

# **Sequential Anytime-Valid Inference (SAVI) using E-Processes (Part II)**

Xiaoqi (Shirley) Liu

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# Overview

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## 1. Recap of part I

- 1.1 Motivation, setup, definitions
- 1.2 Validity: e-processes under  $\mathcal{P}$
- 1.3 Efficiency: e-processes under  $\mathcal{Q}$

## 2. Constructing e-processes (valid for composite $\mathcal{P}$ )

- 2.1 Universal inference (UI)
- 2.2 Betting via sequential e-variables
- 2.3 Mixing fixed-sample-size e-variables

## 3. Further discussions

- 3.1 Caveat: pathologies of e-power
- 3.2 E-processes avoid reasoning about hypothetical worlds

## 4. Summary

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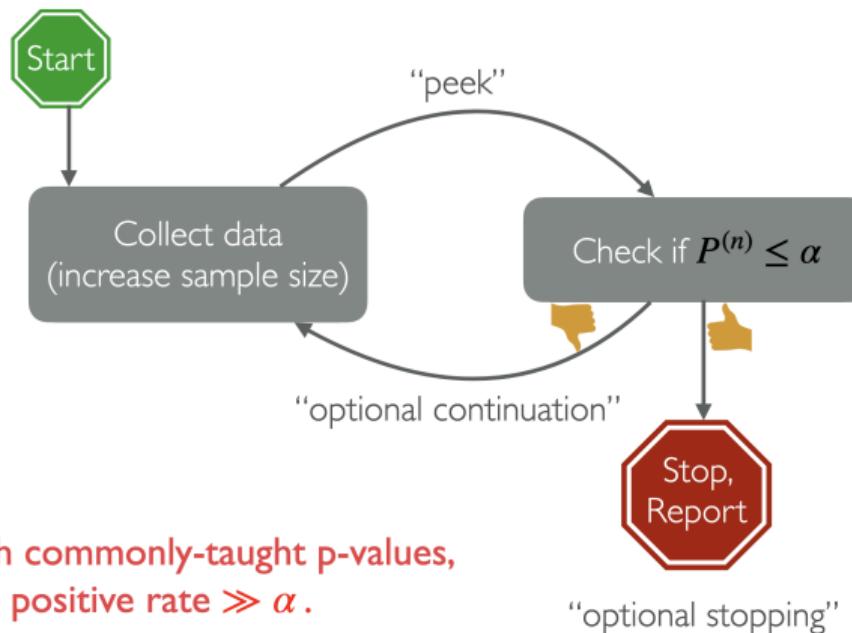
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# Motivation for SAVI (recap)

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What is the problem with continuous monitoring?



## Setup (recap)

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**Natural filtration:** for  $t \in \{0, 1, \dots\}$ ,  $\mathcal{F}_t = \sigma(X_1, \dots, X_t)$ ,  $\mathcal{F}_0 = \{\emptyset, \Omega\}$ ,  $\mathcal{F} := (\mathcal{F}_t)_{t \geq 0}$

**Process:** sequence of r.v.s  $Y = (Y_t)_{t \geq 0}$  **adapted to  $\mathcal{F}$** , i.e., every  $Y_t$  is  $\mathcal{F}_t$ -measurable.

$Y_t$  is **predictable** if  $Y_t$  is  $\mathcal{F}_{t-1}$ -measurable ( $t \geq 1$ )

**Stopping time (or rule)**  $\tau$ : nonnegative integer-valued r.v. s.t.  $\{\tau \leq t\} \in \mathcal{F}_t$  for each  $t \geq 0$ .

**Level- $\alpha$  sequential test for  $\mathcal{P}$ :** a binary process  $\phi = (\phi_t)_{t \geq 1}$  with

$$\sup_{\mathbb{P} \in \mathcal{P}} \mathbb{P}(\exists t \geq 1 : \phi_t = 1) \leq \alpha$$

## Test supermartingales & e-processes (recap)

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**Test supermartingale for  $\mathcal{P}$ .** For every  $\mathbb{P} \in \mathcal{P}$ ,  $M = (M_t)_{t \geq 0}$  satisfies

- $M_t \geq 0$   $\mathbb{P}$ -a.s.,
- $\mathbb{E}^{\mathbb{P}}[M_t | \mathcal{F}_{t-1}] \leq M_{t-1}$ , and
- $\mathbb{E}^{\mathbb{P}}[M_0] \leq 1$ .

