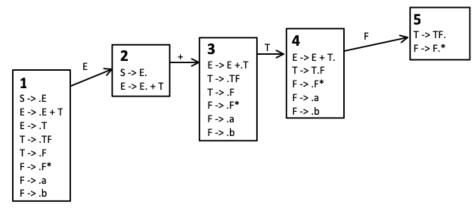
1.

a. $first(H) = \{h\}$

first(E) = {e,
$$\varepsilon$$
}
first(B) = {b, g}
first(F) = {f} + first(E) + first(H) = {f, e, h, ε }
first(A) = first(B) + first(F) = {b, g, f, e, h, ε }
follow(A) = {\$, c}
follow(B) = first(A) = {b, g, f, e, h}
follow(E) = follow(A) + first(F) + first(H) = {\$, c, f, e, h}
follow(F) = first(E) + follow(B) + follow(A) = {\$, e, b, g, f, h, c}
follow(H) = follow(F) = {\$, e, b, g, f, h, c}
b. first(C) = {h, ε }
first(B) = {g, ε }
first(A) = {d} + first(B) + first(C) = {d, g, h, ε }
first(S) = first(A) + first(C) + {b} + first(B) + {a} = {d, g, h, ε , b, a}
follow(S) = {\$}
follow(B) = follow(S) + {a} + first(C) + follow(A) = {\$, e, h, g}
follow(C) = first(B) + follow(S) + {b} + follow(A) = {\$, e, h, g}

2.

a.

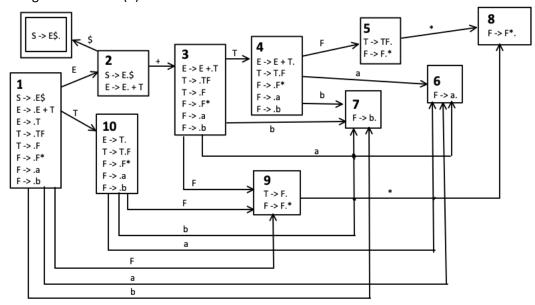


IR(0):

this grammar is not LR(0). As shown by the partial DFA, there occurs a shift-reduce conflict in state 5

SLR(1):

this grammar is SLR(1)



- 1. $E \rightarrow E+T$
- 2. $E \rightarrow T$
- 3. $T \rightarrow TF$
- 4. $T \rightarrow F$
- 5. $F \rightarrow F^*$
- 6. $F \rightarrow a$
- 7. $F \rightarrow b$

Follow(E) = $\{+, \$\}$

Follow(T) = $\{\$, +, a, b\}$

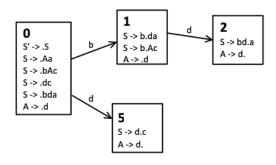
Follow(F) = $\{*, \$, +, a, b\}$

		Action				Go	oto	
	*	+	a	b	\$	E	T	F
1			S9	S7		I2	I10	I9
2		S 3			accept			
3			S6	S7			IT	
4		R1	S6	S7	R1			I5
5	S8	R3	R3	R3	R3			
6	R6	R6	R6	R6	R6			
7	R7	R7	R7	R7	R7			
8	R5	R5	R5	R5	R5			
9	S8	R4	R4	R4	R4			
10		R2	S6	S7	R2			I9

b.

LR(0):

this grammar is not LR(0), at states 5 and 2 we have a shift-reduce conflict



SLR(1):

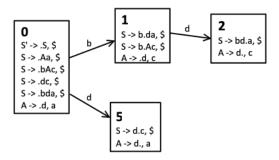
$$follow(S) = \{\$\}$$

$$follow(A) = \{a, c\}$$

this grammar is not SLR(1) either. The shift-reduce conflict in state 5 is between shifting c and reducing to non-terminal S, in state is between shifting a and reducing to S. Since the follow set of S contains a and c, we cannot reduce this conflict with SLR(1).

LALR(1):

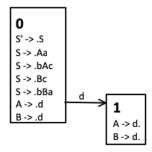
This grammar is LALR(1).



