

Today's Content:

- Binary Tree
- Traversal
- Specs
- Iterative traversals
- level order
- Invert BT

} 45-50

Tree Basics:

○ : node

→ : Edge, connects node

$\text{parent}(F)$: B

$\text{Ancestor}(D)$: E C R

F G E arc: nodes at same level

$\text{height}(\text{Node})$: max length from node to any of its leaf nodes

length calculated based on edges

$\text{height}(\text{Node}) = \max(\text{height of child nodes})$

$\text{height}(\text{leaf}) = 0$

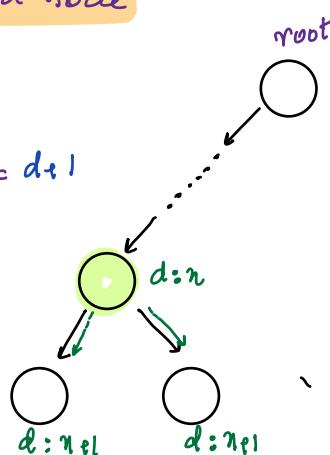
$\text{Depth}(\text{Node})$: length of path from root node to given node

$\text{Depth}(\text{root}) = 0$

$\text{Depth}(\text{Node}) = \text{level of a node}$

$\text{Depth}(\text{node}) = d$

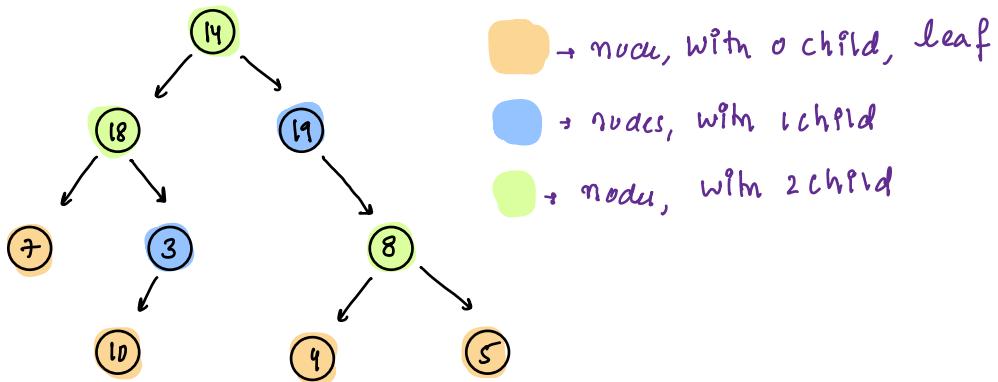
Depth of its child node = $d+1$



Obs: Root node is only node without a parent

Binary Tree: Any node can at max have 2 child nodes

0 1 2 child



Class Node

```

int data
Node left
Node right
  
```

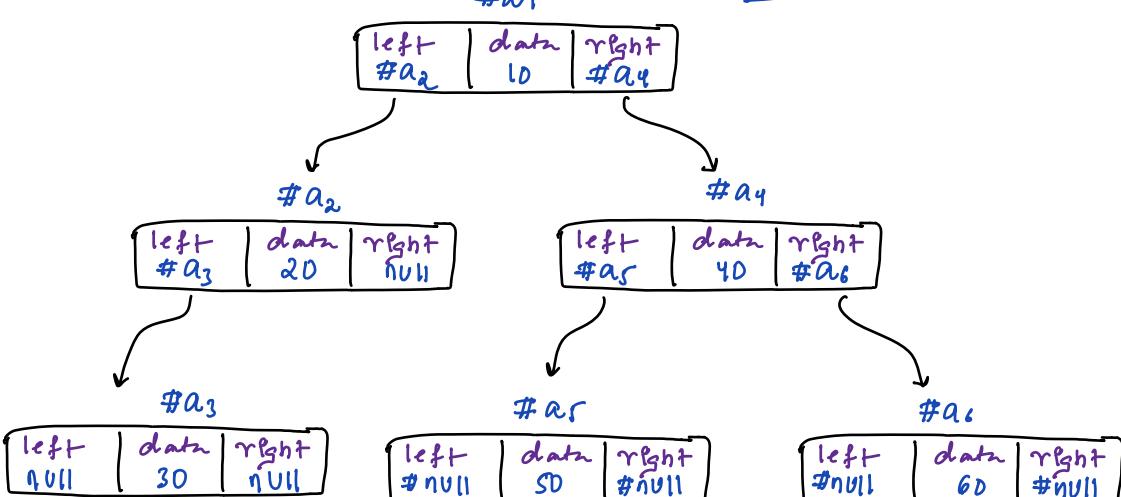
↳ obj reference can hold address of node object

