

Data Structures (Sorting) CSE-207

Bubble sort

- Compare each element (except the last one) with its neighbor to the right
 - If they are out of order, swap them
 - This puts the largest element at the very end
 - The last element is now in the correct and final place
- Compare each element (except the last *two*) with its neighbor to the right
 - If they are out of order, swap them
 - This puts the second largest element next to last
 - The last two elements are now in their correct and final places
- Compare each element (except the last three) with its neighbor to the right
 - Continue as above until you have no unsorted elements on the left



Sorting: Bubble sort

- Bubble sort compares the numbers in pairs from left to right exchanging when necessary.
- Suppose we have array A[1],A[2],.....A[N]
- Pass 1
 - Compare A[1] and A[2] arrange them as A[1]<A[2]
 - Then A[2] and A[3] So that A[2]<A[3]
 - ... Continue so that A[N-1]<A[N].
- Pass 2
 - Repeat step 1 with one less comparison. So A[N-1] in its right place
- Pass N-1
 - Compare A[1] and A[2] so thatA[1] < A[2].



Pass-1

```
6, 2, 12, 11, 9, 3
6, 2, 9, 11,
6, 2, 9, 11, 9,
6, 2, 9, 11, 9, 3, 12
```



Pass-2

6,	2,	9,	11	, 9,	3,	12
2,	6,	9,	1	, 9,	3,	12
2,	6,	9,	11	, 9,	3,	12
2,	6,	9,	11	, 9,	3,	12
					, 3,	
2,	6,	9,	9,	3,	11,	12
2,	6,	9,	9,	3,	11,	12



Pass-3

2,	6,	9,	9,	3,	11,	12
					11,	
2,	6,	9,	9,	3,	11,	12
2,	6,	9,	9,	3,	1),	12
2,	6,	9,	3,	9,	11,	12
					1 1	1.0
2,	6,	9,	3,	9,	11,	12



Pass-4

 2, 6, 9, 3, 9, 11, 12

 2, 6, 9, 3, 9, 11, 12

 2, 6, 9, 3, 9, 11, 12

 2, 6, 9, 3, 9, 11, 12

 2, 6, 3, 9, 9, 11, 12

2, 6, 3, 9, 9, 11, 12



Pass-5

 2, 6, 3, 9, 9, 11, 12

 2, 6, 3, 9, 9, 11, 12

 2, 3, 6, 9, 9, 11, 12

2, 3, 6, 9, 9, 11, 12



Pass-6

2, 3, 6, 9, 9, 11, 12 2, 3, 6, 9, 9, 11, 12 2, 3, 6, 9, 9, 11, 12 So we have done!!!!

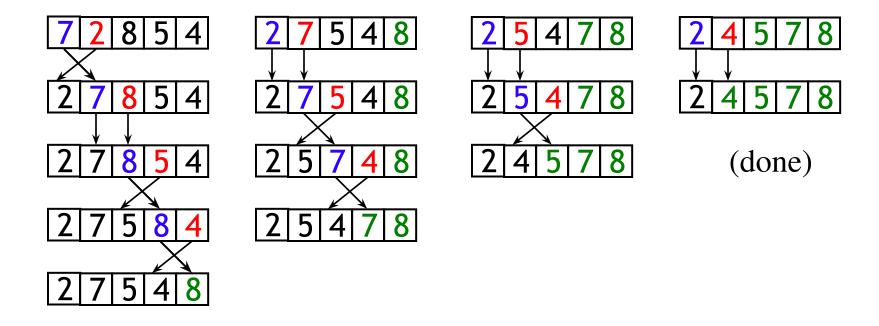
Algorithm: Bubble_Sort(List, N)

Here List is the collection of items and N is the total no. of items.

- Repeat steps 2 and 3 for $I = 1 \dots N$
- Repeat step 3 for J = 1....N
- If List[J] > List[J+1] then swap(List[J], List[J+1]).
- 4. End.



Example of bubble sort





Before we go.....

