

Dynamic Programming: Preliminaries

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Preliminary topics

- Reminders, notation
- The curse of dimensionality
- (General state) Markov chains
- Nonlinear functional equations

References

- Stokey and Lucas (1989)
- Stachurski (2009)

Reminder 1: Distributions

Let S be a nonempty set

A **distribution** ϕ on S is a function that assigns probabilities to subsets of S :

$$\phi(B) = \text{probability mass assigned to } B \subset S$$

I'll often use notation such as

$$\int g(x)\phi(\mathrm{d}x)$$

Think of this as $\mathbb{E} g(X)$ when $X \sim \phi$

Example. If $S = \mathbb{R}$ and ϕ has a density f , then

$$\int g(x)\phi(\mathrm{d}x) = \int_{-\infty}^{\infty} g(x)f(x) \mathrm{d}x$$

Example. If $S = \mathbb{R}^2$ and ϕ has a density f , then

$$\int g(x)\phi(\mathrm{d}x) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} g(x_1, x_2)f(x_1, x_2) \mathrm{d}x_1 \mathrm{d}x_2$$

Example. If $S = \{1, 2, \dots, n\}$, then

$$\int g(x)\phi(\mathrm{d}x) = \sum_{i=1}^n g(x_i)\phi(x_i)$$

Reminder 2: Metric Spaces

Let \mathcal{G} be a nonempty set and let ρ map $\mathcal{G} \times \mathcal{G}$ to \mathbb{R}

The pair (\mathcal{G}, ρ) is called a **metric space** if, for any x, y, z in \mathcal{G} ,

- $\rho(x, y) = 0$ if and only if $x = y$
- $\rho(x, y) = \rho(y, x)$
- $\rho(x, z) \leq \rho(x, y) + \rho(y, z)$

Example. $\mathcal{G} = \mathbb{R}^n$ and $\rho(x, y) = \|x - y\|$

Example. $S \subset \mathbb{R}^n$ and \mathcal{C} is all continuous bounded functions from S to \mathbb{R} ,

$$\rho(f, g) := \sup_{x \in S} |f(x) - g(x)|$$

The three axioms hold for (\mathcal{C}, ρ)

For example, if f, g and h are in \mathcal{C} and $x \in S$, then

$$\begin{aligned} |f(x) - h(x)| &= |f(x) - g(x) - (g(x) - h(x))| \\ &\leq |f(x) - g(x)| + |g(x) - h(x)| \\ &\leq \rho(f, g) + \rho(g, h) \end{aligned}$$

$$\therefore \rho(f, h) \leq \rho(f, g) + \rho(g, h)$$

Optimization and Computers

Some optimization problems are pretty easy

- All functions are differentiable
- Few choice variables (low dimensional)
- Concave (for max) or convex (for min)
- First order / tangency conditions relatively simple

Textbook examples often chosen to have this structure

In reality many problems don't have this structure

- Can't take derivatives
- No analytical solution for FOCs
- Many choice variables (high dimensional)
- Neither concave nor convex — local maxima and minima

Can Computers Save Us?

