## Lambda Semantics

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## 1 TYPING RULES

$$\frac{\mathfrak{n}\in\mathbb{Z}}{\Gamma\vdash \textbf{IntE}\ \mathfrak{n}: \textbf{LInt}}\ \text{TyInt}$$

$$\frac{b \in \mathbb{B}}{\Gamma \vdash \mathbf{BoolE} \ b : \mathbf{LBool}} \ \mathsf{TyBool}$$

$$\frac{}{\Gamma \vdash \textbf{UnitE} : \textbf{LUnit}} \ \, \textbf{TyUnit}$$

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash \textbf{Pair} \; e_1 \; e_2 : \textbf{LPair} \; \tau_1 \; \tau_2} \; \; \text{TyProduct}$$

$$\frac{\Gamma \vdash e : \tau_{1}}{\Gamma \vdash \textbf{LeftE } e : \textbf{LEither } \tau_{1} \ \tau_{2}} \ \text{TyCoProductLeft} \\ \frac{\Gamma \vdash e : \tau_{2}}{\Gamma \vdash \textbf{RightE } e : \textbf{LEither } \tau_{1} \ \tau_{2}} \ \text{TyCoProductRight}$$

$$\frac{x:\tau\in\Gamma}{\Gamma\vdash \textbf{Var}\;x:\tau}\; TyVar$$

$$\frac{\Gamma, x : \tau_1 \vdash e : \tau_2}{\Gamma \vdash \textbf{Lambda} \ [\tau_1] \ e : \textbf{LFun} \ \tau_1 \ \tau_2} \ TyLambda$$

$$\frac{\Gamma \vdash e_1 : \textbf{LFun} \; \tau_1 \; \tau_2 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash \textbf{App} \; e_1 \; e_2 : \tau_2} \; \; \text{TyApp}$$

$$\frac{\Gamma \vdash e : LFun \ \tau \ \tau}{\Gamma \vdash Fix \ e : \tau} \ TyFix$$

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash \textbf{binop} \ e_1 \ e_2 : \tau_2} \ \text{TyPrimBinOp}$$

$$\frac{\Gamma \vdash e : \tau_1}{\Gamma \vdash \textbf{op} \; e : \tau_2} \; \text{TyPrimOp}$$

$$\frac{\Gamma \vdash e_1 : \textbf{LBool} \quad \Gamma \vdash e_2 : \tau \quad \Gamma \vdash e_3 : \tau}{\Gamma \vdash \textbf{Cond} \ e_1 \ e_2 \ e_3 : \tau} \ \text{TyCond}$$

$$\frac{\Gamma \vdash e_1 : \textbf{LEither} \ \tau_1 \ \tau_2 \qquad \Gamma \vdash e_2 : \textbf{LFun} \ \tau_1 \ \tau_3 \qquad \Gamma \vdash e_3 : \textbf{LFun} \ \tau_2 \ \tau_3}{\Gamma \vdash \textbf{Case} \ e_1 \ e_2 \ e_3 : \tau_3} \ \ \text{TyCase}$$

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binop	$\tau_1$	$\tau_2$	op	$\tau_1$	$\tau_2$
PrimAdd	LInt	LInt	PrimNeg	LInt	LInt
PrimSub	LInt	LInt	PrimNot	LBool	LBool
PrimMul	LInt	LInt	PrimFst	LPair a b	a
PrimDiv	LInt	LInt	PrimSnd	LPair a b	b
PrimIntEq	LInt	LBool			
PrimBoolEq	LBool	LBool			
PrimAnd	LBool	LBool			
PrimOr	LBool	LBool			

 Table 1: Arguments and result types of primitive functions.

$$\overline{IntE \ v \ \psi \ v} \quad SemInt$$

$$\overline{BoolE \ v \ \psi \ v} \quad SemBool$$

$$\overline{UnitE \ \psi \ ()} \quad SemUnit$$

$$\frac{e_1 \ \psi \ v_1}{Pair \ e_1 \ e_2 \ \psi \ (v_1, v_2)} \quad SemProduct$$

$$\frac{e \ \psi \ v}{LeftE \ e \ \psi \ Left \ v} \quad SemCoProductLeft$$

$$\frac{e \ \psi \ v}{RightE \ e \ \psi \ Right \ v} \quad SemCoProductRight$$

$$\overline{Lambda \ \ e \ \psi \ v} \quad SemLambda$$

$$\frac{e_1 \ \psi \ f : \alpha \to [\alpha/x]e \quad e_2 \ \psi \ v_2 \quad f \ v_2 \ \psi \ v}{App \ e_1 \ e_2 \ \psi \ v} \quad SemApp}$$

$$\frac{e_1 \ \psi \ f : \alpha \to [\alpha/x]e' \quad f \ (Fix \ e) \ \psi \ v}{Fix \ e \ \psi \ v} \quad SemFix$$

$$\frac{e_1 \ \psi \ v_1 \quad e_2 \ \psi \ v_2 \quad binop \ v_1 \ v_2 = v}{binop \ e_1 \ e_2 \ \psi \ v} \quad SemPrimBinOp$$

$$\frac{e_1 \ \psi \ True \quad e_2 \ \psi \ v}{Op \ e \ \psi \ v'} \quad SemCondTrue$$

$$\frac{e_1 \ \psi \ False \quad e_3 \ \psi \ v}{Cond \ e_1 \ e_2 \ e_3 \ \psi \ v} \quad SemCondFalse$$

$$\frac{e_1 \ \psi \ Left \ v \quad e_2 \ \psi \ f : \alpha \to [\alpha/x]e'_2 \quad f \ v \ \psi \ v'}{Case \ e_1 \ e_2 \ e_3 \ \psi \ v'} \quad SemCaseLeft$$

$$\frac{e_1 \ \psi \ Right \ v \quad e_3 \ \psi \ f : \alpha \to [\alpha/x]e'_3 \quad f \ v \ \psi \ v'}{Case \ e_1 \ e_2 \ e_3 \ \psi \ v'} \quad SemCaseRight$$

Lambda	Haskell		
LInt	Int		
LBool	Bool		
LUnit	()		
LPair a b	(a, b)		
LEither a b	Either a b		
LFun	AST -> AST		

 Table 2: Concrete Haskell representation of Lambda types.