



Information theory



9

Lecture 01

=>INTRO

- 1-Who is considered the father of information theory?
- A) Claude Shannon
- B) Sherif Barakat
- C) Albert Einstein
- D) Newton
- 2- What are the main question that information theory seeks to answer regarding data transmission?
- A) How much can data be compressed?
- B) How fast can data be reliably transmitted over a noisy channel?
- C) Both of them
- D) None of them
- 3- What are the basic "point-to-point" communication theorem?
- A) Source coding theorem
- B) Channel coding theorem
- C) Both of them
- D) None of them
- 4- What are the Applications of Information theory?
- A) Data compression
- B) Error correcting coding
- C) Transmission and modulation
- D) Image processing: texture
- E) Information security
- F) All of them



- 5- is a sequence of symbols that can be interpreted as a message?
- A) Information
- B) Data
- c) uncertainty
- D) Entropy
- 6- is what you get when your uncertainty about something is reduced?
- A) Information
- B) Data
- c) uncertainty
- D) Entropy
- 7- is what you get when your uncertainty about something is reduced?
- A) uncertainty
- B) Data
- c) Information
- D) Entropy

=>AMOUNT OF INFORMATION

-Let X be an information source with M possible outcomes. The amount of information I(X) is given by:

$$I(X) = \log(M)$$

- -The base of logarithm used in this course is base 2
- -The amount of information is measured in bits.



8- When you flip a fair coin, Calculate the amount of information received....

- A) 1 bit
- B) 2 bit
- c) 3 bit
- D) 4 bit



Ans: $I(X) = log_2 2 = 1 bit$

For a fair coin with two possible outcomes, we have one bit of uncertainty. When it lands we lose that uncertainty and gain one bit of information.

9- When you flip an unfair coin, Calculate the amount of information received....

- A) 5 bit
- B) 10 bit
- c) 0 bit
- D) 1 bit



Ans: $I(X) = log_2 1 = 0$ bit

For an unfair coin that always lands on heads, it has only one possible outcome.

10- When you roll a dice, Calculate the amount of information received....

- A) 7.458 bit
- B) 6.789 bit
- c) 0 bit
- D) 2.585 bit



Ans: $I(X) = log_2 6 = 2.585 bits$

Since the last example has more outcomes, the amount of uncertainty increases. Therefore, the amount of information 1s bigger than that of the coin examples.





=>If the outcomes are not equally likely?

Let symbols = $\{a1, a2, ..., aM\}$ be the possible outcomes of an information source, the probabilities associated with these symbols be $\{p1, p2, ..., PM\}$. The amount of information that we receive when we get a specific symbol ai is:

$$I(a_i) = \log(\frac{1}{p_i})$$

11- A fair coin with a set of two probabilities {0.5, 0.5}, What's is I(a1)?

- A) 5 bit
- B) 6 bit
- c) 0 bit
- D) 1 bit



Ans:
$$I(a_1) = \log\left(\frac{1}{0.5}\right) = 1 bit$$

the first probability is the probability of getting head and the second one is the probability of getting tail.

12- An unfair coin that always comes up as heads, what's is I(a1)?

- A) 9 bit
- B) 3 bit
- c) 0 bit
- D) 1 bit

Ans:
$$I(a_1) = \log\left(\frac{1}{1}\right) = 0$$
 bit

13- A biased coin with probabilities of {0.99, 0.01}, what's I(a1), I(a2) respectively?



- ?
- A) 9 bit, 1 bit
- B) 1 bit, 1 bit
- c) .014 bit, 6.644 bit
- D) .145 bit, .4.687 bit

Ans:

$$I(a_1) = \log\left(\frac{1}{0.99}\right) = 0.014 \ bit$$

 $I(a_2) = \log\left(\frac{1}{0.01}\right) = 6.644 \ bit$

=>Entropy

14- the average number of bits required to represent an information source?

- A) Information
- B) Data
- c) uncertainty
- D) Entropy
- The entropy of an information source (X) is called H(X)

$$H(X) = \sum_{i=1}^{M} p_i \log\left(\frac{1}{p_i}\right) = -\sum_{i=1}^{M} p_i \log(p_i)$$

- 15- Fair coin with probabilities of {0.5, 0.5}, What's H(X)?
- A) 9 bit
- B) 3 bit
- c) 0 bit
- D) 1 bit

Ans:
$$H(X) = 0.5 \times \log(\frac{1}{0.5}) + 0.5 \times \log(\frac{1}{0.5}) = 1 \text{ bit.}$$

17- Biased coin with probabilities of {0.75, 0.25}, What's H(X)?



- 9
- A) 1.11 bit
- B) .0147 bit
- c) .811 bit
- D) 4.568 bit

Ans:
$$H(X) = 0.75 \times \log(\frac{1}{0.75}) + 0.25 \times \log(\frac{1}{0.25}) = 0.811 \text{ bits.}$$

- 18- Unbalanced dice with probabilities of {0.1, 0.1, 0.1, 0.5, 0.1, 0.1}, What's H(X)?
- A) 2.161 bit
- B) 1.458 bit
- c) .894 bit
- D) 2.887 bit

Ans:
$$H(X) = 5 \times (0.1 \times \log(\frac{1}{0.1})) + 0.5 \times \log(\frac{1}{0.5}) = 2.161 \text{ bit}$$