For any function we can always try brute force optimization

Here's an example for the following function

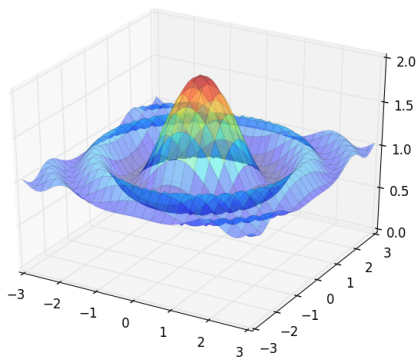


Figure: The function to maximize

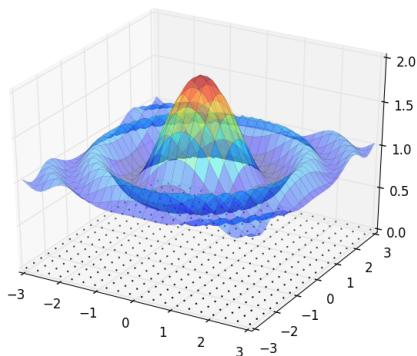


Figure: Grid of points to evaluate the function at

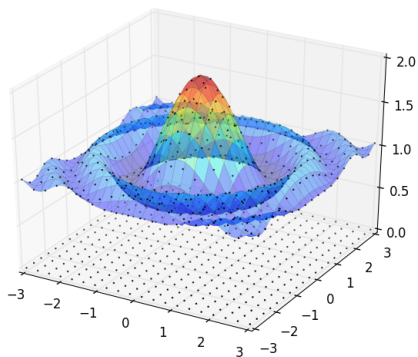


Figure: Evaluations

```
import numpy as np
from numba import jit

grid = np.linspace(-3, 3, 20)
@jit
def compute_max():
    m = -np.inf
    for x in grid:
        for y in grid:
            z = np.cos(x**2 + y**2)/(1 + x**2 + y**2) + 1
            if z > m:
                m = z
    return m

compute_max()
```

Grid size = $20 \times 20 = 400$

Outcomes

- Number of function evaluations = 400
- Time taken = almost zero
- Maximal value recorded = 1.951
- True maximum = 2

Not bad and we can easily do better

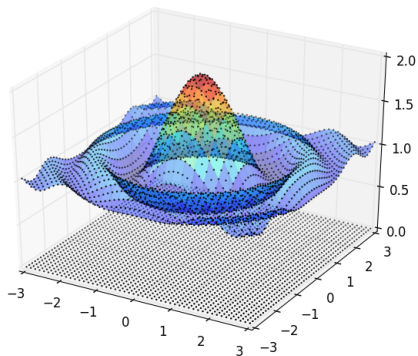


Figure: $50^2 = 2500$ evaluations

- Number of function evaluations = 50^2
- Time taken = $400 \mu s$
- Maximal value recorded = 1.992
- True maximum = 2

So why even study optimization?

The problem is mainly with larger numbers of choice variables

- 3 vars: $\max f(x_1, x_2, x_3)$
- 4 vars: $\max f(x_1, x_2, x_3, x_4)$
- ...

If we have 50 grid points per variable and

- 2 variables then evaluations $= 50^2 = 2500$
- 3 variables then evaluations $= 50^3 = 125,000$
- 4 variables then evaluations $= 50^4 = 6,250,000$
- 5 variables then evaluations $= 50^5 = 312,500,000$
- ...

Example. Recent study: Optimal placement of drinks across vending machines in Tokyo

Approximate dimensions of problem:

- Number of choices for each variable = 2
- Number of choice variables = 1000

Hence number of possibilities = 2^{1000}

How big is that?

```
In [10]: 2**1000
```

```
Out[10]:
```

```
107150860718626732094842504906000181056140481170  
553360744375038837035105112493612249319837881569  
585812759467291755314682518714528569231404359845  
775746985748039345677748242309854210746050623711  
418779541821530464749835819412673987675591655439  
460770629145711964776865421676604298316526243868  
37205668069376
```

Let's say my machine can evaluate about 1 billion possibilities per second

How long would that take?

```
In [16]: (2**1000 / 10**9) / 31556926  # In years
```

```
Out[16]:
```

```
339547840365144349278007955863635707280678989995
899349462539661933596146571733926965255861364854
060286985707326991591901311029244639453805988092
045933072657455119924381235072941549332310199388
301571394569707026437986448403352049168514244509
939816790601568621661265174170019913588941596
```

What about high performance computing?

- more powerful hardware
- faster CPUs
- GPUs
- vector processors
- cloud computing
- massively parallel supercomputers
- ...

