Hashing 2

Sets, Dicts

Problem Solving
Session 2
On Saturday.

Agenda

- Target Pair Sum
- Distinct elements in window of size k

Q1 Target Pair Sum

Adobe Spotify amazon

Given N array elements, check if there exists a pair (i, j) such that

↳ boolean ↑ Indexes

A[i] + A[j] == k && i != j

a + b = k

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5

	i	j	A[i]	A[j]	
k = 11	4	8	4	7	
k = 6	0	3	8	-2	Quiz 1
k = 22	← No pair				

Brute Force

```
for (i=0; i<N; i++) {
    a = A[i]
    b = k-a
    for(j=0; j<i; j++) {
        if( A[j] == b)
            return True
    }
}
return false
```

a + b = k

⇒ b = k - a

TC : O(N²)

SC : O(1)

Idea & Pseudocode

k=6

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5

b = k - a -2

⇒ True

hs = { 8, 9, 1, -2, 4, 5, 11, -6, 7, }

---

k=22

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5
b=k-a	14	13	21	24	18	17	11			

↪ True

hs = { 8, 9, 1, -2, 4, 5, 11, -6, 7, }

Above Hashset approach fails

k=10

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5
b=k-a	2	1								

if (a==b)
check if freq[a] >=2

else
check if freq[b] >=1

Freq

{
8→1, 9→1, 1→1
-2→1, 4→1
5→2, 11→1
-6→1, 7→1
}

k=22

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5
b=k-a	14	13	21	24	18	17	11			

Hashmap <int, int> freq;

// Create a freq map ← TODO

for(i=0; i<N; i++) {

a = A[i]

b = k-a

```
// Search for b
if( a==b and freq[b] >= 2)
        return True
else if( a!=b and freq[b] >= 1)
        return True
}
return false
```

TC: O(N)
SC: O(N)

Java

```java
boolean targetPairSum(int[] A, int target) {
    HashMap<Integer, Integer> hm = new HashMap<Integer, Integer>();
    for (int i = 0; i < A.length; i++) {
        hm.put(A[i], hm.getOrDefault(A[i], 0) + 1);
    }

    for (int i = 0; i < A.length; i++) {
        int b = target - A[i];
        if (hm.containsKey(b)) {
            if (A[i] == b && hm.get(b) >= 2)
                return true;
            else if (A[i] != b)
                return true;
        }
    }

    return false;
}
```

Python

```python
def targetPairSum(A, target):
    hm = {}
    for i in range(len(A)):
        hm[A[i]] = hm.get(A[i], 0) + 1

    for i in range(len(A)):
        b = target - A[i]

        if b in hm:
            if A[i] == b and hm[b] >= 2:
                return True
            elif A[i] != b:
                return True

    return False
```

Using Hashset

k = 12

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5
b = k-a	4	3	11	14	8					

Valid pair - A[0], A[4]
(8, 4)
Ans = True

Hashset
{ 8, 9, 1, -2 }

1) Iterate
2) Check if (k-a) is present
3) Insert A[i] into the set

---

k = 22

	0	1	2	3	4	5	6	7	8	9
A[] =	8	9	1	-2	4	5	11	-6	7	5
b = k-a	14	13	21	24	18	17	11	28		

⇒ Ans = false

Hashset
{ 8, 9, 1, -2, 4, 5 }

At any index i, the hashset should only contain the elements from [0 to i-1]

```
Hashset <int>  s;

for (i=0;  i<N;  i++) {
        a = A[i]
        b = k-a

        // Check if b is present
        if ( b in s)
            return True

        // Insert A[i] into the set
        s.insert (a)
}
return false
```

TC: O(N)
SC: O(N)

Carry forward on our set.

Java

```java
boolean targetPairSum(int[] A, int target) {
    HashSet<Integer> s = new HashSet<Integer>();

    for (int i = 0; i < A.length; i++) {
        int b = target - A[i];
        if (s.contains(b)) {
            return true;
        }
        s.add(A[i]);
    }

    return false;
}
```

Python

```python
def targetPairSum(A, target):
    s = set()
    for i in range(len(A)):
        b = target - A[i]

        if b in s:
            return True

        s.add(A[i])

    return False
```

Break till 10:20 PM

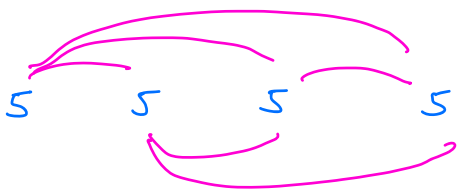
Q Leetcode - Two Sum

Q. What if ... ?

Q Calculate the no of pairs (i, j) such that

$A[i] + A[j] == k$ and $i != j$

$k = 10$



6 pairs

Hashmap

Q Check if there exists a pair (i, j) such that

$A[i] - A[j] == k$, $i != j$

$a - b = k$

$\Rightarrow b = a - k$

Q Same as Two Sum Problem. But return (i, j) pair.

