

Numerical Methods for Computational Science and Engineering

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URL: <http://www.sam.math.ethz.ch/~hiptmair/tmp/NumCSE/NumCSE15.pdf>

IV. Approximation of functions in 1D

Approximation of functions: Generic view

Given: function $f : D \subset \mathbb{R}^n \mapsto \mathbb{R}^d$ (often in procedural form $y = \text{feval}(x)$, Rem. 3.1.4)

Goal: Find a "simple" (*) function $\tilde{f} : D \mapsto \mathbb{R}^d$ such that the approximation error $f - \tilde{f}$ is "small" (♣)

Focus : $n = d = 1$

simple : (piecewise) polynomial

accurate : $\|f - \tilde{f}\|$ small, $\|\cdot\| \hat{=}$ norm on $C^0(D)$,

eg. supremum norm* [uniform approximation] $\|f\|_{L^\infty(D)} := \sup_{x \in D} |f(x)|$ //

L^2 -norm $\|f\|_{L^2(D)}^2 := \int_D |f(x)|^2 dx$

Approximation scheme : $A : X \rightarrow V$
 $\uparrow \quad \quad \uparrow$
spaces of functions

e.g. : $X = C^0(I)$, $I \subset \mathbb{R}$

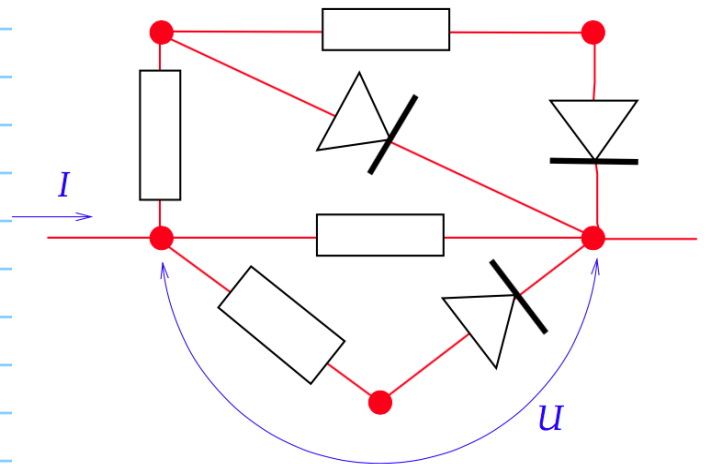
$V = \mathcal{P}_n$ or $V = \mathcal{S}_{3,n}$

Application: Model reduction

Two-port circuit \triangleright

Goal: Model this as a single circuit element described $I = I(U)$

\hookrightarrow computable
 $[\hat{=}$ non-linear solve]



② M.R.: Replace $I(U)$ with $\tilde{I}(U)$, based on a few evaluations of $I(U)$, that can be evaluated fast

Approximation by interpolation

Interpolation scheme + sampling \rightarrow approximation scheme

$$f: I \subset \mathbb{R} \rightarrow \mathbb{K} \xrightarrow{\text{sampling}} (t_i, y_i := f(t_i))_{i=0}^m \xrightarrow{\text{interpolation}} \tilde{f} := I_T \mathbf{y} \quad (\tilde{f}(t_i) = y_i).$$

\uparrow
free choice of nodes $t_i \in I$
 \uparrow
interpolation scheme

4.1 Approximation by global polynomial

Example from analysis: Taylor approximation

$$f(t) = \underbrace{\sum_{j=0}^n \frac{f^{(j)}(t_0)}{j!} (t-t_0)^j}_{\in \mathcal{P}_n} + \underbrace{f^{(n+1)}(\xi) \frac{(t-t_0)^{n+1}}{(n+1)!}}_{\text{remainder, small for } |t-t_0| \leq 1}$$

4.1.1. Polynomial Approximation: Theory

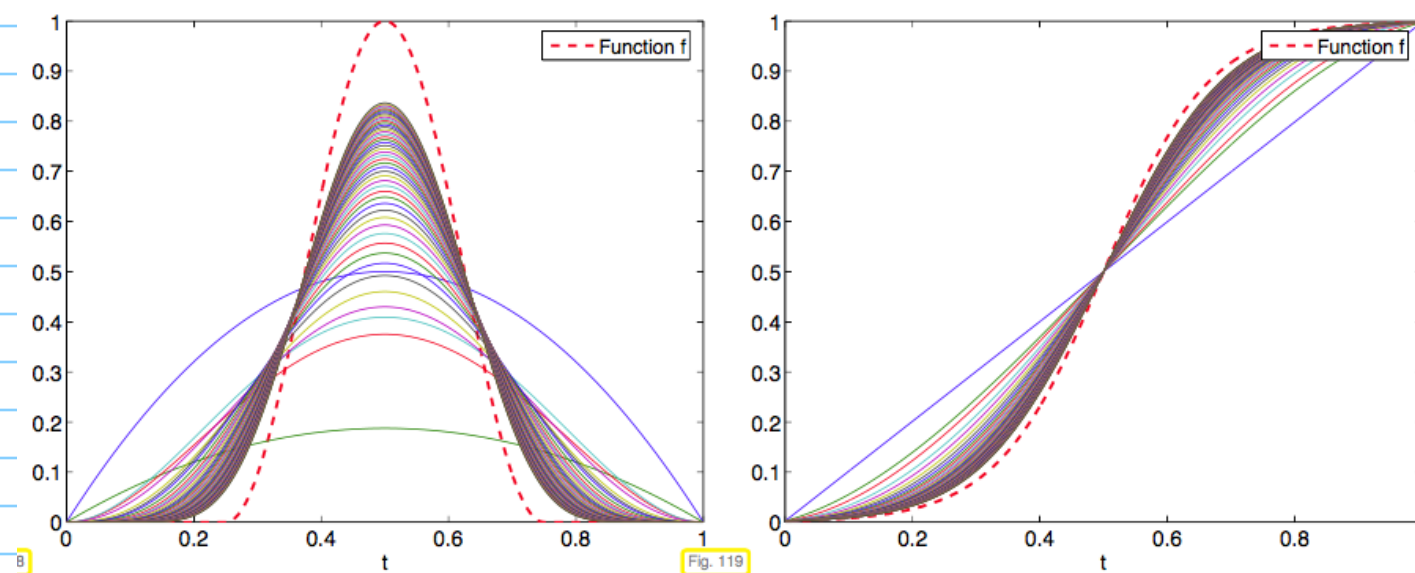
\rightarrow polynomials can **uniformly** approximate any $f \in C^k(I)$, $I \subset \mathbb{R}$ closed, $k \geq 0$, along with their first k derivatives

Theorem 4.1.6. Uniform approximation by polynomials

For $f \in C^0([0, 1])$, define the n -th Bernstein approximant as

$$p_n(t) = \sum_{j=0}^n f(j/n) \binom{n}{j} t^j (1-t)^{n-j}, \quad p_n \in \mathcal{P}_n. \quad (4.1.7)$$

It satisfies $\|f - p_n\|_\infty \rightarrow 0$ for $n \rightarrow \infty$. If $f \in C^m([0, 1])$, then even $\|f^{(k)} - p_n^{(k)}\|_\infty \rightarrow 0$ for $n \rightarrow \infty$ and all $0 \leq k \leq m$.



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Best approximation:

Definition 5.1.10. (Size of) best approximation error

Let $\|\cdot\|$ be a (semi-)norm on a space X of functions $I \mapsto \mathbb{K}$, $I \subset \mathbb{R}$ an interval. The (size of the) **best approximation error** of $f \in X$ in the space \mathcal{P}_k of polynomials of degree $\leq k$ with respect to $\|\cdot\|$ is

$$\text{dist}_{\|\cdot\|}(f, \mathcal{P}_k) := \inf_{p \in \mathcal{P}_k} \|f - p\|.$$

Theorem 4.1.11. L^∞ polynomial best approximation estimate

If $f \in C^r([-1, 1])$ (r times continuously differentiable), $r \in \mathbb{N}$, then

$$\inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([-1, 1])} \leq \underbrace{(1 + \pi^2/2)^r \frac{(n-r)!}{n!}}_{\text{quantitative bound}} \|f^{(r)}\|_{L^\infty([-1, 1])}.$$

crucial role of smoothness

Asymptotics for degree $n \rightarrow \infty$: $O(n^{-r})$

▷ **Algebraic conv. with rate r** in the polynomial degree

In general: No quantitative conv. bounds without smoothness.

Generalization to $[a, b]$:

Tool: Affine hf. $\phi: \begin{cases} [-1, 1] & \longrightarrow [a, b] \\ t & \longrightarrow a + \frac{1}{2}(b-a)(t+1) \end{cases}$

Pullback $\phi^*: \begin{cases} C^0([a, b]) & \longrightarrow C^0([-1, 1]) \\ f & \longrightarrow \{t \mapsto f(\phi(t))\} \end{cases}$

Note: $\phi^*: \mathcal{P}_n \longrightarrow \mathcal{P}_n$ bijective (ii)
 $\|\phi^* f\|_{L^\infty([-1, 1])} = \|f\|_{L^\infty([a, b])}$ (i)
 $\Rightarrow \|\phi^* f^{(r)}\|_{L^\infty([-1, 1])} = \left(\frac{1}{2}(b-a)\right)^r \|f^{(r)}\|_{L^\infty([a, b])}$

(Chain rule $(\phi^* f)' = (f \circ \phi)' = (f' \circ \phi) \phi'$)

$$\begin{aligned} \inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([a, b])} &\stackrel{(i)}{=} \inf_{p \in \mathcal{P}_n} \|\phi^*(f - p)\|_{L^\infty([-1, 1])} \\ &= \inf_{p \in \mathcal{P}_n} \|\phi^* f - \phi^* p\|_{L^\infty([-1, 1])} \end{aligned}$$

$\phi^*(\mathcal{P}_n) = \mathcal{P}_n$!

$$\begin{aligned} &\stackrel{(ii)}{=} \inf_{p \in \mathcal{P}_n} \|\phi^* f - p\|_{L^\infty([-1, 1])} \\ &\stackrel{\text{Thm. 4.1.11}}{\leq} C n^{-r} \|(\phi^* f)^{(r)}\|_{L^\infty([-1, 1])} \\ &\leq C n^{-r} \left(\frac{b-a}{2}\right)^r \|f^{(r)}\|_{L^\infty([a, b])} \end{aligned}$$

4.1.2. Error estimates for polynomial interpolation

Recall : interpolation scheme \Rightarrow approximation scheme

given on $[-1, 1]$, based on nodes $\{\hat{t}_0, \dots, \hat{t}_n\}$

Generalization to $[a, b]$: use nodes $t_j = \Phi(\hat{t}_j)$, where $\Phi: [-1, 1] \rightarrow [a, b]$ affine mapping from above.

