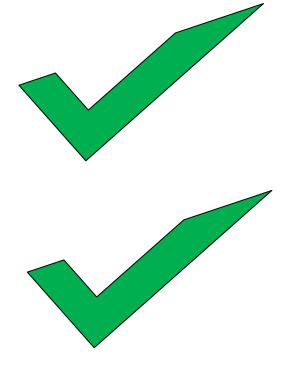
Application of Queues

- First-in first-out (FIFO) inventory control
 - OS scheduling: processes, print jobs, packets, etc.
 - Breadth-first traversal of a graph or level-order traversal of a binary tree (more on these later)
- Real applications
 - iTunes playlist.
 - Data buffers (iPod, TiVo).
 - Asynchronous data transfer (file IO, pipes, sockets).
 - Dispensing requests on a shared resource (printer, processor)

Our Roadmap

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Queue

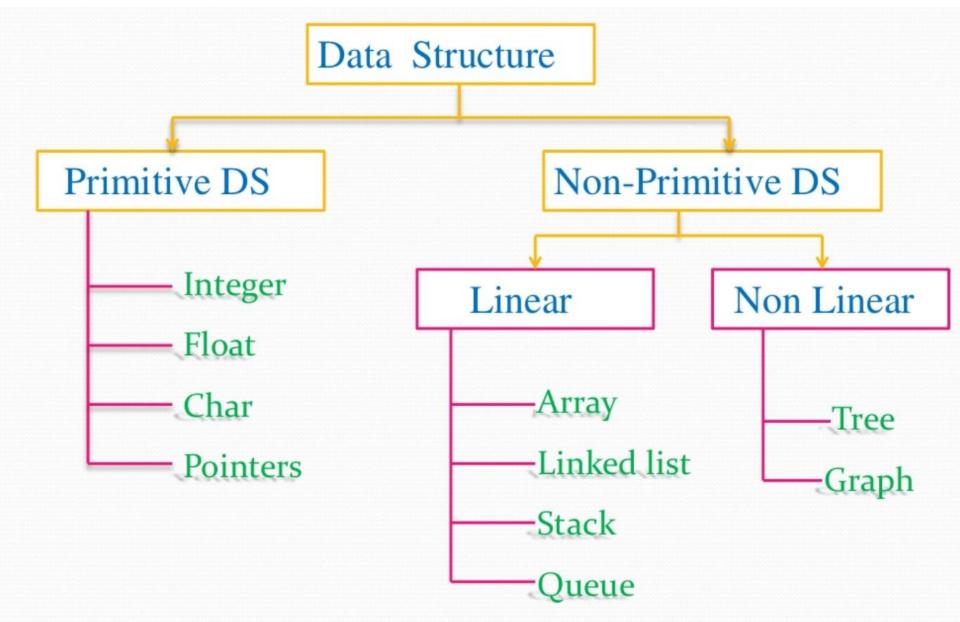


Stack vs. Queue

Stack VS. Queue

	Stack	Queue
In-Out	FILO	FIFO
Application	function runtime	OS scheduling
Operations	push pop	enQueue, deQueue
Ops Time Complexity	0(1)	0(1)
Implementa tion	Array-based, Linked-based	Array-based, Linked-based

Data Structure

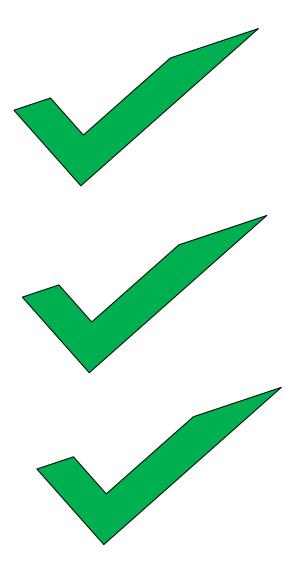


Our Roadmap

Stack

Queue

Stack vs. Queue



Thank You!