Hashmap < num, index >

Q2 Given N array elements, calculate no. of distinct elements in every window of size k.

subarray

amazon

Microsoft

k=4

	0	1	2	3	4	5	6	7	8	9
arr[] =	2	4	3	8	3	9	4	9	4	10

[0 3]	4
[1 4]	3
[2 5]	3
[3 6]	4
[4 7]	3
[5 8]	2
[6 9]	3

First subarray - [0 k-1]

Last subarray - [N-k N-1]

Length	No of subarrays
1	N
2	N-1
3	N-2
4	N-3
⋮	⋮
k	N-k+1

Quiz 2

Brute Force

For every window, get no of distinct elements

```
for(i=0; i<=N-k; i++) {
    // subarray [i  i+k]
    HashSet<int> set
    for(j=i; j<=(i+k); j++) {
        set.insert(A[j])
    }
    print(set.size())
}
```

TC: $(N-k) * k$

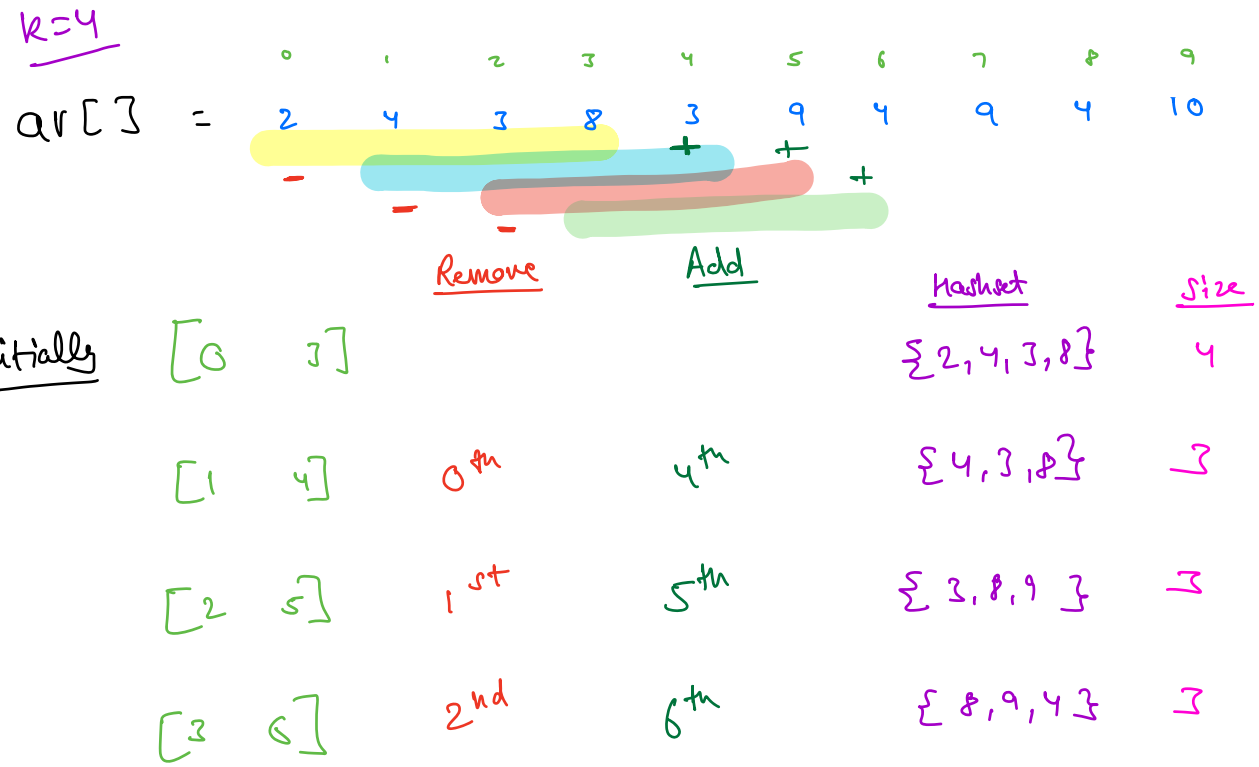
$\Rightarrow \frac{N}{2} * \frac{N}{2}$

$\Rightarrow O(N^2)$

worst case $k = \frac{N}{2}$

SC: $O(k)$

Optimised Approach



Issue:

If we remove an element, all occurrences of that element are removed from the set.

