

华东师范大学软件学院实验报告

实验课程：计算机网络	年级：2023 级本科	实验成绩：
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— 2-1

1 问题与翻译

A noiseless 8-kHz channel is sampled every 1 msec. What is the maximum data rate?

一个无噪声的 8 千赫信道每 1 毫秒采样一次。最大数据传输速率是多少？

2 解答与翻译

Nyquist's theorem states that for an ideal low-pass (no noise, finite bandwidth) channel, the maximum symbol transmission rate is $2W$ symbols per second, where W is the bandwidth of the channel (in Hz). If we use V to represent the number of discrete levels (i.e., how many distinct symbols there are), then the maximum data rate is:

$$\text{Maximum data rate in a noiseless channel} = 2W \log_2(V) \text{ b/s}$$

In this case, the channel bandwidth is 8 kHz (8000 Hz), and if we assume each sample produces 16 bits, then there are 2^{16} possible states. Substituting these values into the formula gives:

$$C = 2 \times 8000 \times \log_2(2^{16}) = 256,000 \text{ bps}$$

Thus, the maximum data rate is approximately 256 Kbps.

So, the maximum data rate of the noiseless channel is decided by how many bits there are in every sample.

奈奎斯特定理规定：在理想低通（没有噪声、带宽有限）信道中，为了避免码间串扰，极限码元传输速率为 $2W$ 波特，其中 W 是信道的频率带宽（单位为 Hz）。若用 V 表示每个码元的离散电平数目（码元的离散电平数目是指有多少种不同的码元），则极限数据传输速率为：

$$\text{理想低通信道下的极限数据传输速率} = 2W \log_2(V) \text{ 比特每秒}$$

在此例中，信道带宽为 8 千赫（8000 赫兹）。

假设每次采样产生 16 比特，则每次采样可以有 2^{16} 种不同的状态。将这些值代入公式，得到：

$$C = 2 \times 8000 \times \log_2(2^{16}) = 256,000 \text{ 比特每秒}$$

因此，无噪声信道的最大数据速率是由每个样本中有多少位决定的。在假设每次采样产生 16 比特时，得到的最大传输速率为 256 Kbps。

参考资料：[StackOverflow](#)

二 2-2

1 问题与翻译

If a binary signal is sent over a 3-kHz channel whose signal-to-noise ratio is 20 dB, what is the maximum achievable data rate?

如果一个二进制信号通过信噪比为 20 分贝的 3 千赫信道传输，最大可实现的数据传输速率是多少？

2 解答与翻译

Using Shannon's theorem, we can calculate the maximum achievable data rate for a noisy channel. The formula is:

$$\text{Maximum Data Rate} = B \times \log_2(1 + SNR)$$

Where B is the bandwidth of the channel, and SNR is the signal-to-noise ratio (in linear terms). First, we convert 20 dB into a linear value:

$$SNR (\text{linear}) = 10^{\frac{20}{10}} = 100$$

Then applying the Shannon formula:

$$\text{Maximum Data Rate} = 3000 \times \log_2(1 + 100) = 3000 \times \log_2(101) \approx 3000 \times 6.6582 = 19974.6 \text{ bps}$$

So, the maximum achievable data rate is approximately 19.97 kbps.

使用香农公式，可以计算带噪信道的最大数据传输速率。公式如下：

$$\text{最大数据速率} = B \times \log_2(1 + SNR)$$

其中，B 是信道带宽，SNR 是信噪比（线性值）。首先将 20 dB 信噪比转换为线性值：

$$SNR (\text{线性}) = 10^{\frac{20}{10}} = 100$$

然后将其代入香农公式：

$$\text{最大数据速率} = 3000 \times \log_2(1 + 100) = 3000 \times \log_2(101) \approx 3000 \times 6.6582 = 19974.6 \text{ 比特每秒}$$

因此，最大可实现的数据传输速率约为 19.97 kbps。

三 2-3

1 问题与翻译

How much bandwidth is there in 0.1 microns of spectrum at a wavelength of 1 micron?