Node(n){ // constructor: To initialize data members of class

```

    data = n
    left = null
    right = null
  
```

→ // Only root node given, with that
we can traverse entire tree



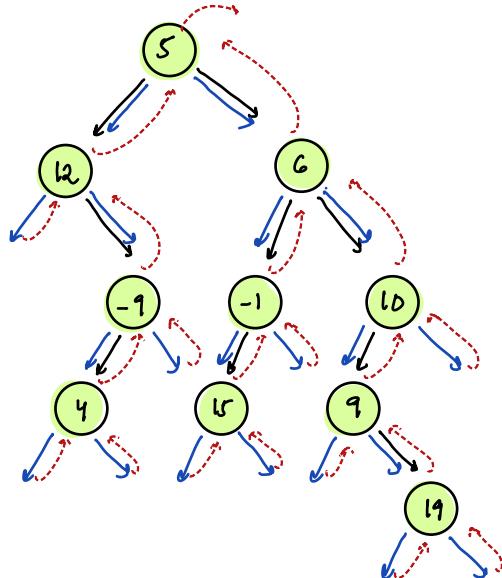
Tree Traversals :

preorder : root Left Right

→ print node data

→ Go to left subtree
print entire left subtree in
preorder

→ Go to right subtree
print entire right subtree in
preorder



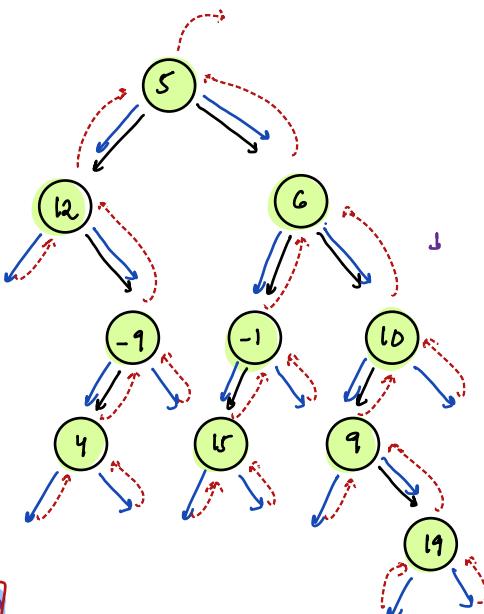
preorder: [root Pre LST] [Pre RST]
[5 | 12 | -9 | 4 | 6 | -1 | 15 | 10 | 9 | 19]

inorder : left data right

→ Go to left subtree
print entire left subtree in
inorder

→ print node data

→ Go to right subtree
print entire right subtree in
inorder



inorder: [In LST] [root] [In RST]
[12 | 4 | -9 | 5 | 15 | -1 | 6 | 9 | 19 | 10]

