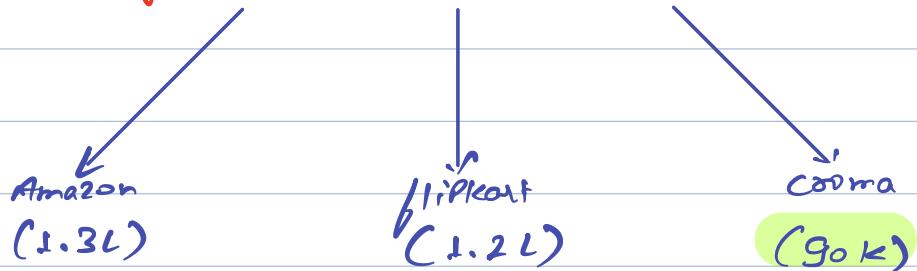


Today's Content

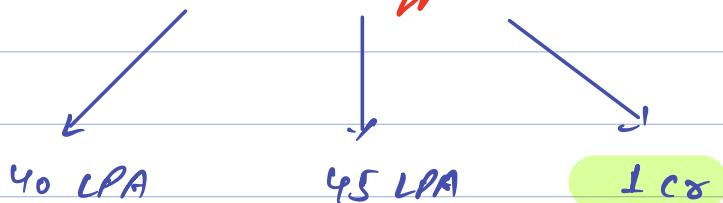
- ↳ Currency exchange
- ↳ fractional knapsack
- ↳ Greedy Properties
- ↳ Activity Selection
- ↳ job Scheduling

Greedy Algorithm

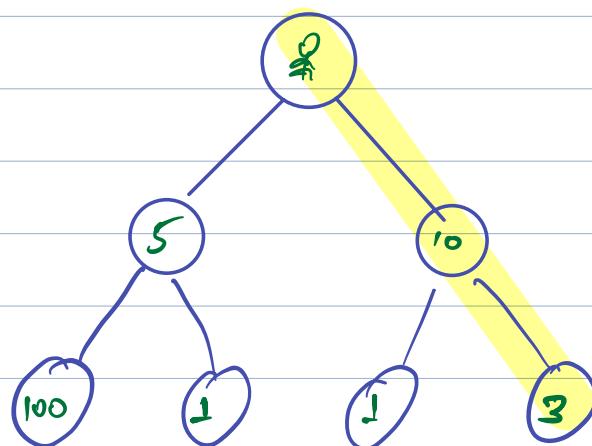
Buy a Phone (iPhone 14 Pro max)



Accept the offer



↳ Greedy algo is an approach to solve problems by making locally optimal choices at each step.



↗ 1 is always there.

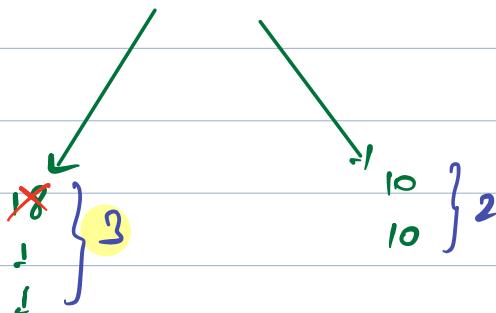
$$\text{curr}(i) \geq \text{curr}(i-1);$$

Indian currency: 1, 2, 5, 10, 20, 50, 100, 200, 500, 2000

Cash: 5548 → min number of coins/notes to get required cash?
↳ multiples of same coins are allowed.

Notes/Coins	Count	left amount.
2000	2	1548
500	3	48
20	2	8
5	1	3
2	1	1
1	1	1
	<u>1</u>	
	10	

Currency: 1 10 18 amount: 20



↳ Need to recheck this example.

