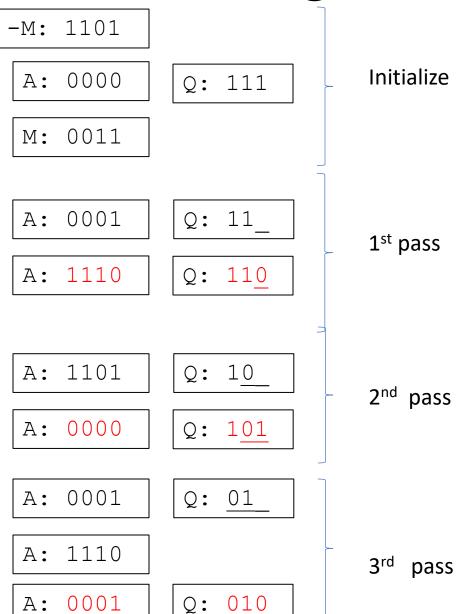
Non-restoring Method

- The non-restoring division defers restoration to reduce operations
 - Addition is performed instead of subtraction in the next step
 - This requires a final restoration, when last operation yielded a negative remainder

Non-Restoring Method

- Initialization
 - Clear A. Requires 1 extra bit for A to be used as a sign bit.
 - Q gets dividend.
 - *M* gets divisor.
- Loop for each bit of the dividend Q
 - If A (previous result) is negative $(A_n = 1)$
 - Shift AQ to the left.
 - Add $(A \leftarrow A + M)$
 - Q_0 gets the complement value of the MSB of A
 - Else
 - Shift AQ to the left.
 - Subtract $(A \leftarrow A M)$
 - Q_0 gets the complement value of the MSB of A
- Perform restoration if last result in A is negative, i.e., if $A_n = 1$ then $A \leftarrow A + M$
- Quotient in Q while remainder in A; adjust sign as needed

Non-Restoring Division



- Initialization
 - Clear A. Requires 1 extra bit for A to be used as a sign bit.
 - Q gets dividend.
 - M gets divisor.
- Loop for each bit of the dividend Q
 - If A (previous result) is negative $(A_n = 1)$
 - Shift AQ to the left.
 - Add $(A \leftarrow A + M)$
 - Q_0 gets the complement value of the MSB of A
 - Else
 - Shift AQ to the left.
 - Subtract $(A \leftarrow A M)$
 - Q_0 gets the complement value of the MSB of A
- Perform restoration if last result in A is negative, i.e., if $A_n = 1$ then $A \leftarrow A + M$
- Quotient in Q while remainder in A; adjust sign as needed



Try: 01101 (Q) / 00101 (M) (using non-restoring division) Show the value of A and Q after the end of each pass

After this pass	A	Q
1 st	111011	11010
2 nd	111100	10100
3 rd	11110	01000
4 th	000001	10001
5th	000011	00010