postorder : left right data

→ Go to left subtree
print entire left subtree in postorder

→ Go to right subtree
print entire right subtree in postorder

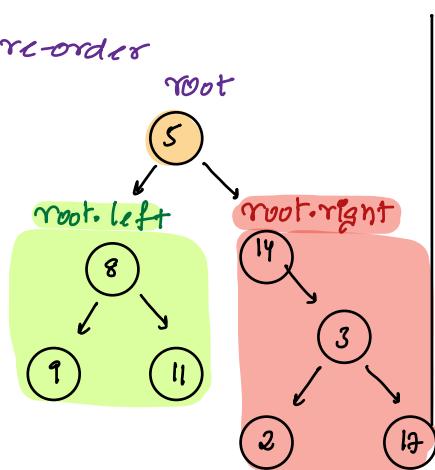
→ print node data

postorder: TODO for above example

Ass: Given root node, print entire tree in pre-order

```
void preorder (Node root) {
```

1. if (root == null) { return; }
2. print (root.data)
3. preorder (root.left)
4. preorder (root.right)



preorder: 1 2 8 4 // Data left Right

inorder: 1 3 2 4 // left Data Right

postorder: 1 3 4 2 // left right Data

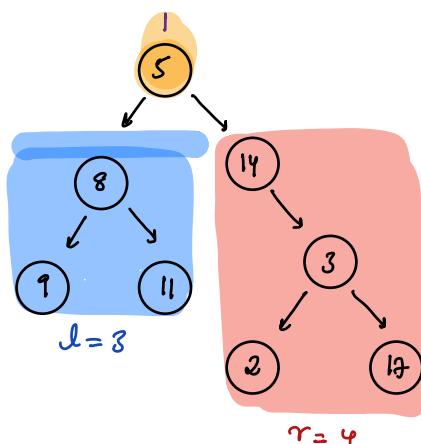
Ass: Given a root node, return no. of nodes in Tree

```
int size (Node root) { Tc: O(N) SC: O(H) } // at max stack size = H Height of Tree
```

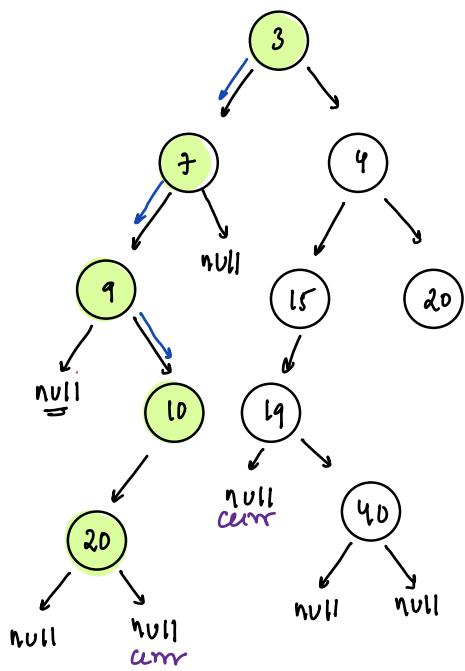
```
if (root == null) { return 0; }

l = size (root.left)
r = size (root.right)

return l+r+1
```



Inorder (Iteratively)



Inorder: 9 20 10 7 3 19 ... soon

TC: O(N) SC: O(H)

void inorder(Node root) {

Node curr = root

Stack<Node> st; // Storing address

while(st.size() > 0 || curr != NULL) {

→ curr = 4

15

4

20

10

19

40

15

4

20

10

19

40

while(curr != null) {

st.push(curr)