Let's say speed up is 10^{12} (wildly optimistic)

```
In [19]: (2**1000 / 10**(9 + 12)) / 31556926
```

```
Out[19]:
```

```
3395478403651443492780079558636357072806789899958
9934946253966193359614657173392696525586136485406
0286985707326991591901311029244639453805988092045
9330726574551199243812350729415493323101993883015
7139456970702643798644840335204916851424450993981
6790601568621661265174170019
```

For comparison:

```
In [20]: 5 * 10**9 # Expected lifespan of sun
```

```
Out[20]: 5000000000
```

Message: There are serious limits to computation

What's required is clever analysis

Exploit what information we have

- without information (oracle) we're stuck
- with information / structure we can do clever things

Examples later on...

Discrete Time Markov Processes

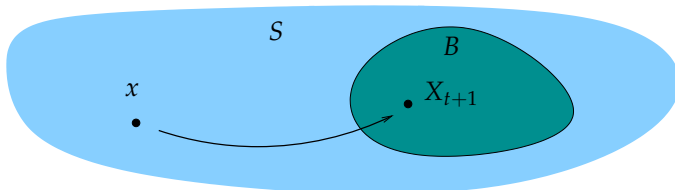
Let

- S be any set (called the **state space**)
- $P(x, dy)$ be a **stochastic kernel** on S — a distribution over S for each $x \in S$

If $\{X_t\}$ is a stochastic process satisfying

$$P(x, B) = \mathbb{P}\{X_{t+1} \in B \mid X_t = x\}$$

then called a **Markov process** with stochastic kernel P

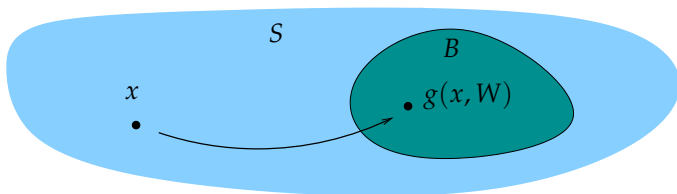


Example. Consider the **stochastic difference equation**

$$X_{t+1} = g(X_t, W_{t+1}) \quad \text{with} \quad \{W_t\} \stackrel{\text{iid}}{\sim} \phi$$

This is a Markov process with stochastic kernel

$$P(x, B) = \phi\{w \in \mathbb{W} : g(x, w) \in B\}$$



Example. Let $S = \mathbb{R}$, let $\{W_t\} \stackrel{\text{iid}}{\sim} N(0, 1)$ and let

$$X_{t+1} = aX_t + b + \sigma W_{t+1}$$

This is a **linear Gaussian** Markov process with kernel

$$P(x, dy) := N(ax + b, \sigma^2)$$

That is, $P(x, B) = \int_B p(x, y) dy$ where

$$p(x, y) = \frac{1}{\sqrt{2\pi}\sigma} \exp \left\{ -\frac{(y - ax - b)^2}{\sigma^2} \right\}$$

Example. Consider the Solow–Swan model

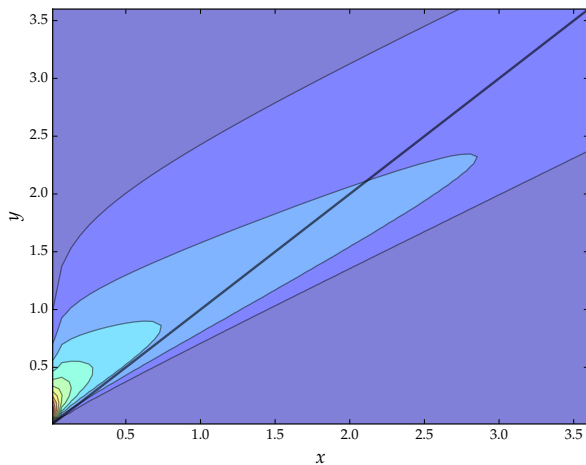
$$k_{t+1} = sf(k_t)W_{t+1} + (1 - \delta)k_t \quad \{W_t\}_{t \geq 1} \stackrel{\text{iid}}{\sim} \phi$$

