# CSE 101: Introduction to Computational Thinking

# Unit 5: Divide-and-Conquer Algorithms; Recursion

# Divide and Conquer

- The strategy for the linear search and insertion sort algorithms was the same: iterate over every location in the list and perform some operation
- We will now look at a different strategy: **divide and conquer**
  - The idea: break a problem into smaller sub-problems and solve the smaller sub-problems
  - Sub-problems are chosen in such a way that their solutions can be combined to provide the solution to the original problem
- It may not seem like that big a deal, but the improvement can be dramatic, as we will see

# Example: Searching a Dictionary

- To get a general sense of how the divide-and-conquer strategy improves search, consider how people find information in a (physical) phone book or dictionary
- Suppose you want to find "janissary" in a dictionary
  - Open the book near the middle
  - The heading on the top left page is "kiwi", so move back a small number of pages
  - Here you find "hypotenuse", so move forward
  - Find "ichthyology", move forward again
- The number of pages you move gets smaller (or at least adjusts in response to the words you find)

# Example: Searching a Dictionary

- A more detailed specification of this process:
  1. The goal is to search for a word *w* in a region of the book.
  2. The initial region is the entire book.
  3. At each step, pick a word *x* in the middle of the current region.
  4. There are now two smaller regions: the part before *x* and the part after *x*.
  5. If *w* comes before *x*, repeat the search on the region before *x*. Otherwise, search the region following *x* (go back to step 3).
- Note: at first a “region” is a group of pages, but eventually a region is a set of words on a single page

# A Note About Data Organization

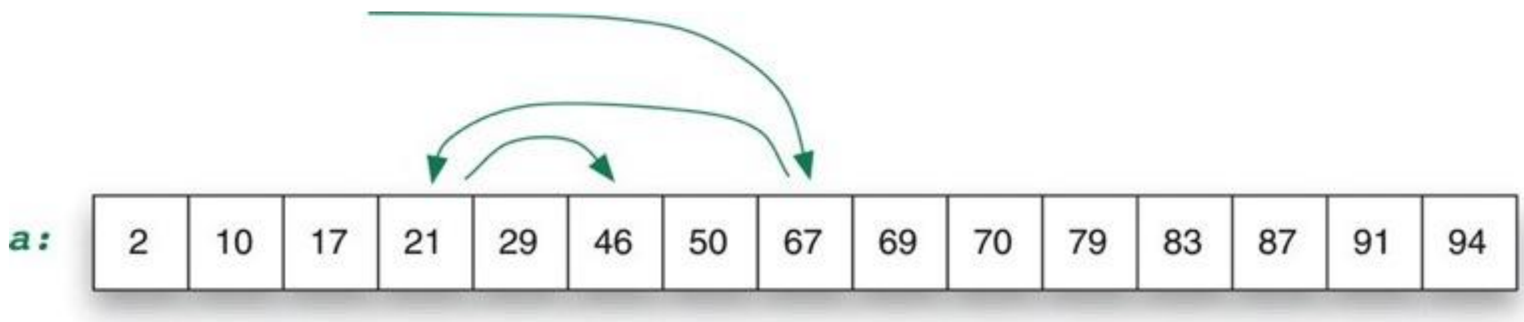
- An important note: an efficient search depends on having the data organized in some fashion
- If books in a library are scattered all over the place we have to do an iterative search
  - Start at one end of the room and progress toward the other
- If books are sorted or carefully cataloged we can try a more efficient method that exploits the sorted nature of the books

# Binary Search: Overview

- The **binary search** algorithm uses the divide-and-conquer strategy to search through a list
- The list must be sorted for this algorithm to work properly
  - The “zeroing in” strategy for looking up a word in the dictionary won't work if the words are not in alphabetical order
  - Similarly, binary search will not work for a list of values unless the list is sorted
- To search a list of *n* items, first look at the item in location $\frac{n}{2}$
- If this is the item we want, then the search has ended successfully
- Otherwise, search either the region from 1 to $\frac{n}{2} - 1$ or the region from $\frac{n}{2} + 1$ to $n$

# Binary Search: Example

- Example: searching for 46 in a sorted list of 15 numbers:




- Note how the search moves backward and forward, quickly finding the **target element**

# Binary Search: the Details

- The algorithm uses two variables to keep track of the boundaries of the region to search
  - `lower`: the index one position below the leftmost item in the region
  - `upper`: the index one position above the rightmost item in the region
- Initial values when searching a list of *n* items:

```
lower = -1
upper = n
```

# Binary Search: the Details

- The algorithm is based on an iteration (loop) that keeps making the search region smaller and smaller
  - The initial region is the complete list
  - The next one is either the upper half or lower half
  - The one after that is one quarter, then one eighth, then…
- The heart of the algorithm contains these operations:
  - Set **`mid`** to a location halfway between **`lower`** and **`upper`**:

```
mid = (lower + upper) // 2
```

  - If the item is at this location, then we're done:

```
if a[mid] == x:
    return mid
```

# Binary Search: the Details

- The heart of the algorithm (continued):
  - Otherwise, move one of the “brackets” to the current mid-point for the next iteration:

```
if x < a[mid]:
    upper = mid
else:
    lower = mid
```

# Binary Search: the Details

- Let's revisit our example from earlier. The star in the figure shows how the value of mid changes with each iteration


```
lower = -1
upper = 15
mid = 7
```

```
lower = -1
upper = 7
mid = 3
```

```
lower = 3
upper = 7
mid = 5
```

# Binary Search: the Details

- How do we handle the case when the target item is not in the list?
  - We have to add a condition that makes sure that **`lower`** is still to the left of **`upper`**
  - If the **`upper`** and **`lower`** pointers meet each other, this means that the search region has no elements in it – the search has failed
- We can now write the complete **`bsearch`** function, which returns:
  - The index of the target item in the list, when the search is successful, or
  - **`None`**, if the target item is not in the list

# Completed `bsearch()` Function

```
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

- See unit05/bsearch.py for fully commented code

# Trace Execution: `bsearch()`

```
a:     [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |

# Trace Execution: `bsearch()`

```
a:     [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index: 0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | -1 |

# Trace Execution: `bsearch()`

```
a:     [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index: 0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| `x` | 8 |
| `lower` | -1 |
| `upper` | 10 |

# Trace Execution: `bsearch()`

```
a:     [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index: 0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1: True
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | -1 |
| upper | 10 |

# Trace Execution: `bsearch()`

```
a:     [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index: 0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | -1 |
| upper | 10 |
| mid | 4 |

# Trace Execution: `bsearch()`

```
a:     [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x: False
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | -1 |
| upper | 10 |
| mid | 4 |
| a[mid] | 6 |

# Trace Execution: `bsearch()`

```
a:      [1, 2, 3, 5, 6, 8, 9, 11, 14, 17]
index:  0   1   2   3   4   5   6    7    8    9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:  False
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | -1 |
| upper | 10 |
| mid | 4 |
| a[mid] | 6 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11, 14, 17]
index: 0   1   2   3   4   5   6    7    8    9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 10 |
| mid | 4 |
| a[mid] | 6 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11, 14, 17]
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1: True
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 10 |
| mid | 4 |
| a[mid] | 6 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11, 14, 17]
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| `x` | 8 |
| `lower` | 4 |
| `upper` | 10 |
| `mid` | 7 |
| `a[mid]` | 11 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11, 14, 17]
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x: False
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 10 |
| mid | 7 |
| a[mid] | 11 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11, 14, 17]
index: 0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:  True
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 10 |
| mid | 7 |
| a[mid] | 11 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11],14, 17
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 7 |
| mid | 7 |
| a[mid] | 11 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11],14, 17
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
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    while upper > lower + 1: True
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 7 |
| mid | 7 |
| a[mid] | 11 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11],14, 17
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
| --- | --- |
| `x` | 8 |
| `lower` | 4 |
| `upper` | 7 |
| `mid` | 5 |
| `a[mid]` | 8 |

# Trace Execution: `bsearch()`

```
a:      1, 2, 3, 5,[6, 8, 9, 11],14, 17
index:  0  1  2  3  4  5  6   7   8   9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x: True
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 7 |
| mid | 5 |
| a[mid] | 8 |

# Trace Execution: `bsearch()`

```
a:       1,  2,  3,  5, [6,  8,  9,  11], 14,  17
index: 0   1   2   3   4   5   6    7    8    9
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

| Variable | Value |
|---|---|
| x | 8 |
| lower | 4 |
| upper | 7 |
| mid | 5 |
| a[mid] | 8 |

# Completed `bsearch()` Function

- In the bsearch.py file there is a function named **`print_bsearch_brackets`** that will let us visualize how the **`lower`** and **`upper`** pointers change as the search progresses
- The call to this function goes near the top of the loop
  - See the code on the next slide

# `bsearch()` with Print-outs

```
def bsearch(a, x):
    lower = -1
    upper = len(a)
    while upper > lower + 1:
        mid = (lower + upper) // 2
        print_bsearch_brackets(a, lower, mid,
                                           upper)
        if a[mid] == x:
            return mid
        if x < a[mid]:
            upper = mid
        else:
            lower = mid
    return None
```

