

CSE 101: Introduction to Computational and Algorithmic Thinking

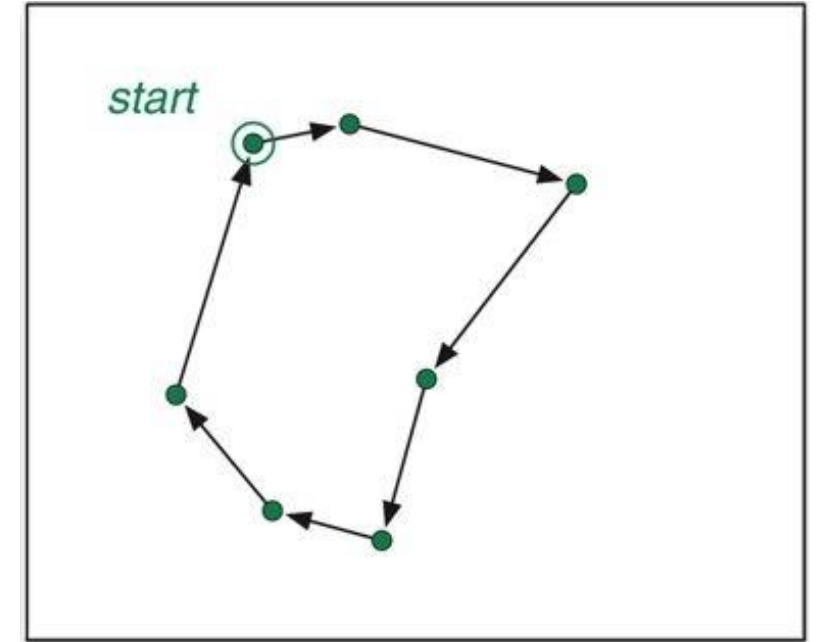
Unit 12: The Limits of Computation

An Intractable Problem

- In computer science we have a broad class of problems which are generally called **intractable**, meaning that there are no known efficient solutions for these problems
- One such famous, intractable problem is called the **Traveling Salesman Problem**, or **TSP** for short
- Imagine a salesman who must travel to n different cities
- He doesn't care about the order in which he visits the cities, as long as he visits each city exactly once, returning to his starting city at the end
- Between every pair of cities we have a travel time or travel cost or some other metric