[Error estimates on $[a, b]$ by transformation techniques]

Definition 4.1.25. Lagrangian (interpolation polynomial) approximation scheme

Given an interval $I \subset \mathbb{R}$, $n \in \mathbb{N}$, a node set $\mathcal{T} = \{t_0, \dots, t_n\} \subset I$, the **Lagrangian (interpolation polynomial) approximation scheme** $L_{\mathcal{T}}: C^0(I) \rightarrow \mathcal{P}_n$ is defined by

$$L_{\mathcal{T}}(f) := l_{\mathcal{T}}(\mathbf{y}) \in \mathcal{P}_n \quad \text{with} \quad \mathbf{y} := (f(t_0), \dots, f(t_n))^T \in \mathbb{K}^{n+1}.$$

Lagrangian interpolation

sampling

New: Given families of node sets $(\mathcal{T}_n)_{n \in \mathbb{N}}$, $\#\mathcal{T}_n = n+1$
 \Rightarrow family of polynomial A.S. $L_{\mathcal{T}_n}: C^0(I) \rightarrow \mathcal{P}_n$

* Example: equispace nodes $\mathcal{T}_n := \{t_j^{(n)} := a + (b-a)\frac{j}{n} : j = 0, \dots, n\} \subset I$

New aspect : **Convergence** $\|f - L_{\mathcal{T}_n} f\| \leq T(n)$

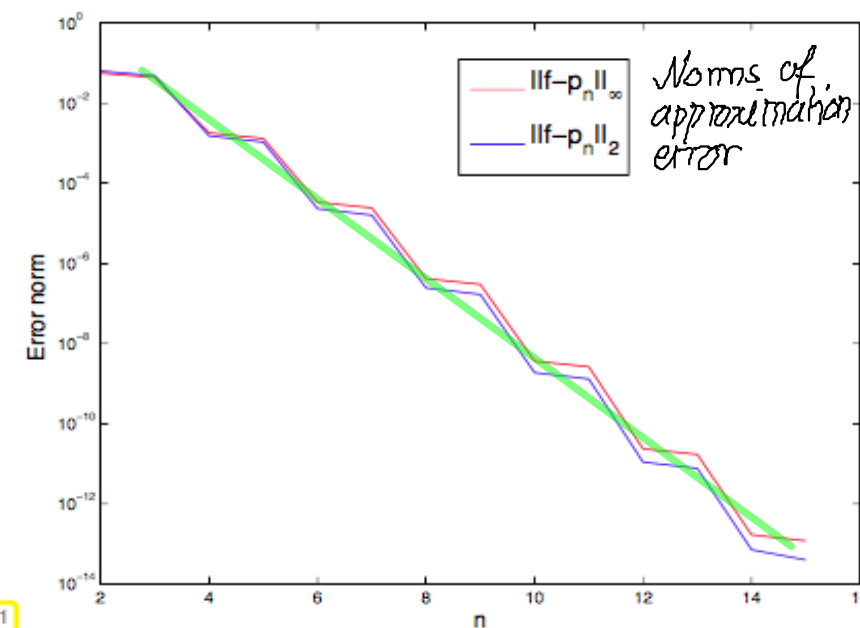
asymptotically for $n \rightarrow \infty$,
 [$n \hat{=}$ polynomial degree]

Example:

$f(t) = \sin t$
 $I = [0, \pi]$
 equispaced nodes

$$E_n := \|f - L_{\mathcal{T}_n} f\|$$

semi-logarithmic scale \triangleright



From plot : $\log E_n \approx C - kn$, $k > 0$
 $E_n \approx C(e^{-k})^n$

Empiric asymptotics : $E_n = O(q^n)$ with $0 < q < 1$
 $\hat{=}$ exponential conv.

$$\exists C \neq C(n) > 0: \|f - L_{\mathcal{T}_n} f\| \leq C T(n) \quad \text{for } n \rightarrow \infty. \quad (4.1.30)$$

Definition 4.1.31. Types of asymptotic convergence of approximation schemes

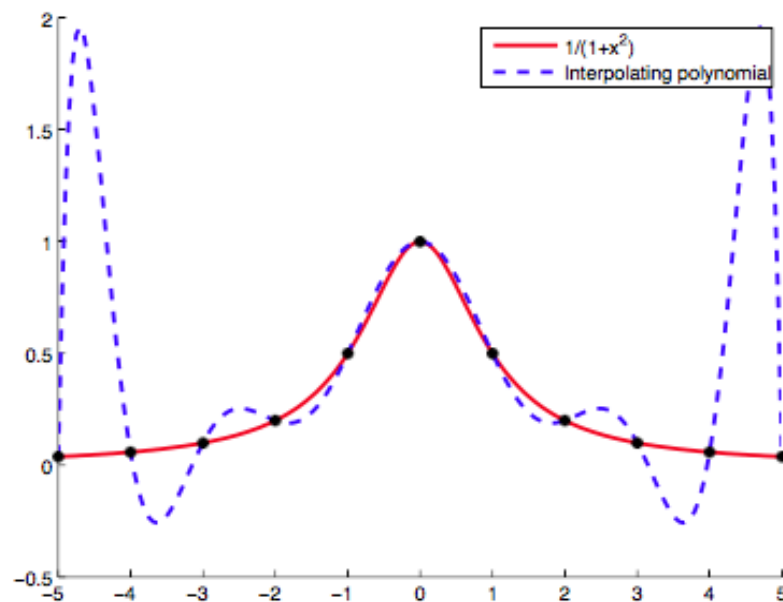
Writing $T(n)$ for the bound of the norm of the interpolation error according to (4.1.30) we distinguish the following types of asymptotic behavior:

$\exists p > 0$: $T(n) \leq n^{-p}$: algebraic convergence, with rate $p > 0$, $\forall n \in \mathbb{N}$.
 $\exists 0 < q < 1$: $T(n) \leq q^n$: exponential convergence,

The bounds are assumed to be sharp in the sense, that no bounds with larger rate p (for algebraic convergence) or smaller q (for exponential convergence) can be found.

\rightarrow see Thm. 4.1.11.

⑤ Will Lagrangian interpolation always converge?



$f(t) = \frac{1}{1+t^2}$
on $[-5, 5]$:
 $\|f - L_{\mathcal{T}} f\|_{\infty} \rightarrow \infty$
for equidistant nodes



Warning + promise

Theory ahead!

The main theorem:

Theorem 4.1.37. Representation of interpolation error [7, Thm. 8.22], [19, Thm. 37.4]

We consider $f \in C^{n+1}(I)$ and the Lagrangian interpolation approximation scheme (\rightarrow Def. 4.1.25) for a node set $\mathcal{T} := \{t_0, \dots, t_n\}$. Then, for every $t \in I$ there exists a $\tau_t \in]\min\{t, t_0, \dots, t_n\}, \max\{t, t_0, \dots, t_n\}[$ such that

$$\underbrace{f(t) - L_{\mathcal{T}}(f)(t)}_{\text{approximation error}} = \frac{f^{(n+1)}(\tau_t)}{(n+1)!} \cdot \prod_{j=0}^n (t - t_j). \quad (4.1.38)$$

Proof: $w_{\mathcal{T}}(t) := \prod_{j=0}^n (t - t_j)$ ["nodal polynomial"] $\in \mathcal{P}_{n+1}$

Fix $t \in I \setminus \mathcal{T}$

$$c \in \mathbb{R}: \quad f(t) - (L_{\mathcal{T}} f)(t) = c \cdot \underbrace{w_{\mathcal{T}}(t)}_{\neq 0}$$

Auxiliary function: $\varphi(x) := f(x) - (L_{\mathcal{T}} f)(x) - c w_{\mathcal{T}}(x)$

$\Rightarrow \begin{matrix} \varphi(t_j) = 0 \\ \varphi(b) = 0 \end{matrix}$ by def. of $c \Rightarrow \varphi$ has $n+2$ distinct zeros

\Rightarrow [Rolle's theorem] φ' has $n+1$ distinct zeros

\vdots
 $\varphi^{(n+1)}$ has at least one zero

$$\exists \tau_t: \varphi^{(n+1)}(\tau) = f^{(n+1)}(\tau) - c(n+1)! = 0$$

$$c = \frac{1}{(n+1)!} f^{(n+1)}(\tau)$$

□

Theorem 4.1.37. Representation of interpolation error [7, Thm. 8.22], [19, Thm. 37.4]

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$$f(t) - L_{\mathcal{T}}(f)(t) = \frac{f^{(n+1)}(\tau_t)}{(n+1)!} \cdot \prod_{j=0}^n (t - t_j). \quad (4.1.38)$$

Estimates from Thm 4.1.37 by taking maximum over t

$$\Rightarrow \|f - L_{\mathcal{T}}f\|_{\infty} \leq \max_{\tau \in I} |f^{(n+1)}(\tau)| \frac{1}{(n+1)!} \|W_{\mathcal{T}}\|_{\infty}$$

$$\|f - L_{\mathcal{T}}f\|_{L^{\infty}(I)} \leq \frac{\|f^{(n+1)}\|_{L^{\infty}(I)}}{(n+1)!} \max_{t \in I} |(t - t_0) \cdots (t - t_n)| \quad (*)$$

smoothness of f is essential!

• Back to $f(t) = \sin t \Rightarrow \|f^{(n)}\|_{\infty} \leq 1 \quad \forall n$
 $f \in C^{\infty}$

For equidistant nodes in $[0, \pi]$

$$(*) \Rightarrow \|\sin - L_{\mathcal{T}_n} \sin\|_{\infty} \leq \frac{1}{(n+1)!} \max_{0 \leq t \leq \pi} \prod_{j=0}^n (t - \frac{\pi}{n+1} j) \leq \frac{1}{(n+1)!} \left(\frac{\pi}{n}\right)^{n+1}$$

attains max for $t = \frac{\pi}{2n}$

\rightarrow more than exponential cvg.

$$f(t) = \frac{1}{1+t^2} \Rightarrow \|f^{(n)}\|_{\infty, [-5,5]} \geq 2^n n!$$

Bound from (*) can blow up as $n \rightarrow \infty$

Interpolation error estimates for analytic functions

Definition 5.1.54. Analyticity of a complex valued function

Let $D \subset \mathbb{C}$ be an open set in the complex plane. A function $f : D \rightarrow \mathbb{C}$ is called **analytic/holomorphic** in D , if $f \in C^{\infty}(D)$ and it possesses a convergent Taylor series at every point $z \in D$.

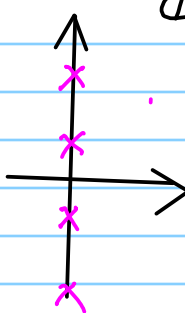
Extension to \mathbb{C} : Reinterpret $f : I \subset \mathbb{R} \rightarrow \mathbb{R}$ as a restriction $f : D \subset \mathbb{C} \rightarrow \mathbb{C}$ to $I \subset D$
 ["just use the same formula"]

Example: $f(t) = \frac{1}{1+t^2} \Rightarrow f(z) = \frac{1}{1+z^2}$

\rightarrow a rational function, analytic in $\mathbb{C} \setminus \{\pm i\}$

Entire functions: $e^z, \sin z, \cos z$: analytic in \mathbb{C}

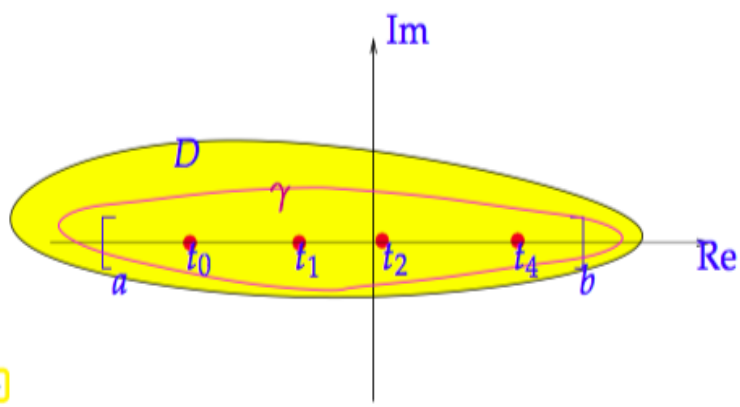
Example: $f(z) = \frac{1}{1+e^z} = \frac{1}{1+e^{x+iy}} = \frac{1}{1+e^x \cdot e^{iy}}$
 $e^z = e^x \cdot e^{iy} \stackrel{!}{=} -1 : |e^x| = 1 \Rightarrow x = 0$
 $e^{-iy} = -1 \Leftrightarrow y = (2k+1)\pi, k \in \mathbb{Z}$
 \Rightarrow analytic in $\mathbb{C} \setminus \{(2k+1)\pi i, k \in \mathbb{Z}\}$



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Example : $f(z) = U^T (A - zI)^{-1} b, z \in \mathbb{C}, A \in \mathbb{R}^{n,n}, U, b \in \mathbb{R}^n$
 Analytic in $\mathbb{C} \setminus \{z : (A - zI) \text{ not invertible}\}$
 $= \mathbb{C} \setminus \{z \text{ is an eigenvalue of } A\}$

Example : $f(z) = \sqrt{z+1}$ is analytic $\mathbb{C} \setminus [-\infty, -1]$
 Note : $\sqrt{\cdot} : \mathbb{C} \setminus [-\infty, 0] \rightarrow \mathbb{C}$ analytic



Assumption 4.1.57. Analyticity of interpoland
 We assume that the interpoland $f : I \rightarrow \mathbb{C}$ can be extended to a function $f : D \subset \mathbb{C} \rightarrow \mathbb{C}$, which is **analytic** (\rightarrow Def. 4.1.54) on an open set $D \subset \mathbb{C}$ with $[a, b] \subset D$.