# `bsearch()` Example

- list: **`[5, 12, 16, 40, 58, 62, 72, 84, 88, 90]`**
- target element: 72
- In the sample visualizations below, the brackets indicate the current search region, and * indicates the middle element

```
[5  12  16  40 *58  62  72  84  88  90]
 5  12  16  40  58 [62  72 *84  88  90]
 5  12  16  40  58 *62  72] 84  88  90
 5  12  16  40  58  62 [72] 84  88  90
```

- Result: 6

# `bsearch()` Example

- list: **`[5, 12, 16, 40, 58, 62, 72, 84, 88, 90]`**
- target element: 65 (not present in list)
- In the sample visualizations below, the brackets indicate the current search region, and * indicates the middle element

```
[5  12  16  40 *58  62  72  84  88  90]
 5  12  16  40  58 [62  72 *84  88  90]
 5  12  16  40  58 *62  72] 84  88  90
```

- Result: **`None`**

# Cutting the Problem Down to Size

- It should be clear why we say that the binary search algorithm uses a divide-and-conquer strategy
  - The problem is to find an item within a given range
  - At each step, the problem is split into two equal sub-problems
  - Focus then turns to one sub-problem for the next step

# Number of Comparisons

- The number of iterations made by this algorithm when it searches a list containing $n$ items is roughly $\log_2 n$
- To see why, consider the question from the other direction
  - Let's start with a list containing 1 item. Suppose each step of an iteration doubles the size of the list.
  - After $n$ steps we will have $2^n$ items in the list
- By definition of logarithm, if $x = 2^y$, then $y = \log_2 x$
- During searching we're reducing a region of size $n$ down to a region of size 1
- A successful search might return after the first comparison
- An unsuccessful search does all $\log_2 n + 1$ comparisons
- Example: 8→4→2→1, or 4 comparisons (note: $\log_2 8 = 3$)

# Searching Long Lists

- Divide-and-conquer might seem like a lot of extra work for such a simple problem (searching)
- For large lists, however, that work leads to a very efficient search
- We would need at most 30 comparisons to find something in a list of 1 billion items
- The worst case for linear search would be 1 billion comparisons!

| $n$ | $\log_2 n$ (rounded up) |
|---:|---:|
| 2 | 1 |
| 4 | 2 |
| 8 | 3 |
| 16 | 4 |
| 1,000 | 10 |
| 2,000 | 11 |
| 4,000 | 12 |
| 1,000,000 | 20 |
| 1,000,000,000 | 30 |
| 1,000,000,000,000 | 40 |

# Divide and Conquer Sorting

- The divide-and-conquer strategy used to make a more efficient search algorithm can also be applied to **sorting**
- Two well-known sorting algorithms:
  - **Merge Sort**: sort subgroups of size 2, merge them into sorted groups of size 4, merge those into sorted groups of size 8, and so on
  - **Quicksort**: divide a list into big values and small values, then sort each part
- Let's first explore merge sort and see how it can be implemented in Python

# Merge Sort

- The merge sort algorithm works from “the bottom up”
  - Start by solving the smallest pieces of the main problem
  - Keep combining their results into larger solutions
  - Eventually the original problem will be solved
- Example: sorting playing cards
  - Divide the cards into groups of two
  - Sort each group, putting the smaller of the two on the top
  - Merge groups of two into groups of four
  - Merge groups of four into groups of eight
  - and so on ...

# Merge Sort: Example

- Let's try an example of merge sort with 7 playing cards:

| | |
|---|---|
| `2 Q J 7 A 10 5` | Original list |
| `2 Q    7 J    10 A    5` | Sorted pairs |
| `2 7 J Q    5 10 A` | Merged pairs into sorted groups of 3 or 4 |
| `2 5 7 10 J Q A` | Merged smaller groups into 1 large sorted group |

# Merge Sort

- What makes this method more effective than simple insertion sort?
  - **Merging** two piles is a very simple operation
  - Only need to look at the two cards currently on the top of each pile
  - No need to look deeper into either group
- In our example, we had these two piles at one point:
  - **`[2 7 J Q]`** and **`[5 10 A]`**
  - Compare 2 with 5, pick up the 2
  - Compare 5 with 7, pick up the 5
  - Compare 7 with 10, pick up the 7
  - and so on...

# Merge Sort Visualization

- See visualgo.net/en/sorting for a visualization of merge sort
- Watching a few animations of merge sort in action will give you a stronger sense of how the algorithm sorts a list of values

# Implementing Merge Sort

- We will now see how to implement merge sort as a function called **`msort`**
- **`msort`** depends on two helper functions:
  - **`merge`**, which merges two sorted lists into one. This function is already implemented in the built-in **`heapq`** module in Python.
  - **`merge_groups`**, which calls **`merge`** and tells it exactly which sub-lists of the original list to merge

# Implementing Merge Sort

- Let's look at an example of the **`merge`** function so we understand how it works

```
import heapq
list1 = [1, 4, 6, 8]
list2 = [2, 5, 7, 9, 10, 13, 19]
merged_list = heapq.merge(list1, list2)
```

- **`merged_list`** will be:

```
[1, 2, 4, 5, 6, 7, 8, 9, 10, 13, 19]
```

# The `merge_groups` Function

- A helper function which we will write ourselves is **`merge_groups`**
- The **`merge_groups`** function takes two arguments: the list and the size of a group, **`group_size`** (e.g., 2, 4, 8, …)
  - It takes adjacent groups of *sorted* values two at a time and merges them into single groups
  - For example, if the group size is 2, this means that **`merge_groups`** will merge adjacent pairs into quartets
- The function depends on Python's **slicing notation**, which works with lists and strings
  - Code like **`nums[i:j]`** means "create a new list containing elements **`i`** through **`j-1`** of **`nums`**"

# Slicing Examples

- Example of slicing:

```
names = ['Abe', 'Barbara', 'Chris',
         'Dave', 'Erin', 'Frank', 'Harry']
print(names[2:6])
```

- Output: `['Chris', 'Dave', 'Erin', 'Frank']`
- Slicing notation can be used to change the contents of a list:

```
names[1:3] = ['Mike', 'Nathan', 'Opal']
```

- `names` becomes:

```
['Abe', 'Mike', 'Nathan', 'Opal', 'Dave',
'Erin', 'Frank', 'Harry']
```

- Note: `'Barbara'` and `'Chris'` have been replaced with the words in red

# The `merge_groups` Function

- To understand how **`merge_groups`** needs to work, consider the task of merging two quartets into one octet ("quartets" means **`group_size = 4`**)
  - The two quartets are adjacent to each other in the list
  - Generally, there are several or many such pairs of quartets we need to merge together, and we have to merge all such pairs of quartets into octets
- We need variables to tell us where each pair of quartets begins
- Call these variables **`i`** and **`j`**
  - **`i`** is the starting index of the first quartet
  - **`j`** is the starting index of the second quartet

# The `merge_groups` Function

- After merging those quartets together, we need to move to the next two quartets
- They can be found at indexes **`i+8`** and **`j+8`** since we need to skip over the octet we just created
- Initially, **`i = 0`** and **`j = 4`**
- For the second iteration, **`i = 8`** and **`j = 12`**
  - Note that **`j = i + 4`** → **`j = i + group_size`**
- Next, **`i = 16`**, **`j = 20`** (**`j = i + group_size`**)
- In general, after merging two groups together, **`i`** will increase by 2 × **`group_size`** and **`j`** will simply become **`j = i + group_size`**

# The `merge_groups` Function

- Now that we have worked out this logic, we can implement the **`merge_groups`** function:

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- The for-loop doubles the group size until the list is just one large group
- See unit05/msort.py for examples of this function in action

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: **`merge_groups(nums, 1)`**
- **`a`**: **`[8, 6, 7, 5, 3, 1, 2, 4] group_size = 1`**
- The loop will iterate as: **`i = 0, 2, 4, 6`**

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: **`merge_groups(nums, 1)`**
- **`a`**: **`[8, 6, 7, 5, 3, 1, 2, 4] group_size = 1`**
- **`i = 0`**
- **`j = i + 1 = 1`**
- **`k = j + 1 = 2`**
- **`a[0:2] = merge(a[0:1], a[1:2])`**
- **`a`**: **`[6, 8, 7, 5, 3, 1, 2, 4]`**

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: `merge_groups(nums, 1)`
- `a`: `[6, 8, 7, 5, 3, 1, 2, 4] group_size = 1`
- `i = 2`
- `j = i + 1 = 3`
- `k = j + 1 = 4`
- `a[2:4] = merge(a[2:3], a[3:4])`
- `a`: `[6, 8, 5, 7, 3, 1, 2, 4]`

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: `merge_groups(nums, 1)`
- `a`: `[6, 8, 5, 7, 3, 1, 2, 4] group_size = 1`
- `i = 4`
- `j = i + 1 = 5`
- `k = j + 1 = 6`
- `a[4:6] = merge(a[4:5], a[5:6])`
- `a`: `[6, 8, 5, 7, 1, 3, 2, 4]`

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: **`merge_groups(nums, 1)`**
- **`a`**: **`[6, 8, 5, 7, 1, 3, 2, 4]`** **`group_size = 1`**
- **`i = 6`**
- **`j = i + 1 = 7`**
- **`k = j + 1 = 8`**
- **`a[6:8] = merge(a[6:7], a[7:8])`**
- **`a`**: **`[6, 8, 5, 7, 1, 3, 2, 4]`**