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**E-process for  $\mathcal{P}$ .** A sequence of e-values  $E = (E_t)_{t \geq 0}$  adapted to  $\mathcal{F}$  and either (equivalently):

- $\mathbb{E}^{\mathbb{P}}[E_\tau] \leq 1$  for **every** stopping time  $\tau$  and  $\mathbb{P} \in \mathcal{P}$ ; or
- $\exists$  a test supermartingale family  $(M^{\mathbb{P}})_{\mathbb{P} \in \mathcal{P}}$  with  $E_t \leq M_t^{\mathbb{P}}$ ,  $\mathbb{P}$ -a.s. for every  $t \geq 0$  and  $\mathbb{P} \in \mathcal{P}$ .

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## Optional stopping & Ville's inequality (recap)

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**Optional stopping.** If  $M$  is a test supermartingale for  $\mathbb{P}$ , then for any stopping time  $\tau$  and  $\sigma \leq \tau$ :

$$\mathbb{E}^{\mathbb{P}}[M_{\tau} \mid \mathcal{F}_{\sigma}] \leq M_{\sigma} \implies \mathbb{E}^{\mathbb{P}}[M_{\tau}] \leq \mathbb{E}^{\mathbb{P}}[M_0] \leq 1.$$

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**Ville's inequality.** If  $M$  is a non-negative supermartingale for  $\mathbb{P}$ , then for any  $x > 0$ ,

$$\mathbb{P}(\exists t \in \mathbb{N} : M_t \geq x) \leq \frac{\mathbb{E}^{\mathbb{P}}[M_0]}{x}$$

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**Ville's inequality.** If  $M$  is an e-process for  $\mathcal{P}$ , then for any  $\alpha \in (0, 1]$ ,

$$\mathbb{P}\left(\exists t \in \mathbb{N} : M_t \geq \frac{1}{\alpha}\right) \leq \alpha$$

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## Proposition

Every level- $\alpha$  sequential test for  $\mathcal{P}$  is a e-process for  $\mathcal{P}$  thresholded at  $1/\alpha$ .

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**Optional continuation.** If  $M^A, M^B$  are e-processes for  $\mathcal{P}$  and  $M_t^B = 1$  on  $\{t \leq \tau\}$ , then  $M = (M_t)$  with  $M_t := M_{\tau \wedge t}^A M_t^B$  is an e-process for  $\mathcal{P}$ .

## Nontrivial e-processes exist

\* : no reference measure

Any  $\mathcal{P}$  that is sequentially testable

Exchangeability\* (in original filtration)

T-test (in original filtration)

## Nontrivial test supermartingales exist

SubGaussian distributions\* (or any bounded MGF)

Robust, heavy-tailed mean estimation\*

## Nontrivial test martingales exist

Testing symmetry\*

Two-sample testing\*

Bounded means\*

T-test (in shrunk filtration)

Exchangeability\* (in shrunk filtration)

Independence testing\* (in shrunk filtration)

Figure: Hierarchy of tools for constructing sequential tests: nontrivial e-processes (outermost), test supermartingales (middle), and test martingales (innermost), with examples of what each can test.

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## Log-optimality (testing $\mathbb{P}$ vs $\mathbb{Q}$ ) (recap)

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The LR process  $M^*$  given by  $M_0^* = 1$  and  $M_t^* = \frac{d\mathbb{Q}|_{\mathcal{F}_t}}{d\mathbb{P}|_{\mathcal{F}_t}}(X_1, \dots, X_t)$  is a test martingale for  $\mathbb{P}$ .

### Theorem (Log-optimality)

For any stopping time  $\tau$  that is finite  $\mathbb{Q}$ -a.s. and any e-process  $M$  for  $\mathbb{P}$ :

$$\mathbb{E}^{\mathbb{Q}}[\log M_{\tau}^*] \geq \mathbb{E}^{\mathbb{Q}}[\log M_{\tau}].$$

⇒ LR processes are useful benchmark for testing against composite alternatives.

# Asymptotic log-optimality (testing $\mathbb{P}$ vs $\mathbb{Q}$ ) (recap)

## Definition

An e-process  $M$  is **asymptotically log-optimal** for  $\mathbb{P}$  against  $\mathbb{Q}$  if for **every**  $\mathbb{Q} \in \mathcal{Q}$ ,

$$\lim_{t \rightarrow \infty} \frac{1}{t} \left( \log M_t - \log M_t^{\mathbb{Q}} \right) \geq 0 \quad \text{in } L^1\text{-convergence under } \mathbb{Q}$$

where  $M^{\mathbb{Q}}$  is the LR process of  $\mathbb{Q}$  to  $\mathbb{P}$ .