Theorem 4.1.55. Residue theorem [30, Ch. 13]

Let $D \subset \mathbb{C}$ be an open set, $\gamma \subset D$ a simple closed smooth curve (in the complex plane), contractible in D , and $\Pi \subset D$ a finite set*. Then for each function f that is analytic in $D \setminus \Pi$ holds

* enclosed by γ

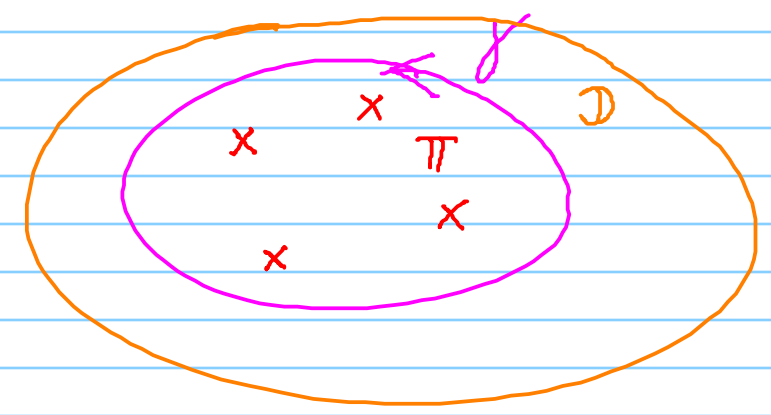
$$\frac{1}{2\pi i} \int_{\gamma} f(z) dz = \sum_{p \in \Pi} \text{res}_p f,$$

where $\text{res}_p f$ is the **residual** of f in $p \in \mathbb{C}$. contour integral

Contour integral $\int_{\gamma} f(z) dz = \int_{\alpha}^{\beta} f(\gamma(\tau)) \cdot \frac{d\gamma}{d\tau} d\tau$
 $\gamma : [\alpha, \beta] \rightarrow \mathbb{C}$

$$\rightarrow \left| \int_{\gamma} f(z) dz \right| \leq |\gamma| \max_{z \in \gamma} |f(z)|$$

\uparrow length



Lemma 4.1.56. Residual formula for quotients

let g and h be complex valued functions that are both analytic in a neighborhood of $p \in \mathbb{C}$, and satisfy $h(p) = 0, h'(p) \neq 0$. Then

$$\text{res}_p \frac{g}{h} = \frac{g(p)}{h'(p)}. \quad [\text{residual in a simple pole}]$$

$$L_j(t) = \prod_{k=0, k \neq j}^n \frac{t - t_k}{t_j - t_k} = \frac{w(t)}{(t - t_j) \prod_{k=0, k \neq j}^n (t_j - t_k)} = \frac{w(t)}{(t - t_j) w'(t_j)}$$

$$L_{\mathcal{T}} f = \sum_{j=0}^n f(t_j) L_j$$

Lagrange polynomials for node set $\{t_0, \dots, t_n\}$

8) $w(z) = (z - t_0) \cdots (z - t_n) \in \mathcal{P}_{n+1}$

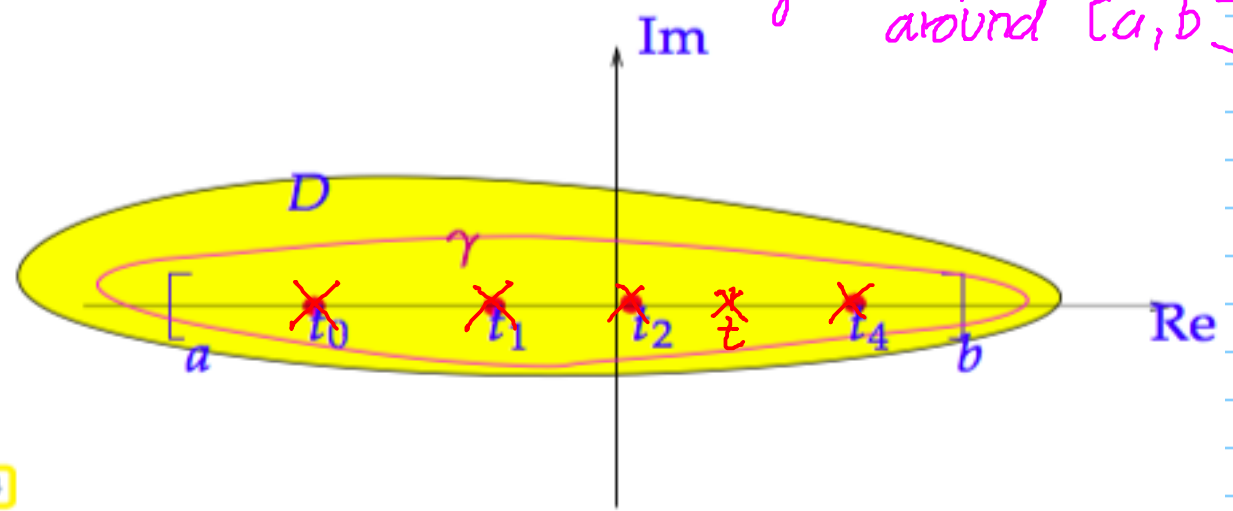
Fix $t \in [a, b] \setminus J$:

Auxiliary function:

$g_t(z) := \frac{f(z)}{(z-t)w(z)}, z \in \mathcal{D} \setminus \{t, t_0, \dots, t_n\}$

simple poles there!

$\gamma \triangleq$ simple closed contour around $[a, b]$



$$\begin{aligned} \frac{1}{2\pi i} \int_{\gamma} g_t(z) dz &= \text{res}_t g_t + \sum_{j=0}^n \text{res}_{t_j} g_t \\ &= \frac{f(t)}{w(t)} + \sum_{j=0}^n \frac{f(t_j)}{(t_j - t)w'(t_j)} \quad [\cdot w(t)] \end{aligned}$$

$L_j(t) = \prod_{k=0, k \neq j}^n \frac{t - t_k}{t_j - t_k} = \frac{w(t)}{(t - t_j) \prod_{k=0, k \neq j}^n (t_j - t_k)} = \frac{w(t)}{(t - t_j)w'(t_j)}$

$w(t) \frac{1}{2\pi i} \int_{\gamma} g_t(z) dz = f(t) - \underbrace{\sum_{j=0}^n f(t_j) L_j(t)}_{\text{interpolation error}}$

$f(t) = \underbrace{- \sum_{j=1}^n f(t_j) \frac{w(t)}{(t_j - t)w'(t_j)}}_{\text{polynomial interpolant!}} + \underbrace{\frac{w(t)}{2\pi i} \int_{\gamma} g_t(z) dz}_{\text{interpolation error!}}$

A representation formula! $|\cdot| \leq \|w\|_{\infty} \frac{|t|}{2\pi} \max_{a \leq t \leq b} \|g_t\|_{\infty, \gamma}$

Towards estimates:

Use $g_t :=$

$|f(t) - L_T f(t)| \leq \left| \frac{w(t)}{2\pi i} \int_{\gamma} \frac{f(z)}{(z-t)w(z)} dz \right| \leq \frac{|\gamma|}{2\pi} \frac{\max_{a \leq \tau \leq b} |w(\tau)|}{\min_{z \in \gamma} |w(z)|} \cdot \frac{\max_{z \in \gamma} |f(z)|}{\text{dist}([a, b], \gamma)}$

$\text{dist}([a, b], \gamma) = \inf_{z \in \gamma, a \leq t \leq b} |z - t|$

The only term depending on n , and $\{t_0, \dots, t_n\}$ only this depends on f

"derivative free bound"

4.1.3. Chebyshev Interpolation

Goal: Optimal placement* of nodes t_j for Lagrange interpolation

Theory :
$$\|f - L_T f\|_{L^\infty(I)} \leq \frac{1}{(n+1)!} \|f^{(n+1)}\|_{L^\infty(I)} \|w\|_{L^\infty(I)},$$

with $w(t) := (t - t_0) \cdots (t - t_n)$

* make $\|w\|_{L^\infty(I)}$ as small as possible \rightarrow f-independent optimal nodes
 nodes enter only here

Optimal choice of interpolation nodes independent of interpoland



Idea: choose nodes t_0, \dots, t_n such that $\|w\|_{L^\infty(I)}$ is minimal!

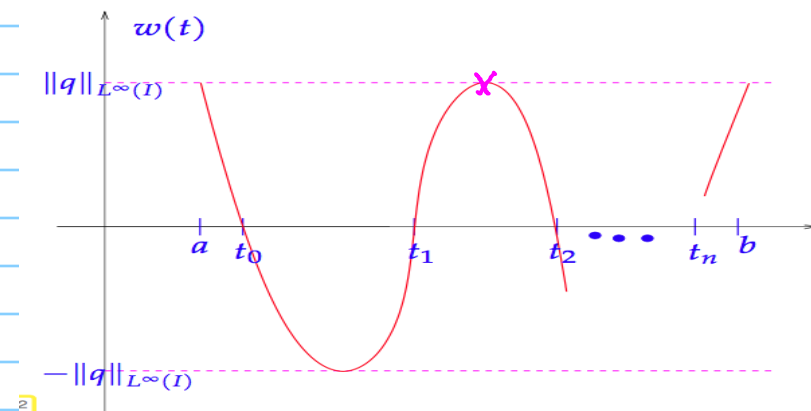
This is equivalent to finding a polynomial $q \in \mathcal{P}_{n+1}$

- * with leading coefficient = 1,
- * such that it minimizes the norm $\|q\|_{L^\infty(I)}$.

Then choose nodes t_0, \dots, t_n as zeros of q (caution: t_j must belong to I).

Heuristics :

- All zeros of w in $[a, b]$ \rightarrow meaningful interp. nodes
- Same extremal value (modulus) in all extrema



The solution

Definition 4.1.66. Chebyshev polynomials \rightarrow [19, Ch. 32]

The n^{th} Chebyshev polynomial is $T_n(t) := \cos(n \arccos t)$, $-1 \leq t \leq 1, n \in \mathbb{N}$.

$$\Rightarrow |T_n(t)| \leq 1$$

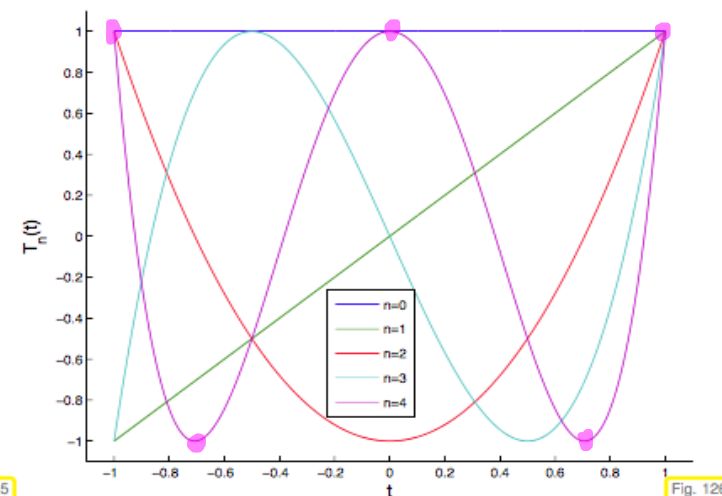
Theorem 4.1.67. 3-term recursion for Chebyshev polynomials \rightarrow [19, (32.2)]

The function T_n defined in Def. 4.1.66 satisfy the 3-term recursion

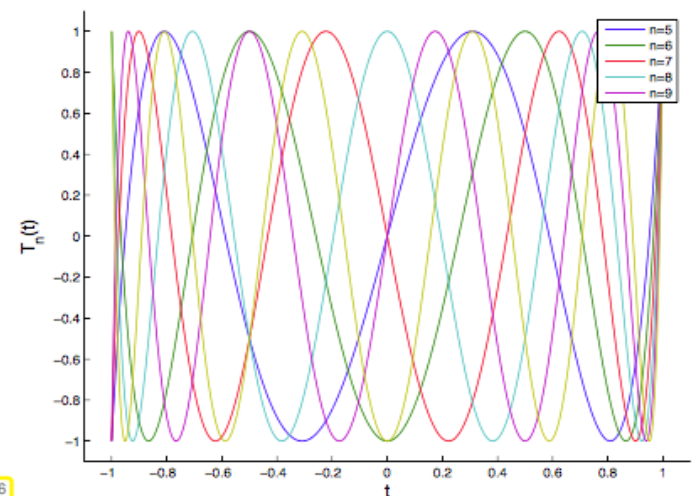
$$T_{n+1}(t) = 2tT_n(t) - T_{n-1}(t), \quad T_0 \equiv 1, \quad T_1(t) = t, \quad n \in \mathbb{N}. \quad (4.1.68)$$

$$\Rightarrow T_n \in \mathcal{P}_n \text{ w/ leading coefficient } 2^{n-1}$$

Proof: $\cos((n+1)t) = 2\cos nt \cdot \cos t - \cos(n-1)t$ \square



Chebyshev polynomials T_0, \dots, T_4



Chebyshev polynomials T_5, \dots, T_9

Zeros : $n \arccos t \in (2j+1)\frac{\pi}{2} \Rightarrow t_j = \cos\left(\frac{2j+1}{n} \frac{\pi}{2}\right)$
 $t_j \hat{=}$ Chebyshev points $\hat{=}$ optimal nodes $j = 0, \dots, n-1$