Here

- k_t takes values in $S = (0, \infty)$
- $s, \delta \in (0, 1)$ and $f(k) > 0$ when $k > 0$

The stochastic kernel is

$$P(k, B) = \phi \{w \in \mathbb{W} \mid sf(k)w + (1 - \delta)k \in B\}$$



Markov Operators

Given stochastic kernel P on S and $h: S \rightarrow \mathbb{R}$, let

$$Ph(x) = \int h(y)P(x, dy)$$

Called the **Markov operator** corresponding to P

Interpretation:

$$Ph(x) = \mathbb{E} [h(X_{t+1}) \mid X_t = x]$$

Solving Equations

What is/are the solution/solutions to these equations?

1. $x = ax + b$

2. $x = x + 1$

3. $x^2 = 1$

Now let x be $n \times 1$ and A be $n \times n$

When does this **vector equation** have a solution ?

$$Ax = b$$

When does this vector equation in \mathbb{R}^n have a unique solution?

$$x = Ax + b$$

When does the **method of successive approximations** converge?

1. pick any $x_0 \in \mathbb{R}^n$
2. set $x_{n+1} = Ax_n + b$ for $n = 0, 1, \dots$

Is there a unique $k > 0$ that solves

$$k = sk^\alpha + (1 - \delta)k$$

Does $k_{n+1} = sk_n^\alpha + (1 - \delta)k_n$ converge to the solution?

Discussion: Consider the **asset price equation**

$$q_t = \beta \mathbb{E}_t[q_{t+1} + \delta(X_{t+1})]$$

Here $\{X_t\}$ is Markov $\sim P$ and $\delta(X_t)$ is current dividend

Guess $q_t = q(X_t)$ and rewrite as

$$q(X_t) = \beta \mathbb{E}_t[q(X_{t+1}) + \delta(X_{t+1})]$$

or as the **functional equation**

$$q(x) = \beta \int q(y)P(x, dy) + \beta \int \delta(y)P(x, dy) \quad (x \in S)$$

- Unique solution? How to solve?

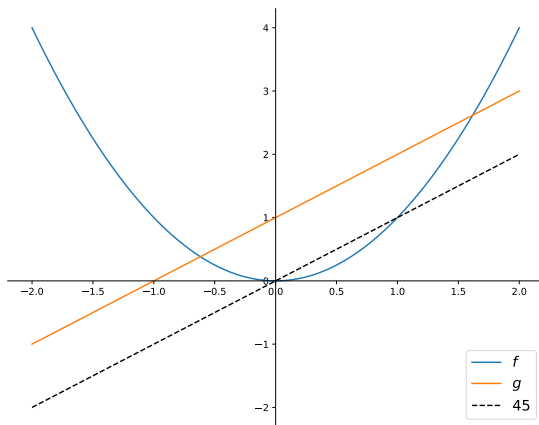
Fixed Points

Let (\mathcal{G}, ρ) be a metric space and let $T: \mathcal{G} \rightarrow \mathcal{G}$

A **fixed point** of T is a point $x^* \in \mathcal{G}$ such that $Tx^* = x^*$

Examples.

- If $f(x) = x^2$ on \mathbb{R} , then 0 and 1 are fixed points
- If $g(x) = x + 1$ on \mathbb{R} , then g has no fixed points on \mathbb{R}



T is called a **contraction map** on (\mathcal{G}, ρ) if

$$\exists \alpha < 1 \quad \text{such that} \quad \rho(Tx, Ty) \leq \alpha \rho(x, y), \quad \forall x, y \in \mathcal{G}$$