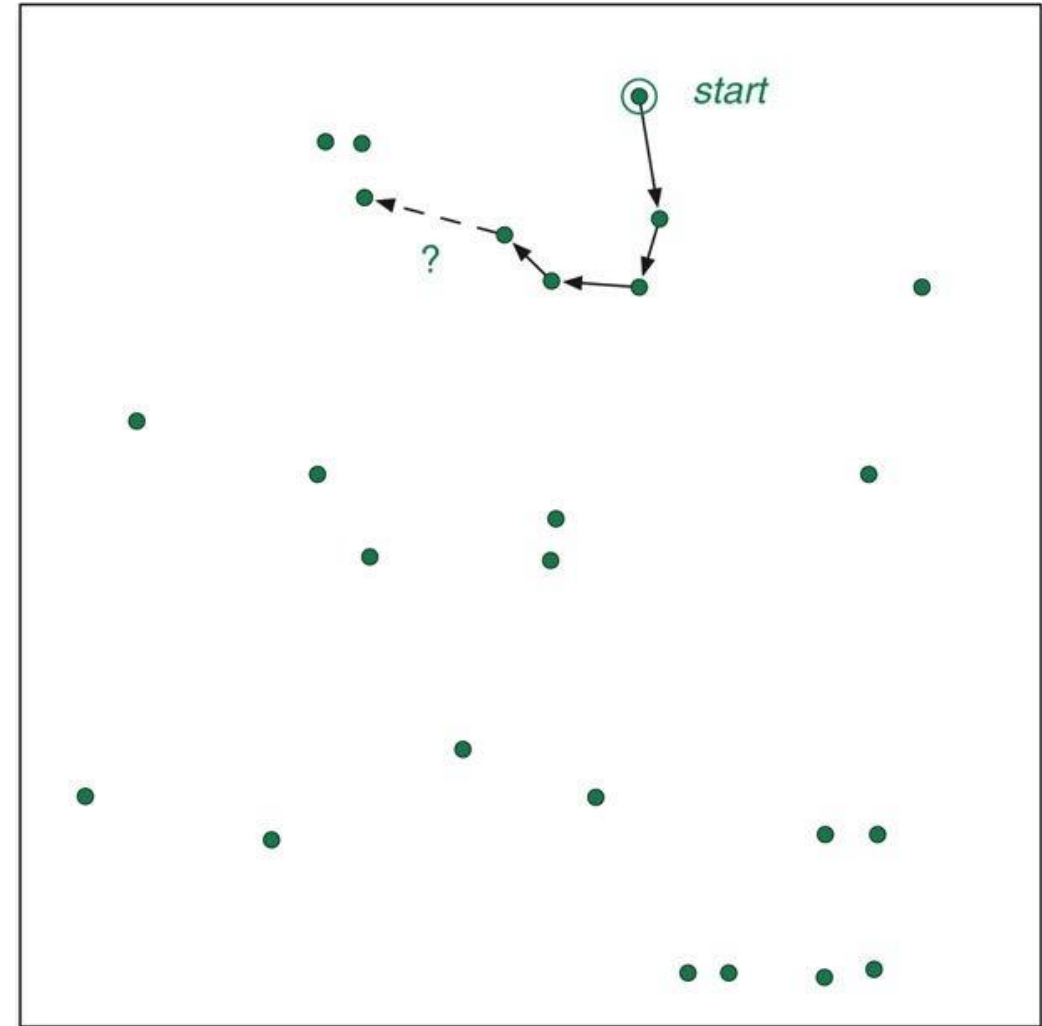
An Intractable Problem

- The goal is to minimize the total cost of the trip while still visiting every city exactly once (except the starting city, which we visit twice)
- TSP arises in many real-world situations, such as routing delivery trucks or school buses and even in CPU manufacturing (to find the fastest way to make the connections between components)
- How can we solve this problem?
- It turns out that simply visiting the nearest unvisited city will not provide the optimal solution, so we need to find another strategy



An Intractable Problem

- As an example, the dashed line in the figure is actually not part of the best tour



Exhaustive Search

- One simple strategy is to list all of the possible tours through the n cities and pick the tour with the smallest cost
- We *exhaustively* list all possible tours and search through them for the one with smallest cost
- There are approximately $(n - 1)!$ such tours
- For very small values of n the number of tours to check is not that large
- But the magnitude of $(n - 1)!$ becomes very large very quickly as n gets bigger
- For $n = 25$ cities there are approximately 3×10^{23} tours
- That's 300,000,000,000,000,000,000,000 possibilities!

Evolutionary Algorithms

- An algorithm based on exhaustive search is usually impractical
- In this Unit we will look at an **evolutionary algorithm**, one based on concepts from biological evolution
- The algorithm creates a set of random tours and picks a few with lowest costs (the “fittest” tours)
- The worst tours are discarded, and new tours are developed based on the “survivors” by “mutating” the survivors a little
- Once again the best tours are retained, the worst are discarded, and another round of “evolution” occurs to search for the best possible solution

Maps and Tours

- To make the problem concrete we can use some driving distances between several cities in Ireland, as organized in the **matrix** shown below:



	Belfast	Cork	Dublin	Galway	Limerick
Belfast	—				
Cork	425	—			
Dublin	167	257	—		
Galway	306	209	219	—	
Limerick	323	105	198	105	—

Maps and Tours

- **TSP.py** has a **Map** class that defines the distance between cities
- We can load a file to get the distances:

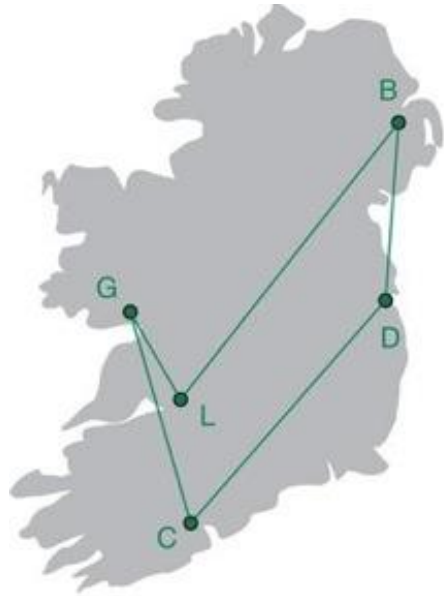
```
from TSP import *  
  
m = Map('tsp/ireland.txt')
```

- A call to **m.display()** generates this output of distances between cities in Ireland:

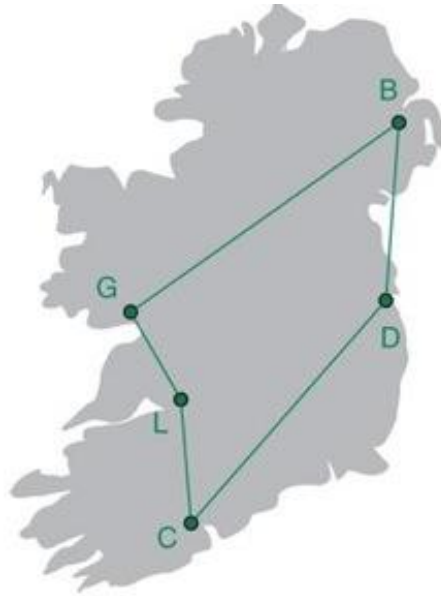
	belfast	cork	dublin	galway	limerick
belfast	0.00				
cork	425.00	0.00			
dublin	167.00	257.00	0.00		
galway	306.00	209.00	219.00	0.00	
limerick	323.00	105.00	198.00	105.00	0.00

Exhaustive Search

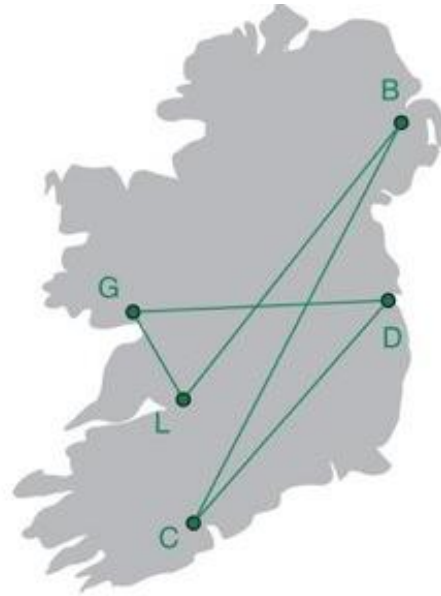
- In mathematics, a **permutation** is an ordering of the items in a list or the characters in a string
- To solve TSP we could list all permutations (tours) of the n cities and pick the tour with lowest cost
- Three example permutations of five cities:



DCGLB



DCLGB



BCDGL

Exhaustive Search

- To see an example of how this might work, **TSP.py** contains a **generator** function that will create a list of all permutations of a string
- A generator is a function that produces a sequence of values *on demand*, instead of all at once

```
from TSP import each_permutation
for s in each_permutation('ABC') :
    print(s)
```

- Since there are 3 letters in “ABC”, the function produces $3! = 6$ permutations

Exhaustive Search

- The function **xsearch** will take a **Map** object and perform an exhaustive search for the best tour

```
def xsearch(m):  
    best = m.make_tour()  
    for t in m.each_tour():  
        if t.cost() < best.cost():  
            best = t  
    return best
```

- Example: for the five-city map of Ireland (**m**), **xsearch(m)** returns this **Tour** object: `['belfast', 'galway', 'limerick', 'cork', 'dublin']`
940.000 (940 is the cost of the tour)
- See **TSP_xsearch.py**

Exhaustive Search

- Earlier in the course we studied **big-Oh notation**, a mathematical technique for analyzing algorithm efficiency
- The exhaustive algorithm for finding the optimal solution to TSP has **factorial time complexity**, meaning the algorithm takes a very long time to find the solution

Random Search

- Evolutionary algorithms are based in part on randomization of data
- For TSP, imagine if we generate a few random permutations of the cities and make small adjustments to the tours until we find the optimal one
- As a first step towards implementing algorithm we need some code that will generate random permutations and pick the one with the lowest cost
- The steps to create one permutation would be this:
 1. **`a = m.cities()`** to get a list of the cities in the map
 2. **`permute(a)`** to shuffle the cities
 3. **`t = m.make_tour(a)`** to compute the tour cost

Random Search

- These three steps can be performed by **make_tour** itself if we pass it the value '**random**' as its argument
- Consider the function **rsearch**, which generates a set of n random tours and picks the best one:

```
def rsearch(m, n):  
    best = m.make_tour('random')  
    for i in range(n-1):  
        t = m.make_tour('random')  
        if t.cost() < best.cost():  
            best = t  
    return best
```

- See **TSP_rsearch.py**

Random Search

- An example of how to call this function:

```
m = Map('tsp/ireland.txt')  
rsearch(m, 10) # generate 10 random tours
```

- Possible return value:

```
['belfast', 'cork', 'limerick', 'galway', 'dublin']  
1021.0
```

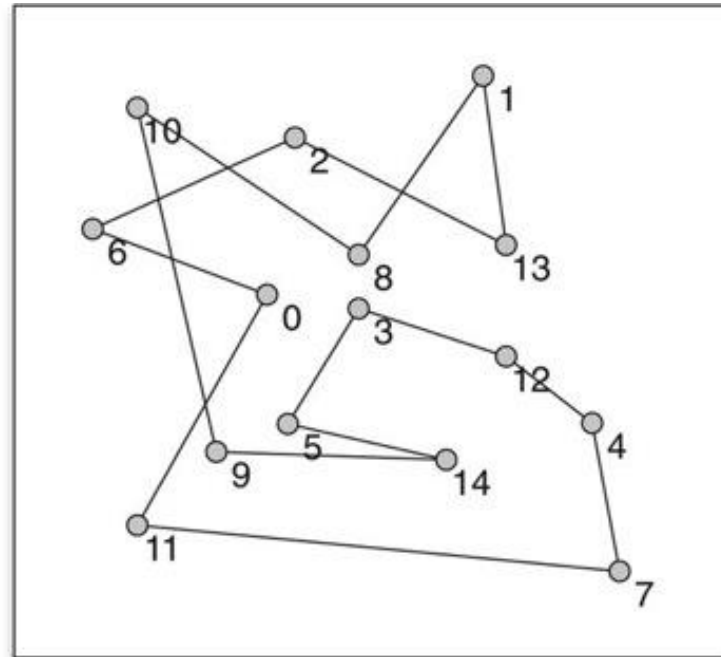
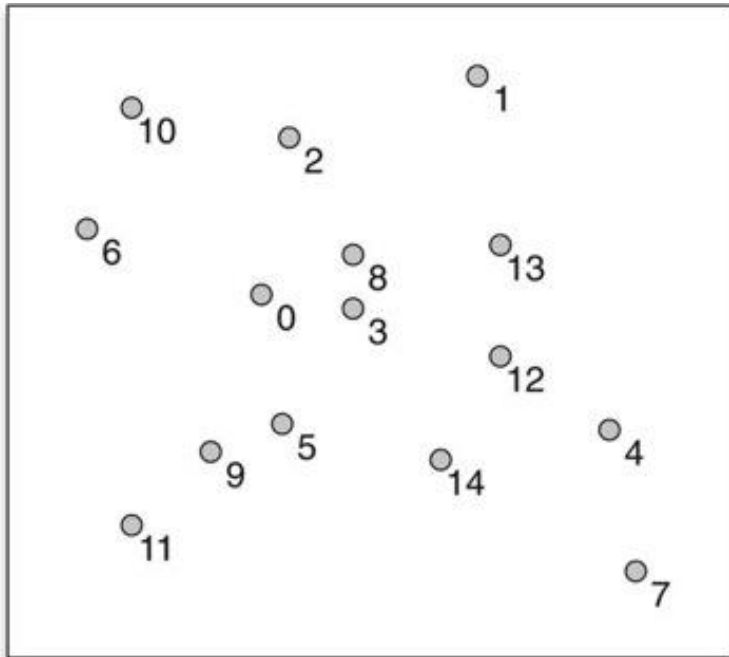
- Note that the tour generated may not be the optimal one (as in this case)
- For large numbers of cities, the exhaustive search approach is impractical
- We could try random search with a value of n like 1000, which will give an approximate solution pretty quickly