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: **`merge_groups(nums, 2)`**
- **`a`**: **`[6, 8, 5, 7, 1, 3, 2, 4] group_size = 2`**
- The loop will iterate as: **`i = 0, 4`**

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: `merge_groups(nums, 2)`
- `a`: `[6, 8, 5, 7, 1, 3, 2, 4]` `group_size = 2`
- `i = 0`
- `j = i + 2 = 2`
- `k = j + 2 = 4`
- `a[0:4] = merge(a[0:2], a[2:4])`
- `a`: `[5, 6, 7, 8, 1, 3, 2, 4]`

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: **`merge_groups(nums, 2)`**
- **`a`**: **`[5, 6, 7, 8, 1, 3, 2, 4]`** **`group_size = 2`**
- **`i = 4`**
- **`j = i + 2 = 6`**
- **`k = j + 2 = 8`**
- **`a[4:8] = merge(a[4:6], a[6:8])`**
- **`a`**: **`[5, 6, 7, 8, 1, 2, 3, 4]`**

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: **`merge_groups(nums, 4)`**
- **`a`**: **`[5, 6, 7, 8, 1, 2, 3, 4] group_size = 4`**
- The loop will iterate as: **`i = 0, 8`** (iterate once only)

# Trace Execution: `merge_groups`

```
def merge_groups(a, group_size):
    for i in range(0, len(a), 2*group_size):
        j = i + group_size
        k = j + group_size
        a[i:k] = list(heapq.merge(a[i:j], a[j:k]))
```

- Function call: `merge_groups(nums, 4)`
- `a`: `[5, 6, 7, 8, 1, 2, 3, 4]` `group_size = 4`
- `i = 0`
- `j = i + 4 = 4`
- `k = j + 4 = 8`
- `a[0:8] = merge(a[0:4], a[4:8])`
- `a`: `[1, 2, 3, 4, 5, 6, 7, 8]`

# Completed `msort` Function

- All that remains now is to write **`msort`**, which will be straightforward with the help of **`merge_groups`**
- The main thing that **`msort`** needs to do is tell **`merge_groups`** how large each group is. But that's easy:
  - First we take single elements and merge them into sorted pairs
  - Then merge all the sorted pairs into sorted quartets
  - Next merge all the sorted quartets into sorted octets
  - and so on ...
- To help visualize the progress of **`msort`**, we can call a function named **`print_msort_brackets`**

# Completed `msort` Function

```
def msort(a):
    size = 1
    while size < len(a):
        print_msort_brackets(a, size) # optional
        merge_groups(a, size)
        size *= 2
        print_msort_brackets(a, len(a)) # optional
```

- See unit05/msort.py for a test run of the **`msort`** function

# Completed `msort` Function

- Example run of **`msort`**, with **`print_msort_brackets`**:

```
nums: [33, 93, 7, 15, 50, 11, 65, 43]
[33] [93] [7] [15] [50] [11] [65] [43]
[33 93] [7 15] [11 50] [43 65]
[7 15 33 93] [11 43 50 65]
[7 11 15 33 43 50 65 93]
```

# Comparisons in Merge Sort

- To completely sort a list with $n$ items requires $\log_2 n$ iterations
  - Why? The group size starts at 1 and doubles with each iteration. The group size equals or exceeds $n$ after $\log_2 n$ rounds of doubling
- During each iteration of **`msort`** there are at most $n$ comparisons. Why?
  - Comparisons occur in the built-in **`merge`** method
  - It compares values at the front of each group
  - It may have to work all the way to the end of each group, but might stop early
  - So, the total number of comparisons is roughly $n \log_2 n$

# Scalability of Merge Sort

- So, merge sort is a $O(n \log_2 n)$ algorithm
- Is this new formula that much better than the comparisons made by insertion sort?
- Not that big of a difference for small lists
- But as the length of the list increases, the savings start to add up

| $\boldsymbol{n}$ | $\boldsymbol{n^2/2}$ | $\boldsymbol{n \log n}$ |
|---:|---:|---:|
| 8 | 32 | 24 |
| 16 | 128 | 64 |
| 32 | 512 | 160 |
| 1,000 | 500,000 | 10,000 |
| 5,000 | 12,500,000 | 65,000 |
| 10,000 | 50,000,000 | 140,000 |

# Recursion

- An algorithm that uses divide and conquer can be written using iteration or **recursion**
- A recursive solution to a problem features “self-similarity”, meaning that a function that solves a problem *calls itself*
- You actually already have familiarity with this concept
  - Consider the factorial operation in mathematics
  - $n! = n \times (n-1)!$ for integers $n \geq 1$, where $0! = 1$
  - Note how factorial is defined in terms of itself (i.e., the ! symbol appears on both sides of the equals sign)
  - This is a **recursive definition** of factorial
  - The simplest case of a recursive definition is called the **base case**

# Recursion Example: Factorial

- Writing a recursive Python function that implements factorial is very straightforward
- We need to define both the recursive part (which is when the factorial function calls itself), and the base case

```
def factorial(n):
    if n == 0:      # base case
        return 1
    else:
        return n * factorial(n-1)
```

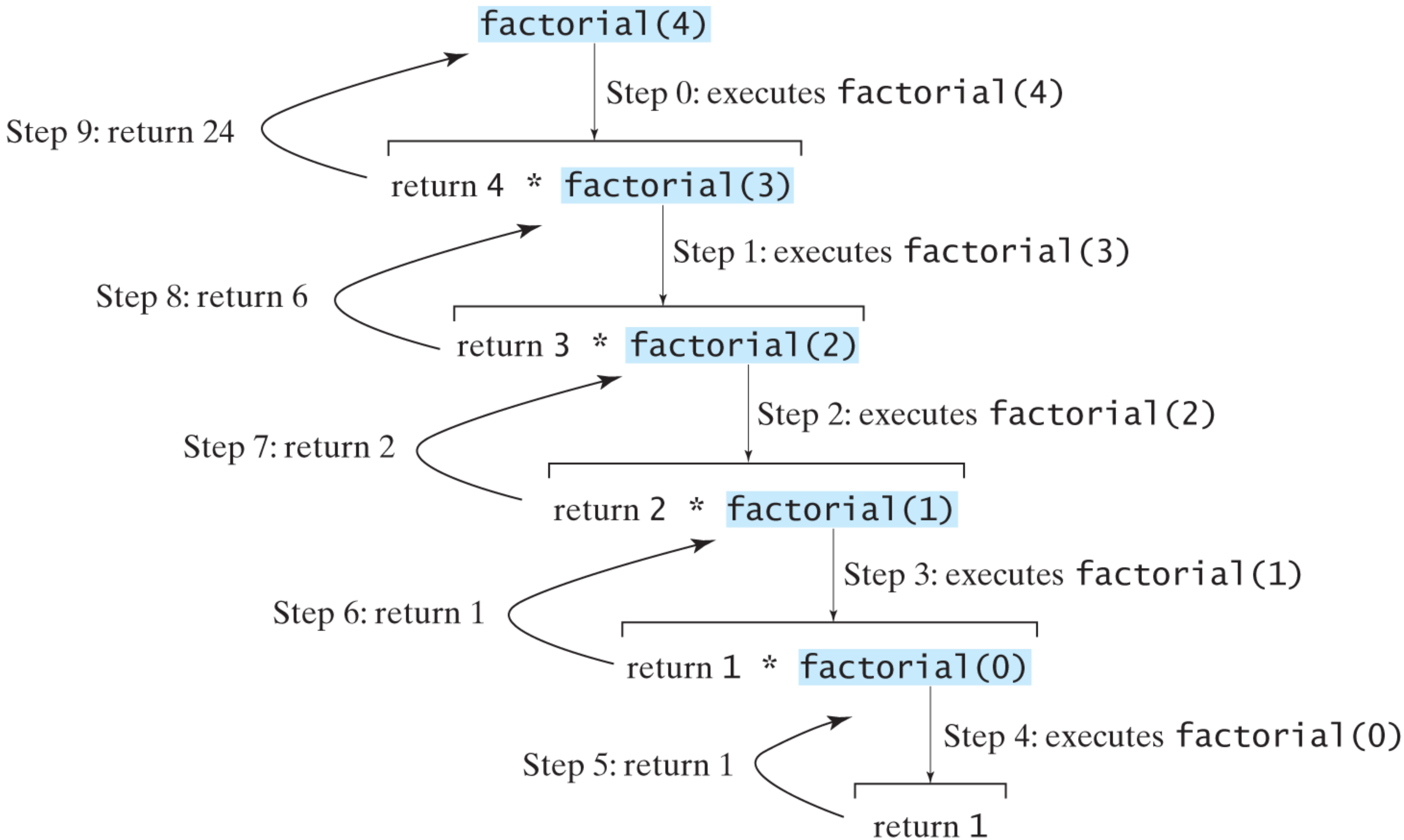
Recursive call to **`factorial()`**

- See unit05/recursion_examples.py for code for many of the example recursive functions from these notes

# Recursion

- All recursive functions have the following characteristics:
  - One or more **base cases** (the simplest cases) are used to stop recursion
  - One or more **recursive calls** that reduce the original problem in size, bringing it increasingly closer to a base case until it becomes that case
  - A recursive call can result in many more recursive calls, because the method keeps on dividing a sub-problem into new sub-problems *that are of smaller size than the original*
  - These sub-problems are of the same nature as the original
- Please note: *solutions* can be recursive, not problems!