1 10 21 amount = 51
amount = 30

a) Fractional Knapsack

If you can consume K Kg of vegetables, you can eat any integral amount of item. Max Protein you Can get?

vegetables	Eating complete item Protein gained	Max Protein/kg	$K = 70 \text{ kg}$
Tomato $\rightarrow 20 \text{ kg}$	$200 \text{ P} \rightarrow 10 \text{ P/kg} * 20 = 200 \text{ P}$		
Apples $\rightarrow 15 \text{ kg}$	$180 \text{ P} \rightarrow 12 \text{ P/kg} * 15 = 180 \text{ P}$		
onion $\rightarrow 50 \text{ kg}$	$250 \text{ P} \rightarrow 5 \text{ P/kg}$		
Chicken $\rightarrow 10 \text{ kg}$	$150 \text{ P} \rightarrow 15 \text{ P/kg} * 10 = 150 \text{ P}$		
Potato $\rightarrow 25 \text{ kg}$	$200 \text{ P} \rightarrow 8 \text{ P/kg} * 25 = 200 \text{ P}$		
Mango $\rightarrow 12 \text{ kg}$	$132 \text{ P} \rightarrow 11 \text{ P/kg} * 12 = 132 \text{ P}$		
Seafood $\rightarrow 5 \text{ kg}$	$100 \text{ P} \rightarrow 20 \text{ P/kg} * 5 = 100 \text{ P}$		
			eat on the basis of max Protein.
			onion $\rightarrow 50 \text{ kg} \rightarrow 250 \text{ P}$
			Tomato $\rightarrow 20 \text{ kg} \rightarrow \frac{200 \text{ P}}{450 \text{ P}}$

given $K = 8$

$25 \leq K$

Class Pair {

int w;

int p;

double PPV;

//Pseudo code

fractional Knapsack (int w[n], int Protein[n], int k) {

Pair [] items = new Pair [n];

T.C: $O(n \log n)$

S.C: $O(n)$

```
for (int i=0; i<n; i++) {  
    Pair P = new Pair (w[i], Protein[i],  $\frac{\text{Protein}[i]}{w[i]+1.0}$ );  
    items[i] = P;  
}
```

Arrays.sort (items);

double ans=0

```
for (int i=N-1; i>=0; i--) {  
    Pair xem = items[i];  
    if (xem.w <= k) {  
        ans = ans + xem.p;  
        k = k - xem.w;  
    } else {  
        ans = ans + (k * xem.ppu);  
        break;  
    }  
}
```

return ans;

}

Knapsack

double ans = 0;

```

for (int i=N-1; i>=0; i--) {
    if (item[i].w <= k) {
        ans = ans + item[i].v;
        k = k - item[i].w;
    } else {
        ans = ans + (k * item[i].ppu);
        break;
    }
}
    
```

3

return ans;

+

0 1 2 3 4 5 6

item[7]: {50, 250} {25, 200} {20, 200, 10} {12, 132, 11} {15, 180} {10, 150, 15} {5, 100, 20}

WIP: 0 1 2 3 4 5
20 15 50 10 25 12 5

PPU: 200 180 250 150 200 132 100
10

$$K = 70 - 5 = 65 - 10 = 55 - 15$$

$$= 40 - 12 = 28 - 20 = 8$$

Pair item: item[6]

$$ans = 0 + 100 + 150 + 180 + 132 + 200 + 64$$

* Greedy Properties

binary search
greedy
dp

6(i) for min/max related Problems.

(ii) based on what Parameter we want to apply greedy.

(iii) either Prove it logically or discard it with One counter example.