Random Search

- Another way to test the **rsearch** function is to have the **Map** class generate a map with randomly placed cities:

```
m = Map(15) # cities numbered 0 to 14
```

- **make_tour** would generate output like the figure on the right



Point Mutations

- The evolutionary algorithm we seek to implement is “biologically inspired” (hence the name "evolutionary")
- Specifically, it is a kind of **genetic algorithm**, wherein our list of city names serve as the “DNA” of a single tour
- Each slight change we will make to a tour is called a **point mutation**
- A point mutation for us will be to exchange a city with its neighbor in a tour
- We start with a tour:

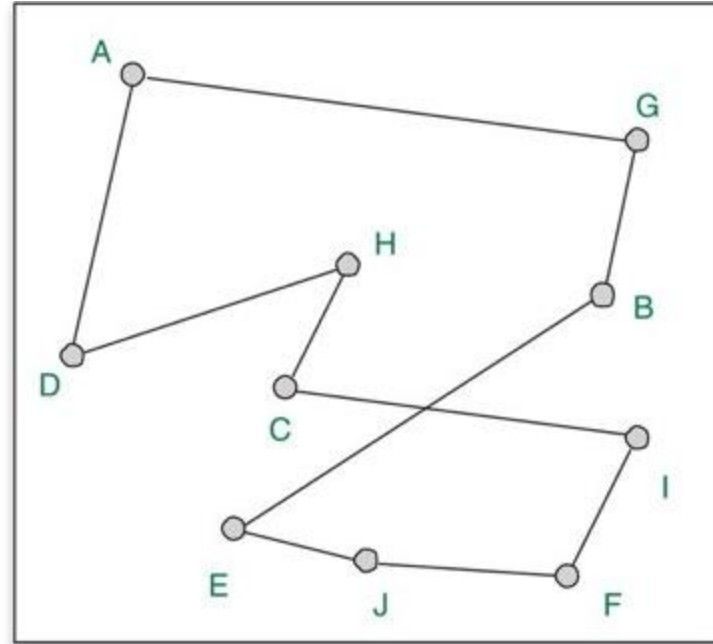
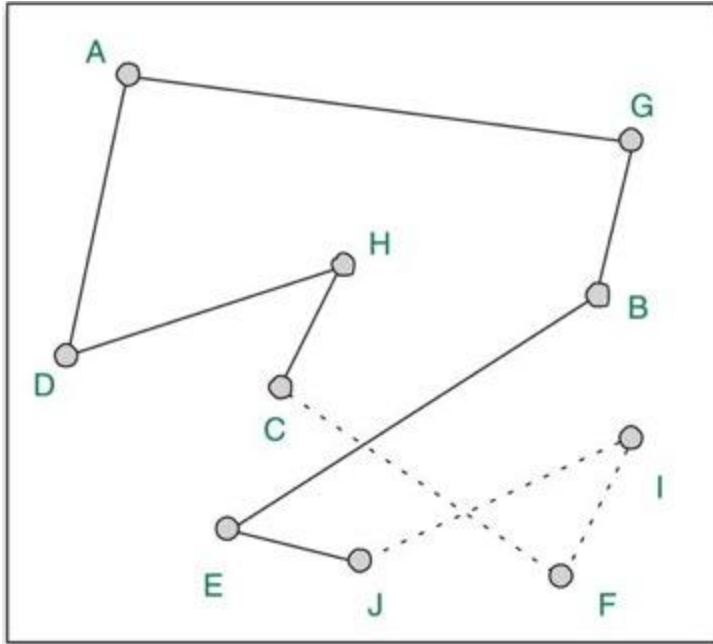
```
m = Map('tsp/ireland.txt')  
t = m.make_tour()
```

Point Mutations

- To mutate the cities we can call the **mutate** method and indicate at what index we want to make a mutation: **t.mutate(2)**
- For example, suppose we start with this tour:
['belfast', 'cork', 'dublin', 'galway', 'limerick'] 329.0
- Perhaps it will mutate into this:
['belfast', 'cork', 'galway', 'dublin', 'limerick'] 1374.0
- We see that the cities at indexes 2 and 3 have been exchanged, resulting in a different tour with a higher cost

Point Mutations

- Another example, with randomized input:



- Exchanging cities I and F in the tour shortened the tour length by “untwisting” it in one place