在波长为 1 微米的频谱中，0.1 微米的频谱带宽是多少？

2 解答与翻译

The bandwidth Δf can be calculated using the formula:

$$\Delta f = \frac{c \cdot \Delta \lambda}{\lambda^2}$$

Assuming:

- Speed of light $c = 3 \times 10^8 \text{ m/s}$,
- Wavelength range $\Delta \lambda = 0.1 \times 10^{-6} \text{ m}$,

- Wavelength $\lambda = 1 \times 10^{-6} \text{ m}$.

Substituting these values into the formula:

$$\Delta f = \frac{3 \times 10^8 \text{ m/s} \times 0.1 \times 10^{-6} \text{ m}}{(1 \times 10^{-6} \text{ m})^2} = 3 \times 10^{13} \text{ Hz}$$

Thus, the bandwidth is approximately 30 THz.

频谱带宽 Δf 可以通过以下公式计算:

$$\Delta f = \frac{c \cdot \Delta \lambda}{\lambda^2}$$

假设:

- 光速 $c = 3 \times 10^8 \text{ 米/秒}$,
- 波长范围 $\Delta \lambda = 0.1 \times 10^{-6} \text{ 米}$,
- 波长 $\lambda = 1 \times 10^{-6} \text{ 米}$.

将这些值代入公式:

$$\Delta f = \frac{3 \times 10^8 \text{ 米/秒} \times 0.1 \times 10^{-6} \text{ 米}}{(1 \times 10^{-6} \text{ 米})^2} = 3 \times 10^{13} \text{ 赫兹}$$

因此, 带宽约为 30 太赫兹 (THz)。

四 2-4

1 问题与翻译

It is desired to send a sequence of computer screen images over an optical fiber. The screen is 1920x1200 pixels, each pixel being 24 bits. There are 50 screen images per second. How much bandwidth is needed?

希望通过光纤传送一系列计算机屏幕图像。屏幕分辨率为 1920x1200 像素, 每个像素为 24 比特。每秒传输 50 幅图像。所需的带宽是多少?

2 解答与翻译

Each image contains:

$$1920 \times 1200 \times 24 = 55296000 \text{ bits/image}$$

With 50 images transmitted per second, the data rate is:

$$55296000 \times 50 = 2764800000 \text{ bps} = 2.7648 \text{ Gbps}$$

Thus, the bandwidth required is 2.7648 Gbps.

每张图像的比特数为:

$$1920 \times 1200 \times 24 = 55296000 \text{ 比特/图像}$$

每秒 50 张图像的传输速率为:

$$55296000 \times 50 = 2764800000 \text{ 比特/秒} = 2.7648 \text{ 千兆比特每秒}$$

因此, 所需的带宽为 2.7648 Gbps。

五 2-5

1 问题与翻译

Radio antennas often work best when the diameter of the antenna is equal to the wavelength of the radio wave. Reasonable antennas range from 1 cm to 5 meters in diameter. What frequency range does this cover?

无线电天线的最佳工作状态往往是天线的直径等于无线电波的波长。合理的天线直径范围为 1 厘米到 5 米。这涵盖了什么频率范围？

2 解答与翻译

The frequency is related to the wavelength by the equation:

$$f = \frac{c}{\lambda}$$

For a wavelength of 1 cm (0.01 m):

$$f = \frac{3 \times 10^8}{0.01} = 3 \times 10^{10} \text{ Hz} = 30 \text{ GHz}$$

For a wavelength of 5 meters:

$$f = \frac{3 \times 10^8}{5} = 6 \times 10^7 \text{ Hz} = 60 \text{ MHz}$$

Thus, the frequency range is from 60 MHz to 30 GHz.