Swap two variable

```
if (a[j] > a[j+1])
{
    temp = a[j+1];
    a[j+1] = a[j];
    a[j] = temp;
}
```



Code for bubble sort

```
for (j = 0; j < array_size - 1; j++ )</pre>
      if (a[j] > a[j+1])
     temp = a[j+1];
     a[j+1] = a[j];
     a[j] = temp;
```

```
    7
    2
    8
    5
    4

    2
    7
    8
    5
    4

    2
    7
    8
    5
    4

    2
    7
    5
    8
    4

    2
    7
    5
    4
    8
```



Code for bubble sort

```
for (i = 0; i < (array size - 1); ++i)
   for (j = 0; j < array_size - 1; j++ )</pre>
         if (a[j] > a[j+1])
        temp = a[j+1];
        a[j+1] = a[j];
        a[j] = temp;
```



Code for bubble sort

Efficient

```
for (i = 0; i < (array size - 1); ++i)
   for (j = 0; j < array_size - 1-i; j++)
         if (a[j] > a[j+1])
        temp = a[j+1];
        a[j+1] = a[j];
        a[j] = temp;
```



- Selection is a simple sorting algorithm.
- It works by first finding the smallest element using a linear scan and swapping it into the first position in the list.
- Then finding the second smallest element by scanning the remaining elements, and so on.

7.

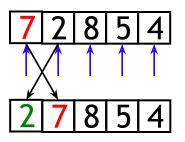
```
Algorithm: Selection_Sort (List, N)
     Repeat steps 2 to 6 for I = 1 to N
1.
        Set Min := I
2.
        Repeat steps 4 and 5 for J = I+1 to N
3.
           If List [J] < List [Min] then
4.
              Set Min := J
5.
        swap( List[ I ], List[Min] )
6.
     End
```

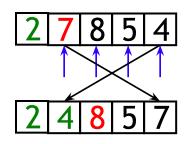


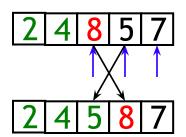
- Given an array of length n,
 - Search elements 0 through n-1 and select the smallest
 - Swap it with the element in location 0
 - Search elements 1 through n-1 and select the smallest
 - Swap it with the element in location 1
 - Search elements 2 through n-1 and select the smallest
 - Swap it with the element in location 2
 - Search elements 3 through n-1 and select the smallest
 - Swap it with the element in location 3
 - Continue in this fashion until there's nothing left to search

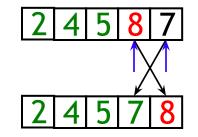


Example and analysis of selection sort









Before we go.....

Find the min value and its position

```
min = 0;
for (j = min+1; j < array_size; j++)
{
   if (a[min]>a[j])
   min = j;
   }
   printf("%d %d", min, a[min]);
```



Before we go.....

- Find the min value and its position
- swap it with first position

```
min = 0;
for (j = min+1; j < array size; j++)
  if (a[min]>a[j])
  min = j;
 printf("%d %d", min, a[min]);
temp = a[0];
a[0] = a[min];
a[min] = temp;
```



```
for (i = 0; i < array size - 1; ++i)
    min = i; //0
    for (j = min+1; j < array size; j++)
     if (a[min]>a[j])
     min = j;
     temp = a[i]; //a[0]
        a[i] = a[min];
        a[min] = temp;
```