# Trace: `factorial(4)`



# Trace: `factorial(4)`

```
factorial(4) = 4*factorial(3)
  factorial(3) = 3*factorial(2)
    factorial(2) = 2*factorial(1)
      factorial(1) = 1*factorial(0)
        factorial(0) = 1
      factorial(1) = 1*factorial(0) = 1*1 = 1
    factorial(2) = 2*factorial(1) = 2*1 = 2
  factorial(3) = 3*factorial(2) = 3*2 = 6
factorial(4) = 4*factorial(3) = 4*6 = 24
```

recursive function calls

functions returning values

# A Disclaimer

- The true benefit of *recursive thinking* is not realized until one starts trying to solve challenging problems that are more complicated than what we will explore in CSE 101
- Some (but not all) of the problems described in these lecture notes would be better solved using iterative, non-recursive functions
  - One notable exception is sorting, which *can* be solved efficiently using recursive algorithms like Quicksort
- The purpose of these examples, therefore, is to help you understand *how to think recursively* when solving problems, not necessarily how to solve the stated problems in the most efficient manner

# Example: Fibonacci Numbers

- Suppose we have one pair of rabbits (male and female) at the beginning of a year
- Rabbit pairs are not fertile during their first month of life but thereafter give birth to one new male and female pair at the end of every month
- Also, these are immortal rabbits and never die

# Example: Fibonacci Numbers

- So we can now compute how many rabbit *pairs* will be alive at the end of month *k*:

$$\left[\begin{array}{l}\text{the number}\\ \text{of rabbit}\\ \text{pairs alive}\\ \text{at the end}\\ \text{of month } k\end{array}\right] = \left[\begin{array}{l}\text{the number}\\ \text{of rabbit}\\ \text{pairs alive}\\ \text{at the end}\\ \text{of month } k-1\end{array}\right] + \left[\begin{array}{l}\text{the number}\\ \text{of rabbit}\\ \text{pairs born}\\ \text{at the end}\\ \text{of month } k\end{array}\right]$$

$$= \left[\begin{array}{l}\text{the number}\\ \text{of rabbit}\\ \text{pairs alive}\\ \text{at the end}\\ \text{of month } k-1\end{array}\right] + \left[\begin{array}{l}\text{the number}\\ \text{of rabbit}\\ \text{pairs alive}\\ \text{at the end}\\ \text{of month } k-2\end{array}\right]$$

# Example: Fibonacci Numbers

- At the start of the year (after 0 months), we have $F_0 = 1$ *pair* of rabbits
- At the end of the first month we still have only $F_1 = 1$ pair of rabbits because they haven't reproduced yet
- At the end of *k* months there will be $F_k = F_{k-1} + F_{k-2}$ pairs of rabbits
    - $F_{k-1}$ is how many rabbits were alive the previous month
    - $F_{k-2}$ is how many rabbits were alive two months ago, which equals how many rabbits will be born in month *k*
- By now you have probably guessed that *F* is the Fibonacci sequence $(1, 1, 2, 3, 5, 8, 13, 21, \ldots)$
- Let's see a function that returns the $n^{\text{th}}$ Fibonacci number

# Example: Fibonacci Numbers

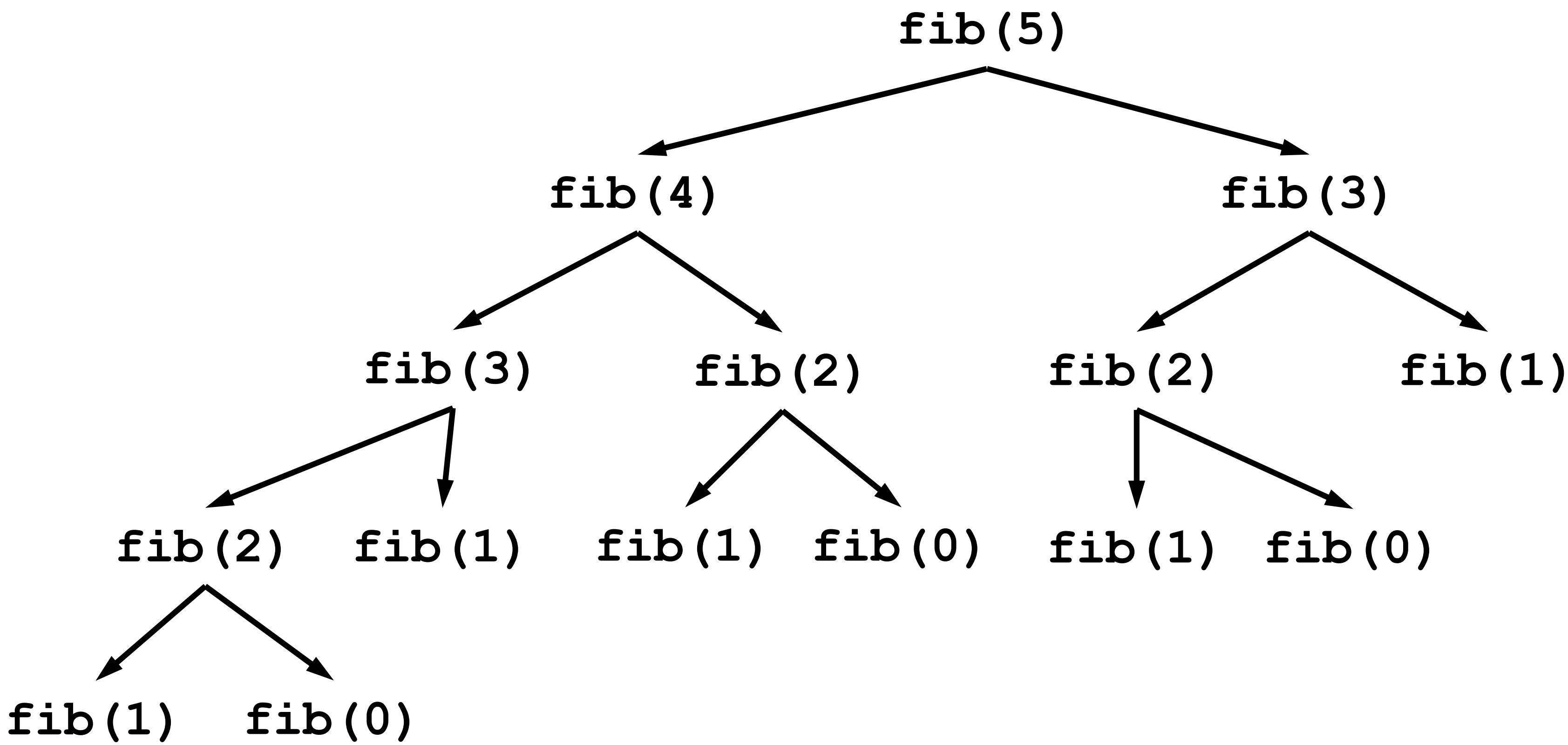
```
def fib(n):
    if n == 0 or n == 1: # two base cases
        return 1
    return fib(n - 1) + fib(n - 2)
```

- Examples:

```
fib(0) = 1
fib(1) = 1
fib(2) = fib(1) + fib(0) = 1 + 1 = 2
fib(3) = fib(2) + fib(1) = 2 + 1 = 3
fib(4) = fib(3) + fib(2) = 3 + 2 = 5
```