Once we add the lookaheads to each LR(0) item, the shift-reduce conflicts are reduced. In state 2, the shift symbol a is different from A's lookahead symbol c and in state 5 the shift symbol c is different from A's lookahead symbol a. We can deduce from the grammar rules that no other items will have the same grammar rules as state 2 or 5, so these two items will not be merged by any other items in LALR(1). Therefore, it is safe to say that this grammar LALR(1).

c.

LR(0):



This grammar is not LR(0), in state 1 there is a reduce-reduce conflict. Thus LR(0) fails.

SLR(1):

This grammar is not LR(1) either, in the partial DFA of LR(0), state 1 has two reduce rules deriving from non-terminals A and B. If we check the follow set of non-terminals we will find out:

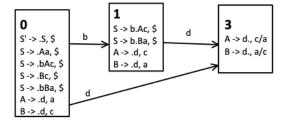
 $Follow(S) = \{\$\}$

 $Follow(A) = \{a, c\}$

 $Follow(B) = \{a, c\}$

among which the follow sets of A and B are identical. Thus, SLR(1) cannot resolve the reduce-reduce conflict in state 1.

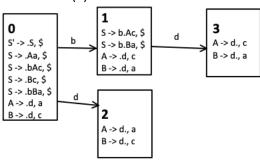
LALR(1):



This grammar is not LALR(1) either. In state 0 and 1, we see there is an identical set of grammar rules derived from non-terminal A and B. They have different lookahead symbols in state 0 and state 1. Yet in state 3, after being dynamically merged, these two rules face reduce-reduce conflict again because their sets of lookahead symbols are merged as well.

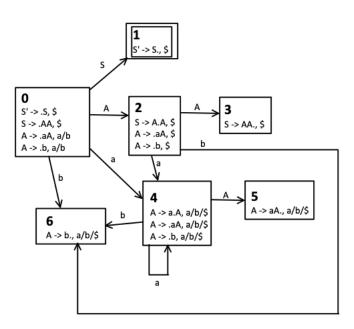
CLR(1)

This grammar is CLR(1).



CLR(1) solves the reduce-reduce conflict in LALR(1) by splitting the old state 3 into two states with identical grammatical rules but different sets of lookaheads. Unlike in LALR(1), state 1 and 0 herein branch to two separates states, leading to more items than LR(0)/LALR(1). Other branches in this CLR(1) DFA all lead a single-rule items and thus shouldn't have any conflicts. Therefore, this grammar is CLR(1).

3.



- 1. $S \rightarrow AA$
- 2. $A \rightarrow aA$
- 3. $A \rightarrow b$

	Ac	tion		Goto	
	а	b	\$	S	Α
0	S4	S6		I1	12
1			accept		
2	S4	S6			13
3			R1		15
4	S4	S6			15
5	R2	R2	R2		
6	R3	R3	R3		

Input: "abaab\$" action: shift 4

0 a Input: "abaab\$" action: shift 6

4	b
0	а

Input: "abaab\$" action: reduce 3

6	
4	D
0	а

Input: "abaab\$" action: goto 5

4	Α
0	а

Input: "abaab\$" Action: reduce 2

5	
4	A
0	a

Input: "abaab\$" action: goto 2

	_
0	Α

Input: "abaab\$" action: shift 4

2	a
0	Α

Input: "abaab\$" action: shift 4

4	а
2	а
0	Α

Input: "abaa<u>b</u>\$" Action: shift 6

6	
4	/ مر
4	а

Input: "abaab\$" Action: goto 5

2	а
0	Α

4	Α
4	а
2	а
0	Α

Input: "abaab\$" Action: reduce 2

5	A
4	/ /a
2	а
0	Α

Input: "abaab\$" action: goto 3

2	Α
0	Α

Input: "abaab\$" Action: reduce 1

3	A
0	A

Input: "abaab\$" Action: goto 1

•	
0	S

Input: "abaab\$" Action: shift \$

1	0

Input: "abaab\$" action: accept

	\$
1	S