频率与波长的关系为：

$$f = \frac{c}{\lambda}$$

对于波长 1 厘米 (0.01 米)：

$$f = \frac{3 \times 10^8}{0.01} = 3 \times 10^{10} \text{ 赫兹} = 30 \text{ 千兆赫兹}$$

对于波长 5 米：

$$f = \frac{3 \times 10^8}{5} = 6 \times 10^7 \text{ 赫兹} = 60 \text{ 兆赫兹}$$

因此，频率范围为 60 MHz 至 30 GHz。

六 2-6

1 问题与翻译

Ten signals, each requiring 4000 Hz, are multiplexed onto a single channel using FDM. What is the minimum bandwidth required for the multiplexed channel? Assume that the guard bands are 400 Hz wide.

十个信号，每个需要 4000 赫兹，通过频分复用 (FDM) 复用到单一信道。复用信道所需的最小带宽是多少？假设保护带宽为 400 赫兹。

2 解答与翻译

Each signal requires 4000 Hz, and there are 10 signals. In addition, 400 Hz of guard band is needed between adjacent signals. There are 9 guard bands because only 9 are required to separate 10 signals.

The total bandwidth required is the sum of the bandwidth of the 10 signals and the 9 guard bands:

$$\text{Total Bandwidth} = (10 \times 4000) + (9 \times 400) = 40000 + 3600 = 43600 \text{ Hz}$$

Thus, the minimum bandwidth required for the multiplexed channel is 43.6 kHz.

每个信号需要 4000 赫兹，有 10 个信号。相邻信号之间的保护带宽为 400 赫兹。

一共有 10 个信号，因此需要为这 10 个信号分配带宽。但是信号本身占用带宽之外，信号之间还需要设置保护带宽。保护带宽的个数是 9 个，因为 10 个信号只需要 9 段保护带宽（类似于 10 个房间只需要 9 面墙来分隔）。

总带宽由两部分组成：

10 个信号，每个信号需要 4000 Hz，所以信号的总带宽是：

$$10 \times 4000 = 40000 \text{ Hz}$$

9 段保护带宽，每段为 400 Hz，所以保护带宽的总宽度是：

$$9 \times 400 = 3600 \text{ Hz}$$

将信号的带宽和保护带宽相加，得到所需的总带宽：

$$\text{Total Bandwidth} = 40000 \text{ Hz} + 3600 \text{ Hz} = 43600 \text{ Hz}$$

即：43.6 kHz。

七 2-7

1 问题与翻译

Why has the PCM sampling time been set at 125 sec?

为什么 PCM 采样时间被设置为 125 微秒？

2 解答与翻译

PCM (Pulse Code Modulation) is typically used in telephone communication, where the standard voice communication bandwidth is about 4 kHz. According to the Nyquist theorem, the sampling rate must be at least twice the bandwidth, which means sampling at a rate of 8000 samples per second. The sampling interval is:

$$\text{Sampling Interval} = \frac{1}{8000} = 125 \text{ microseconds}$$

Thus, the PCM sampling time is set at 125 microseconds to meet the Nyquist sampling requirement for voice communication.

PCM（脉冲编码调制）通常用于电话通信中，标准语音通信带宽约为 4 kHz。根据奈奎斯特定理，采样率必须至少是带宽的两倍，即每秒 8000 次采样。采样间隔为：

$$\text{SamplingInterval} = \frac{1}{8000} = 125 \text{ 微秒}$$

所以显然，PCM 采样时间设置为 125 微秒，数据速率不会超过 8KHz，因此没有必要更频繁地采样。以满足语音通信的奈奎斯特采样要求。

八 2-8

1 问题与翻译

Compare the maximum data rate of a noiseless 4-kHz channel using (a) Analog encoding (e.g., QPSK) with 2 bits per sample. (b) The T1 PCM system.

比较无噪声的 4 千赫信道在以下两种情况下的最大数据速率：

(a) 模拟编码（例如 QPSK）使用每个采样 2 比特。

(b) T1 PCM 系统。

2 解答与翻译

For (a), using QPSK (Quadrature Phase Shift Keying), each symbol represents 2 bits, and there are 4 distinct phase shifts. The Nyquist formula for maximum data rate is:

$$C = 2B \log_2 M$$

For a 4-kHz bandwidth, and $M = 4$ (since each symbol encodes 2 bits), the maximum data rate is:

$$C = 2 \times 4000 \times \log_2 4 = 16,000 \text{ bps}$$

For (b), in the T1 PCM system, the sampling rate is twice the bandwidth, and each sample is encoded using 8 bits. Thus, the maximum data rate is:

$$C = 8000 \times 8 = 64,000 \text{ bps}$$

Thus, the maximum data rate for QPSK is 16 kbps, and for the T1 PCM system, it is 64 kbps.