⇒ covers any e-process  $M$  that grows an  $e^{o(t)}$  factor slower than  $M^{\mathbb{Q}}$ .

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**Testing  $\mathbb{P}$  against  $\mathcal{Q} = \{\mathbb{Q}_{\theta} : \theta \in \Theta_1\}$  with iid data:**  $M_0 = 1$ , and with predictable  $\hat{\theta}_{i-1}$ , use

- **Plug-in LR:**  $M_t = \prod_{i=1}^t \frac{q_{\hat{\theta}_{i-1}}(X_i)}{p(X_i)}$ , or
- **Mixture LR:**  $M_t = \int_{\Theta_1} \prod_{i=1}^t \frac{q_{\theta}(X_i)}{p(X_i)} \nu(d\theta)$ ,  $\nu$  a prior on  $\Theta_1$ .

Plug-in is asymptotically log-optimal when  $\theta_i \rightarrow \theta$  under  $\mathbb{Q}_{\theta}$  in a suitable sense, given log-LR is concave, score function has bounded variance.

**Example:** Given iid data from  $N(\theta^\dagger, 1)$ , goal is to test  $H_0 : \theta^\dagger = 0$  vs  $H_1 : \theta^\dagger > 0$ . For illustration, take  $\theta^\dagger = 0.3$ .

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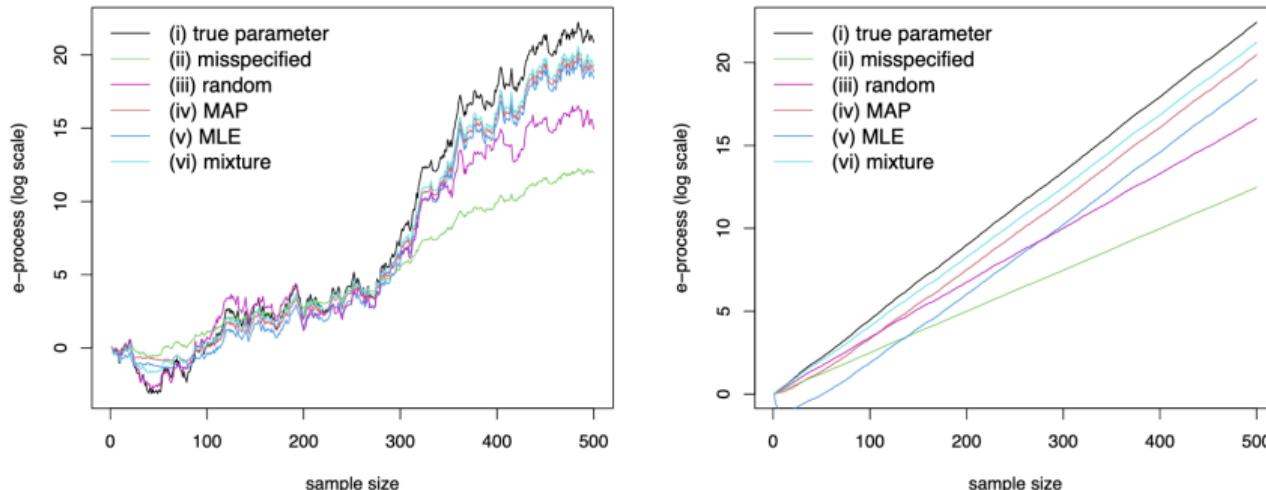


Figure: few ways of constructing e-values from LR processes. Left: one run. Right: average of 1000 runs.

- (i) true parameter: choose  $\theta_i = \theta^\dagger = 0.3$
- (ii) misspecified: choose  $\theta_i = 0.1$
- (iii) random: take iid  $\theta_i$  from  $U[0, 0.5]$
- (iv) MAP: choose  $\theta_i$  by the MAP estimator with prior  $\theta \sim N(0.1, 0.2^2)$
- (v) MLE: choose  $\theta_i$  with  $\theta_1 := 0.1$  and  $\theta_i$  the MLE of  $\theta$  based on  $X_1, \dots, X_{i-1}$
- (vi) mixture: compute a mixture of  $M$  each with a fixed  $\theta_i \in [0, 0.6]$ , uniformly weighted via a discrete grid