Theorem 4.1.72. Minimax property of the Chebychev polynomials [9, Section 7.1.4.], [19, Thm. 32.2]

The polynomials T_n from Def. 4.1.66 minimize the supremum norm in the following sense:

$$\|T_n\|_{L^\infty([-1,1])} = \inf\{\|p\|_{L^\infty([-1,1])} : p \in \mathcal{P}_n, p(t) = 2^{n-1}t^n + \dots\}, \quad \forall n \in \mathbb{N}.$$

Proof: (indirect proof), Assume $\exists q(t) = 2^{n-1}t^n + \dots \in \mathcal{P}_n$

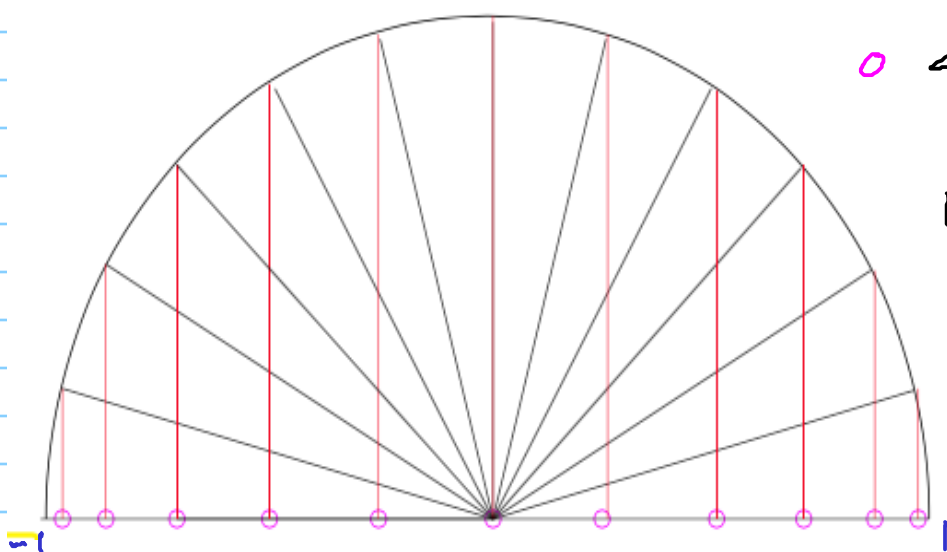
$$\|q\|_\infty < \|T_n\|_\infty \Rightarrow (T_n - q)(\bar{t}) > 0 \text{ in maxima } \bar{t} \text{ of } T_n$$

$$(T_n - q)(\bar{t}) < 0 \text{ in minima } \bar{t} \text{ of } T_n$$

↑
n+1 sign changes

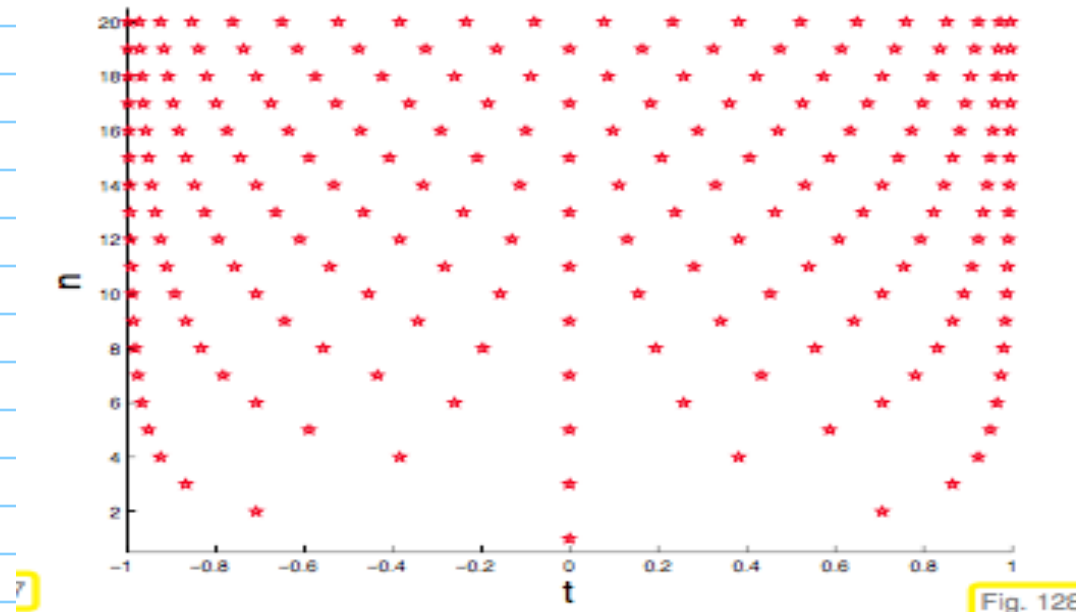
▷ $T_n - q \in \mathcal{P}_{n-1}$ has n zeros $\Rightarrow T_n - q = 0$ \square

▷ Optimal choice $W = 2^{-n} T_{n+1}$



◦ ← Chebyshev nodes on $[-1, 1]$

▷ Chebyshev interpolation



C.N. cluster at endpoints of $[-1, 1]$

The Chebyshev nodes in the interval $I = [a, b]$ are

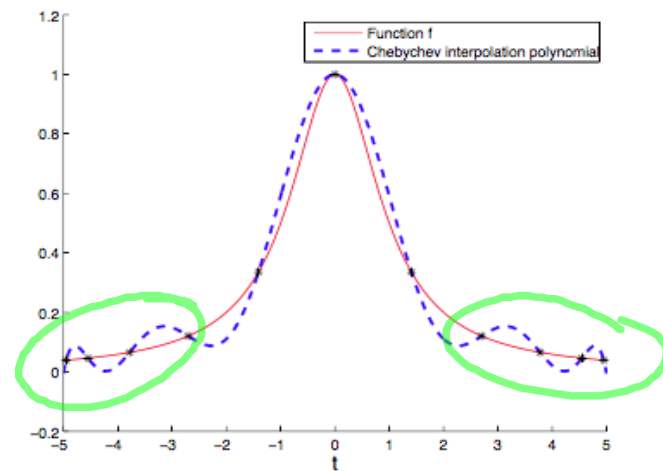
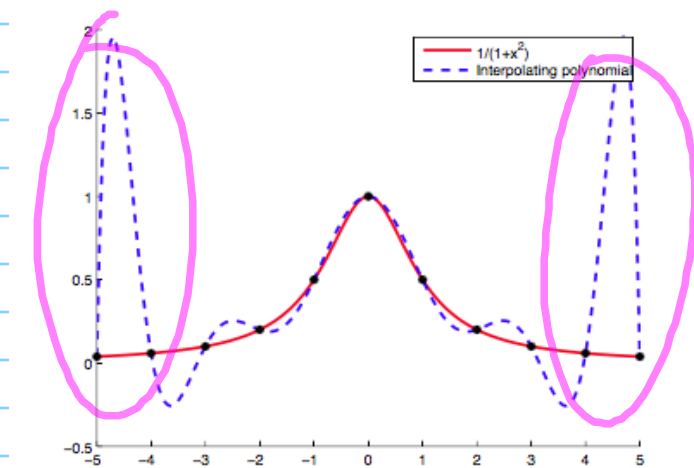
$$t_k := a + \frac{1}{2}(b-a) \left(\cos\left(\frac{2k+1}{2(n+1)}\pi\right) + 1 \right), \quad k = 0, \dots, n. \quad (4.1.77)$$

$$\|W\|_{L^\infty([a,b])} = 2^{-n}$$

▷ For Chebyshev interpolation on $[a, b]$
 $f \in C^{n+1}([a, b])$:

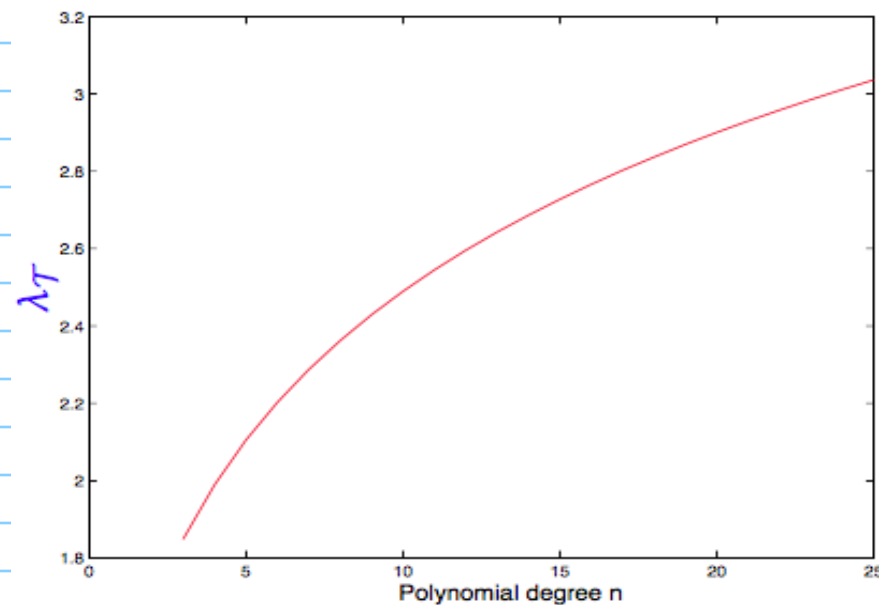
$$\|f - I_T(f)\|_{L^\infty([a,b])} \leq \frac{2^{-n}}{(n+1)!} \|f^{(n+1)}\|_{L^\infty([a,b])} \left(\frac{b-a}{2}\right)^{n+1}$$

(1) $f(t) = \frac{1}{1+t^2}$, $I = [-5, 5]$



Sensitivity of Chebyshev interpolation: Lebesgue constant λ_T

exponential blowup for
equidistant nodes



Theoretical estimate

$$\lambda_{\text{cheb}} \leq \frac{2}{\pi} \log(1+n) + 1$$

λ_T connects interpolation & best approximation

$$\|f - L_T f\|_{\infty} = \|f - I_T[f(t_j)]\|_{\infty} =$$

[Note $L_T p = p \quad \forall p \in \mathcal{P}_n$]

$$p \in \mathcal{P}_n$$

↑
arbitrary!

$$= \|(f-p) - I_T[(f-p)(t_j)]\|_{\infty}$$

$$\leq \|f-p\|_{\infty} + \|I_T[(f-p)(t_j)]\|_{\infty}$$

$$\leq \|f-p\|_{\infty} + \lambda_T \|(f-p)(t_j)\|_{\infty}$$

$$\leq (1 + \lambda_T) \|f-p\|_{\infty}$$

$$\Rightarrow \underbrace{\|f - L_T f\|}_{\text{interpolation}} \leq (1 + \lambda_T) \underbrace{\inf_{p \in \mathcal{P}_n} \|f-p\|_{\infty}}_{\| \cdot \|_{\infty} \text{ of best approx. error}}$$

$$\|f - L_T f\|_{L^{\infty}(I)} \leq \frac{(2/\pi \log(1+n) + 2)(1 + \pi^2/2)^r (n-r)!}{n!} \|f^{(r)}\|_{L^{\infty}([-1,1])}$$

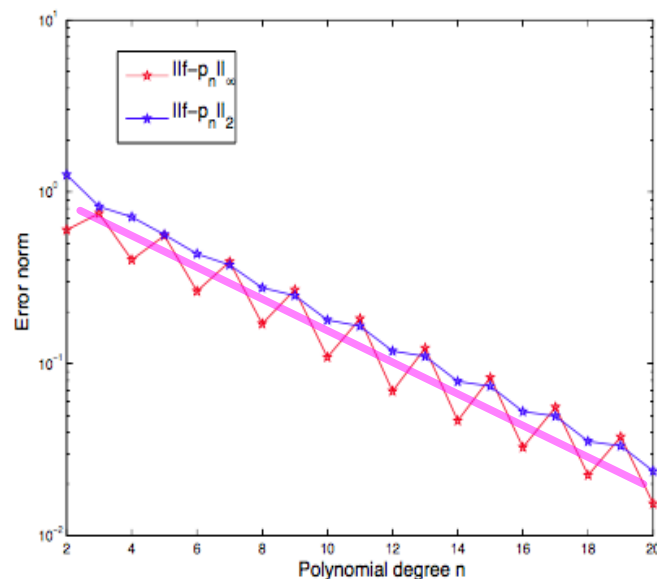
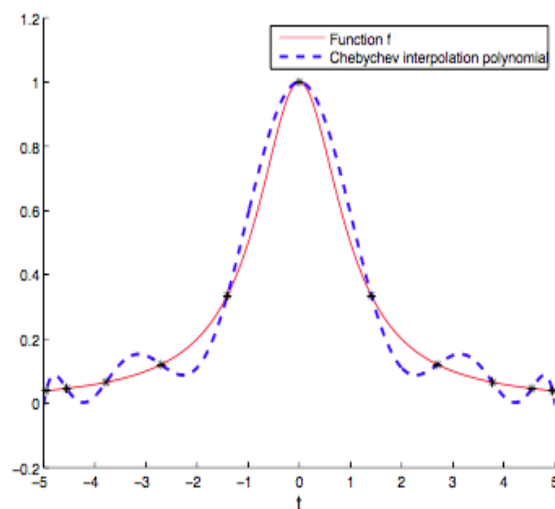
$\rightarrow O(n^{-r})$ algebraic cug.