curr = curr.left

curr = st.top() → curr = 19

st.pop() // Deleting from stack

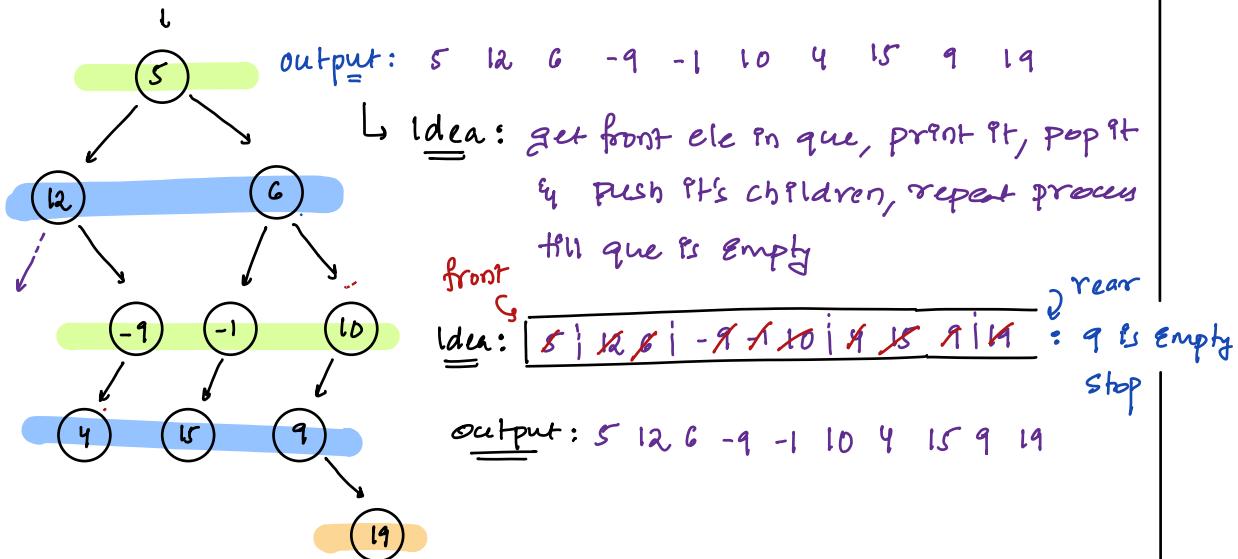
print(curr.data) -

curr = curr.right → curr = 40

}

preorder / postorder : TODO / Iterative with Stack

level order traversal : level by level, left to right



void levelorder(Node root) { TC: O(N) SC: max nodes in a level?

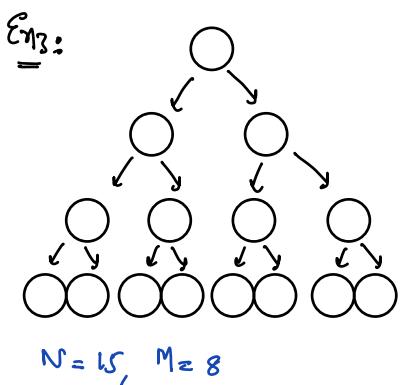
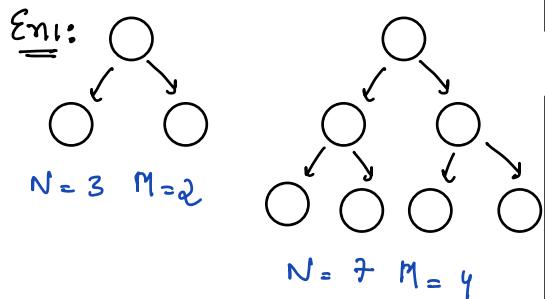
```
if (root == null) { return; }

Queue<Node> q
q.enqueue(root)

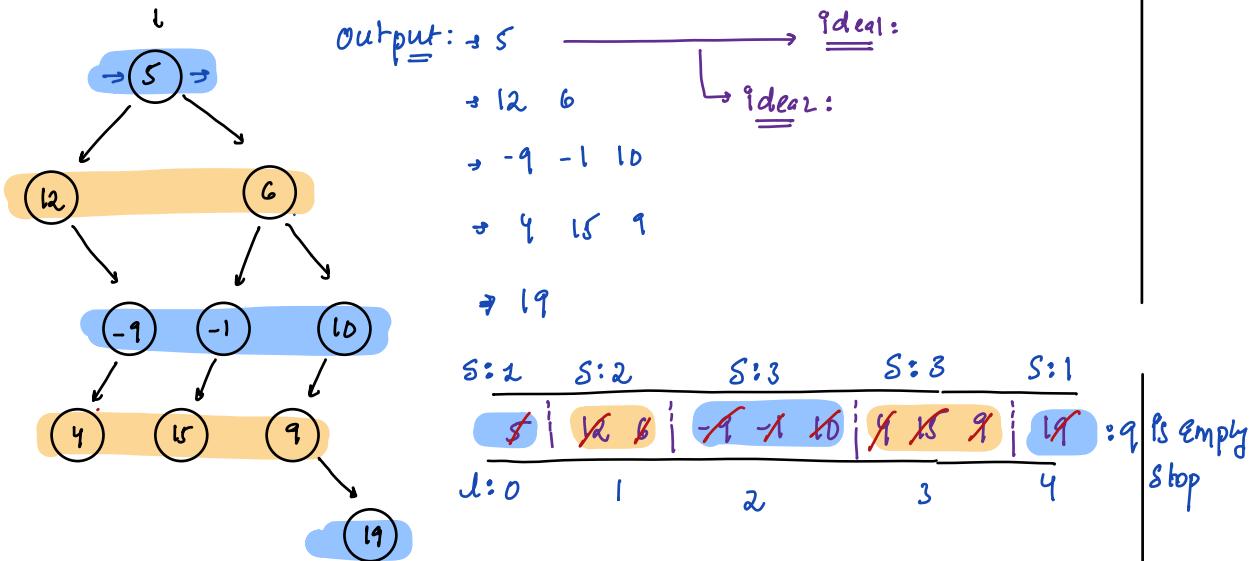
while (q.size() > 0) {
    Node f = q.front(),
    q.dequeue()
    print(f.data);

    if (f.left != null) {
        q.enqueue(f.left)
    }
    if (f.right != null) {
        q.enqueue(f.right)
    }
}
```