The Genetic Algorithm

- We can now combine the ideas of random generation and mutation to derive the genetic algorithm
- To select the best tour(s) that will “evolve” into the next generation of tours, we need to have a “population” of tours to draw from
- We will write the function **esearch**, which takes these arguments:
 - a map, which is a list of cities in the tour
 - the number of generations to simulate
 - the number of tours in the population
- A helper function **init_population** will handle generating a list of tours (the “population”)

The Genetic Algorithm

```
def init_population(m, popsize):  
    pop = [m.make_tour('random')  
           for i in range(popsize)]  
    return pop
```

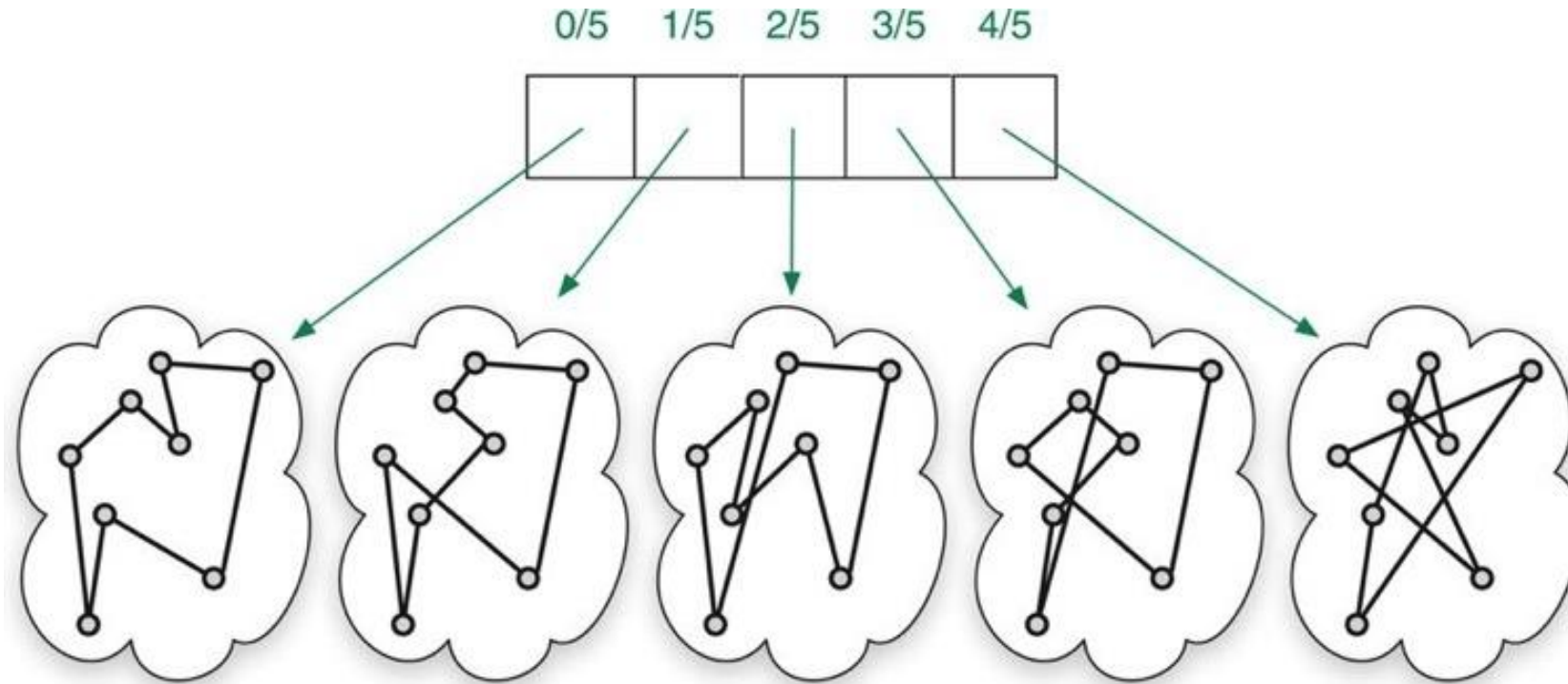
- The function generates **popsize** random tours from the map **m**
- Note that a list comprehension is being used
- The list of tours is then sorted by tour length, putting the shortest tour is at the front of the **pop** list
- A different helper function called **select_survivors** performs “random selection,” similar to how the process is described for biological evolution
- See **TSP_esearch.py**

The Genetic Algorithm

- The **select_survivors** function removes tours from the list with a certain probability
 - Remember that the shortest tours are at the front of the list
 - Specifically, the tour in index i is deleted with probability i/p , where p is the size of the population
 - Note that the best tour has no chance of being deleted because its probability is $0/p$
- This gives some of the worst paths a chance to “survive” and perhaps evolve into better solutions
- On average, half of the tours will be “deleted” by being replaced with the value **None** in the list