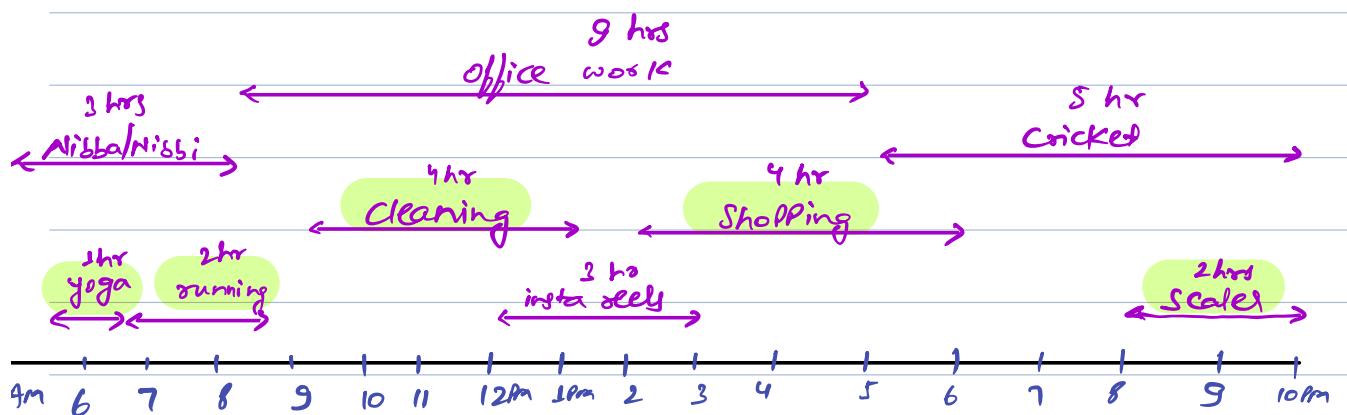
Balak till 10:25pm

Q2) Activity Selection

↳ max Count of task you can do.

Note: ① Start a task, we need to complete.

② At any point of time, single task.



① min duration task first XX

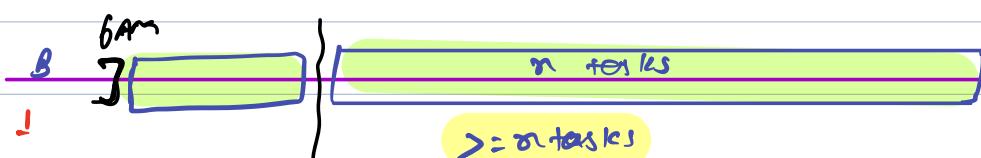
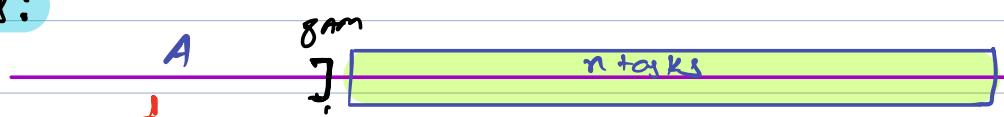
↳ yoga, running, Scalper, insta reels

② ending early

↳ yoga, running, Cleaning, Shopping, Scalper

↳ ans = 5

Correctness:



// Pseudo code

class Pair {

 int s;
 int e;

int ActivitySelection (int Start[n], int end[n]) {

 Pair [] arr = new Pair[n];

 for (int i=0; i<n; i++) {

 Pair P = new Pair [start[i], end[i]);

 arr[i] = P

}

 // sort on the basis of e.

(end time)

 Arrays.sort (arr);

T.C: O(n log n)

S.C: O(n)

// Pick all the Non-overlapping tasks.

int ans = 1;

int endtime = arr[0].e;

for (int i=1; i<n; i++) {

 if (arr[i].s >= endtime) {

 ans++;

 endtime = arr[i].e;

3

3

return ans;

3

yoga

(5, 7)

(5, 8)

(7, 9)

(9, 13)

nibolnissi

(12, 15)

(8, 17)

(14, 18)

dr

Job Scheduling

Given n tasks to complete

→ Deadline assigned for each task, day on or before we can do task.

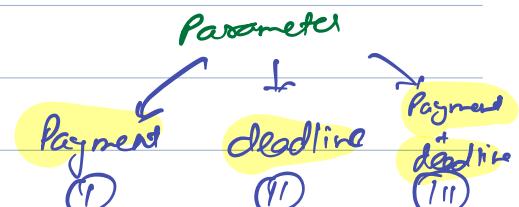
→ Payment assigned to each task.

→ On any given day we can perform only 1 task and each task takes 1 day to finish.

→ find max Payment we can get.

Ex1:

Job	deadline day	Payment
a	3	100
b	1	19
c	2	27
d	1	25
e	3	30

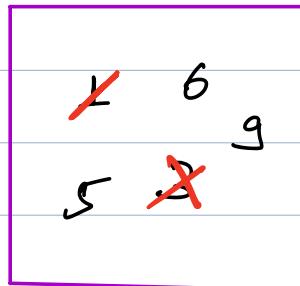


Ex2:

Job	deadline day	Payment	→ Sort on the basis of deadline
a	3	5	Jobs: b d a c e
b	1	1	deadline: 1 2 3 3 3
c	3	6	Payment: 1 3 5 6 9
d	2	3	task at greater index can replace task at lower index.
e	3	9	

Jobs: b d a c e
 deadline: 1 2 3 3 3
 Payment: 1 3 5 6 9

$$\frac{c}{1} \quad \frac{d}{2} \quad \frac{a}{3}$$



Ans = 20

Class Pair {

int d;
 int p;

}

int jobscheduling (int deadline[n], int Payment[n]) {

Pair [] arr = new Pair [n];

for (int i=0; i<n; i++) {

Pair P: new Pair (deadline[i], Payment[i]);

arr[i] = P

}

on the basis of deadline.