QPSK (正交相移键控) 是一种相位调制技术, 通过改变参考信号 (载波) 的相位来传输数据。在 QPSK 中, 每个符号代表 2 个比特, 因此有 4 种不同的相位变化 (因为 $2^2 = 4$)。

对于无噪声信道, Nyquist 最大数据速率公式为:

$$C = 2B \log_2 M$$

$$\begin{aligned} C &= 2 \times 4000 \times \log_2 4 \\ &= 2 \times 4000 \times 2 \quad (\text{因为 } \log_2 4 = 2) \\ &= 16,000 \text{ 比特每秒 (bps)} \end{aligned}$$

最大数据速率: 16,000 bps (或 16 kbps)

T1 PCM (脉冲编码调制) 系统是一种数字传输方法, 常用于电信。它通过对模拟信号进行采样并将其转换为数字比特流。T1 系统以信号最高频率的两倍进行采样 (根据 Nyquist 定理), 并用 8 位来表示每个采样。

根据 Nyquist 定理, 最低采样率 (f_s) 必须至少是信号最大频率 (f_{\max}) 的两倍, 以避免混叠:

$$f_s = 2f_{\max}$$

数据速率 = 采样率 \times 每个样本的比特数

$$= 8000 \text{ 样本/秒} \times 8 \text{ 比特/样本} = 64,000 \text{ 比特每秒 (bps)}$$

最大数据速率: 64,000 bps (或 64 kbps)

九 2-9

1 问题与翻译

A CDMA receiver gets the following chips: $(-1 +1 -3 +1 -1 -3 +1 +1)$. Assuming the chip sequences defined in Fig. 2-28(a), which stations transmitted, and which bits did each one send?

一个 CDMA 接收机接收到如下码片序列: $(-1 +1 -3 +1 -1 -3 +1 +1)$ 。假设码片序列如图 2-28(a) 所示, 哪些站点发送了数据, 每个站点发送了哪些比特?

2 解答与翻译

To determine which stations transmitted, we correlate the received chips with each station's chip sequence using the dot product. For each station, we compute the dot product of the received chips with the station's chip sequence. A positive result indicates a transmitted bit of 1, while a negative result indicates a transmitted bit of 0.

For example:

$$S.A = (-1 + 1 - 3 + 1 - 1 - 3 + 1 + 1) \cdot (+1 + 1 + 1 - 1 - 1 + 1 - 1 - 1) = 1$$

Thus, station A transmitted a bit of 1. Repeating this process for the other stations gives:

$$S.B = (-1 + 1 - 3 + 1 - 1 - 3 + 1 + 1) \cdot (+1 + 1 - 1 + 1 - 1 - 1 + 1 - 1) = -1$$

Thus, station B transmitted a bit of 0.

$$S.C = (-1 + 1 - 3 + 1 - 1 - 3 + 1 + 1) \cdot (+1 - 1 + 1 - 1 + 1 + 1 - 1 - 1) = 0$$

Thus, station C did not transmit.

So, the transmitted bits are: A: 1, B: 0, C: no transmission.

