

On a system using Round Robin Scheduling, let s represent the time required to perform a process switch, q represent the RR time quantum, and r represent the average time a process runs before blocking on I/O. Give a formula for CPU efficiency given the following:

- i) $q = \infty$ ii) $q > r$ iii) $s < q < r$ iv) $s = q < r$ v) q nearly 0

- 1) Assuming r is the total processing time, and s is the overhead for each context switch.
 Q is the time quantum, then the CPU efficiency can be represented by the equation:

$$r / ([r/Q] - 1) * s + r$$

Solution:

- (a) When, $Q = \infty$, $s = 0$

It becomes $T/T = 1$.

So, CPU efficiency is 100%.

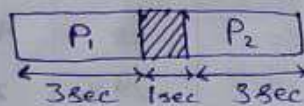
- (b) When, $Q > r$, $s = 0$

It becomes $T/T = 1$.

So, CPU efficiency is 100%.

- (c) $s < Q < r$

Example:



$$P_1 = 3 \text{ sec} \quad P_2 = 3 \text{ sec}$$

Switching = 1 Sec.

$$\text{Efficiency Formula} = \frac{r}{([r/Q] - 1) * s + r}$$

Let, $Q = 2$

$$u = \frac{6}{([6/2] - 1) * 1 + 6} = \frac{6}{2 + 6} = \frac{6}{8} = 75\%$$

Let, $Q = 4$

$$u = \frac{6}{([6/4] - 1) * 1 + 6} = \frac{6}{(1.5 - 1) + 6} = 85.71\%$$

2
∴ The CPU efficiency is $\eta / ([\eta/Q] - 1) * S + \eta$.

(d) When, $S = Q < \eta$

$$\begin{aligned}\eta / (\eta/Q * S - S) + \eta &= \eta / (\eta - S) + \eta \\ &= \eta / 2\eta - S \\ &= \eta / 2\eta - Q\end{aligned}$$

(e) Q nearly zero

When, Q is nearly zero.

It becomes $\eta / \text{infinity} = 0$.