Example. $f(x) = \alpha x + b$ on metric space $(\mathbb{R}, |\cdot|)$ with $|\alpha| < 1$, since

$$|f(x) - f(y)| = |\alpha x - \alpha y| = |\alpha| |x - y|$$

Fact. Every contraction T is continuous on \mathcal{G}

Proof: If $x_n \rightarrow x$ in (\mathcal{G}, ρ) , then

$$\rho(Tx_n, Tx) \leq \alpha \rho(x_n, x) \rightarrow 0$$

Fact. If T is a contraction map on (\mathcal{G}, ρ) and $x \in \mathcal{G}$, then $\{T^k x\}$ is Cauchy

Sketch of proof: Along the trajectory $\{T^k x\}$ from x , we have

$$\begin{aligned}\rho(T^{k+1}x, T^k x) &\leq \alpha \rho(T^k x, T^{k-1}x) \\ &\leq \alpha^2 \rho(T^{k-1}x, T^{k-2}x) \\ &\vdots \\ &\leq \alpha^k \rho(Tx, x)\end{aligned}$$

Banach's Fixed Point Theorem

Theorem. If (\mathcal{G}, ρ) is complete and T is a contraction, then T has a unique fixed point x^* in \mathcal{G} and, for all $x \in \mathcal{G}$,

$$\lim_{k \rightarrow \infty} \rho(T^k x, x^*) = 0$$

Proof: Pick any $x \in \mathcal{G}$

The sequence $\{T^k x\}$ is Cauchy and hence converges to some x^*

The point x^* is a fixed point, since

$$Tx^* = T(\lim_k T^k x) = \lim_k T(T^k x) = \lim_k T^{k+1} x = x^*$$

Regarding uniqueness, if x^* and x^{**} are fixed points of T , then

$$\rho(x^*, x^{**}) = \rho(Tx^*, Tx^{**}) \leq \alpha \rho(x^*, x^{**})$$

$$\therefore \rho(x^*, x^{**}) = 0$$

$$\therefore x^* = x^{**}$$

Application: Asset Pricing

Recall the asset pricing equation

$$q(x) = \beta Pq(x) + \beta P\delta(x)$$

Let \mathcal{C} be all continuous bounded functions on S with metric

$$\rho(f, g) := \sup_{x \in S} |f(x) - g(x)|$$

Let P have the **Feller property**, which is to say that

$$h \in \mathcal{C} \implies Ph \in \mathcal{C}$$

Let $\delta \in \mathcal{C}$ and let $\beta \in (0, 1)$

Claim: The asset pricing equation

$$q(x) = \beta Pq(x) + \beta P\delta(x) \quad (x \in S)$$

has a unique solution $q^* \in \mathcal{C}$

We often write this as $q = \beta Pq + \beta P\delta$

Remarks:

- Equivalent: the operator $T: \mathcal{C} \rightarrow \mathcal{C}$ defined by

$$Tq = \beta Pq + \beta P\delta$$

has a unique fixed point in \mathcal{C}

To prove this we need to show that

1. $Tq = \beta Pq + \beta P\delta$ is in \mathcal{C} when $q \in \mathcal{C}$
2. the pair (\mathcal{C}, ρ) forms a complete metric space
3. T is a contraction map on (\mathcal{C}, ρ)

Here (1) follows from assumption and the proof of (2) is omitted

Regarding (3), fix any q, q' in \mathcal{C} and any $x \in S$

We have,

$$\begin{aligned} |Tq(x) - Tq'(x)| &= |\beta Pq(x) + \beta P\delta(x) - \beta Pq'(x) - \beta P\delta(x)| \\ &= \beta \left| \int q(y)P(x, dy) - \int q'(y)P(x, dy) \right| \\ &= \beta \left| \int [q(y) - q'(y)]P(x, dy) \right| \\ &\leq \beta \int |q(y) - q'(y)|P(x, dy) \\ &\leq \beta \sup_y |q(y) - q'(y)| = \beta \rho(q, q') \end{aligned}$$

Taking the supremum with respect to x completes the proof