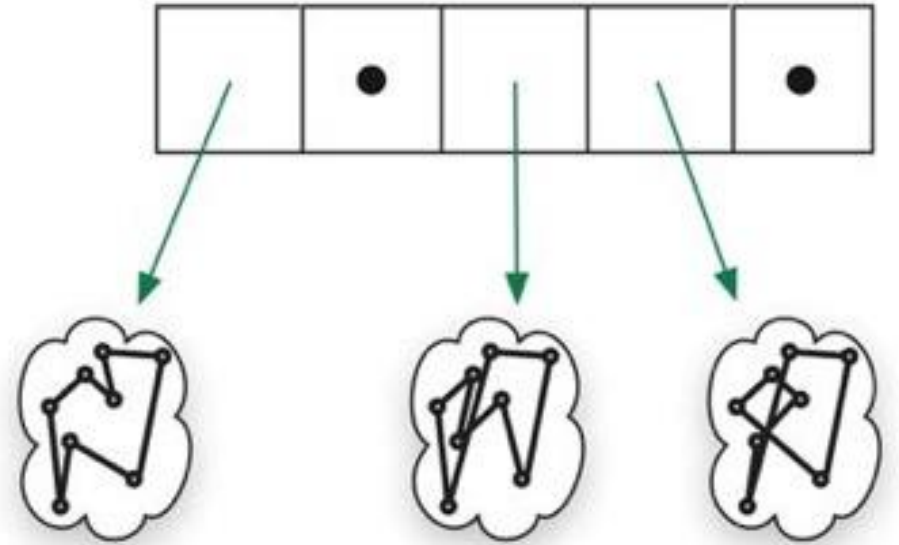
The Genetic Algorithm

- The figure below shows the sorted tours before **select_survivors** has been called



The Genetic Algorithm

```
import random
def select_survivors(population):
    n = len(population)
    for i in range(1,n):
        if random.random() < i/n:
            population[i] = None
```

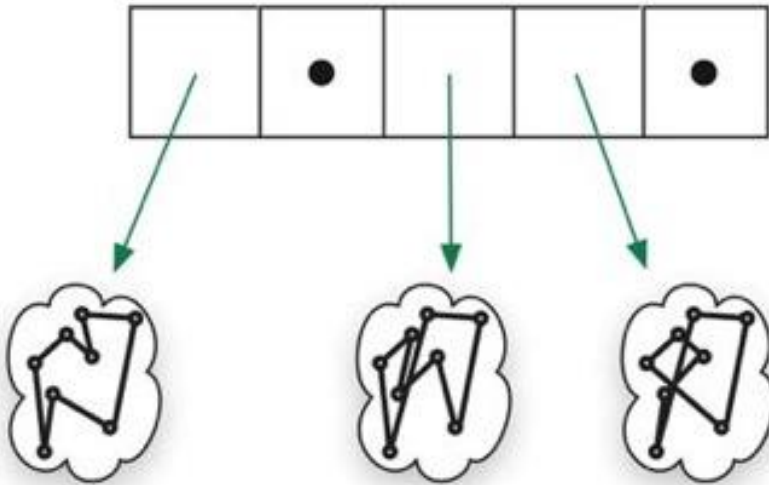


- This algorithm will introduce **None** objects, as shown in the example above
- See `TSP_eseach.py`

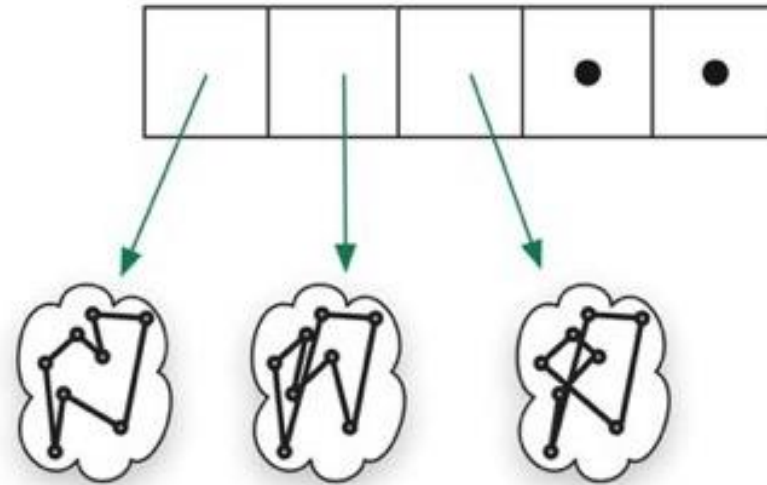
The Genetic Algorithm

- A helper function named **compact_population** moves the survivors to the front of the lists and moves None objects to the end
- The figure below shows the list of **Tour** objects before and after a call to **compact_population**:

before:



after:



The Genetic Algorithm

```
def compact_population(population):  
    d = 0  
    for i in range(1, len(population)):  
        if population[i] is None:  
            d += 1  
        elif d > 0:  
            population[i-d], population[i] =  
                population[i], population[i-d]  
    return len(population) - d
```

- **d** is a running count of how many **None** objects there are
- Inside the for-loop **d** also tells us how many positions *backwards* we to move an element so that the **None** objects are shuffled to the end.
- See **TSP_esearch.py**

The Genetic Algorithm

- To see why this works, let's call the function with a list of integers and None objects

```
a = [1, 2, None, 3, 4, None, None, 5]
```

```
compact_population(a)
```

- Here is how the list is updated with each swap operation:

```
[1, 2, None, 3, 4, None, None, 5]
```

- `d = 1`, so move the 3 back 1 position:

```
[1, 2, 3, None, 4, None, None, 5]
```

- `d = 1`, so move the 4 back 1 position:

```
[1, 2, 3, 4, None, None, None, 5]
```

- `d = 3`, so move the 5 back 3 positions:

```
[1, 2, 3, 4, 5, None, None, None]
```