通过将接收到的码片序列与每个站点的码片序列进行点积运算，可以确定哪些站点发送了数据。正值表示发送了比特 1，负值表示发送了比特 0。

根据题目，每个站点在发送 0 比特时对应的码片序列为：

- A 发送 0: $+1 + 1 + 1 - 1 - 1 + 1 - 1 - 1$
- B 发送 0: $+1 + 1 - 1 + 1 - 1 - 1 + 1 - 1$
- C 发送 0: $+1 - 1 + 1 - 1 + 1 + 1 - 1 - 1$
- D 不发送: $+0 + 0 + 0 + 0 + 0 + 0 + 0 + 0$

总的信号为各站点传输的信号叠加，接收端的信号为：

$$(+3 + 1 + 1 - 1 - 3 - 1 + 1 + 1)$$

这表示 A、B、C、D 发送信号的组合。

通过计算接收到的信号和每个站点码片序列的内积，可以解码出各站点传输的比特。

$$S.A = (+3 + 1 + 1 - 1 - 3 - 1 + 1 + 1) \cdot (+1 + 1 + 1 - 1 - 1 + 1 - 1 - 1)$$

计算得到：

$$S.A = (1 + 1 + 1 + 1 + 3 + 1 + 1 + 1)/8 = 1$$

因此，A 发送了比特 1。

$$S.B = (+3 + 1 + 1 - 1 - 3 - 1 + 1 + 1) \cdot (+1 + 1 - 1 + 1 - 1 - 1 + 1 - 1)$$

计算得到：

$$S.B = (1 + 1 - 1 - 1 + 3 + 1 + 1 - 1)/8 = -1$$

因此，B 发送了比特 0。

$$S.C = (+3 + 1 + 1 - 1 - 3 - 1 + 1 + 1) \cdot (+1 - 1 + 1 - 1 + 1 + 1 - 1 - 1)$$

计算得到：

$$S.C = (1 - 1 + 1 + 1 - 3 - 1 - 1 + 1)/8 = 0$$

因此，C 没有传输。

$$S.D = (+3 + 1 + 1 - 1 - 3 - 1 + 1 + 1) \cdot (0 + 0 + 0 + 0 + 0 + 0 + 0 + 0)$$

计算得到：

$$S.D = 1$$

因此，D 发送了比特 1。

十 2-10

1 问题与翻译

A cable company decides to provide Internet access over cable in a neighborhood consisting of 5000 houses. The company uses a coaxial cable and spectrum allocation allowing 100 Mbps downstream bandwidth per cable. To attract customers, the company decides to guarantee at least 2 Mbps downstream bandwidth to each house at any time. Describe what the cable company needs to do to provide this guarantee.

一家有线公司决定在包含 5000 户家庭的社区提供有线互联网接入。公司使用同轴电缆和频谱分配，每条电缆允许 100 Mbps 的下行带宽。为了吸引客户，公司决定保证每户家庭在任何时候至少有 2 Mbps 的下行带宽。描述公司需要采取哪些措施来提供这一保证。

2 解答与翻译

To guarantee 2 Mbps per household, the total downstream bandwidth required is:

$$5000 \times 2 \text{ Mbps} = 10000 \text{ Mbps} = 10 \text{ Gbps}$$

Since each cable provides 100 Mbps, the company will need:

$$\frac{10000 \text{ Mbps}}{100 \text{ Mbps/cable}} = 100 \text{ cables}$$

The company must divide the neighborhood into smaller groups, with each group served by one cable. To guarantee 2 Mbps per house, each group should consist of no more than:

$$\frac{100 \text{ Mbps}}{2 \text{ Mbps/house}} = 50 \text{ houses}$$

Thus, the company needs to install 100 separate cables or split the network accordingly to provide sufficient bandwidth.

为了保证 5000 户每户都有 2 Mbps 的带宽，总所需的下行带宽为：

$$5000 \times 2 \text{ Mbps} = 10000 \text{ Mbps} = 10 \text{ Gbps}$$

由于每条电缆提供 100 Mbps，公司的需求为：

$$\frac{10000 \text{ Mbps}}{100 \text{ Mbps/cable}} = 100 \text{ cables}$$

公司需要将社区划分为更小的组，每组由一条电缆提供服务。为了保证 2 Mbps 的带宽，每组的家庭数量不应超过：

$$\frac{100 \text{ Mbps}}{2 \text{ Mbps/house}} = 50 \text{ households}$$

公司应安装 100 条独立电缆或分割网络，以为每户提供足够的带宽。同时，电缆公司需将每根电缆直接连接到光纤节点。