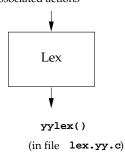
Lex and Yacc: A Brisk Tutorial

Saumya K. Debray Department of Computer Science The University of Arizona Tucson, AZ 85721

Lex: A Scanner Generator

 Helps write programs whose control flow is directed by instances of regular expressions in the input stream.

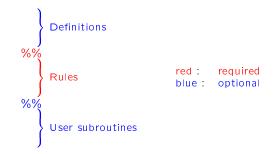
Table of regular expressions + associated actions



- yylex():
 - matches the input stream against the table of regular expressions supplied
 - carries out the associated action when a match is found.

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Structure of Lex Specification File



Rules: line oriented:

 $\langle reg. exp \rangle \langle whitespace \rangle \langle action \rangle$

 $\langle reg. exp \rangle$: starts at beginning of line, continues upto first unescaped whitespace

(action) : a single C statement
 (multiple statements: enclose in braces { }).

unmatched input characters : copied to stdout.

Lex Regular Expressions

Similar to egrep:

- <u>operators</u>: " \ [] ^ ? . * | () \$ / { } % < >
- letters and digits match themselves
- period '.' matches any character (except newline)
- brackets [] enclose a sequence of characters, termed a *character class*. This matches:
 - any character in the sequence
 - a '-' in a character class denotes an inclusive range, e.g.: [0-9] matches any digit.
 - a $\hat{}$ at the beginning denotes <u>negation</u>: [^0-9] matches any character that is <u>not a digit</u>.
- a quoted character " " matches that character. operators can be escaped via \.
- \n, \t match newline, tab.

parentheses () grouping
 bar | alternatives
 star * zero or more occurrences
 one or more occurrence
 zero or one occurrence

Examples of Lex Rules

- int printf("keyword: INTEGER\n");
- [0-9]+ printf("number\n");
- "-"?[0-9]+("."[0-9]+)? printf("number\n");

Choosing between different possible matches:

When more than one pattern can match the input, lex chooses as follows:

- 1. The *longest* match is preferred.
- Among rules that match the same number of characters, the rule that occurs earliest in the list is preferred.

Example: the pattern

(intended to match multi-line comments) may consume all the input!

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Communicating with the user program

 ${\tt yytext}$: a character array that contains the actual string that matched a pattern.

yyleng: the no. of characters matched.

Example:

- [a-z][a-z0-9_]* printf("ident: %s\n", yytext);
- Counting the number of words in a file and their total size:

```
[a-zA-Z]+ {nwords += 1; size += yyleng;}
```

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Lex source definitions

- Any source not intercepted by lex is copied into the generated program:
 - a line that is not part of a lex rule or action, which begins with a blank or tab, is copied out as above (useful for, e.g., global declarations)
 - anything included between lines containing only %{ and %} is copied out as above (useful, e.g., for preprocessor statements that must start in col.1)
 - anything after the second %% delimiter is copied out after the lex output (useful for local function definitions).
- Definitions intended for lex are given before the first %%. Any line in this section that does not begin with a blank or tab, or is not enclosed by %{...%}, is assumed to be defining a lex substitution string of the form

```
name translation
```

E.g.:

```
letter [a-zA-Z]
```

An Example

```
#include "tokdefs.h"
#include <strings.h>
static int id_or_keywd(char *s);
letter
              [a-zA-Z]
digit
              [0-9]
              [a-zA-Z0-9_]
alfa
whitesp
              [ \t\n]
%%
{whitesp}*
{comment}
{letter}{alfa}
                   REPORT(id_or_keywd(yytext), yytext);
%%
static struct {
  char *name;
  int val;
} keywd_entry,
  keywd_table[] = {
                      CHAR
    "char",
    "int",
                      INT.
                      WHILE.
    "while",
static int id_or_keywd(s)
char *s:
}
```

Left Context Sensitivity: Start Conditions

Start conditions are a mechanism for conditionally activating patterns. This is useful for handling

- conceptually different components of an input; or
- situations where the lex defaults (e.g., "longest possible match") don't work well, e.g., comments or quoted strings.

Basic Idea:

- Declare a set of start condition names using

%Start name₁ name₂ ...

- If scn is a start condition name, then a pattern prefixed with <scn> will only be active when the scanner is in start condition scn.
- The scanner begins in start condition INITIAL, of which all non-<scn>-prefixed rules are members.
- Start conditions such as these are <u>inclusive</u>: i.e., being in that start condition adds appropriately prefixed rules to the active rule set.

flex also allows $\underline{exclusive}$ start conditions (declared using x), which are sometimes more convenient.

