

Proof $p \rightarrow q$:

1 direct: assume p is true, show q is true

2: contraposition: assume q is false, show p is false

3: if and only if: first prove $p \rightarrow q$, then prove $q \rightarrow p$

4: contradiction: assume p is false, find a contradiction that must be false, so p is true

Proof by induction:

Proof Template VII: Induction to prove $\forall n \geq 1 : P(n)$.

Proof. We use induction to prove $\forall n \geq 1 : P(n)$.

1. Show that $P(1)$ is T. (Usually a simple verification.) [base case]
2. Show $P(n) \rightarrow P(n+1)$ for $n \geq 1$ [induction step]

Prove the implication using direct proof or contraposition.

Direct	Contraposition
Assume $P(n)$ is T. (valid derivations) must show for any $n \geq 1$ must use $P(n)$ here Show $P(n+1)$ is T.	Assume $P(n+1)$ is F. (valid derivations) must show for any $n \geq 1$ must use $\neg P(n+1)$ here Show $P(n)$ is F.

3. Conclude: by induction, $\forall n \geq 1 : P(n)$.

Recursive definition of Rooted Binary Trees (RBT).

- ① The empty tree ϵ is an RBT.

- ② If T_1, T_2 are disjoint RBTs with roots r_1 and r_2 , then linking r_1 and r_2 to a new root r gives a new RBT with root r .



- ③ A single root-node \bullet is an RBT.

- ④ If T_1, T_2 are disjoint RBTs with roots r_1 and r_2 , then linking r_1 and r_2 to a new root r gives a new RBT with root r .

Sum calculation:

$$\sum_{i=k}^n 1 = n+1-k$$

$$\sum_{i=1}^n i = \frac{1}{2}n(n+1)$$

$$\sum_{i=0}^n 2^i = 2^{n+1} - 1$$

$$\sum_{i=1}^n f(x) = nf(x)$$

$$\sum_{i=1}^n i^2 = \frac{1}{6}n(n+1)(2n+1)$$

$$\sum_{i=0}^n \frac{1}{2^i} = 2 - \frac{1}{2^n}$$

$$\sum_{i=0}^n r^i = \frac{1-r^{n+1}}{1-r} \quad (r \neq 1)$$

$$\sum_{i=1}^n i^3 = \frac{1}{4}n^2(n+1)^2$$

$$\sum_{i=1}^n \log i = \log n!$$

$$\sum_{i=0}^n i2^{i+1} = 0 + \sum_{i=1}^n i2^{i+1}$$

$$= (n-1)2^{n+1} + 2$$

$$\frac{T(n)}{f(n)} \xrightarrow{n \rightarrow \infty} \begin{cases} \infty & T \in \omega(f); \\ \text{constant} > 0 & T \in \Theta(f); \\ 0 & T \in o(f). \end{cases}$$

(More formally, $T \in \Theta(f)$ if $cf(n) \leq T(n) \leq Cf(n)$.)

$T \in o(f)$	$T \in O(f)$	$T \in \Theta(f)$	$T \in \Omega(f)$	$T \in \omega(f)$
" $T < f$ "	" $T \leq f$ "	" $T = f$ "	" $T \geq f$ "	" $T > f$ "

Approximation via Integration

Theorem 9.3. Let f be a monotonically increasing function. Then,

$$\int_{m-1}^n dx f(x) \leq \sum_{i=m}^n f(i) \leq \int_m^{n+1} dx f(x).$$

(If f is monotonically decreasing instead, the inequalities are reversed.)

$\gcd(6,15) = 15 - 2 \times 6 = 3$

$= \gcd(3,6) = 6 - 2 \times (15 - 2 \times 6) = 5 \times 6 - 2 \times 15 = 0$

$= \gcd(0,3) = 15 - 2 \times 6 - (5 \times 6 - 2 \times 15) = 3$

$= 3 = 3 \times 15 - 7 \times 6 = 3$

So $\gcd(6,15) = 3 = 3 \times 15 - 7 \times 6$

Modular Equivalence Properties:

$$a \equiv b \pmod{d} \text{ and } r \equiv s \pmod{d}$$

$$\textcircled{1} ar \equiv bs \pmod{d}$$

$$\textcircled{2} a + r \equiv b + s \pmod{d}$$

$$\textcircled{3} a^n \equiv b^n \pmod{d}$$

Set 1. **Associative:** $(A \cap B) \cap C = A \cap (B \cap C);$
 $(A \cup B) \cup C = A \cup (B \cup C).$

2. **Commutative:** $A \cap B = B \cap A;$
 $A \cup B = B \cup A.$

3. **Complements:** $\overline{\overline{A}} = A;$
 $\overline{A \cap B} = \overline{A} \cup \overline{B};$
 $\overline{A \cup B} = \overline{A} \cap \overline{B}.$