See unit05/recursion_examples.py

# Trace: `fib(5)`



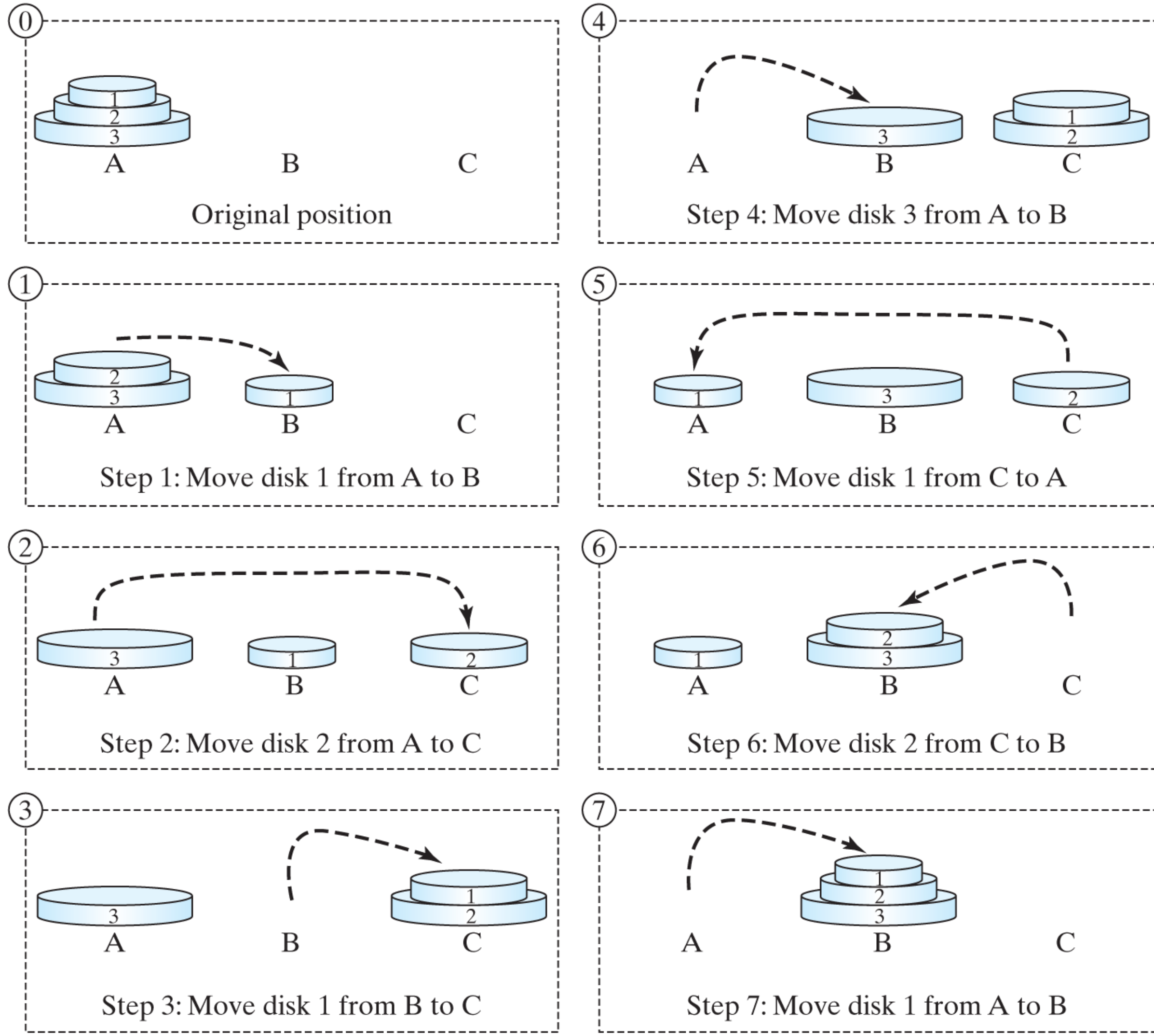

- This *call tree* diagram illustrates how the initial call to **`fib(n)`** generates a large number of recursive calls, even for a small value for **`n`**

# Example: The Towers of Hanoi

- The Towers of Hanoi problem is a classic problem that can be solved more easily with recursion than iteration
- Goal: move a set number of disks of distinct sizes from one tower to another while observing the following rules:
  - There are *n* disks labeled 1, 2, 3, ……, *n* and three towers labeled A, B, and C
  - No disk can be on top of a smaller disk at any time.
  - All the disks are initially placed on tower A
  - Only one disk can be moved at a time, and it must be the smallest disk on a tower.
  - Move all the disks from A to B with the assistance of C
- See [goo.gl/8hjRyP](goo.gl/8hjRyP) for animations

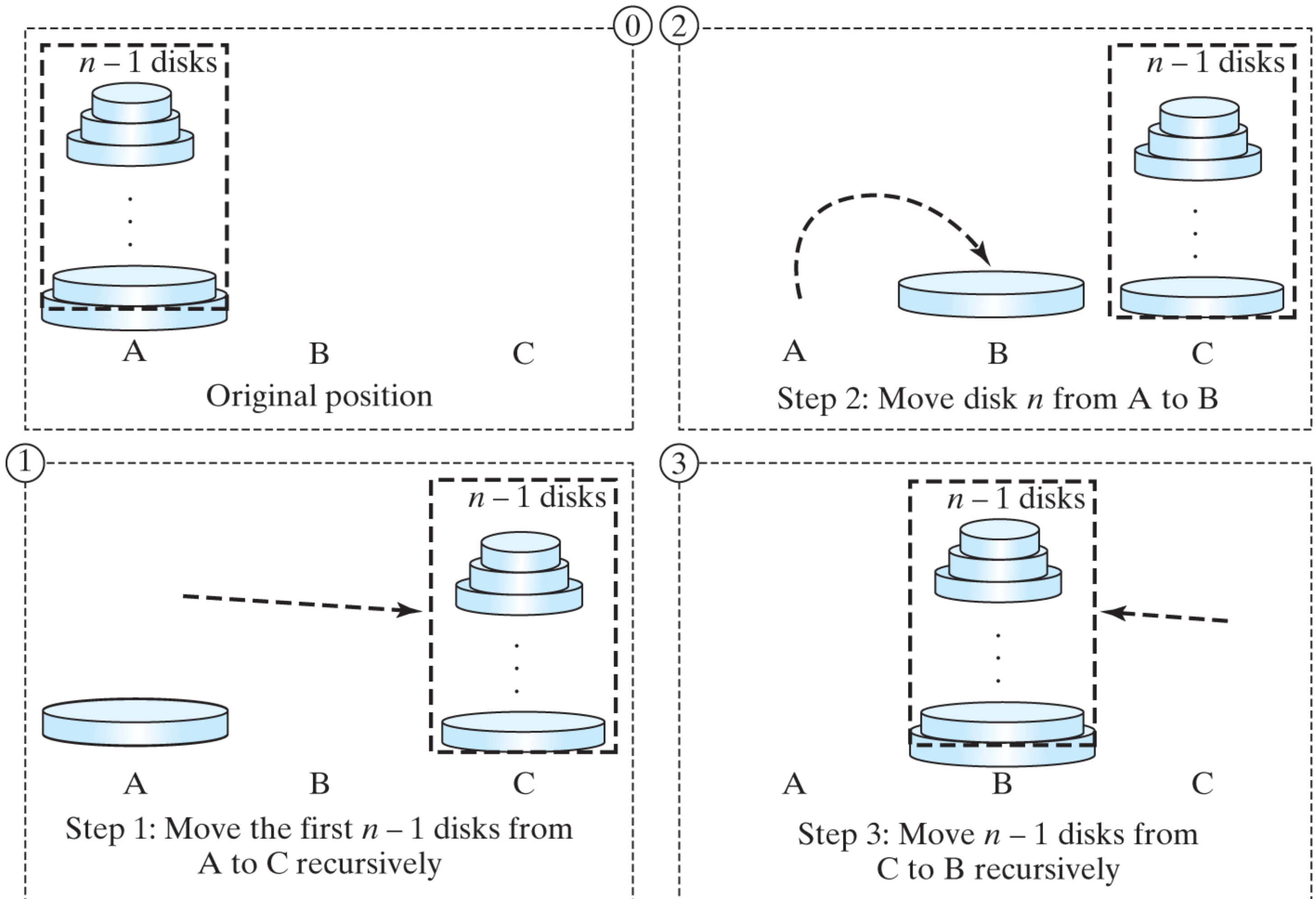
# Example: The Towers of Hanoi

- Towers of Hanoi with $n = 3$
- What about the recursive step?
- Base case?
- Well, the base case is when there is only $n = 1$ disk



# Example: The Towers of Hanoi

- The recursive solution for $n - 1$ disks is shown below



# Example: The Towers of Hanoi

- The key idea of the recursive step is this:
  - We want to move $n$ disks from A to B
  - Let's break this down into three phases:

  1. *Recursively* move the top $n - 1$ disks from A to C
  2. Move the largest disk (disk # $n$) from A to B
  3. *Recursively* move the top $n - 1$ disks from C to B
- Moving $n - 1$ involves recursively moving $n - 2$ disks, then moving disk # $n - 1$, then recursively moving those $n - 2$ disks again. And likewise for $n - 2$, $n - 3$, etc.
- We see that the problem of moving $n$ disks can be expressed in terms of a solution to the problem of moving $n - 1$ disks, which is what makes the solution recursive

# Example: The Towers of Hanoi

```
def move_disks(n, from_tower, to_tower, aux_tower):
   if (n == 1):
      print('Move disk ' + str(n) + ' from ' +
            from_tower + ' to ' + to_tower)
   else:
      move_disks(n-1, from_tower, aux_tower,
                 to_tower)
      print('Move disk ' + str(n) + ' from ' +
           from_tower + ' to ' + to_tower)
      move_disks(n-1, aux_tower, to_tower,
                 from_tower)
```

- See unit05/recursion_examples.py

# Example: The Towers of Hanoi

- You may be surprised that the **`move_disks`** function has no return statements
- The computer will execute the function until its last statement (which is the second recursive call to **`move_disks`**) and, when that call completes, simply return to the caller
  - It's as though the word **`return`** were typed as the last line of the function

# Example: The Towers of Hanoi

- Example: **`move_disks(3, 'A', 'B', 'C')`**
- Output:

```
Move disk 1 from A to B
Move disk 2 from A to C
Move disk 1 from B to C
Move disk 3 from A to B
Move disk 1 from C to A
Move disk 2 from C to B
Move disk 1 from A to B
```