↕ compare!

$$\|f - I_T(f)\|_{L^{\infty}([-1,1])} \leq \frac{2^{-n}}{(n+1)!} \|f^{(n+1)}\|_{L^{\infty}([-1,1])}$$

① $f(t) = (1+t^2)^{-1}$, $I = [-5, 5]$ (see Ex. 4.1.34): analytic in a neighborhood of I .

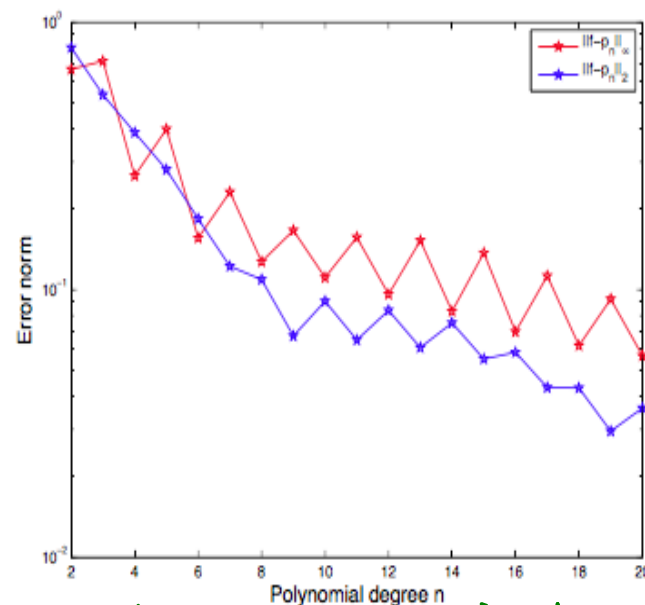
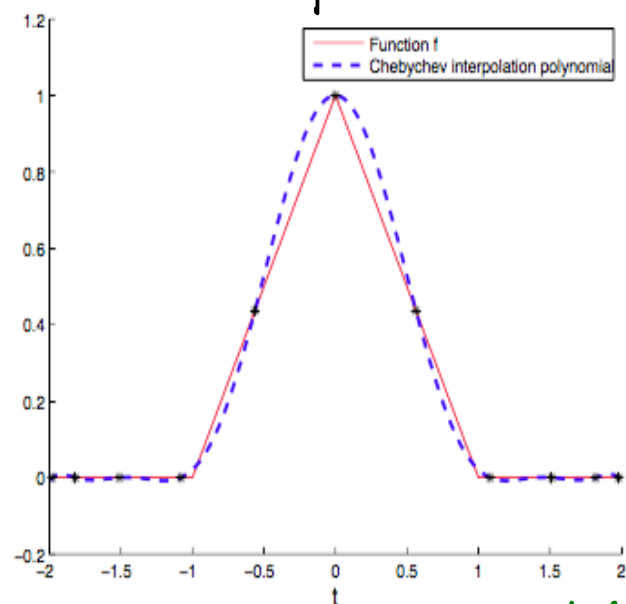
Interpolation with $n = 10$ Chebychev nodes (plot on the left).



↓

Exponential conv $O(q^n)$

$f \equiv$ test function $\notin C^1$



→ Estimates for analytic interpolands cannot be applied

Chebyshev interpolation: analytic interpolands

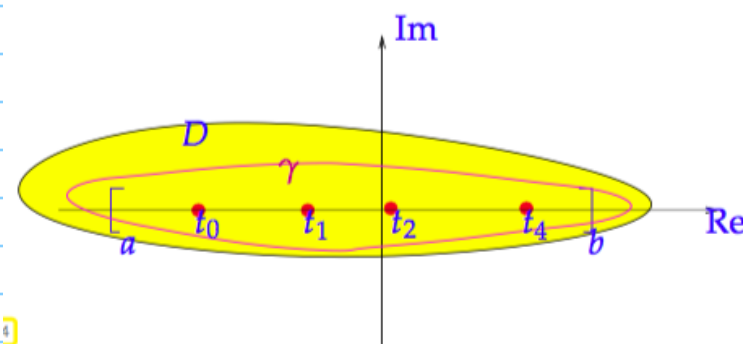
Focus: n -asymptotics

$[I = [-1, 1]]$

$$|f(t) - L_n f(t)| \leq \left| \frac{w(t)}{2\pi i} \int_{\gamma} \frac{f(z)}{(z-t)w(z)} dz \right| \leq \frac{|\gamma|}{2\pi} \frac{\max_{a \leq \tau \leq b} |w(\tau)|}{\min_{z \in \gamma} |w(z)|} \cdot \frac{\max_{z \in \gamma} |f(z)|}{\text{dist}([a, b], \gamma)}$$

$$w = 2^{-n} T_{n+1} \Rightarrow \|w\|_{L^\infty[-1, 1]} \leq 2^{-n}$$

Choice of integration contour $\gamma \subset \mathbb{C}$ around $[-1, 1]$



choice of γ should cancel arcs

Definition 4.1.66. Chebyshev polynomials → [19, Ch. 32]

The n^{th} Chebyshev polynomial is $T_n(t) := \cos(n \arccos t)$, $-1 \leq t \leq 1, n \in \mathbb{N}$.

$$\gamma := \{z = \cos(\theta - i \log \rho), 0 \leq \theta \leq 2\pi\}$$

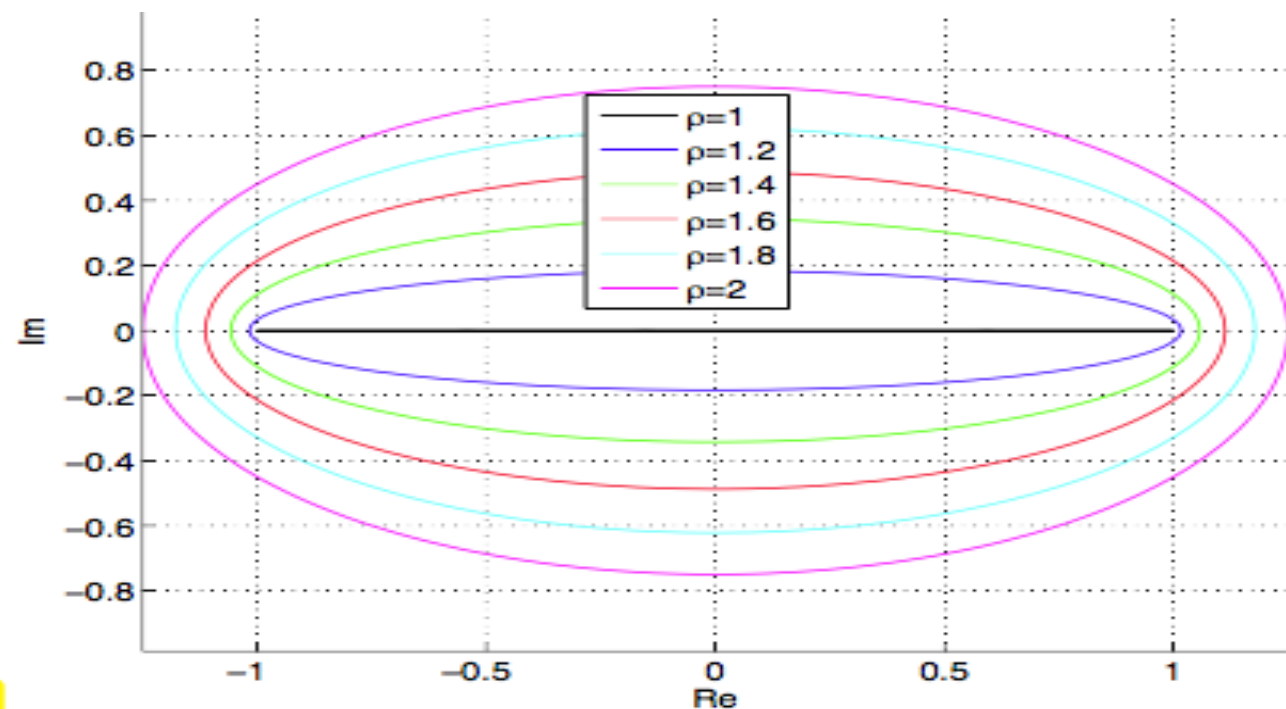
$$= \left\{ z = \frac{1}{2} (\exp(i(\theta - i \log \rho)) + \exp(-i(\theta - i \log \rho))) \right\}$$

$$= \left\{ z = \frac{1}{2} (\rho e^{i\theta} + \rho^{-1} e^{-i\theta}) \right\}$$

$$= \left\{ z = \underbrace{\frac{1}{2}(\rho + \rho^{-1}) \cos \theta}_{\text{Re}} + i \underbrace{\frac{1}{2}(\rho - \rho^{-1}) \sin \theta}_{\text{Im}} \right\}$$

$\theta \triangleq$ curve parameter

$$y(\theta) = A \cos \theta + B \sin \theta, \theta \in [0, 2\pi]: \text{ellipse}$$



\Rightarrow Larger $\rho \Rightarrow$ larger ellipses

Lower bound of $|w| \Leftrightarrow |T_{n+1}|$ on γ

$$|T_s(\cos(\theta - i \log \rho))|^2 = |\cos(s(\theta - i \log \rho))|^2$$

$$s := n+1$$

$$= \cos(s(\theta - i \log \rho)) \cdot \overline{\cos(s(\theta - i \log \rho))}$$

$$= \frac{1}{4} (\rho^s e^{is\theta} + \rho^{-s} e^{-is\theta}) (\rho^s e^{-is\theta} + \rho^{-s} e^{is\theta})$$

$$= \frac{1}{4} (\rho^{2s} + \rho^{-2s} + e^{2is\theta} + e^{-2is\theta})$$

$$= \frac{1}{4} (\rho^s - \underbrace{\rho^{-s}}_{<1})^2 + \underbrace{\frac{1}{4} (e^{is\theta} + e^{-is\theta})^2}_{=\cos^2(s\theta) \geq 0} \geq \frac{1}{4} (\rho^s - 1)^2,$$

\downarrow Plug into

$$\|f - L_T f\|_{L^\infty([-1,1])} \leq \frac{4|\gamma|}{\pi} \underbrace{\frac{1}{(\rho^{n+1} - 1)(\rho + \rho^{-1} - 2)}}_{O(\rho^{-n})} \cdot \max_{z \in \gamma} |f(z)|$$