$$S = \frac{N+1}{2} = O(N)$$



level order traversal: 2 → { After every level goto new line }



void levelorder(Node root) { TC: O(N) SC: max nodes in a level ?

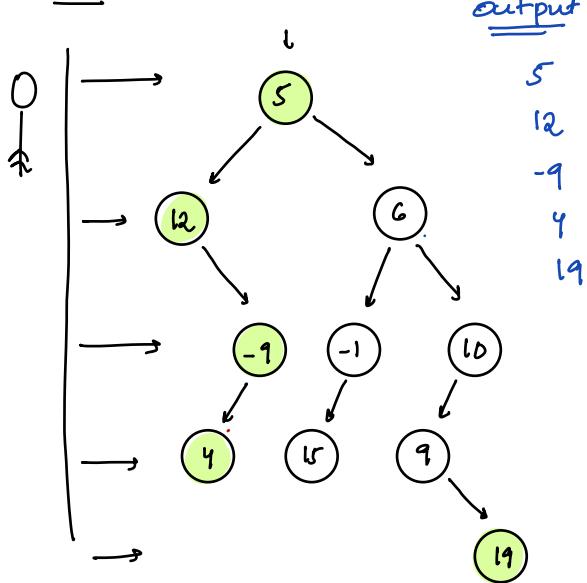
```

if (root == null) { return }
Queue Node > q
q.enqueue(root)
while ( q.size() > 0 ) {
    int n = q.size()
    for ( i = 1; i <= n; i++ ) {
        Node f = q.front()
        q.dequeue()
        print(f.data)
        if ( f.left != null ) {
            q.enqueue(f.left)
        }
        if ( f.right != null ) {
            q.enqueue(f.right)
        }
    }
    print("\n") // new line
}