# Recursive Binary Search

- For recursive binary search (**`rsearch`**), the idea is basically the same as iterative binary search (**`bsearch`**)
- But, the while-loop is replaced with a recursive call to the function
- The algorithm checks the middle element to see if it equals the target
- If not, the function calls itself on the first half or second half, depending on whether the middle element is greater than or less than the target (respectively)

# Completed `rsearch` Function

```
def rsearch(a, x, lower, upper):
    if upper == lower + 1:
        return None
    mid = (lower + upper) // 2
    if a[mid] == x:
        return mid
    if x < a[mid]:
        return rsearch(a, x, lower, mid)
    else:
        return rsearch(a, x, mid, upper)
```

- See unit05/rsearch.py

# Binary Search Algorithms

Iterative version:

```
def bsearch(a, x):
  lower = -1
  upper = len(a)
  while upper > lower + 1:
    mid = (lower + upper) // 2
    if a[mid] == x:
      return mid
    if x < a[mid]:
      upper = mid
    else:
      lower = mid
  return None
```

Recursive version:

```
def rsearch(a, x, lower, upper):
  if upper == lower + 1:
    return None

  mid = (lower + upper) // 2
  if a[mid] == x:
    return mid
  if x < a[mid]:
    return rsearch(a, x, lower, mid)
  else:
    return rsearch(a, x, mid, upper)
```

# Quicksort

- Quicksort is a recursive sorting algorithm
- Like merge sort, quicksort breaks a list into smaller sub-lists and sorts the smaller lists
  - It divides the list into sub-lists in a different manner, however
  - The first element in a region to be sorted is chosen as the **pivot element**
  - The region is then **partitioned** into two sub-regions with a helper function called **`partition`**
- The full code is available in unit05/qsort.py

# Quicksort

- The **`partition`** function performs this work:
  - Elements less than the pivot element are put in the left sub-region
  - Elements greater than the pivot element are put in the right sub-region
  - The pivot element is placed between the two sub-regions
  - The pivot element is now in its final position
- Quicksort works in a "top-down" approach by repeatedly splitting largest lists into smaller ones, whereas merge sort works in a "bottom-up" manner to recombine smaller lists into larger ones.
- See [visualgo.net/en/sorting](visualgo.net/en/sorting) for animations

# Quicksort

- The **`partition`** function partitions only a portion of a list
- The function takes three arguments:
  - A list of numbers
  - The starting index of the region to partition
  - The ending index of the region to partition
- For instance, **`partition(nums, 4, 15)`** means that **`nums[4]`** is the pivot element and that we want to partition elements in the range **`nums[4:16]`**
- Example: **`nums = [62 88 6 85 39 19 82 23]`**
- Function call: **`partition(nums, 0, 7)`**
- Pivot element: **`62`** (element **`[0]`** is always the pivot element)
- After partition: **`[23 6 39 19 62 85 82 88]`**

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

- The function returns the index of where the pivot element eventually winds up in **`a[]`**. That number also happens to be the number of elements ≤ the pivot element.

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

**x** stores a copy of the pivot element

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

`i` will eventually store the final position of the pivot element

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

For each element in the region, except the pivot element…

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

Compare each element in the region with **x**

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

**a[i]** ≤ **x**, so we found another element that will go in the first half of the region

The **i** variable essentially counts the number of elements that are ≤ the pivot element

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

Move the small element to the front half of the region

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

Move the pivot element into its final position

# Completed `partition` Function

```
def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

Return final index of where pivot element was moved to

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 0 |

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 0 |
| j | 1 |

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:  False
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 0 |
| j | 1 |
| a[j] | 8 |

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 0 |
| j | 2 |
| a[j] | 1 |

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index: 0   1   2   3   4   5   6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:  True
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 0 |
| j | 2 |
| a[j] | 1 |

# Trace Execution: `partition()`

```
a:     [5, 8, 1, 6, 3, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| `p` | 0 |
| `r` | 6 |
| `x` | 5 |
| `i` | 1 |
| `j` | 2 |
| `a[j]` | 1 |

# Trace Execution: `partition()`

```
a:     [5, 1, 8, 6, 3, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 1 |
| j | 2 |
| a[j] | 8 |

# Trace Execution: `partition()`

```
a:     [5, 1, 8, 6, 3, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 1 |
| j | 3 |
| a[j] | 6 |

# Trace Execution: `partition()`

```
a:     [5, 1, 8, 6, 3, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:  False
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
| --- | --- |
| p | 0 |
| r | 6 |
| x | 5 |
| i | 1 |
| j | 3 |
| a[j] | 6 |

# Trace Execution: `partition()`

```
a:     [5, 1, 8, 6, 3, 7, 2]
index: 0   1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 1 |
| j | 4 |
| a[j] | 3 |

# Trace Execution: `partition()`

```
a:     [5, 1, 8, 6, 3, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:  True
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
| --- | --- |
| p | 0 |
| r | 6 |
| x | 5 |
| i | 1 |
| j | 4 |
| a[j] | 3 |

# Trace Execution: `partition()`

```
a:     [5, 1, 8, 6, 3, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 2 |
| j | 4 |
| a[j] | 3 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 6, 8, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 2 |
| j | 4 |
| a[j] | 8 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 6, 8, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 2 |
| j | 5 |
| a[j] | 7 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 6, 8, 7, 2]
index: 0   1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:  False
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 2 |
| j | 5 |
| a[j] | 7 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 6, 8, 7, 2]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 2 |
| j | 6 |
| a[j] | 2 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 6, 8, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:  True
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 2 |
| j | 6 |
| a[j] | 2 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 6, 8, 7, 2]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
| --- | --- |
| p | 0 |
| r | 6 |
| x | 5 |
| i | 3 |
| j | 6 |
| a[j] | 2 |

# Trace Execution: `partition()`

```
a:     [5, 1, 3, 2, 8, 7, 6]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| `p` | 0 |
| `r` | 6 |
| `x` | 5 |
| `i` | 3 |
| `j` | 6 |
| `a[j]` | 2 |

# Trace Execution: `partition()`

```
a:     [2, 1, 3, 5, 8, 7, 6]
index: 0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
| --- | --- |
| p | 0 |
| r | 6 |
| x | 5 |
| i | 3 |

# Trace Execution: `partition()`

```
a:     [2, 1, 3, 5, 8, 7, 6]
index:  0  1  2  3  4  5  6

def partition(a, p, r):
    x = a[p]
    i = p
    for j in range(p+1, r+1):
        if a[j] <= x:
            i += 1
            a[i], a[j] = a[j], a[i]
    a[p], a[i] = a[i], a[p]
    return i
```

| Variable | Value |
|---|---|
| p | 0 |
| r | 6 |
| x | 5 |
| i | 3 |

# Quicksort

- The `partition` function will do most of the work in the **quicksort algorithm**:
  - First, partition the entire list. The first pivot element will now be at its final position.
  - Take the first sub-region of the list and partition it, and likewise for the second sub-region
  - By now, 3 elements (the 3 pivot elements) are in their final positions and we have 4 small regions
  - We partition those 4 sub-regions, causing 4 more pivot elements to be finally positioned (7 total pivots)
- This process continues until a region is so small that there is nothing to partition (zero elements in the region)

# Completed `qsort` Function

- The top-level function is **`qsort`**, which depends on a helper function called **`qs`**, which in turn calls **`partition`**:

```
def qsort(a):
    qs(a, 0, len(a)-1)  # sort the entire list

def qs(a, p, r):
    if p < r:  # base case: region has 0 elements
        q = partition(a, p, r)
        qs(a, p, q-1) # recursively sort first
        qs(a, q+1, r) # and second sub-regions

regions: [... p p+1 ... q-1 q q+1 ... r r + 1 ...]
```

# Trace Execution: `qsort()`

- Input: **`[55, 46, 89, 64, 93, 45, 15, 96]`**
- Below, red indicates a pivot element
- Values in brackets are sub-regions that are being partitioned
- A red value outside brackets is a pivot element that was positioned during an earlier round of partitioning

| | |
|---|---|
| **`[15, 46, 45, 55, 93, 89, 64, 96]`** | 1st partition |
| **`[15, 46, 45] 55 [64, 89, 93, 96]`** | 2nd partitions |
| **`15 [45, 46] 55 [64, 89] 93 [96]`** | 3rd partitions |
| **`15 45 [46] 55 64 [89] 93 96`** | 4th partitions |
| **`[15, 45, 46, 55, 64, 89, 93, 96]`** | Done! |

# Trace Execution: `qsort()`

```
                         qs(a,0,7)
[15, 46, 45, 55, 93, 89, 64, 96]

  qs(a,0,2)                qs(a,4,7)
[15, 46, 45] 55 [64, 89, 93, 96]

      qs(a,1,2)       qs(4,5)     qs(7,7)
 15 [45, 46] 55 [64, 89] 93 [96]

          qs(3,2)          qs(5,5)
 15   45 [46] 55   64 [89] 93   96

[15, 45, 46, 55, 64, 89, 93, 96]
```

Done!