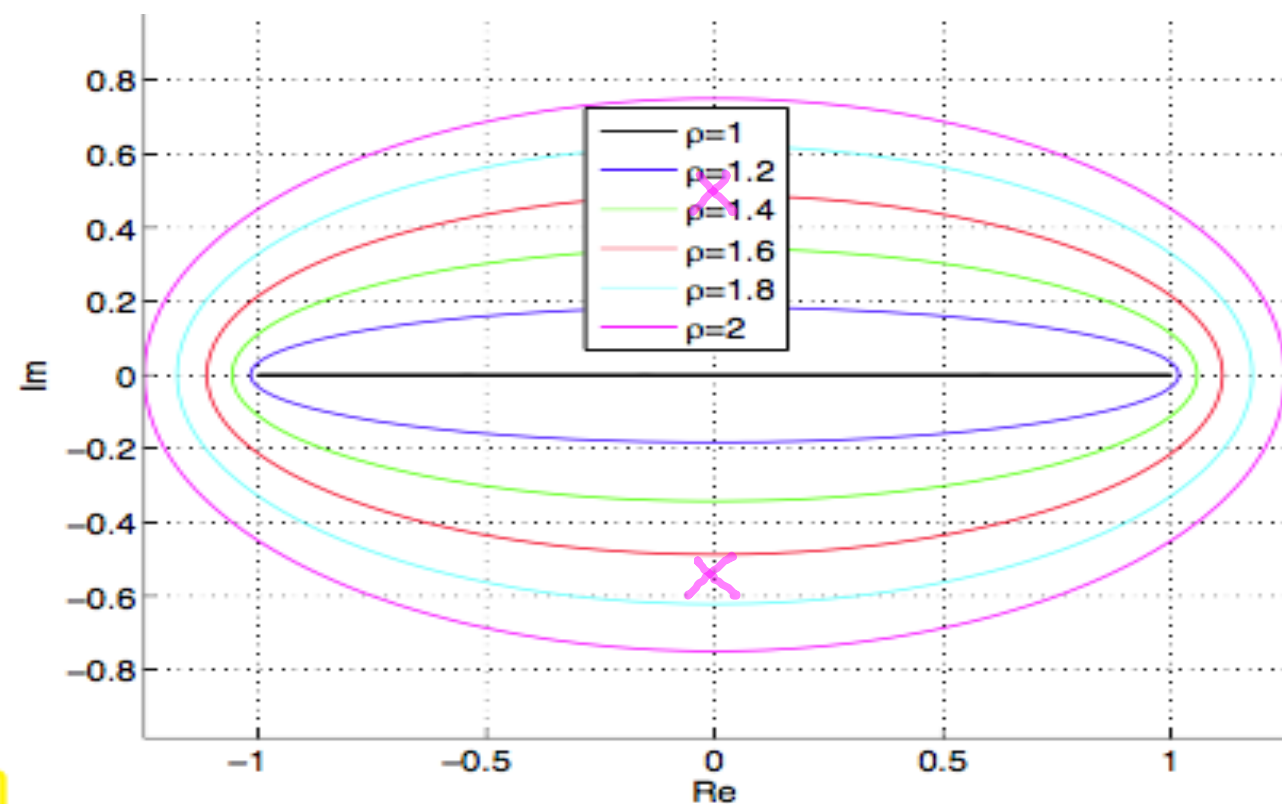
$\rho > 1$:

$O(\rho^{-n})$

The larger ρ , the faster conv.

Admissible ρ : $\gamma_\rho \subset D$ [domain of analyticity]

14 Ex: $f(t) = \frac{1}{1+(2t)^2} \Rightarrow D = \mathbb{C} \setminus \{\pm \frac{1}{2}\}$



Interior of f_ρ must be contained in $D \Rightarrow \rho \lesssim 1.65$

4.2. Mean square best approximation

Goal: compute best approximants

4.2.1. Abstract theory

$X \cong \mathbb{R}$ -vector space \leftrightarrow function space

$V \subset X$ finite-dim. subspace \leftrightarrow space of polynomials
with basis $\{b_1, \dots, b_N\}$

$(\cdot, \cdot)_X \triangleq$ inner product on $X \Rightarrow$ mean square norm $\|u\|_X = (u, u)_X^{1/2}$

Sought: $q := \inf_{p \in V} \|f - p\|_X^2$ for $f \in X$

$$\|f - p\|_X^2 = \|f\|^2 - 2(f, p)_X + \|p\|_X^2$$

Basis rep.: $p = \sum_{j=1}^N f_j b_j$, $f_j \in \mathbb{R}$

$$\|f - p\|_X^2 = \|f\|^2 - 2 \sum_j f_j (f, b_j)_X + \sum_j \sum_k f_j f_k (b_j, b_k)_X$$

$$= \|f\|^2 - 2 \underline{d}^T \underline{c} + \underline{c}^T \underline{M} \underline{c} =: \phi(\underline{c})$$

$$\underline{d} := [(f, b_j)_X]_{j=1}^N \in \mathbb{R}^N \quad \phi: \mathbb{R}^N \rightarrow \mathbb{R}$$

$$\underline{M} := [(b_j, b_k)_X]_{j,k=1}^N \in \mathbb{R}^{N,N} \quad \phi \in C^\infty$$

$$\underline{c} := [f_j]_{j=1}^N$$

16

\Rightarrow Find $\underline{c} \in \mathbb{R}^N$: $\Phi(\underline{c}) \rightarrow \min$

$$\text{grad } \Phi(\underline{c}) = -2\underline{d} + 2M\underline{c} \stackrel{!}{=} \underline{0} \quad (*)$$

$$\text{Hess } \Phi(\underline{c}) = M$$

$$X := \sum \xi_i b_i \in X : \underline{x}^T M \underline{x} = \|\sum \xi_i b_i\|_X^2 > 0 \Leftrightarrow \underline{x} \neq 0$$

$\Rightarrow M$ s.p.d. !

$$(*) \text{ LSE : } \underline{c} = M^{-1} \underline{d} : \text{A formula!}$$

\hookrightarrow normal equations

Theorem 4.2.7. Mean square best approximation through normal equations

Given any $f \in X$ there is a unique $q \in V$ such that

$$\|f - q\|_X = \inf_{p \in V} \|f - p\|_X.$$

Its coefficients γ_j , $j = 1, \dots, N$, with respect to the basis $\mathfrak{B}_V := \{b_1, \dots, b_N\}$ of V ($q = \sum_{j=1}^N \gamma_j b_j$) are the unique solution of the normal equations

$$M[\gamma_j]_{j=1}^N = \left[(f, b_j)_X \right]_{j=1}^N, \quad M := \begin{bmatrix} (b_1, b_1)_X & \dots & (b_1, b_N)_X \\ \vdots & & \vdots \\ (b_N, b_1)_X & \dots & (b_N, b_N)_X \end{bmatrix} \in \mathbb{K}^{N,N}. \quad (4.2.8)$$

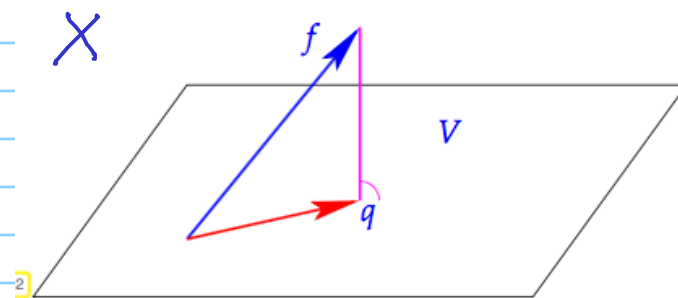
"Geometric" view:

Corollary 4.2.12. Best approximant by orthogonal projection

If q is the best approximant of f in V , then $f - q$ is **orthogonal** to every $p \in V$:

$$(f - q, p)_X = 0 \quad \forall p \in V \Leftrightarrow f - q \perp q.$$

\uparrow
best approximation error



$$\begin{aligned} \xi &\rightarrow \|f - q + \xi p\|_X^2 = \\ &= \|f - q\|_X^2 - 2\xi (f - q, p) + \xi^2 \|p\|_X^2 \end{aligned}$$

a parabola with apex in
 $\xi = 0 \Rightarrow (f - q, p) = 0$

$$\forall p \in V$$

Aiming for simple N.E. (diagonal M)

Definition 4.2.13. Orthonormal basis

A subset $\{b_1, \dots, b_N\}$ of an N -dimensional vector space V with inner product (\rightarrow Def. 4.2.1) $(\cdot, \cdot)_V$ is an **orthonormal basis (ONB)**, if $(b_k, b_j)_V = \delta_{kj}$.

A basis of $\{b_1, \dots, b_N\}$ of V is called **orthogonal**, if $(b_k, b_j)_V = 0$ for $k \neq j$.

$\{b_1, \dots, b_N\}$ ONB \Rightarrow

$$q = \sum_{j=1}^N (f, b_j)_X b_j \quad (*)$$

We already know how to compute ONBs!

Gram-Schmidt orthonormalization

```

1:  $b_1 := \frac{p_1}{\|p_1\|_V}$  % 1st output vector
2: for  $j = 2, \dots, k$  do
  { % Orthogonal projection
3:    $b_j := p_j$ 
4:   for  $\ell = 1, 2, \dots, j-1$  do (4.2.17)
5:     {  $b_j \leftarrow b_j - (p_j, b_\ell)_V b_\ell$  }
6:   if  $(b_j = 0)$  then STOP
7:   else {  $b_j \leftarrow \frac{b_j}{\|b_j\|_V}$  }
  }

```

Theorem 4.2.18. Gram-Schmidt orthonormalization

When supplied with $k \in \mathbb{N}$ linearly independent vectors $p_1, \dots, p_k \in V$ in a vector space with inner product $(\cdot, \cdot)_V$, Algorithm (4.2.17) computes vectors b_1, \dots, b_k with

$$(b_\ell, b_j)_V = \delta_{\ell j}, \quad \ell, j \in \{1, \dots, k\},$$

$$\text{Span}\{b_1, \dots, b_\ell\} = \text{Span}\{p_1, \dots, p_\ell\} \quad //$$

for all $\ell \in \{1, \dots, k\}$.

→ orthogonal projections

Alg. for computing q :
 (i) G.S. → $\{b_1, \dots, b_n\}$
 (ii) (*)

4.2.2. Polynomial mean square best approximation

$$[X = C^0([a, b]), \quad V = \mathcal{P}_m]$$

inner product: $(f, g)_X = \int_a^b f(t)g(t)dt$: L^2 -I.P.

Definition 4.2.24. Orthonormal polynomials → Def. 4.2.13

Let $(\cdot, \cdot)_X$ be an inner product on \mathcal{P}_m . A sequence r_0, r_1, \dots, r_m provides orthonormal polynomials with respect to $(\cdot, \cdot)_X$, if

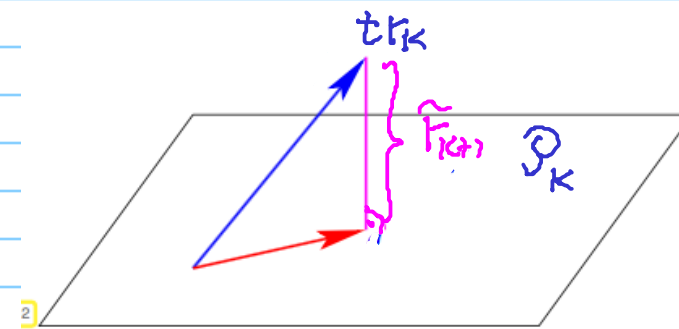
$$r_\ell \in \mathcal{P}_\ell, \quad (r_k, r_\ell)_X = \delta_{k\ell}, \quad \ell, k \in \{0, \dots, m\}. \quad (4.2.25)$$

The polynomials are just orthogonal, if $(r_k, r_\ell)_X = 0$ for $k \neq \ell$.