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## Universal Inference (UI) e-variable & e-process

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**Goal:** build an e-variable (or an e-process) to test a composite null  $\mathcal{P}$  against composite alternative  $\mathcal{Q}$  **under no regularity conditions.**

**Recipe:** a reduction from hypothesis testing to MLE under  $\mathcal{P}$ .

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## Theorem (UI/Split LR e-variable)

Split data  $X_{1:n}$  into  $D_0 \perp D_1$ , define

$$E = \prod_{i \in D_0} \frac{\hat{q}_1(X_i)}{\hat{p}_0(X_i)}, \quad \hat{q}_1 \in \mathcal{Q} \text{ learned from } D_1, \quad \hat{p}_0 \in \mathcal{P} \text{ is the MLE on } D_0,$$

then  $E$  is an e-variable for  $\mathcal{P}$ .

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## Proposition (UI e-process)

Define  $E_0 = 1$  and

$$E_t = \prod_{i=1}^t \frac{\hat{q}_{i-1}(X_i)}{\hat{p}_t(X_i)} \text{ for } t \geq 1, \quad \hat{q}_{i-1} \in \mathcal{Q} \text{ learned from } X_{1:i-1}, \quad \hat{p}_t \in \mathcal{P} \text{ is the MLE on } X_{1:t},$$

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then  $E = (E_t)$  is an e-process for  $\mathcal{P}$ .

**Proof.** For every  $\mathbb{P} \in \mathcal{P}$ , let  $M_0^\mathbb{P} = E_0 = 1$ , and

$$M_t^\mathbb{P} := \prod_{i=1}^t \frac{\hat{q}_{i-1}(X_i)}{p(X_i)} \geq \prod_{i=1}^t \frac{\hat{q}_{i-1}(X_i)}{\hat{p}_t(X_i)} = E_t \text{ for } t \geq 1.$$

Then  $M_t^\mathbb{P}$  is a test martingale for  $\mathbb{P}$ .

$\Rightarrow$  optional stopping:  $\mathbb{E}^\mathbb{P}[M_\tau^\mathbb{P}] \leq \mathbb{E}^\mathbb{P}[M_0^\mathbb{P}] = 1 \Rightarrow \mathbb{E}^\mathbb{P}[E_\tau] \leq \mathbb{E}^\mathbb{P}[M_\tau^\mathbb{P}] \leq 1$  for every stopping time  $\tau$ .

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then  $E = (E_t)$  is an e-process for  $\mathcal{P}$ .

## Remarks.

- Learn alternative out-of-sample; fit null in-sample.
- The derived sequential test  $\mathbb{1}\{E_t \geq \frac{1}{\alpha}\}$  is computationally expensive: while numerator can be updated in an online fashion, all terms in denominator need to be recalculated after observing each new data point costing  $O(t)$  at step  $t$ .
- Can use mixture instead of plug-in:  $E_t = \int_{\mathcal{Q}} \prod_{i=1}^t \frac{q(X_i)}{\hat{p}_t(X_i)} \nu(dq)$  for any distribution  $\nu$  over  $\mathcal{Q}$ .

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## Testing by betting

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Initialize wealth  $M_0 = 1$ .

For  $t = 1, 2, \dots$ :

- Declare a bet  $E_t : \mathcal{X} \rightarrow [0, \infty)$  with  $\mathbb{E}^{\mathbb{P}}[E_t(X_t) | \mathcal{F}_{t-1}] \leq 1 \quad \forall \mathbb{P} \in \mathcal{P}$ .
- Observe data  $X_t$ .
- Update wealth:  $M_t = M_{t-1} \cdot E_t(X_t) = \prod_{s=1}^t E_s(X_s)$ .

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## Definition (Sequential e-variables)

The e-variables  $E_t$  with  $t \geq 1$  for  $\mathcal{P}$ , adapted to the filtration  $\mathcal{F}$ , are *sequential* if

$$\mathbb{E}^{\mathbb{P}}[E_t | \underbrace{E_1, \dots, E_{t-1}}_{\mathcal{F}_{t-1}}] \leq 1, \quad \forall \mathbb{P} \in \mathcal{P}, \quad \forall t \geq 1.$$

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## Proposition

If  $(E_t)_{t \geq 1}$  are sequential e-variables for  $\mathcal{P}$ ,

$$M_t = \prod_{s=1}^t E_s \quad \text{for } t \geq 1, \quad M_0 = 1,$$

is a test supermartingale (hence e-process) for  $\mathcal{P}$ .