- Comparison
- Data Movement
- Sorted

 5
 1
 3
 4
 6
 2

- Comparison
- Data Movement
- Sorted

5 1 3 4 6 2

Large st

- Comparison
- Data Movement
- Sorted

- Comparison
- Data Movement
- Sorted

- Comparison
- Data Movement
- Sorted

 5
 1
 3
 4
 2
 6

- Comparison
- Data Movement
- Sorted

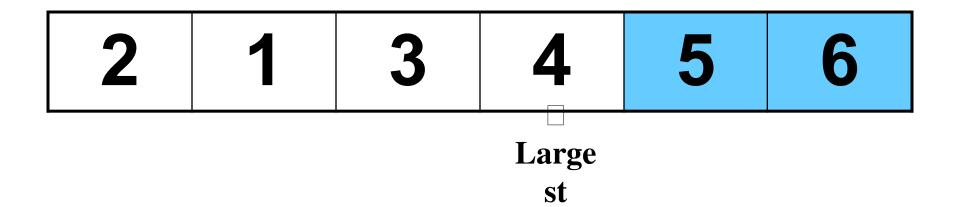
5 1 3 4 2 6

- Comparison
- Data Movement
- Sorted

5 1 3 4 2 6

Largest

- Comparison
- Data Movement
- Sorted



- Comparison
- Data Movement
- Sorted

2 1 3 4 5 6

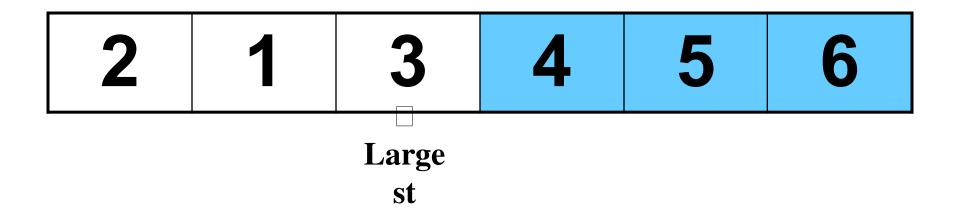
Comparison

__ Data

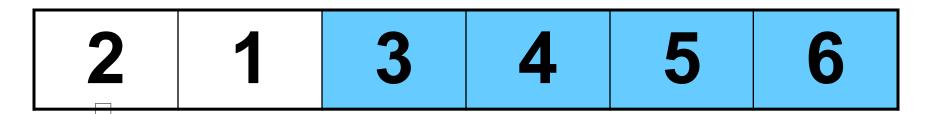
Movement

~ 1

- Comparison
- Data Movement
- Sorted



- Comparison
- Data Movement
- Sorted



Large st

- Comparison
- Data Movement
- Sorted

1 2 3 4 5 6

- Comparison
- Data Movement
- Sorted

1 2 3 4 5 6

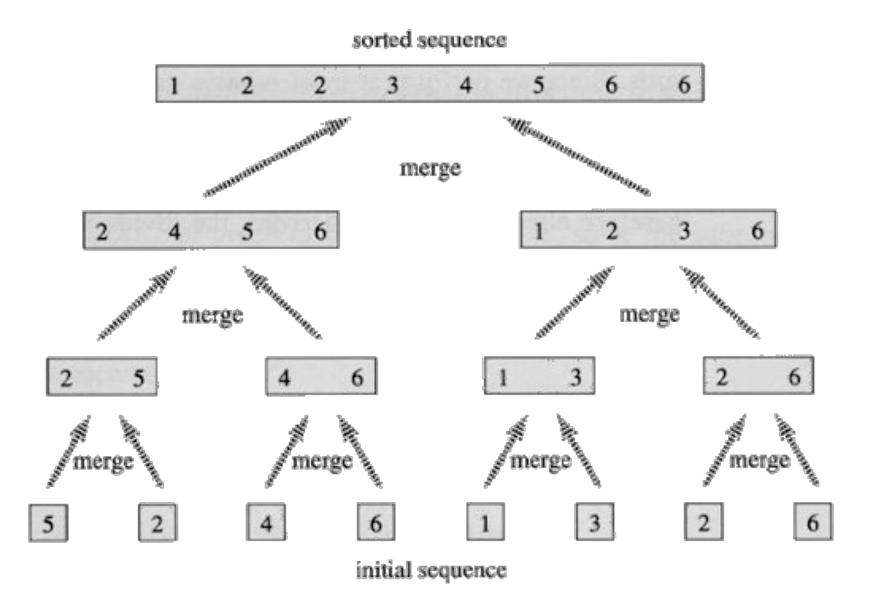
DONE!

- Comparison
- Data Movement
- Sorted

Merge Sort

- Merge sort is a sorting technique based on divide and conquer technique. With worstcase time complexity being O(n log n), it is one of the most respected algorithms.
- Merge sort first divides the array into equal halves and then combines them in a sorted manner.

Merge Sort



Merge Sort

```
MERGE(A, p, q, r)
    n_1 = q - p + 1
2. \quad n_2 = r - q
3. let L[1..n_1 + 1] and R[1..n_2 + 1] be new arrays
4. for i = 1 to n_1
5.
         L[i] = A[p+i-1]
6. for j = 1 to n_2
7. 	 R[j] = A[q+j]
8. L[n_1+1]=\infty
9. R[n_2 + 1] = \infty
10. i = 1
11. j = 1
12. for k = p to r
13.
      if L[i] \leq R[j]
             A[k] = L[i]
14.
15.
             i = i + 1
        else A[k] = R[j]
16.
17.
             j = j + 1
```

Merge Sort

MERGE-SORT(A, p, r)

- 1. if p < r
- $2. q = \lfloor (p+r)/2 \rfloor$
- 3. MERGE-SORT(A, p, q)
- 4. MERGE-SORT(A, q+1, r)
- 5. MERGE(A, p, q, r)