→ increasing degree

$\{r_0, \dots, r_m\}$ ONPs: $tr_k(t) \in \mathcal{P}_{k+1} \setminus \mathcal{P}_k$
 linearly indep ← leading coeff. $\neq 0$

$$\begin{cases} r_{k+1} \in \text{Span}\{tr_k, r_k, r_{k-1}, \dots, r_0\} \\ r_{k+1} \perp \text{Span}\{r_0, \dots, r_k\} \end{cases}$$



orthogonal projection

$$r_{k+1} = \frac{\hat{r}_{k+1}}{\|\hat{r}_{k+1}\|} \quad ; \quad \hat{r}_{k+1} = tr_k - \sum_{j=1}^k (tr_k, r_j)_X r_j \quad (*)$$

Assumption: $(tf, g)_X = (f, tg)_X \quad \forall f, g$

(17)

$$(tr_k, r_j)_x = (r_k, \underbrace{tr_j}_{\in \mathcal{P}_{j+1}})_x = 0 \quad \forall j \leq k-2 \text{ by orthogonality!}$$

→ Sum (+) collapses to a **3-term recursion**,
 \rightarrow not normalized

Theorem 4.2.31. 3-term recursion for orthogonal polynomials

Given an inner product $(\cdot, \cdot)_X$ on \mathcal{P}_m , $m \in \mathbb{N}$, define $p_{-1} := 0$, $p_0 = 1$, and

$$p_{k+1}(t) := (t - \alpha_{k+1})p_k(t) - \beta_k p_{k-1}(t), \quad k = 0, 1, \dots, m-1,$$

with $\alpha_{k+1} := \frac{(\{t \mapsto tp_k(t)\}, p_k)_X}{\|p_k\|_X^2}$, $\beta_k := \frac{\|p_k\|_X^2}{\|p_{k-1}\|_X^2}$. (4.2.32)

Then $p_k \in \mathcal{P}_k$ with leading coefficient = 1, and $\{p_0, p_1, \dots, p_m\}$ is an **orthogonal basis** of \mathcal{P}_m .

```

1 function [alpha, beta] = coeffortho(t,n)
2 % Vector t passes the points in the definition of the discrete L^2-inner
3 % product, n the maximal index desired
4 m = numel(t); % Maximal degree of orthogonal polynomial
5 alpha(1) = sum(t)/m;
6 % Initialization of recursion; we store only the values of
7 % the polynomials at the points in T.
8 p1 = ones(size(t));
9 p2 = t-alpha(1);
10 % Main loop
11 for k=1:min(n-1,m-2)
12     p0 = p1; p1 = p2;
13     % 3-term recursion (4.2.32),
14     alpha(k+1) = dot(p1, (t.*p1))/norm(p1)^2;
15     beta(k) = (norm(p1)/norm(p0))^2;
16     p2 = (t-alpha(k+1)).*p1-beta(k)*p0;
17 end

```

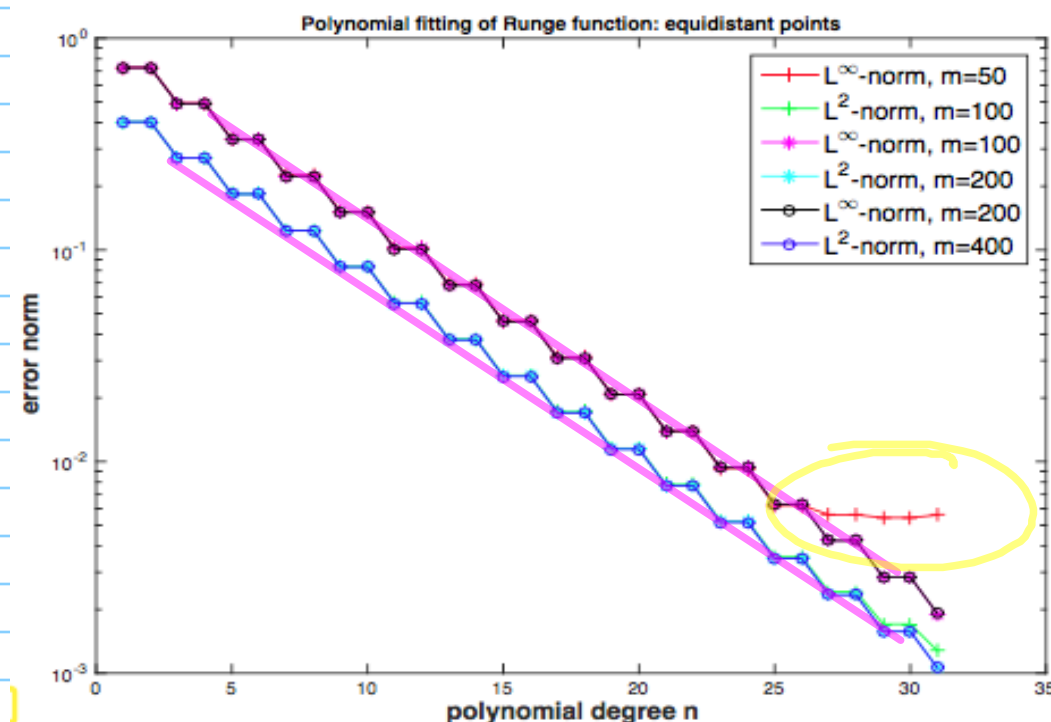
Practical (semi) inner products: **discrete L^2 -inner products**

Given $\mathcal{T} = \{t_j\}_{j=0}^m$: $(f, g)_\mathcal{T} := \sum_{j=0}^m f(t_j)g(t_j)$

→ At least definite on \mathcal{P}_k , $k \leq m$.

Experiment: equidistant t_j in $[-1, 1]$

$$f(t) = \frac{1}{1+(5t)^2}$$



→ **exponential convergence**

$$\|f - p\|_{\mathcal{T}}^2 = \sum_{j=0}^m ((f - p)(t_j))^2$$

(18)

4.3. Uniform Best-Approximation

$f \in C^0([a, b])$, find $q = \operatorname{argmin}_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([a, b])}$

Theorem 4.3.2. Chebychev alternation theorem

Given $f \in C[a, b]$, $a < b$, and a polynomial degree $n \in \mathbb{N}$, a polynomial $q \in \mathcal{P}_n$ satisfies

$$q = \operatorname{argmin}_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty(I)}$$

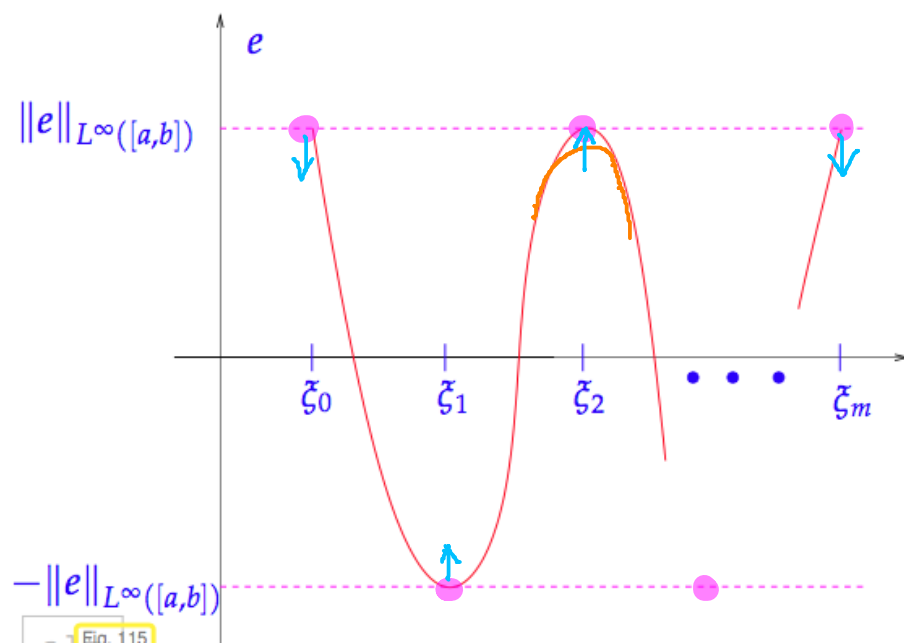
if and only if there exist $n + 2$ points $a \leq \xi_0 < \xi_1 < \dots < \xi_{n+1} \leq b$ such that

$$|e(\xi_j)| = \|e\|_{L^\infty([a, b])}, \quad j = 0, \dots, n+1,$$

$$e(\xi_j) = -e(\xi_{j+1}), \quad j = 0, \dots, n,$$

where $e := f - q$ denotes the approximation error.

$\xi_j \hat{=}$ Chebychev alternants



Remez algorithm:

$$= \|e\|_\infty$$

$$CAT \Rightarrow (+) \quad q(\xi_k) + (-1)^k \delta = f(\xi_k), \quad k=0, \dots, n+1$$

After choosing a basis for $\mathcal{P}_n \rightarrow (n+2) \times (n+2)$ LSE,
unknowns = expansion coeffs of q & δ

• Start from approximate alternants $\xi_j^{(0)}$, $j=0, \dots, n+1$

- Solve (+) $\rightarrow q$
- Check (sampling) $\|f - q\|_\infty \rightarrow \text{STOP, if small enough}$
- New set of alternants $\hat{=}$ extrema of $f - q$

18 4.5. Approximation by piecewise^{*} polynomials

Recall: cubic splines, cubic Hermite, p.w. linear

Advantage: Cost $O(N)$, $N \triangleq$ no. of data points

* w.r.t. to a **mesh** of $[a, b]$: $\mathcal{M} = \{a = x_0 < x_1 < \dots < x_m = b\}$
 mesh cell / interval: $[x_{j-1}, x_j]$ ↑
node

4.5.1. Piecewise Lagrange Interpolation

↳ Separate Lagrange intp. on cells

General local Lagrange interpolation on a mesh

- 1 Choose **local degree** $n_j \in \mathbb{N}_0$ for each cell of the mesh, $j = 1, \dots, m$.
- 2 Choose set of **local** interpolation points (nodes)

$$\mathcal{T}^j := \{t_0^j, \dots, t_{n_j}^j\} \subset I_j := [x_{j-1}, x_j], \quad j = 1, \dots, m,$$

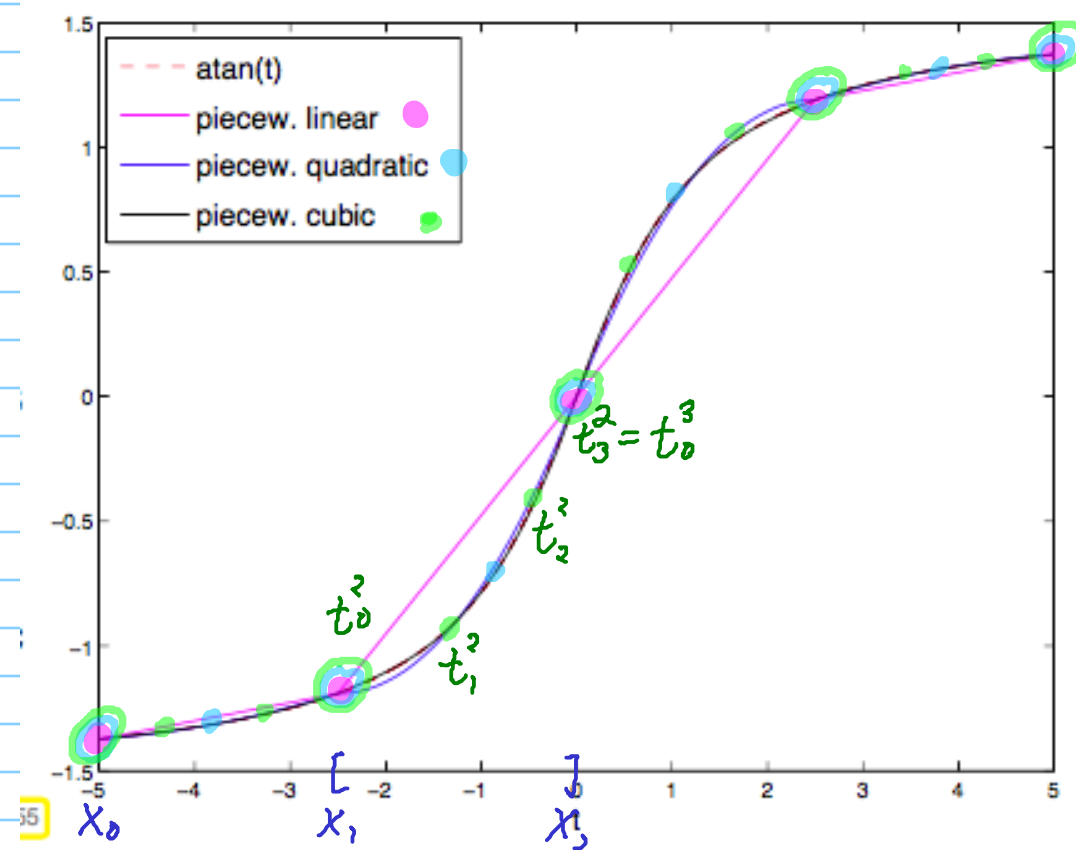
for each mesh cell/grid interval I_j .

- 3 Define **piecewise polynomial** interpolant $s : [x_0, x_m] \rightarrow \mathbb{K}$:

$$s_j := s|_{I_j} \in \mathcal{P}_{n_j} \quad \text{and} \quad s_j(t_i^j) = f(t_i^j) \quad i = 0, \dots, n_j, \quad j = 1, \dots, m. \quad (4.5.5)$$

Owing to Thm. 3.2.14, s_j is well defined.