**Proof.**  $\mathbb{E}^{\mathbb{P}}[M_t | \mathcal{F}_{t-1}] = M_{t-1} \mathbb{E}^{\mathbb{P}}[E_t | \mathcal{F}_{t-1}] \leq M_{t-1}$  for every  $\mathbb{P} \in \mathcal{P}$ .

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**Question:** What are the optimal bets?

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**Question: What are the optimal bets?**

- For simple  $\mathcal{P} = \{\mathbb{P}\}$  and  $\mathcal{Q} = \{\mathbb{Q}\}$ ,  $E_t(X_t) = \frac{q(X_t | \mathcal{F}_{t-1})}{p(X_t | \mathcal{F}_{t-1})}$  ensures  $(M_t)$  is log-optimal.

# Testing by betting

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- Update wealth:  $M_t = M_{t-1} \cdot E_t(X_t) = \prod_{s=1}^t E_s(X_s)$ .

**Question: What are the optimal bets?**

- For simple  $\mathcal{P} = \{\mathbb{P}\}$  and  $\mathcal{Q} = \{\mathbb{Q}\}$ ,  $E_t(X_t) = \frac{q(X_t | \mathcal{F}_{t-1})}{p(X_t | \mathcal{F}_{t-1})}$  ensures  $(M_t)$  is log-optimal.
- For composite  $\mathcal{P}$ ,
  - (i) No known analogue of the LR increments that makes  $(M_t)$  log-optimal;
  - (ii) While **numeraires** are log-optimal, they are designed for fixed sample size,  
 $\implies E_t = E^{(t)*}/E^{(t-1)*}$  need not satisfy  $\mathbb{E}^{\mathbb{P}}[E_t | \mathcal{F}_{t-1}] \leq 1$  and  $(E^{(t)*})_{t \geq 1}$  is valid for the “least-favourable”  $\mathbb{P}^* \in \mathcal{P}^{\circ\circ}$  via RIPr but not for all  $\mathbb{P} \in \mathcal{P}$ .

# Testing by betting

---

Initialize wealth  $M_0 = 1$ .

For  $t = 1, 2, \dots$ :

- Declare a bet  $E_t : \mathcal{X} \rightarrow [0, \infty)$  with  $\mathbb{E}^{\mathbb{P}}[E_t(X_t) | \mathcal{F}_{t-1}] \leq 1 \quad \forall \mathbb{P} \in \mathcal{P}$ .
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  - (iii) Possible solution: Avoid all-in; pick  $\lambda_t \in [0, 1]$  to hedge misspecification.

## Testing by betting (generalised)

---

Initialize wealth  $M_0 = 1$ .

For  $t = 1, 2, \dots$ :

- Declare a bet  $E_t : \mathcal{X} \rightarrow [0, \infty)$  with  $\mathbb{E}^{\mathbb{P}}[E_t(X_t) | \mathcal{F}_{t-1}] \leq 1 \quad \forall \mathbb{P} \in \mathcal{P}$ .
- Choose stake  $\lambda_t \in [0, 1]$ .
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- Update wealth:  $M_t = \underbrace{(1 - \lambda_t)M_{t-1} \cdot 1}_{\text{guaranteed wealth}} + \underbrace{\lambda_t M_{t-1} \cdot E_t}_{\text{risky payoff}} = \prod_{s=1}^t ((1 - \lambda_s) + \lambda_s E_s)$

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## Proposition

$(M_t)_{t \geq 0}$  is a test supermartingale (hence e-process) for  $\mathcal{P}$ .

**Proof.**  $\mathbb{E}^{\mathbb{P}}[(1 - \lambda_t) + \lambda_t E_t | \mathcal{F}_{t-1}] \leq (1 - \lambda_t) + \lambda_t \cdot 1 = 1$  for every  $\mathbb{P} \in \mathcal{P}$ .

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## Definition

$(M_t)_{t \geq 0}$  is called an e-process built on  $(E_t)_{t \geq 1}$ .

