

Design and Analysis Of Algorithm

Tutorial 2

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Section - C

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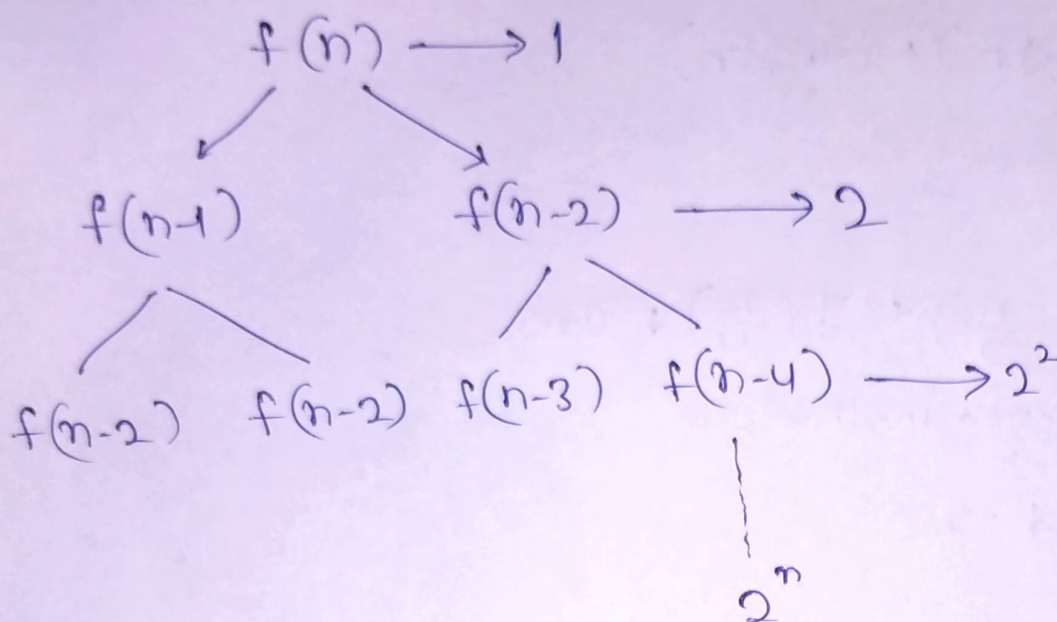
Q.1 void fun (int n)
{
 int j=1, i=0;
 while (i < n) {
 i = i + j;
 j++;
 }
}

j=1	i=0
1	1
2	3
3	6
4	10

Series of i nearly dependent on j as 2^j

Time Complexity = $O(2^n)$

Q.2 Space Complexity = $O(n)$ as clear call of $f(n-1)$



Time Complexity = $O(2^n)$

Q.3 i) $n(\log n)$

```

for(i=0; i<n; i++)
  for(j=0; j<n; j=j*2)
    c++;
  
```

ii) n^3

```

for(i=0; i<n; i++)
  for(j=0; j<n; j++)
    for(k=0; k<n; k++)
      c++;
  
```

iii) $\log(\log n)$

```

int fun(int n)
{
  if(n==1)
    return n;
  }
  
```


else

return fun(\sqrt{n}) + fun(\sqrt{n})

}

Q4. $T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{4}\right) + cn^2$

using master's method:

$$a=2, b=2$$

$$c=1$$

$$f(n) > n^2 \quad f(n/2) > 1$$

$$\text{Time Complexity} = O(n^2)$$

Q5. int fun(int n)

for(int i=1; i<=n; i++) {

for(int j=1; j<=n; j=j+i)

{

// some O(1) task #

}

} }

$$\text{Time Complexity} : O(n\sqrt{n})$$

Q6. for(int i=2; i<=n; i=pow(i,k))

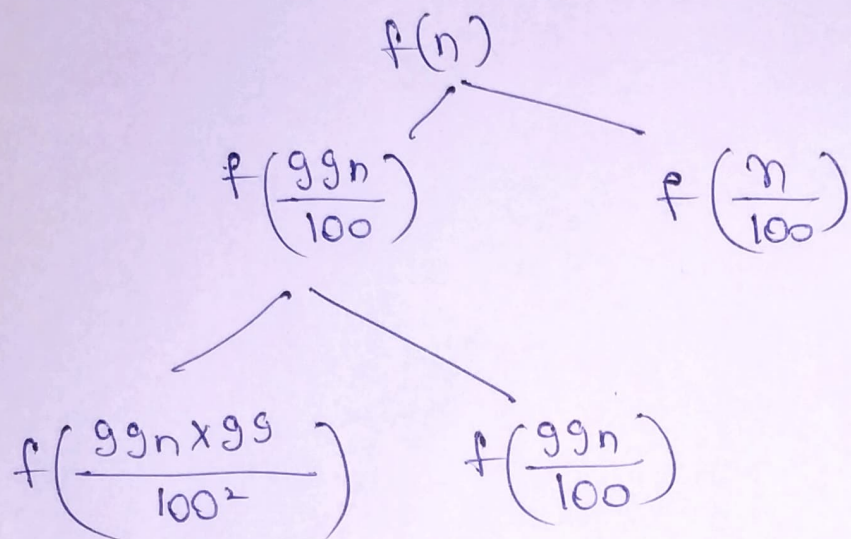
{

// some O(1) statements or expressions

}

where k is a constant (Time Complexity - $O(\log(\log(n)))$)

Q.7. $T(n) = T\left(\frac{99n}{100}\right) + T\left(\frac{n}{100}\right)$



Time Complexity - $O(\log n)$

Q.8 a) $100 < \log(\log(n)) < \log(n) < \sqrt{n} < n < \log(n!) < n \log n < n^2 < 2^n < 2^{2^n} < 4^n < n!$

b) $1 < \log(\log n) < \sqrt{\log n} < \log 2^n < \log n < 2 \log(n) < n \log(n) < n < 2n < 4n < n^2 < \log(n!) < 2(2^n) < n!$

c) $96 < \log_2 n < \log_3 n < \log_5 n < \log n < n \log_e(n) < n \log_2(n) < 8n^2 < 7n^3 < 3^{2^n} < n!$

Shikha