Hashmap with Sliding Window

	0	1	2	3	4	5	6	7	8	9
arr[] =	2	4	3	8	3	9	4	9	4	10

K=4

	Remove $(i-1)^{th}$	Add	Hashmap	Size
[0 3]			{2→1, 4→1, 3→1, 8→1}	4
[1 4]	0^{th}	4^{th}	{4→1, 3→2, 8→1}	3
[2 5]	1^{st}	5^{th}	{3→2, 8→1, 9→1}	3
[3 6]	2^{nd}	6^{th}	{3→1, 8→1, 9→1, 4→1}	4
[4 7]	3^{rd}	7^{th}	{3→1, 9→2, 4→1}	3
[5 8]	4^{th}	8^{th}	{9→2, 4→2}	2
[6 9]	5^{th}	9^{th}	{4→2, 9→1, 10→1}	3

if freq == 0, remove it.

Pseudocode

```
Hashmap <int, int>  freq

// Prepare the first window
for(i=0; i<k; i++) {
        if( arr[i] in freq )
                freq[arr[i]] ++
        else
                freq[arr[i]] = 1
}
print( freq.size())
i=1
j=k
while (i<=n-k OR j<n) {
        // Subarray [i  j]

        // Remove (i-1)th index element
        freq[ arr[i-1]] --;
```

```
if ( freq[arr[i-1]] == 0)
    freq.remove ( arr[i-1])

// Add jth element

if ( arr[j] in freq)
    freq[arr[j]] ++

else
    freq[arr[j]] = 1

print( freq.size())

i ++

j ++

}
```

Java

```java
void distinctElementsInWindow(int[] arr, int k) {
    int n = arr.length;
    HashMap<Integer, Integer> freq = new HashMap<Integer, Integer>();

    for (int i = 0; i < k; i++) {
        // freq.put(arr[i], freq.getOrDefault(arr[i], 0) + 1);
        if (freq.containsKey(arr[i]))
            freq.put(arr[i], freq.get(arr[i]) + 1);
        else
            freq.put(arr[i], 1);
    }

    System.out.println(freq.size());

    int i = 1;
    int j = k;

    while (i <= (n - k)) {
        // Subarray - [i j]

        // Remove (i-1)th element
        freq.put(arr[i - 1], freq.get(arr[i - 1]) - 1);

        if (freq.get(arr[i - 1]) == 0)
            freq.remove(arr[i - 1]);

        // Add (j)th element
        // freq.put(arr[j], freq.getOrDefault(arr[j], 0) + 1);
        if (freq.containsKey(arr[j]))
            freq.put(arr[j], freq.get(arr[j]) + 1);
        else
            freq.put(arr[j], 1);

        System.out.println(freq.size());

        i++;
        j++;
    }
}
```

Python

```python
def distinctElementsInWindow(arr, k):
    n = len(arr)
    freq = {}

    for i in range(k):
        # freq[arr[i]] = freq.get(arr[i], 0) + 1
        if arr[i] in freq:
            freq[arr[i]] += 1
        else:
            freq[arr[i]] = 1

    print(len(freq))

    i = 1
    j = k

    while i <= (n - k):
        # Subarray - [i j]

        # Remove (i-1)th element
        freq[arr[i - 1]] -= 1

        if freq[arr[i - 1]] == 0:
            freq.pop(arr[i - 1])

        # Add (j)th element
        # freq[arr[j]] = freq.get(arr[j], 0) + 1

        if arr[j] in freq:
            freq[arr[j]] += 1
        else:
            freq[arr[j]] = 1

        print(len(freq))

        i += 1
        j += 1
```

Doubts

Thank You

Wednesday

Python dict - .get( )

getOrDefault

d.get(2,0) → 9

d.get(10,0) → 0

d.get(50, 500) → 50

d = { 2 → 9, 4 → 5 }

Good Night

Happy Rakhi

Happy Independence Day

Happy Pavi New Year

Happy Long Weekend

Thank you

Wednesday