# Trace Execution: `qsort()`

- Another way to visualize the partitioning steps:

```
[15, 46, 45, 55, 93, 89, 64, 96]

[15, 46, 45] 55 [64, 89, 93, 96]

 15 [45, 46] 55 [64, 89] 93 [96]

 15  45 [46] 55  64 [89] 93  96
```

# Additional Examples of Recursive Solutions to Problems

# Example: Sum of Fractions

- Although a sum is computed most efficiently with a loop, it is a simple problem to understand, which makes it a good candidate for solving with recursion
- Consider the problem of trying the compute the following sum, where $n$ is a positive integer: $1 + \frac{1}{2} + \frac{1}{3} + \cdots + \frac{1}{n}$
- Let's consider a function **`sum_fracs()`** that computes and returns this sum
- The simplest case (base case) is when $n = 1$
- For $n > 1$ we can compute the sum as $\frac{1}{n}$ plus the sum of $1 + \frac{1}{2} + \frac{1}{3} + \cdots + \frac{1}{n-1}$, *which we will compute recursively*

# Example: Sum of Fractions

```
def sum_fracs(n):
    if n == 1:
        return 1
    return 1/n + sum_fracs(n - 1)
```

- See unit05/recursion_examples.py

# Trace: `sum_fracs(4)`

```
sum_fracs(4) = 1/4 + sum_fracs(3)
  sum_fracs(3) = 1/3 + sum_fracs(2)
    sum_fracs(2) = 1/2 + sum_fracs(1)
      sum_fracs(1) = 1
```

recursive function calls

```
    sum_fracs(2) = 1/2 + sum_fracs(1)
                 = 1/2 + 1 = 1.5

  sum_fracs(3) = 1/3 + sum_fracs(2)
               = 1/3 + 1.5 = 1.833…

sum_fracs(4) = 1/4 + sum_fracs(3)
             = 1/4 + 1.833… = 2.0833…
```

functions returning values

# Example: Sum a List

- Suppose we want to write a function **`rsum`** that computes the sum of the values in the list **`nums`**
- If **`nums`** has just one item, then the sum is just the value of **`nums[0]`**
- Otherwise, the sum is **`nums[0]`** plus the sum of the rest of the values, which is computed by a recursive call to the function

# Example: Sum a List

```
def rsum(nums):
    if len(nums) == 1:
        return nums[0]
    return nums[0] + rsum(nums[1:len(nums)])
```

- See unit05/recursion_examples.py

# Trace: `rsum([8,1,4,5])`

```
rsum([8,1,4,5]) = 8 + rsum([1,4,5])

  rsum([1,4,5]) = 1 + rsum([4,5])

    rsum([4,5]) = 4 + rsum([5])

      rsum([5]) = 5

    rsum([4,5]) = 4 + rsum([5]) = 4 + 5 = 9

  rsum([1,4,5]) = 1 + rsum([4,5]) = 1 + 9 = 10

rsum([8,1,4,5]) = 8 + rsum([1,4,5]) = 8 + 10 = 18
```

# Example: Exponentiation

- One way to compute $a^n$ for integer $n$ is to multiply $a$ by itself $n$ times: $a^n = a \cdot a \cdot \cdots \cdot a$
  - This is easy to implement using a loop, but it is somewhat inefficient
  - A more efficient approach uses recursion
- Example: suppose we want to compute $2^8$
  - From the laws of exponents we know $2^8 = 2^4 \cdot 2^4$
  - If we compute the value of $2^4$ once, we can simply multiply the value of $2^4$ by itself
  - Likewise, $2^4 = 2^2 \cdot 2^2$
- In general, $a^n = a^{n/2} \cdot a^{n/2}$ when $n$ is even and
  $a^n = a^{(n-1)/2} \cdot a^{(n-1)/2} \cdot a$ when $n$ is odd

# Example: Exponentiation

- Let's use these formulas to write a function that recursively computes the $n^{\text{th}}$ power of any nonzero integer
- For the base case we will use the fact that any non-zero value raised to the 0th power is 1

# Example: Exponentiation

```
def power(base, exponent):
   if exponent == 0:
      return 1                   # base case
   elif exponent % 2 == 0:       # even exponent
      temp = power(base, exponent // 2)
      return temp * temp
   else:                         # odd exponent
      temp = power(base, exponent // 2)
      return temp * temp * base
```

- See unit05/recursion_examples.py

# Trace: `power(3,5)`

```
power(3,5) = power(3,2) * power(3,2) * 3

   power(3,2) = power(3,1) * power(3,1)

      power(3,1) = power(3,0) * power(3,0) * 3

         power(3,0) = 1
```

**power(3,0)** is computed only once and stored in **temp**

```
      power(3,1) = power(3,0) * power(3,0) * 3
                 = 1 * 1 * 3 = 3
```

**power(3,1)** is computed only once and stored in **temp**

```
   power(3,2) = power(3,1) * power(3,1)
              = 3 * 3 = 9
```

**power(3,2)** is computed only once and stored in **temp**

```
power(3,5) = power(3,2) * power(3,2) * 3
           = 9 * 9 * 3 = 243
```

# Example: Reverse a String

- Consider the problem of taking a string and reversing its characters
  - Example: convert **`'stony'`** to **`'ynots'`**
- Let's explore a recursive function **`rev`** that solves this problem
- Let $n$ be the length of a string **`s`**
- In the base case $n = 1$, since **`s`** has only one character, just return **`s`** ("reversing" a single letter requires no work)
- Otherwise, when $n > 1$, return a string consisting of the last letter in **`s`**, followed by the reverse of the first $n - 1$ characters
- Doing this will require a recursive call to the **`rev`** function.

# Example: Reverse a String

- Some slicing notation that will help us:
  - **`s[-1]`** means "get the last character of string **`s`**"
  - **`s[:-1]`** means "get all but the last character of string **`s`**"
  - The same syntax can be used to slice lists and strings

```
def rev(s):
    if len(s) == 1:
        return s
    return s[-1] + rev(s[:-1])
```

# Trace: `rev('stony')`

```
rev('stony') = 'y' + rev('ston')
  rev('ston') = 'n' + rev('sto')
     rev('sto') = 'o' + rev('st')
       rev('st') = 't' + rev('s')
         rev('s') = 's'
       rev('st') = 't' + rev('s') = 't'+'s'= 'ts'
     rev('sto') = 'o' + rev('st') = 'o'+'ts' = 'ots'
  rev('ston') = 'n' + rev('sto') = 'n'+'ots' = 'nots'
rev('stony') = 'y' + rev('ston') = 'y'+'nots' = 'ynots'
```

# Example: Count Occurrences

- Python has a method named **`count()`**, which counts the number of times a target character appears in a string
- For example, **`'stonybrook'.count('o')`** is 3 because there are three lowercase o's in **`'stonybrook'`**
- How might we implement a recursive function that solves the same problem?
- First, inspect the first character of the string
  - If the character matches the target, we need to add 1 to the number of matches in the *remainder* of the string
  - Otherwise, we simply continue by counting the number of matches in the remainder of the string

# Example: Count Occurrences

- But, how do we know how many times the target character appears in the remainder of the string?
  - We perform a recursive call to the function!
- So here's our algorithm:

  If the string has at least one character in it then:

  If the first character matches, then return
  (1 + the # of matches of the target in the rest of the string)

  Otherwise, return the # of matches in the rest of the string

  Otherwise, return 0

# Example: Count Occurrences

```
def count_occurrences(string, ch):
   if len(string) > 0:
      if string[0] == ch:
         return count_occurrences(string[1:], ch) + 1
      return count_occurrences(string[1:], ch)
   return 0
```

- See unit05/recursion_examples.py

# Trace: `count_occurrences()`

- Example: **`count_occurrences('stat', 't')`**
- Abbreviating **`count_occurrences`** as **`count`**:

```
count('stat') = count('tat')
   count('tat') = 1 + count('at')
      count('at') = count('t')
         count('t') = 1 + count('')
            count('') = 0
         count('t') = 1 + count('') = 1 + 0 = 1
      count('at') = count('t') = 1
   count('tat') = 1 + count('at') = 1 + 1 = 2
count('stat') = count('tat') = 2
```

# Example: Find Palindromes

- A palindrome is a word or phrase that can be read backwards and forwards
- Examples: radar, dad, toot, e
- Let's consider a function, **`is_palindrome`**, which returns **`True`** if its argument is a palindrome, and **`False`** if not
- How could we formulate a recursive solution to this problem?
  - We need to consider the base case(s) and recursive step(s)

# Example: Find Palindromes

- The simplest case (base case) would be a string with exactly one character, which, by definition, would be a palindrome
- For the more general case we have two sub-problems:
  1. Verify that the first character and the last character of the string are equal
  2. If they match, ignore the two end characters and check whether the rest of the substring is a palindrome
  - If the first and last characters don't match, then the string (or sub-string) is not a palindrome
- The notation to slice out the first and last elements of string **`s`** and keep the remaining characters is **`s[1:-1]`**

# Example: Find Palindromes

```
def is_palindrome(s):
    if len(s) <= 1:   # a string of 0 or 1 characters
        return True   # is a palindrome
    elif s[0] != s[-1]: # the first and last
        return False    # characters don't match
    else:
        return is_palindrome(s[1:-1])
```