# Optimising predictable stakes $(\lambda_t)_{t \geq 1}$

## Definitions

- (i) For an alternative measure  $\mathbb{Q}$ , the  $\mathbb{Q}$ -log-optimal e-process built on  $(E_t)$  is  $(M_t)$  with

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- (ii) For  $\gamma \in (0, 1]$ , the empirically adaptive e-process is  $(M_t)$  with

$$\lambda_t \in \arg \max_{\lambda \in [0, \gamma]} \frac{1}{t-1} \sum_{s=1}^{t-1} \log ((1 - \lambda) + \lambda E_s), \quad \lambda_1 = 0.$$

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  - Theorem 3.14: for  $E \geq 0$ ,  $\mathbb{E}^{\mathbb{Q}}[E] > 1 \Leftrightarrow \exists \lambda \in [0, 1]$  s.t.  $\mathbb{E}^{\mathbb{Q}}[\log ((1 - \lambda) + \lambda E)] > 0$   
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- $(M_t)$  generated from method (i) is **log-optimal among e-processes built on  $(E_t)$**

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- $(M_t)$  generated from method (ii), **GRAPA**, has good e-power if  $(E_t)$  are roughly iid under  $\mathbb{Q}$ .

# Empirically adaptive e-processes

## Theorem

Let  $(E_t)_{t \geq 1}$  be iid under the alternative distribution  $\mathbb{Q}$  such that  $\mathbb{E}^{\mathbb{Q}}[\log E_1]$  is finite. The **empirically adaptive** e-process  $(M_t)_{t \geq 0}$  with  $\gamma = 1$  satisfies the following:

- (i) **Asymptotic log-optimality** in the sense that

$$\lim_{t \rightarrow \infty} \frac{1}{t} \left( \log M_t - \log M_t^{\mathbb{Q}} \right) \geq 0 \quad \text{in } L^1\text{-convergence under } \mathbb{Q}$$

for the  **$\mathbb{Q}$ -log-optimal** e-process  $(M_t^{\mathbb{Q}})_{t \geq 0}$  built on  $(E_t)_{t \geq 1}$ .

- (ii) **Consistency**, i.e., if  $\mathbb{E}^{\mathbb{Q}}[E_1] > 1$ , then  $M_t \rightarrow \infty$   $\mathbb{Q}$ -a.s. as  $t \rightarrow \infty$ .

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## Proof.

- (i) follows from LLN.

- (ii) applies Theorem 3.14: for  $E \geq 0$ ,  $\mathbb{E}^{\mathbb{Q}}[E] > 1 \Leftrightarrow \exists \lambda \in [0, 1] \text{ s.t. } \mathbb{E}^{\mathbb{Q}}[\log((1 - \lambda) + \lambda E)] > 0$ .

