

# Assignment One

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## Direction:

Please answer all the questions below and hand in your answers before the due day. All work, must be handed in **on time**.

## Due day:

April. 12, 2021

Please hand it in by the class time.

## Questions:

1. For each of the following functions, indicate how much the function's value will change if its argument is increased fourfold.

a.  $\log_2 n$    b.  $\sqrt{n}$    c.  $n$    d.  $n^2$    e.  $n^3$    f.  $2^n$

1. a.  $\log_2 n$  answer: let  $f(n) = \log_2 n$ , then  $f(4n) = \log_2(4n) = 2 + \log_2 n$

b.  $\sqrt{n}$  answer: let  $f(n) = \sqrt{n}$ , then  $f(4n) = \sqrt{4n} = 2\sqrt{n}$

c.  $n$  answer: let  $f(n) = n$ , then  $f(4n) = 4n$

d.  $n^2$  answer: let  $f(n) = n^2$ , then  $f(4n) = (4n)^2 = 16n^2$

e.  $n^3$  answer: let  $f(n) = n^3$ , then  $f(4n) = (4n)^3 = 64n^3$

f.  $2^n$  answer: let  $f(n) = 2^n$ , then  $f(4n) = 2^{(4n)} = 16 \cdot 2^n$

2. Prove (by using the definitions of the notations involved) or disprove (by giving a specific counterexample) the following assertions.

a. If  $t(n) \in O(g(n))$ , then  $g(n) \in \Omega(t(n))$ .

b.  $\Theta(\alpha g(n)) = \Theta(g(n))$ , where  $\alpha > 0$ .

c.  $\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$ .

d. For any two nonnegative functions  $t(n)$  and  $g(n)$  defined on the set of nonnegative integers, either  $t(n) \in O(g(n))$ , or  $t(n) \in \Omega(g(n))$ , or both.

2. a. if  $t(n) \in O(g(n))$ , then  $g(n) \in \Omega(t(n))$

proof:  $t(n) \in O(g(n)) \Leftrightarrow \exists c > 0, \forall n \geq n_0, t(n) \leq c g(n)$

$$\Leftrightarrow \exists c > 0, \forall n \geq n_0, g(n) \geq \frac{1}{c} t(n)$$

$$\text{let } c_2 = \frac{1}{c_1}, \exists c_2 > 0, \forall n \geq n_0, g(n) \geq c_2 t(n)$$

$$\Leftrightarrow g(n) \in \Omega(t(n))$$

b.  $\Theta(\alpha g(n)) = \Theta(g(n))$ , where  $\alpha > 0$

proof:  $\lim_{n \rightarrow \infty} \frac{\alpha g(n)}{g(n)} = \alpha > 0$ , then  $\alpha g(n) \in \Theta(g(n))$

$$\lim_{n \rightarrow \infty} \frac{g(n)}{\alpha g(n)} = \frac{1}{\alpha} > 0, \text{ then } g(n) \in \Theta(\alpha g(n))$$

$$\therefore \Theta(g(n)) = \Theta(\alpha g(n))$$

c.  $\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$

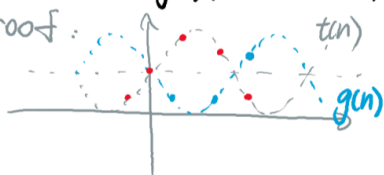
proof: let  $t(n) \in O(g(n)) \cap \Omega(g(n))$ , then  $\exists c_1, c_2 > 0, \forall n \geq n_0$ ,  
 $c_1 g(n) \leq t(n) \leq c_2 g(n)$

$$\therefore t(n) \in \Theta(g(n))$$

$$\therefore \Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$$

d.  $t(n) \geq 0, g(n) \geq 0, n \in \mathbb{N}^*$ , either  $t(n) \in O(g(n))$ , or  $t(n) \in \Omega(g(n))$ , or both.

disproof:



$$\text{let } t(n) = \sin x + 1, g(n) = -\sin x + 1$$

$$\text{then } t(n) \notin O(g(n))$$

$$t(n) \notin \Omega(g(n))$$

$$t(n) \notin O(g(n)) \cap \Omega(g(n))$$

3. Solve the following recurrence relations.

a.  $x(n) = 3x(n-1)$  for  $n > 1, x(1) = 4$

b.  $x(n) = x(n-1) + n$  for  $n > 0, x(0) = 0$

c.  $x(n) = x(n/2) + n$  for  $n > 1, x(1) = 1$  (solve for  $n = 2^k$ )

3. a.  $x(n) = 3x(n-1)$ , for  $n > 1, x(1) = 4$

$$\frac{x(n)}{x(n-1)} = 3 \Rightarrow \frac{x(n)}{x(n-1)} \cdot \frac{x(n-1)}{x(n-2)} \cdots \frac{x(2)}{x(1)} = 3^{n-1}$$

$$\therefore x(n) = 4 \cdot 3^{n-1}$$

b.  $x(n) = x(n-1) + n$  for  $n > 0, x(0) = 0$

$$x(n) - x(n-1) = n$$

$$x(n-1) - x(n-2) = n-1$$

$$\vdots$$

$$x(2) - x(1) = 2$$

$$x(1) - x(0) = 1$$

$$\Rightarrow x(n) - x(0) = \frac{(1+n)n}{2}$$

$$\therefore x(n) = \frac{n(n+1)}{2}$$

C.  $\chi(n) = \chi(n/2) + n$  for  $n > 1$ ,  $\chi(1) = 1$  (solve for  $n = 2^k$ )

$$\chi(n) = \chi\left(\frac{n}{2}\right) + n$$

$$\chi(2^k) = \chi(2^{k-1}) + 2^k \quad (k \in \mathbb{Z}^+) \quad \chi(1) = \chi(2^0) = 1$$

$$\left. \begin{array}{l} \chi(2^k) - \chi(2^{k-1}) = 2^k \\ \chi(2^{k-1}) - \chi(2^{k-2}) = 2^{k-1} \\ \vdots \\ \chi(2) - \chi(2^0) = 2 \end{array} \right\} \Rightarrow \chi(2^k) - \chi(2^0) = \frac{2(1-2^k)}{1-2}$$

$$\Rightarrow \chi(n) = \chi(2^k) = 2^{k+1} - 1$$