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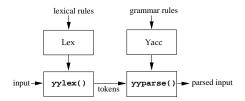
Example of use of start conditions

```
%Start comment0 comment1
#include "tokens.h"
%}
whitespace
               [ \t\n]
               ΓO-9]
digit
intcon
               {digit}+
               {digit}+"."{digit}+
floatcon
start_comment "/""*"
<INITIAL>{start_comment}
                           BEGIN(comment0);
<comment0>"*"
                           BEGIN(comment1);
<comment0>[^*]
<comment1>"*"
<comment1>"/"
                           BEGIN(INITIAL);
<comment1>[^*/]
                           BEGIN(comment0);
                           return(INTCON);
{intcon}
{floatcon}
                           return(FLOATCON);
%%
  . . .
```

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Yacc: A Parser Generator

Takes a specification for a CFG, produces an LALR parser



• Form of a yacc specification file:

```
Declarations

%%

Grammar rules

red: required blue: optional

Programs
```

Yacc: Grammar Rules

Terminals (tokens): Names must be declared:

```
%token name1 name2 ...
```

Any name not declared as a token in the declarations section is assumed to be a nonterminal.

Start symbol:

- may be declared, via: %start name
- if not declared explicitly, defaults to the nonterminal on the LHS of the first grammar rule listed

Productions : A grammar production $A \to B_1 B_2 \cdots B_n$ is written as

$$A: B_1B_2\cdots B_n$$
;

Note: Left-recursion is preferred to right-recursion for efficiency reasons.

Example:

```
stmt : KEYWD_IF '(' expr ')' stmt ;
```

Communication between Scanner and Parser

- The user must supply an integer-valued function <u>yylex()</u> that implements the lexical analyzer (scanner).
- If there is a value associated with the token, it should be assigned to the external variable yylval.
- The token error is reserved for error handling.
- <u>Token numbers</u>: These may be chosen by the user if desired. The default is:
 - chosen by yacc
 - the token no. for a literal is its ASCII value
 - other tokens are assigned numbers starting at 257
 - the endmarker must have a number zero or negative.

Using Yacc

• Suppose the grammar spec is in a file foo.y. Then,

yacc foo.y

yields a file ${\tt y.tab.c}$ containing the parser constructed by yacc.

The command

yacc -v foo.y

additionally constructs a file y.output containing a description of the parser (useful for debugging).

• The uer needs to supply a function main() to driver, and a function yyerror() that will be called by the parser if there is an error in the input.

Conflicts and Ambiguities

- Conflicts may be either *shift/reduce* or *reduce/reduce*:
 - In a shift/reduce conflict, the default is to shift.
 - In a reduce/reduce conflict, the default is to reduce using the first applicable grammar rule.
- Arithmetic Operators: associativity and precedence can be specified:

Associativity: use %left, %right, %nonassoc Precedence (Binary Operators):

- Specify associativity using %left etc.
- Operators within a group have same precedence.
 Between groups, precedence increases going down.

Precedence (Unary Operators): use %prec keyword. This changes the precedence of a rule to be that of the following token.

Example:

Yacc: Error Handling

- The token error is reserved for error handling. This can be used in grammar rules, to indicate where error might occur and recovery take place.
- When an error is detected :
 - If an error token is specified, the parser pops its stack until it finds a state where the error token is legal.

It then behaves as if error is the current lookahead token, and performs the action encountered.

- If there is no rule using the error token, processing halts when an error is encountered.
- To prevent cascading error messages, the parser remains in an error state after detecting an error until 3 tokens have been successfully read and shifted.

If an error is encountered when in the error state, no error message is given, and the input token is discarded.

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Yacc Error Handling: (cont'd)

• A rule of the form

```
stmt : error
```

means that on syntax error, the parser would attempt to skip over the offending statement, looking for 3 tokens that can legally follow ${\tt stmt}$.

• A rule of the form

```
stmt : error ';'
```

causes the parser to skip to the next ';' after ${\tt stmt}$: all intervening tokens are deleted.

 Actions may be associated with these special error rules: these might attempt to (re)initialize tables, reclaim space, turn off code generation, etc.

Adding error symbols

Their placement is guided by the following (conflicting!) goals:

as close as possible to the start symbol of the grammar

(to allow recovery without discarding the entire program)

• as close as possible to each terminal symbol

(to allow only a small amount of input to be discarded on an error)

• without introducing conflicts

(this may be difficult;

shift/reduce conflicts may be acceptable if they serve to lengthen strings, i.e., delay reporting of errors)

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Error Messages

The use should provide a function yyerror() that is called when a syntax error is detected:

More informative error messages :

```
- line no. in source program : yylineno
- token no. causing error : yychar
```

Example:

Controlling error actions

Sometimes we may want to stop discarding tokens, if a certain (synchronizing) token is seen: for this, attach an action {yyerrok;}

Example:

Special-purpose error handling:

- set a global flag to indicate the problem;
- use this flag in yyerror() to give better error messages.

Example:

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