# Empirically adaptive e-processes

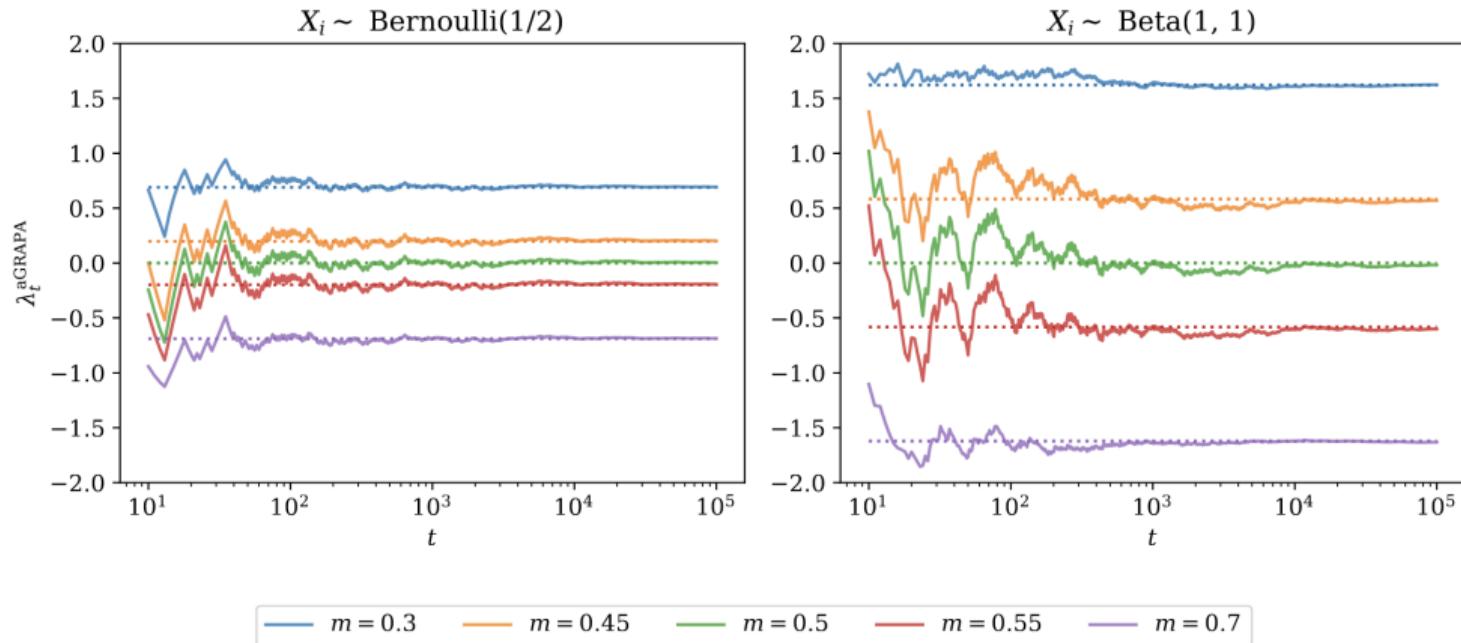


Figure:  $\lambda_t$  chosen using approximated GRAPA under 2 data distributions; dotted lines show oracle stakes.

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## 4. Summary

## Mixing fixed-sample-size e-variables into e-process

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For all  $t \geq 1$ , let  $E^{(t)}$  be e-variable for  $\mathcal{P}$  based on  $X_{1:t}$ . With any pmf  $w$  on  $\mathbb{N}$ , define

$$M_t = \sum_{j=1}^t w(j) E^{(j)}, \quad M_0 = 1.$$

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## Proposition

$(M_t)_{t \geq 1}$  is an e-process for  $\mathcal{P}$ .

**Proof.** For any  $\mathbb{P} \in \mathcal{P}$  and any random time  $\tau$ ,

$$\mathbb{E}^{\mathbb{P}}[M_{\tau}] = \sum_{j \geq 1} w(j) \mathbb{E}^{\mathbb{P}}[E^{(j)} \cdot \mathbb{1}\{\tau \geq j\}] \leq \sum_{j \geq 1} w(j) = 1.$$

**Remark.**  $(M_t)_{t \geq 1}$  is an increasing process, which is unusual for e-processes.

## Mixing fixed-sample-size **numeraires** into e-process

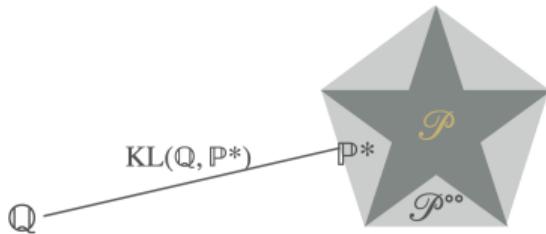
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# Mixing fixed-sample-size numeraires into e-process

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Recall:

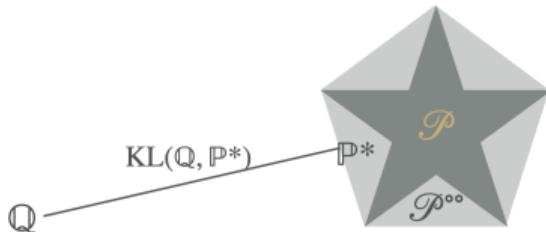
- A numeraire is a  $\mathbb{Q}$ -a.s. strictly positive e-variable  $E^*$  s.t.  $\mathbb{E}^{\mathbb{Q}}[E/E^*] \leq 1$  for every e-variable  $E$ .  $\implies \mathbb{E}^{\mathbb{Q}}[\log(E/E^*)] \leq 0$ .
- “Composite LR”:  $E^*$  is LR between  $\mathbb{Q}$  and some element  $\mathbb{P}^* \in \mathcal{P}^{\circ\circ}$  (i.e., RIPr of  $\mathbb{Q}$  onto  $\mathcal{P}$ ).