Arrays-sort (arr);

minheap <int> mhi;

```
for (int i=0; i<N; i++) {  
    Pair elem = arr[i];  
    if (elem.d > mh.size()) {  
        mh.insert(elem.p);  
    } else {  
        if (elem.p > mh.peek()) {  
            mh.remove();  
            mh.insert(elem.p);  
        }  
    }  
}
```

// Sum all the elements of the heap

3

Tracing

0 1 2 3 4 5 6 7 8 9 ✓

Head: 1 1 1 2 2 4 4 5 5 5

for (int i=0; i<N; i++) {
 Pair sem = arr[i];

 if (sem.d > mh.size()) {

 mh.insert(sem.P);

 } else {

 if (sem.P > mh.PeekC()) {

 mh.removeC();

 mh.insert(sem.P);

}

// Sum all the elements of the heap

350 40
250 200 300 200 150 300 250 100 600 400
1 2 3 4 5

250
600 350
200 400
300 250
↑ mh

WHP: 0 1 2 3 4 5

WHP: 20 15 50 10 25 12 5

Protocols: 200 180 250 150 200 132 100

10

```

void bubblesort (int arr[], int n) {
    for (int i=0; i<n-1; i++) { // no. of iter.
        for (int j=0; j<=n-2-i; j++) {
            if (arr[j] > arr[j+1]) {
                swap (arr[j], arr[j+1]);
            }
        }
    }
}

Pair<int, int> wt = new Pair[n];

```

Pair¹
int val1;
int val2;

Quick Sort ~~Bubble Sort~~ Arrays-Sort (0 to 6) Comparator {

0 1 2 3 4 5 6

WTF: (2,2) (1,7) (5,5) (0,0) (5,3) (1,1) (5,6)

↳ sorting on basis of val1 or val2 ??

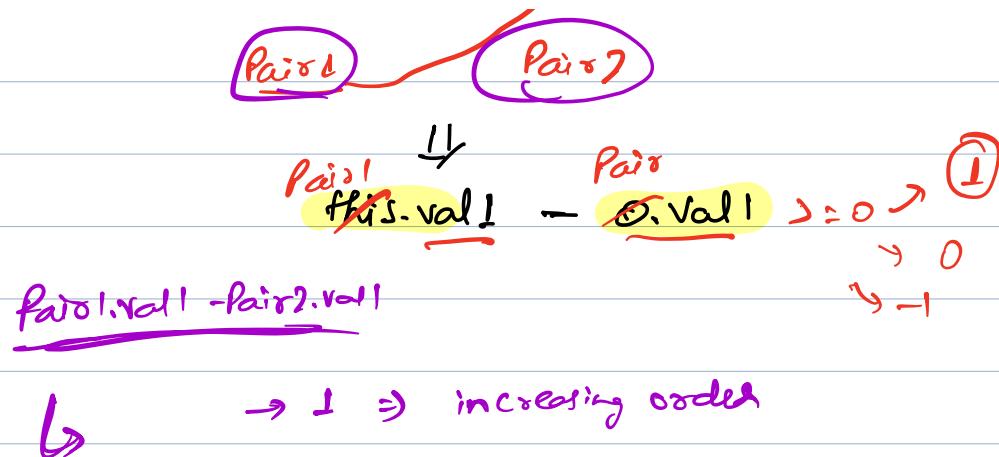
$i=0 \rightarrow$

$$arr[j] > arr[j+1]$$

$$arr[j] - arr[j+1] > 0$$

Comparator ~~(arr[j], arr[j+1])~~ > 0





$$1 + 2 + 3 + \dots + N \approx \frac{n(n+1)}{2} \sim O(n^2)$$

return $(P1.val1 + P1.val2) - (P2.val1 + P2.val2)$

$$\log(a) + \log(b) = \log(a+b)$$

$$\log(1) + \log(2) + \dots + \log(N)$$

$$\log((1+2+\dots+N)) \sim n \log n$$