↳ $I_m \triangleq$ interpolation operator



$$s \in C^0([a, b]) \text{ , if } t_{n_j}^j = t_0^{j+1} \quad \forall j$$

[Assume $f \in C^0([a, b])$]

In general $s \notin C^1$

Focus: Asymptotic conv in maximum norm:

h-convergence: $n_j = n \quad \forall j$ [fixed local degree]
 Consider sequence of meshes \mathcal{M}_k with
 meshwidth $h_k := \max_j |x_j^{(k)} - x_{j-1}^{(k)}| \xrightarrow{k \rightarrow \infty} 0$

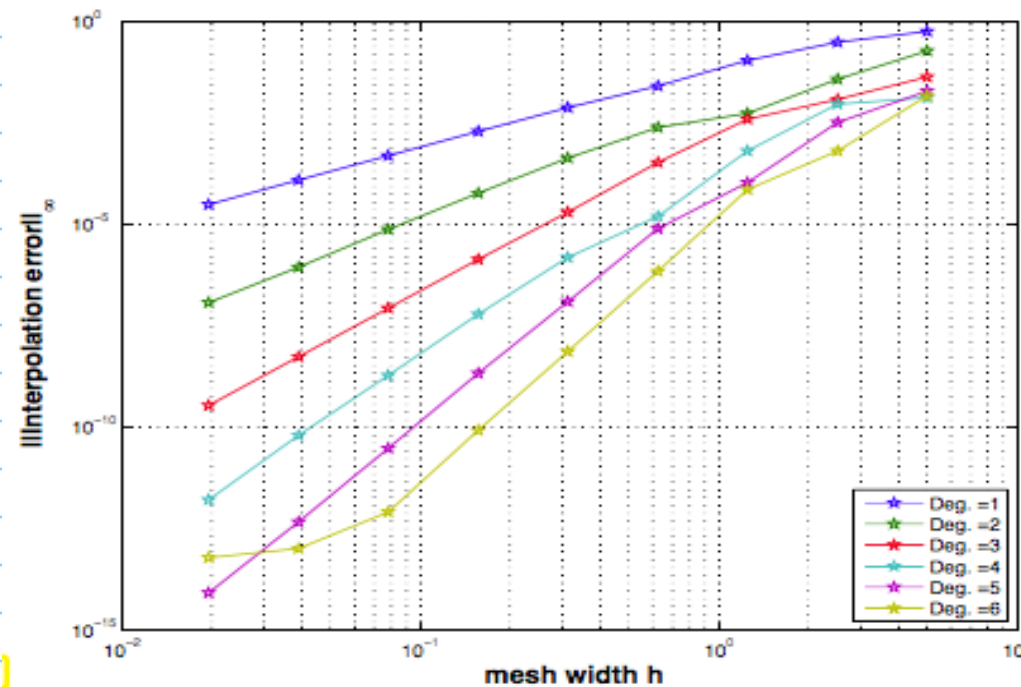
(19)

Find bound $\|f - I_{M_k} f\|_{L^\infty([a,b])} \leq CT(h_k)$
 study asymptotics of T for $h_k \rightarrow 0$

Example: $f(t) = \arctan(t)$ on $[-5, 5]$

$M_k = \{-5 + j_{h_k} 10; j = 0, \dots, 2^k\}$, $h_k = 2^{-k}$
 ["equidistant mesh"]

equidistant local interpolation points
 (endpoints included)



Δ Alg. conv.
 $\|error\|_\infty = O(h^\alpha)$

$\alpha \hat{=}$ rate

Empiric

$\alpha \approx n + 1$

$n \hat{=}$ local pol. degree

Theory:

Known estimate
 $I = [x_{j-1}, x_j]$

$$\|f - L_T f\|_{L^\infty(I)} \leq \frac{\|f^{(n+1)}\|_{L^\infty(I)}}{(n+1)!} \max_{t \in I} |(t - t_0) \cdots (t - t_n)| \quad (4.1.43)$$

$$\max_j |x_j - x_{j-1}|^{n+1} \leq h_m^{n+1} \quad [h_m \hat{=} \text{meshwidth } h]$$

[Uniform local degree n]

$$(4.1.43) \Rightarrow \|f - s\|_{L^\infty([x_0, x_m])} \leq \frac{h^{n+1}}{(n+1)!} \|f^{(n+1)}\|_{L^\infty([x_0, x_m])}$$

$$h := h_M = \max\{|x_j - x_{j-1}| : j = 1, \dots, m\}.$$

smoothness requirement

Remark: p-convergence: M fixed, raise local degree

4.5.2. Cubic Hermite interpolation

Definition 4.5.14. Piecewise cubic Hermite interpolant (with exact slopes) \rightarrow Def. 3.4.1

[Piecewise cubic Hermite interpolant (with exact slopes) Given $f \in C^1([a, b])$ and a mesh $M := \{a = x_0 < x_1 < \dots < x_{m-1} < x_m = b\}$ the piecewise cubic Hermite interpolant (with exact slopes) $s : [a, b] \rightarrow \mathbb{R}$ is defined as

$$s|_{[x_{j-1}, x_j]} \in \mathcal{P}_3, \quad j = 1, \dots, m, \quad s(x_j) = f(x_j), \quad s'(x_j) = f'(x_j), \quad j = 0, \dots, m.$$

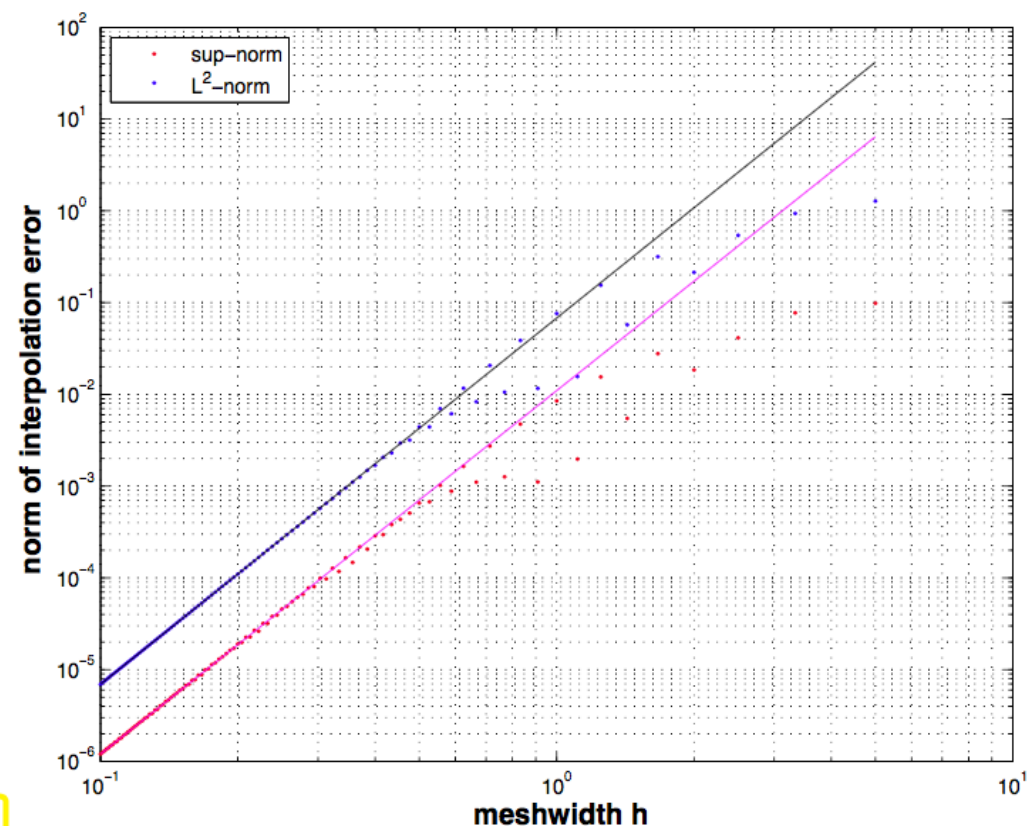
(20)

$$\Rightarrow s \in C^1 !$$

h -convergence :

Experiment : $f(t) = \frac{1}{1+t^2}$ on $[-5, 5]$

• equidistant mesh



Alg. conv.
rate 4

$$\| \text{error} \|_0 = O(h^4)$$

Theorem 4.5.17. Convergence of approximation by cubic Hermite interpolation

Let s be the cubic Hermite interpolant of $f \in C^4([a, b])$ on a mesh $\mathcal{M} := \{a = x_0 < x_1 < \dots < x_{m-1} < x_m = b\}$ according to Def. 4.5.14. Then

$$\|f - s\|_{L^\infty([a, b])} \leq \frac{1}{4!} h_{\mathcal{M}}^4 \|f^{(4)}\|_{L^\infty([a, b])},$$

with the meshwidth $h_{\mathcal{M}} := \max_j |x_j - x_{j-1}|$.

4.5.3. Cubic spline interpolation

NCSE15

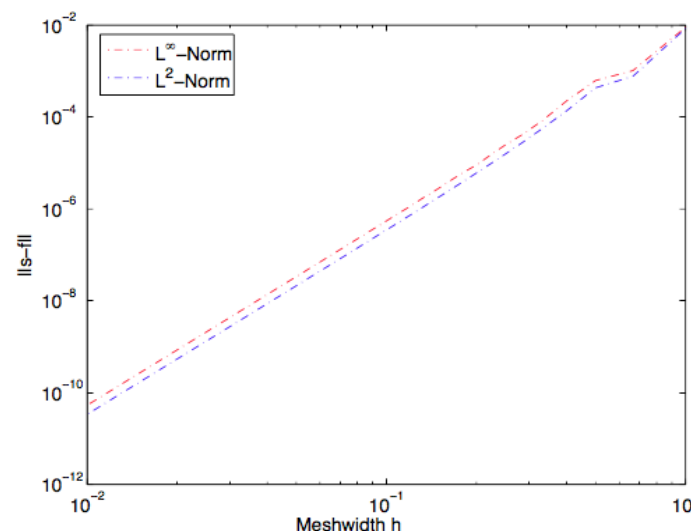
Exp.: h -convergence for complete cubic spline interpolant s on $[a, b]$:

$$\begin{aligned} s(a) &= f(a) \\ s(b) &= f(b) \end{aligned}$$

$$\begin{aligned} s'(a) &= f'(a) \\ s'(b) &= f'(b) \end{aligned}$$

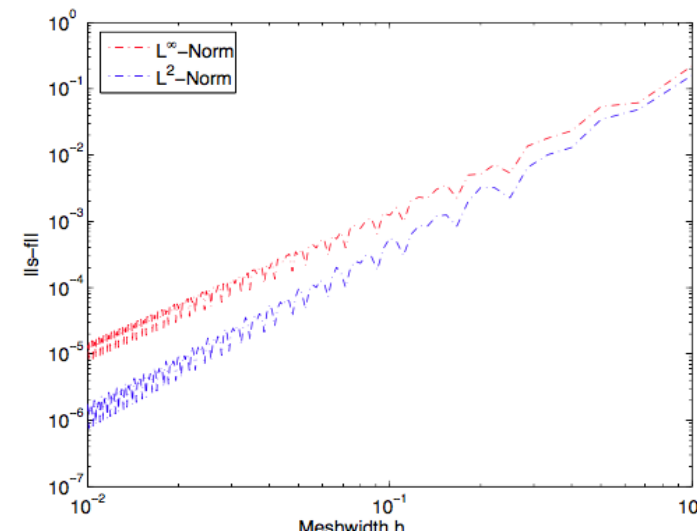
$$f_1(t) = \frac{1}{1+e^{-2t}} \in C^\infty(I), \quad f_2(t) = \begin{cases} 0 & \text{if } t < -\frac{2}{3}, \\ \frac{1}{2}(1 + \cos(\pi(t - \frac{3}{5}))) & \text{if } -\frac{2}{3} < t < \frac{3}{5}, \\ 1 & \text{otherwise.} \end{cases} \in C^1(I).$$

$\nexists C^2$



$$\|f_1 - s\|_{L^\infty([-1, 1])} = O(h^4)$$

Alg. conv., rate 4



$$\|f_2 - s\|_{L^\infty([-1, 1])} = O(h^2)$$

Alg. conv., rate 2

(21)

Rule of thumb :

p.w. approximation w/ polynomials of degree p ,

h -convergence : algebraic, maximal rate $p+1$

↑
smoothness of f required!