```

size	
1	pop que, push children 1 time, print(new line)
2	pop que, push children 2 time print(new line)
3	pop que, push children 3 time print(new line)
3	pop que, push children 3 time print(new line)
1	pop que, push children 1 time print(new line)

left view:



output: 1st node of every level

5
12
-9
4
19

void leftview (Node root) { TC: O(N) Sc: max nodes in a level?

```

if (root == NULL) { return }

Queue Node q
q.enqueue (root)

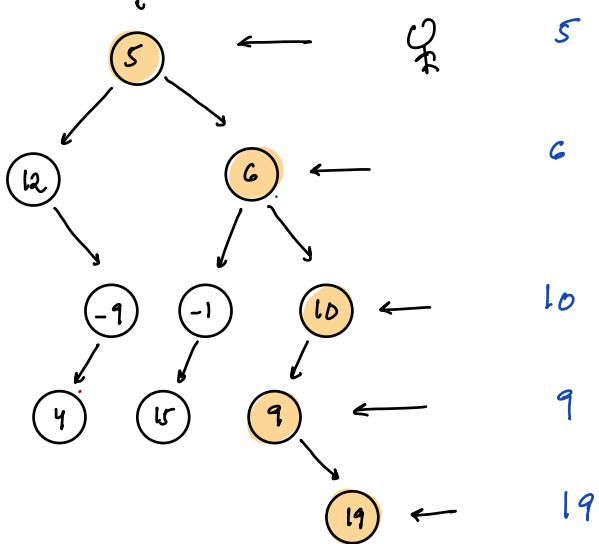
while ( q.size() > 0 ) {
    int n = q.size()

    i=1; i<=n; i++ {
        Node f = q.front()
        q.dequeue() // node will be 1st node
        if (i==1) { print(f.data); }

        if (f.left != NULL) {
            q.enqueue (f.left)
        }
        if (f.right != NULL) {
            q.enqueue (f.right)
        }
    }
    print("\n") // new line
}

```

right view:



output : last node of every level

5

6

10

9

19

```
void rightview (Node root){
```

```
if (root == null) { return; }
```

```
Queue Node > q
```

```
q.enqueue (root)
```

```
while ( q.size() > 0 ) {
```

```
int n = q.size()
```

```
i = 1; i <= n; i++ ) {
```

```
Node f = q.front();
```

```
q.dequeue() // node will be last node
```

```
if (i == n) { print(f.data); }
```

```
if (f.left != null) {
```

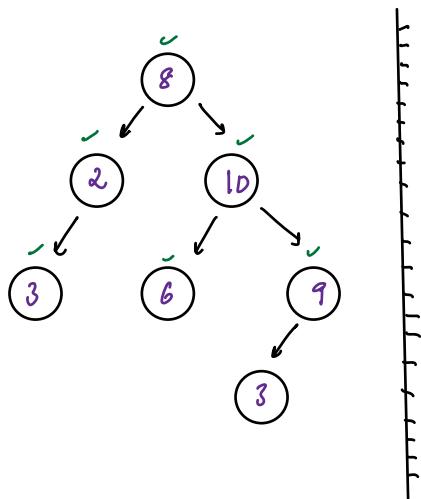
```
q.enqueue (f.left)
```

```
if (f.right != null) {
```

```
q.enqueue (f.right)
```

```
print("\n") // newline
```

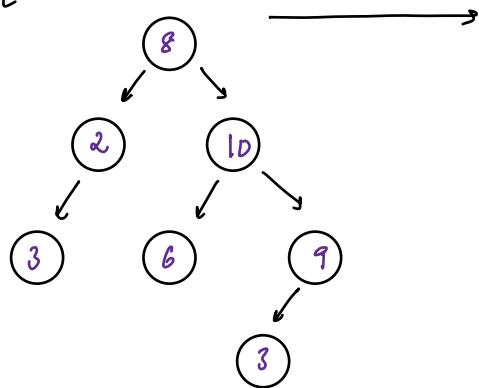
Invert BT / Mirror view of BT : TODO?



Ideal:

Ass:

```
void mirror(Node root){
```



}

—

—

—

—

func (Int arr) L

i = 0; i < n; i++ {

| if ~ e