The Genetic Algorithm

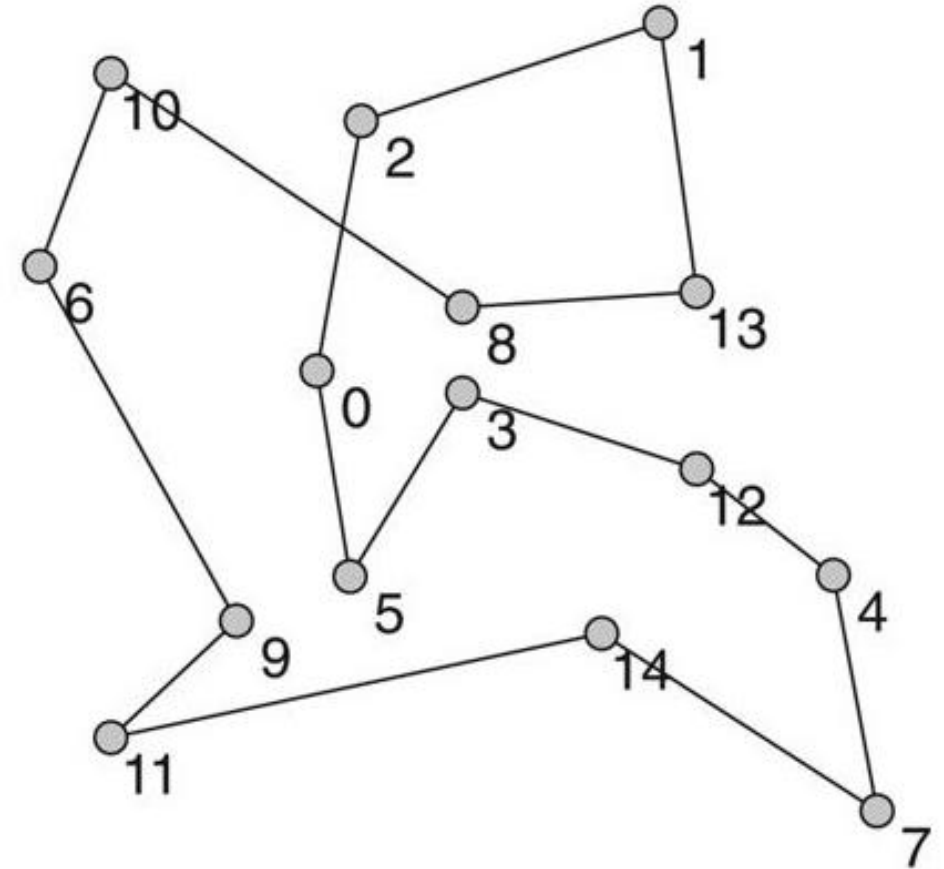
- A helper function `rebuild_population` replaces all the **None** objects by new tours
 - See `TSP.py` for the implementation if you are interested in the details, but it's pretty heavy stuff!

```
def esearch(m, ngen, popsize):  
    pop = init_population(m, popsize)  
    for i in range(ngen):  
        pop.sort(key = Tour.cost)  
        select_survivors(pop)  
        ns = compact_population(pop)  
        rebuild_population(pop, m, ns)  
    return pop[0]
```

See `TSP_esearch.py`

Crossovers

- While **esearch** is running, it can generate tours with twists or kinks that are hard to remove by exchanging only two neighboring cities
- The figure on the right shows an example (2, 1, 13, 8)