# Mixing fixed-sample-size numeraires into e-process

Recall:

- A numeraire is a  $\mathbb{Q}$ -a.s. strictly positive e-variable  $E^*$  s.t.  $\mathbb{E}^{\mathbb{Q}}[E/E^*] \leq 1$  for every e-variable  $E$ .  $\Rightarrow \mathbb{E}^{\mathbb{Q}}[\log(E/E^*)] \leq 0$ .
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Can mix fixed-sample-size numeraires to obtain an e-process:

Popular choice of mixing pmf:  $w(t) = \frac{c}{t(\log t)^2}$  with  $\sum_{t \in \mathbb{N}} w(t) = 1$  and constant  $c > 0$ .  
 $\Rightarrow \log M_t \geq \log(w(t)E^{(t)*}) = \log E^{(t)*} - \log t - 2 \log \log t + \log c$   
 $\Rightarrow \lim_{t \rightarrow \infty} \frac{1}{t} (\log M_t - \log E^{(t)*}) = 0$ .

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Heavy tailed  $\mathbb{Q}$  can make finite-sample e-power misleading:

**Example:** let  $Y_1, Y_2, \dots$  be iid Pareto(1) under  $\mathbb{Q}$  (i.e.,  $\mathbb{Q}(Y > x) = 1/x$  for  $x \geq 1$ ), define  $E_t = \exp(t^2 - Y_t)$  and set  $M_t = \prod_{s=1}^t E_s$  for  $t \in \mathbb{N}$ . Then

$\Rightarrow \mathbb{E}^{\mathbb{Q}}[\log E_t] = -\infty$ , but  $\log E_t$  is  $\mathbb{Q}$ -a.s. finite

$\Rightarrow \mathbb{E}^{\mathbb{Q}}[\log M_t] = -\infty$  for every  $t$ , while  $M_t \rightarrow \infty$  in probability under  $\mathbb{Q}$  (so  $\mathbb{Q}(M_t \geq 1/\alpha) \rightarrow 1$ )

**Takeaway:** Extra caution needed when infinity is involved in the calculation of e-power. Use consistency diagnostics alongside e-power.

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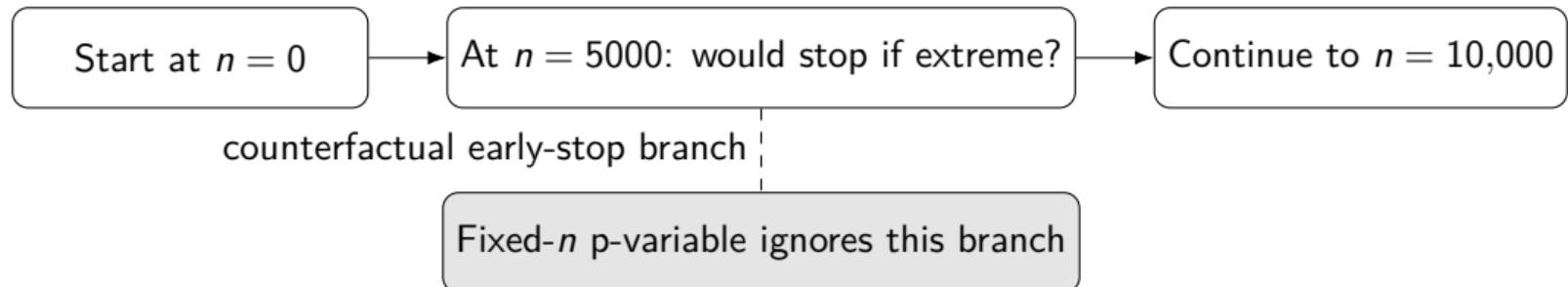
## E-processes avoid reasoning about hypothetical worlds

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Even if an analyst **actually** stops at  $n = 10,000$ , the mere **willingness** to stop earlier if results had been extreme makes the stopping time data-dependent.

**Consequence.** Reasoning about the validity of p-variables requires reasoning about all possible “hypothetical worlds”.

**E-process fix.** Report  $E_\tau$  (or threshold at  $1/\alpha$ ); valid for **any** stopping time  $\tau$ , including adaptive or hypothetical ones.



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**Anytime validity:** E-process controls type-I error at all stopping times; threshold at  $1/\alpha$  via Ville.

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**Simple null:**

- LR process is log-optimal
- plug-in or mixture LR for composite alternative can be asymptotically log-optimal.

**Composite null:**

- **UI e-process:** learn alt out-of-sample, fit null in-sample; applicable to irregular problems.
- **Betting via sequential e-variables ( $E_t$ ):** Tune stakes  $\lambda_t \in [0, 1]$  to form empirically adaptive e-process, matching asymptotic growth rate of its oracle counterpart if  $(E_t)$  are iid.
- **Mixing fixed-sample-size numeraires:** resulting e-process matches the numeraires' asymptotic growth rate.