set error flag

else

1.a.

Search OPTAB for OPCODE if found then Add 3 to LOCCTR else if opcode = BYTE then Add Begin find length of constant in bytes. Add length to LOCCTR End else if OPLODE = WORD then Add 3 to the Add 3 to LOCCTR else if OPCODE=' RESB' then Add # Coperand] to LOCKTR else if OPLODE = 'RESW' then Add 3* # Coperand J to LOCCTP else set error flag End. write line to intermediate till Read next input line End work tast input tine Writ line to intermediate file Read next input Save LLOCCTR-Starting address & to program length End.

1. b. SIC /XE MACHINE ARCHITECTURE:

SICIXE stands for simple Instructional computer (SIC) / Extra equipment or extra expensive (XE).

i) MEMORY: The total memory used or ofited by SIC(XE is about 200 bytes. Number of addrers lines ou 20.

ii> REGISTERS:

NUMBERI	FUNCTION
0	sumulated
1	Index Register
2	Lenkage Register Bare register
3	y general purpor
6	Register
7	floating hoint integer
8	Program Counter
9	status word.
	(000

iii> DATA FORMAT

a> Integers - 3 bytes

b> wharater - 1 byte

er floating point - 6 byte

```
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SIGN Bourabl
IVY INSTRUCTION FORMAT
A) FORMAT 1:
                   -> 8bits
            opiode
 b> FORMAT-2:
                    平平子生661
             opwde | RI | RZ
                6 111111 12 -> 246.20
  X) FORMAT-3:
             opcode n'ix bpe Dispidala
  d> FORMAT- 4:
                                -> 32 bits
                              20
            6 11111
           opude nix bpc Duploata
 flere
      n> Inducet bil
       i> Immediate bit
       X-) Index bil
       base bil
        PD PL relative bit
        e-> extended bil
   n=0, i=0 il represents simple sic instruction
when,
   n=0, i=1 it represents immediate addressing
   n=1, i=0 it refreunt inderect addressing
    n=1, i=1 il represents direct addressing
    x = 0 it represent direct addressing.
     = 1 il represent undered addressing
    6-0, P=0 neither bare not fic relative addressing
```

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> b=1, p=0 at represents bare relative addressing b=0, P=0 it represent PC relative addressing e=0 it represents format 3 instruction e=1 it represents format 4 instruction

V> ADDRESING MODES;

ar Bare relative

67 Program counter relative.

2) Direct

d> Index

er &mmediate

& & Indirect

& g> Extended addressing mother.

VIY INSTRUCTION SET:

as Load and stone instructions: LDB, LDA, STB,

6> Arithematu instructions. ADD, SUR, MUL.

er Register authematic instructions: ADDR, SUBR, MULR

VIIY INPUT AND OUTPUT !

at Test Drive (TD): It is to Test whether The device is ready or not. The wondition wide ((1) is used to test, if cal less then device is ready or else device is not ready.

by Read Data (RD): Reads the byte from device and store it in register A.

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e) write data: Writer byte from register 1

to device.

27 SIO - Starts IIO Operation.

er 410 - fealts IIO Operation.

1> TO - Test IO operations.

1.C. SIC/XE

LOCATION	LENGTH	LABEL	MMEMONIC	OPERAND	OBJECT
		STAPP	START	0	
0000	3	FIRST	LDX	#0	050000
0003	3		LDA	#0	010000
0006	4		+LDB	#TABLE 2	69101790
			BASE	TABLE 2	
000A	3	LOOP	ADD	TABLEIX	18A013
000D	3		ADD	TABLEZIX	136000
00010	3		TKK	COUNT	2F200 A
00 13	3		JLT	LOOP	3B2FF4
0016	4		+57A	TOTAL	0F102F00
001A	3		RSUB		4F0000
00LD	1×3=3	COUNT	RESW	1	
0020	1770	TABLE	REJW	2000	
1790	1770	TABLEZ	RESW	2000	
2 F 0 0	1×3=3	TOTAL	RESW	1	
21-03			END	FIRST	
A CONTRACTOR OF THE PARTY OF TH			DOMESTIC OF THE PARTY OF THE PA		

SIGH: Nourable

OBJECT PROGRAM:

H1 SUM 10000000 1002F03

T1 000000 1 10 050000 0 010000 69101790 1 1B COOO 1 2 F200 A 1 382 FF4 1 0 F102 F00 14 F0000 M1 000007 05 M1 000017 105 E1 000000

PART-B

4.a. BUFFER PAIRS:

is Buffering technique is used to reduce the amount of overhead required to prover beingle buffering when The program has to be loaded and proversed.

WORKING

a) The figure above is a buffer carray) of me tength, n is number of sharaster on disk.

It is divided into two n-charaster halves.

b) Two pointers are used here sexual Begin and forward.

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1) Initially both pointees point to the first sharaster of lexeme to be found. d> Now the foeward pointer increments and scan until a moth for hattien is obtained. et once next beaune is detremined, peoursed and then both the pointer ore set to the character unne diately. f? The steing within learne Begin and. focused is the resume to be persent. LOOK-AHEAD CODE: Switch (* forward tt) 2 case est: if Forward is at end of first buffer) } reload second buffer; forward = beginning of second buffer, else if (Forward is at end of second buffer)? reload first buffer; forward = beginning of first buffer,

> else terminate lexical analysis; break;

TOKEN:

a) A token is a Turninal symbol in grammas for the source language.

b) A token is a frais of token name and.

an oftenal alleibute value

1) The Token name is abstract symbol representing a lexical unit like keyword or identifier

PATTERN:

ar A hattun is a rule discibing set of lexemes that van represent a token in the source fresquams.

b) It is a description of form that lexenes of a token may take

LEXEME:

a) It is a requerce of characters in source pergram that matcher the frattern with a token. b> It is identified by resided analyzer as an instance of token.

NO. OF PAGES - 10 QUESTIONS ATTEMPTED: 1.0,1.6,1.6 H.a, 4.b.