4. **Distributive:** $A \cup (B \cap C) = (A \cup B) \cap (A \cup C);$
 $A \cap (B \cup C) = (A \cap B) \cup (A \cap C).$

Ordinary (Weak) Induction	Base Case: $P(1)$ Induction: $P(n) \rightarrow P(n+1)$	$P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow \dots$
Strong Induction	Base Case: $P(1)$ Induction: $P(1) \wedge \dots \wedge P(n) \rightarrow P(n+1)$	$P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow \dots$ $P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow \dots$ $P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow \dots$
k-Leaping Induction	Base Cases: $P(1), P(2), \dots, P(k)$ Induction: $P(n) \rightarrow P(n+k)$	$P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow \dots$
Exponential Induction	Base Case: $P(1)$ Induction: $P(n) \rightarrow P(2n) \wedge P(2n+1)$	$P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow P(6) \rightarrow P(7) \rightarrow \dots$
Backward Induction	Base Case: $P(1)$ Induction: $P(n) \rightarrow P(2n) \wedge P(n-1)$	$P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow P(4) \rightarrow P(5) \rightarrow P(6) \rightarrow P(7) \rightarrow P(8) \rightarrow \dots$

1. **Associative:** $(p \wedge q) \wedge r \stackrel{\text{eqv}}{=} p \wedge (q \wedge r);$
 $(p \vee q) \vee r \stackrel{\text{eqv}}{=} p \vee (q \vee r).$

2. **Commutative:** $p \wedge q \stackrel{\text{eqv}}{=} q \wedge p;$
 $p \vee q \stackrel{\text{eqv}}{=} q \vee p.$

3. **Negations:** $\neg(\neg p) \stackrel{\text{eqv}}{=} p;$
 $\neg(p \wedge q) \stackrel{\text{eqv}}{=} \neg p \vee \neg q;$
 $\neg(p \vee q) \stackrel{\text{eqv}}{=} \neg p \wedge \neg q.$

4. **Distributive:** $p \vee (q \wedge r) \stackrel{\text{eqv}}{=} (p \vee q) \wedge (p \vee r);$
 $p \wedge (q \vee r) \stackrel{\text{eqv}}{=} (p \wedge q) \vee (p \wedge r).$

5. **Implication:** $p \rightarrow q \stackrel{\text{eqv}}{=} \neg q \rightarrow \neg p;$
 $p \rightarrow q \stackrel{\text{eqv}}{=} \neg p \vee q.$

Harmonic Sum: $H_n = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} = \sum_{i=1}^n \frac{1}{i}.$ Use $f(x) = 1/x$ to get that

$$H_n = 1 + \sum_{i=2}^n \frac{1}{i} \leq 1 + \int_1^n dx \frac{1}{x} = 1 + \ln n.$$

GCD facts:

- (i) $\gcd(m, n) = \gcd(m, \text{rem}(n, m)).$
- (ii) Every common divisor of m, n divides $\gcd(m, n)$.
- (iii) For $k \in \mathbb{N}$, $\gcd(km, kn) = k \cdot \gcd(m, n)$.
- (iv) IF $\gcd(l, m) = 1$ AND $\gcd(l, n) = 1$, THEN $\gcd(l, mn) = 1$.
- (v) IF $d|mn$ AND $\gcd(d, m) = 1$, THEN $d|n$.

Big Q is all rational numbers
 Big N is all positive integers
 Big Z is all integers
 E is all positive evens
 R is all real numbers

Binomial Distribution. Let \mathbf{X} be the number of successful trials in an experiment with n independent trials, each trial having success probability p ; $\mathbf{X} = \mathbf{X}_1 + \cdots + \mathbf{X}_n$ is a sum of n independent Bernoullis, with $\mathbb{P}[\mathbf{X}_i = 1] = p$. The PDF of \mathbf{X} is a Binomial distribution, $P_{\mathbf{X}}(k) = B(k; n, p)$, where

$$B(k; n, p) = \binom{n}{k} p^k (1 - p)^{n-k}.$$

$$b^y = x$$

Then the base b logarithm of x is equal to y:

$$\log_b(x) = y$$

For example when:

$$2^4 = 16$$

Then

$$\log_2(16) = 4$$

$A \subseteq B$	<i>A is a subset of B</i>	<i>Every element of A is in B</i>
$A \subset B$	<i>A is a proper subset of B</i>	<i>Every element of A is in B and at least 1 element of B is not in A</i>
$A = B$	<i>A is equals to B</i>	<i>All elements in A are in B, same with B</i>

	No repetition	With repetition
K sequence (order matters)	$N!/(n-k)!$	N^k
K subset (order does not matter)	$\binom{n}{k}$	$\binom{k+n-1}{n-1}$