- See unit05/recursion_examples.py

# Trace: `is_palindrome()`

- Example: **`is_palindrome('racecar')`**
- Abbreviating **`is_palindrome`** as **`is_pal`**:

```
is_pal('racecar') = is_pal('aceca')
  is_pal('aceca') = is_pal('cec')
    is_pal('cec') = is_pal('e')
      is_pal('e') = True
    is_pal('cec') = is_pal('e') = True
  is_pal('aceca') = is_pal('cec') = True
is_pal('racecar') = is_pal('aceca') = True
```

Keep making recursive calls while the first and last characters match

# Trace: `is_palindrome()`

- Example: **`is_palindrome('hannah')`**
- Abbreviating **`is_palindrome`** as **`is_pal`**:

```
is_pal('hannah') = is_pal('anna')
   is_pal('anna') = is_pal('nn')
      is_pal('nn') = is_pal('')
         is_pal('') = True
      is_pal('nn') = is_pal('') = True
   is_pal('anna') = is_pal('cec') = True
is_pal('hannah') = is_pal('anna') = True
```

Keep making recursive calls while the first and last characters match

# Trace: `is_palindrome()`

- Example: **`is_palindrome('struts')`**
- Abbreviating **`is_palindrome`** as **`is_pal`**:

```
is_pal('struts') = is_pal('trut')
  is_pal('trut') = is_pal('ru')
    is_pal('ru') = False
  is_pal('trut') = is_pal('ru') = False
is_pal('struts') = is_pal('trut') = False
```

Keep making recursive calls while the first and last characters match

# Example: Replace Multiples of 5

- Consider a peculiar function named **`replace_mult5`** that takes a list of numbers and replaces all multiples of 5 with a substitute number
  - The list and the substitute are passed as arguments
- Here's an example:

```
nums = [5,3,15,50,2,4,6,60]
replace_mult5(nums, 77)
```

- **`nums`** becomes: **`[77,3,77,77,2,4,6,77]`**
- Since this function does not return a value, it's not entirely clear how to write it recursively
  - Consider: how do we keep track of what part of the list we have processed so far?

# Example: Replace Multiples of 5

- We can implement **`replace_mult5`** more easily if we use a **helper function**
  - Recall the **`qs`** helper function that helped us implement the **`qsort`** function earlier in this Unit
- Our helper function, **`replace_mult5_helper`**, will take the same two arguments as **`replace_mult5`**, plus a third argument that tracks what part of the list we have already processed:

```
def replace_mult5(nums, sub)
def replace_mult5_helper(nums, sub, i)
```

- In a certain sense, the helper function will simulate the behavior of a loop, as we can see in the implementation on the next slide

# Example: Replace Multiples of 5

```
def replace_mult5(nums, sub):
   replace_mult5_helper(nums, sub, 0)

def replace_mult5_helper(nums, sub, i):
   if i == len(nums):  # base case
      return
   if nums[i] % 5 == 0:
      nums[i] = sub
   replace_mult5_helper(nums, sub, i+1)
```

# Example: Replace Multiples of 5

- The recursive helper function could be written iteratively using the code below

```
def replace_mult5_helper(nums, sub, i):
    for i in range(len(nums)):
        if nums[i] % 5 == 0:
            nums[i] = sub
```

- Compare this code with the recursive version. Do you see how the recursive version is essentially simulating a for-loop?

# Trace: `replace_mult5_helper`

- Example: **`nums = [4,10,2,5]`**

  **`replace_mult5_helper(nums, 8, 0)`**
- Abbreviating **`replace_mult5_helper`** as **`rmh`**

```
rmh([4,10,2,5],8,0) → rmh([4,10,2,5],8,1)
   rmh([4,10,2,5],8,1) → rmh([4,8,2,5],8,2)
      rmh([4,8,2,5],8,2) → rmh([4,8,2,5],8,3)
         rmh([4,8,2,5],8,3) → rmh([4,8,2,8],8,3)
            rmh([4,8,2,8],8,3) → do nothing & return
```

- Since the recursive call is the last statement in the function, the four recursive calls now simply return to each other, in sequence, performing no additional work
- The final contents of **`nums`** is **`[4,8,2,8]`**

# Example: Find Index of Character

- Python has a built-in string method called **`index`** that returns the index of the first occurrence of a character (or substring, actually) in a string
- Example:

```
school = 'stony brook'
pos = school.index('o')   # pos will be 2
```

- If the target character or substring does not appear in the string, the program crashes
- Let's consider a recursive solution to this problem and implement it in a function **`rindex`**
- In cases where the target string is not found, the **`rindex`** function will simply return **`None`** instead of crashing the program

# Example: Find Index of Character

- One challenge we face is that somehow we need to keep track of what part of the string we have searched so far
- We will write a helper function, **`rindex_helper`**, that will assist with this task
- The helper function will ultimately solve the problem
- All that **`rindex`** will need to do is call **`rindex_helper`** with the correct arguments

# Example: Find Index of Character

```
def rindex(string, target):
    return rindex_helper(string, target, 0)

def rindex_helper(string, target, i):
    if i >= len(string):
        return None
    elif string[i] == target:
        return i
    else:
        return rindex_helper(string, target, i+1)
```

# Trace: `rindex_helper`

- Example: **`rindex_helper('stony','n',0)`**
- Abbreviating **`rindex_helper`** as **`rh`**

```
rh('stony','n',0) = rh('stony','n',1)
  rh('stony','n',1) = rh('stony','n',2)
    rh('stony','n',2) = rh('stony','n',3)
      rh('stony','n',3) = 3  # found match!
    rh('stony','n',2) = rh('stony','n',3) = 3
  rh('stony','n',1) = rh('stony','n',2) = 3
rh('stony','n',0) = rh('stony','n',1) = 3
```

# Trace: `rindex_helper`

- Example: **`rindex_helper('stop','z',0)`**
- Abbreviating **`rindex_helper`** as **`rh`**

```
rh('stop','z',0) = rh('stop','z',1)
  rh('stop','z',1) = rh('stop','z',2)
    rh('stop','z',2) = rh('stop','z',3)
      rh('stop','z',3) = rh('stop','z',4)
        rh('stop','z',4) = None
      rh('stop','z',3) = rh('stony','n',4) = None
    rh('stop','z',2) = rh('stony','n',3) = None
  rh('stop','z',1) = rh('stop','z',2) = None
rh('stop','z',0) = rh('stop','z',1) = None
```

# Recursive Helper Functions

- The following material is based on notes by Jayesh Ranjan, a computer science major who served as a TA for CSE 101 many times during his studies at Stony Book University
- The main focus on this guide is on understanding the relationship between iteration and recursion: both are forms of repetition, but each implements the repetition in a different way
- Suppose you wanted to use a loop find the sum of all integers from 0 through *n*, inclusive
- One possible solution is given on the next slide
  - A while-loop is used because it will match up more closely with the recursive version
- See unit05/iter_to_rec.py

# Recursive Helper Functions

```
def iter_sum(n):
    index = 1
    total = 1
    while index <= n:
        total += index
        index += 1
    return total
```

- Somehow we need to map this iterative algorithm to a recursive implementation, specifically:
  - the **`index`** and **`total`** variables
  - the while-loop condition and body
  - the return statement

# Recursive Helper Functions

- Ultimately, we want a function **`rec_sum(n)`** we can call that will return the correct value
- We will use a recursive helper function to keep track of the **`index`** variable by taking it as an argument to the helper
- The **`total`** variable will be implemented as the return value of the helper function
- The while-loop's condition in the iterative solution will need to be replaced with a condition for an if-statement that will terminate the recursion
  - So, in both the iterative and recursive solutions we need a carefully-written condition to stop the repetition
- The recursive implementation is given on the next slide

# Recursive Helper Functions

```
def rec_sum(n):
   return rec_sum_helper(n, 1)

def rec_sum_helper(n, index):
   if index == n:
      return n
   return index + rec_sum_helper(n, index+1)
```

- Let's try to understand now how this code matches up with the iterative solution

# Recursive Helper Functions

Iterative version:

```
index = 1
```

Recursive version:

```
rec_sum_helper(n, 1)
```

- Initializing **`index`** to 1 in the iterative version is akin to calling the recursive helper function with an **`index`** argument of 1.

# Recursive Helper Functions

Iterative version:

**`while index <= n`**

Recursive version:

**`if index == n`**

- The while-loop will stop iterating (repeating) once **`index > n`**. Similarly, the recursive version will stop making function calls once **`index == n`**.
- When **`index == n`** in the recursive version, we return **`n`** itself. This means that **`n`** will be added to the running total that the function is recursively computing.

# Recursive Helper Functions

Iterative version:

```
total += index
index += 1
```

Recursive version:

```
return index +
    rec_sum_helper(n, index+1)
```

- The two **+=** statements from the iterative version are captured in the single line from the recursive version.
- Color is used to show the connection between versions. There is no **total** variable for the recursive function. Rather, the function's return value serves this purpose.

# Recursive Helper Functions

```
def iter_sum(n):
    index = 1
    total = 0
    while index <= n:
        total += index
        index += 1
    return total
```

```
def rec_sum(n):
    return rec_sum_helper(n, 1)

def rec_sum_helper(n, index):
    if index == n:
        return n
    return index +
      rec_sum_helper(n, index+1)
```

- It's not really possible to make a perfect one-to-one matching between the code in both versions, but I've attempted to do so here using color.