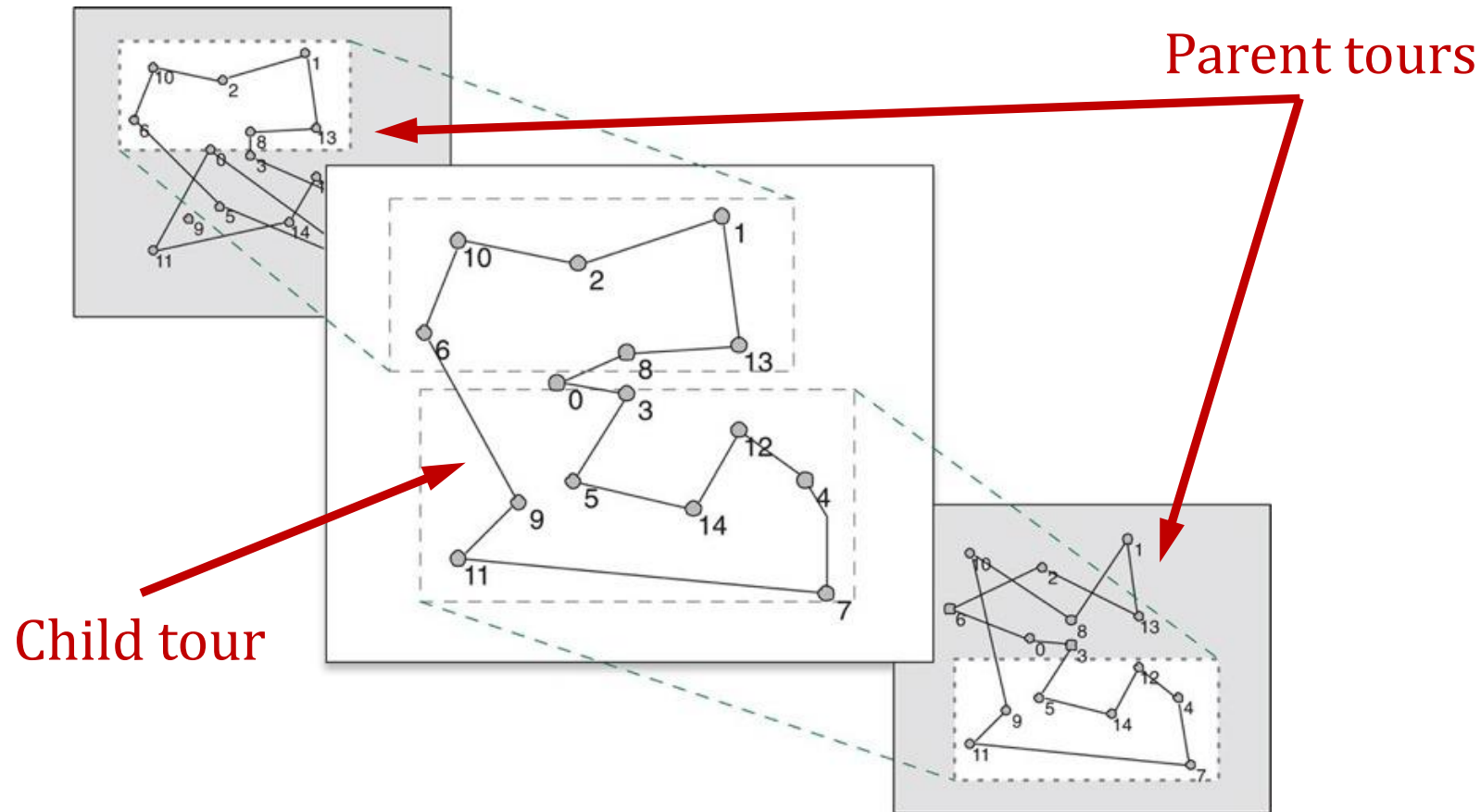


Crossovers

- In genetics, a **crossover mutation** occurs when two chromosomes (very long strands of DNA) break apart
 - And when they come back together there is some mixing, and a new combination is formed
- The idea for us in solving TSP is to take two existing tours and splice them together to form a new tour
- Ideally, the new tour will combine the best portions of its “parent” tours to generate the “child” tour
- We perform a kind of “cut-and-paste” from the parent tours
- Part of the list of city names from one tour is appended to part of the list from a second tour, resulting in a third tour that has large pieces from each of the original tours

Crossovers

- The implementation of this idea is a bit complicated
- See **TSP.py** if you are interested in the details



Crossovers

- Now we have two kinds of mutations: point mutations and crossover mutations
- A modified version of the **rebuild_population** function chooses randomly between the two types for each new tour it adds to the list of tours
- The **esearch** function is otherwise unchanged

The Halting Problem

- Alan Turing in 1936 discovered the **Halting Problem**, which is an **unsolvable problem**
- We have seen in programming that it is possible to create an infinite loop
- In Python it would look like this:

```
while True:  
    # do something
```
- So a program that contained such a loop would never actually stop running. It would never halt (terminate).
- Infinite loops are (usually) bugs in code. Wouldn't it be nice if we could write a program to tell us if our code contains such bugs?

The Halting Problem

- Imagine we could write a program called **HaltChecker**
 - The **HaltChecker** program takes as its input the source code of another program (call it ProgramX)
 - **HaltChecker** examines **ProgramX**'s source and determines whether **ProgramX** will halt eventually
 - It prints either "**ProgramX** will halt" or "**ProgramX** will not halt"
- Let's say we had our infinite loop from the previous slide in a program called **InterestingProgram**

The Halting Problem

- Specifically, our **InterestingProgram** looks like this:
 if **HaltChecker** says "**InterestingProgram** will halt" then:
 while **True**:
 do something
 else:
 halt
- If **HaltChecker** determines that **InterestingProgram** will halt, then **InterestingProgram** starts executing an infinite loop
- Otherwise, **HaltChecker** determines that **InterestingProgram** will run forever, and so **InterestingProgram** terminates

The Halting Problem

```
if HaltChecker says "InterestingProgram will halt" then:  
    while True:  
        do something  
else:  
    halt
```

- **InterestingProgram** makes reference to itself
- **InterestingProgram** begins by telling **HaltChecker** to inspect **InterestingProgram** (i.e., itself)
- We know that **HaltChecker** will give one of two answers: “InterestingProgram will halt” or “InterestingProgram will not halt”

The Halting Problem

```
if HaltChecker says "InterestingProgram will halt" then:  
    while True:  
        do something  
else:  
    halt
```

- If **HaltChecker** says “InterestingProgram will halt”, then the if-statement in **InterestingProgram** will execute an infinite loop, and so **InterestingProgram** will not halt. (A contradiction!)
- Well, what if **HaltChecker** says “InterestingProgram will not halt”. Then the if-statement in **InterestingProgram** says that **InterestingProgram** will halt. (Also a contradiction!)

The Halting Problem

- This means that it is truly impossible to write a program that, for any possible input program, determine whether the input program will halt or not
- So **HaltChecker** is an “impossible algorithm” and cannot exist in the world as we understand it
- Now of course we could write a program like **HaltChecker** that would work some or most of the time, but not all of the time