

Lab 9 – OJ

Divide and Conquer (p2)

CS208 Algorithm Design and Analysis
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Q1: FFT

Description

Given a polynomial of degree n called $F(x) = a_0 + a_1x + \cdots + a_nx^n$, and a polynomial of degree m called $G(x) = b_0 + b_1x^1 + \cdots + b_mx^m$.

Calculate the coefficients of the resulting polynomial by convolving $F(x)$ and $G(x)$, $F(x) * G(x) = c_0 + c_1x^1 + c_2x^2 + \cdots + c_{n+m}x^{n+m}$.

Input Format

The first line contains two integers n, m .

The second line contains $n + 1$ integers, representing the coefficients of $F(x)$ from low to high.

The third line contains $m + 1$ integers, representing the coefficients of $G(x)$ from low to high.

Output Format

One line with $n + m + 1$ integers, representing the coefficients of $F(x) * G(x)$ from low to high.

Q1: FFT

Output Format

One line with $n + m + 1$ integers, representing the coefficients of $F(x) * G(x)$ from low to high.

Sample Input 1

```
1 2
2 5
2 4 3
```

Sample Output 1

```
4 18 26 15
```

Explanation

$$(5x + 2)(3x^2 + 4x + 2) = 15x^3 + 20x^2 + 10x + 6x^2 + 8x + 4 = 15x^3 + 26x^2 + 18x + 4$$

So the output will be 4 18 26 15.

Question: How to deal
with the unequal lengths?

$$m \neq n$$

Pad to equal lengths (of 2
to the power of some
number)

Q1: FFT

Can add ω_n to the argument

RECURSIVE-FFT(a)

```

1   $n = a.length$ 
2  if  $n == 1$ 
3      return  $a$ 
4   $\omega_n = e^{2\pi i/n}$ 
5   $\omega = 1$ 
6   $a^{[0]} = (a_0, a_2, \dots, a_{n-2})$ 
7   $a^{[1]} = (a_1, a_3, \dots, a_{n-1})$ 
8   $y^{[0]} = \text{RECURSIVE-FFT}(a^{[0]})$ 
9   $y^{[1]} = \text{RECURSIVE-FFT}(a^{[1]})$ 
10 for  $k = 0$  to  $n/2 - 1$ 
11      $y_k = y_k^{[0]} + \omega y_k^{[1]}$ 
12      $y_{k+(n/2)} = y_k^{[0]} - \omega y_k^{[1]}$ 
13      $\omega = \omega \omega_n$ 
14 return  $y$ 
```

- Let $F_A = \text{FFT}(A, \omega_n)$ // time $O(n \log n)$
- Let $F_B = \text{FFT}(B, \omega_n)$ // time $O(n \log n)$
- For $i=1$ to m , let $F_C[i] = F_A[i] * F_B[i]$ // time $O(n)$
- Output $C = 1/m * \text{FFT}(F_C, \omega_n^{-1})$. // time $O(n \log n)$

Handle complex numbers:

<https://introcs.cs.princeton.edu/java/32class/Complex.java.html>

Q2: Urban Construction

Sjkmmost persuaded Justin to force his citizens to trip with zipline. However, the citizens are not strong enough that they often fell from the zipline. Therefore, they decided to provide the citizens with some cable cars.



↑ zipline

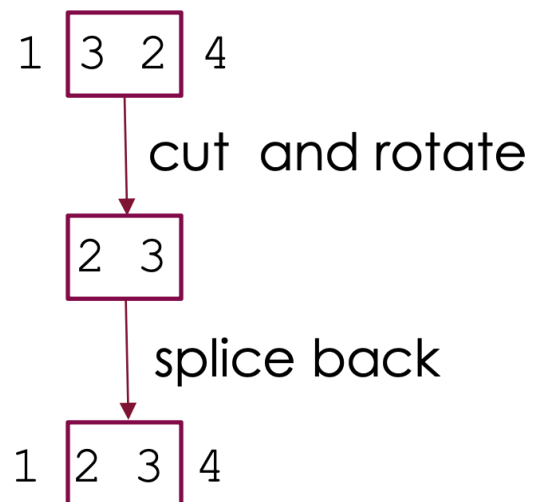
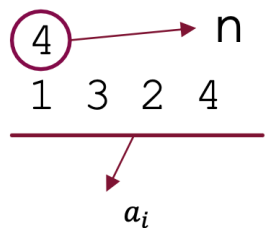
There are n cable cars on the cable, each has an index a_i . Sijmost is trying to put them in order.

You can spend c *justin* (a type of currency) to cut down a segment of rope of length c , rotate it and splice back. That is, he can spent $r - l + 1$ *justin* to rotate the cable cars in an interval.

Sjkmmost has a budget of 2×10^7 *justin*. He should sort the cable cars with some operations with a total cost no more than 2×10^7 *justin*. Can you help him?

Q2: Urban Construction

Sample Input1:



spend $3 - 2 + 1 = 2$ justin

Total cost: $2 \leq 2 \times 10^7$

Left and right ends of the rotation:

-1 -1 denoting end of output:

An empty line:

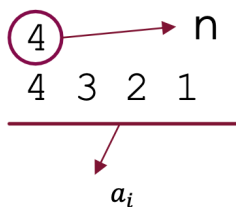
Sample output 1:

2 3

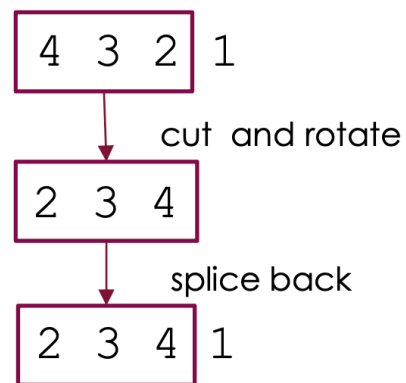
-1 -1

Q2: Urban Construction

Sample Input2:

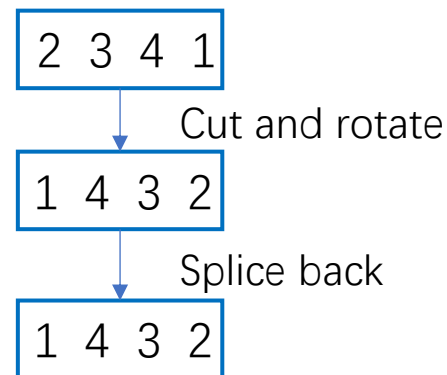


Step 1



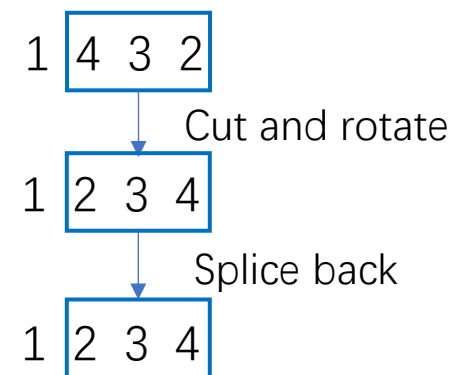
spend $3 - 1 + 1 = 3$ justin

Step 2



Spend $4 - 1 + 1 = 4$

Step 3



Spend $4 - 2 + 1 = 3$

Sample output 2:

1 3
1 4
2 4
-1 -1

Total cost: $3 + 4 + 3 = 10 < 2 \times 10^7$