Hash Length Extension Attack Lab

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 $Instruction: \ https://seedsecuritylabs.org/Labs_20.04/Files/Crypto_Hash_Length_Ext/Crypt$

Lab Environment

Set up the container and run it (www-10.9.0.80) in the background:

```
1 curl
          https://seedsecuritylabs.org/Labs_20.04/Files/Crypto_Hash_Length_Ext/Labsetup.zip
          -o Labsetup.zip
2 unzip Labsetup.zip
3 cd Labsetup
4 dcbuild
5 dcup -d
```

If necessary, get the running container id by dockps and use docksh \leq id \geq to start a shell on this container.

Add the following entry in /etc/hosts (root privilege required, try sudo vi /etc/hosts):

1 10.9.0.80 www.seedlab-hashlen.com

Task 1

Construct and send a benign request to the server:

- Pick up a uid with its key value from Labsetup/image_flask/app/LabHome/key.txt instead of using a real name, for example, I choose the entry 1001:123456 in this task.
- 2. Calculate the MAC of the key concatenated with request content R, that is
- 1 Key:R = 123456:myname=koji&uid=1001&lstcmd=1

Suppose that the name used here is "koji" and it requests for listing all the files in LabHome folder.

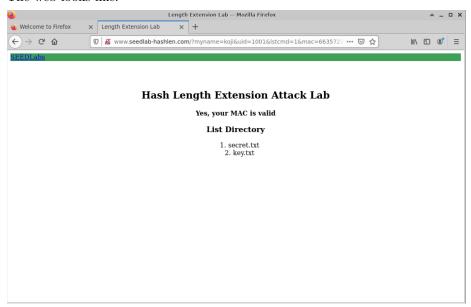
So the MAC is calculated as:

Thus the complete request is:

1 http://www.seedlab-hashlen.com/?myname=koji&uid=1001&lstcmd=1&mac=66357225216e2e9d1eb27b44fc

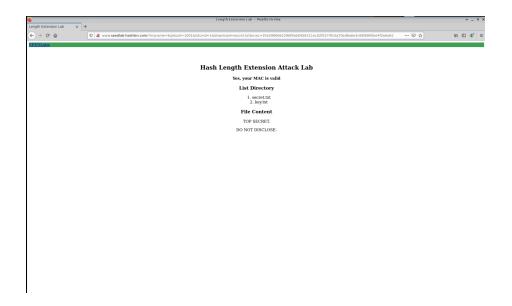
Don't use curl or wget, it doesn't support. Just open a Firefox browser via VNC client and visit the url link above.

The web looks like:



For a download request, we take a similar strategy to construct:

1 http://www.seedlab-hashlen.com/?myname=koji&uid=1001&lstcmd=1&download=secret.txt&mac=35e59



Task 2

Construct the padding for

```
1 123456:myname=koji&uid=1001&lstcmd=1
```

Use Python REPL to complete this work:

Task 3

Compile and run calculate_mac.c, in which SHA256_Update takes the padding bytes we obtained in previous task followed by &download=secret.txt as the second argument.

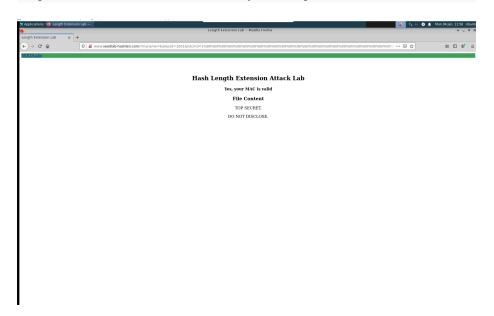
```
gcc calculate_mac.c -o calculate_mac -lcrypto
2 ./calculate_mac

It gives:
1 14797c6db7ca0309d20e0b3c54ac19df60861a83fe64b2713a45e18469b5f3fc

If it reports an error as:
1 gcc: error: calculate_mac.c: No such file or directory

try:
1 sudo apt install libssl-dev
```

Then, visit



Task 4

Alternatively, to distinguish from the existing work, we turn to apply the 1002:983abe as mackey-uid and "MUR" as current username.

A legitimate request to list files without MAC value:

Assume that we have already observed the full request URL as

1 http://www.seedlab-hashlen.com/?myname=MUR&uid=1002&lstcmd=1&mac=3a286321c4cb101ce172c1377a

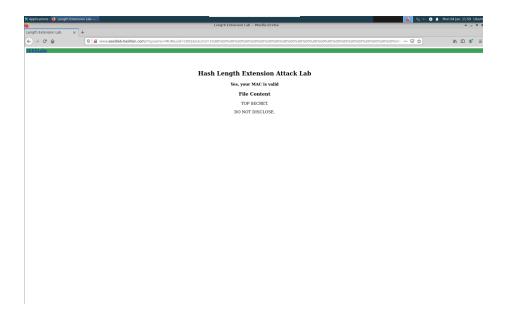
But we do not know the mac key of it. So we use length_ext.c to obtain the MAC after appending "&download=secret.txt" argument. Compile and run:

```
1 gcc length_ext.c -o length_ext -lcrypto
2 ./length_ext
3 #
    bcea031dd94604d9b84e4886aab8e083b6ba2ec66f50316555cf1cb451bf4aed
```

Then, construct the padding of the original message as task-2, recall that we don't know what the mac key exactly is but we know the length of keys are fixed, so we can easily calculate the padding:

So the full request is:

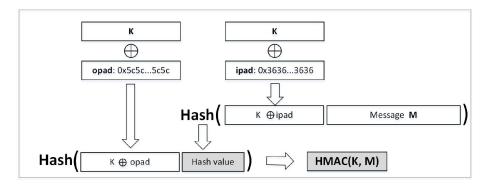
1 http://www.seedlab-hashlen.com/?myname=MUR&uid=1002&lstcmd=1%80%00%00%00%00%00%00%00%00%00%00%00%



Task 5

Keyed-hash mesaage authentication code (HMAC) can be used as the follwing example:

HMAC works as the figure below shows:



H is a hash function and K is a secret key, which could be of any length. B denotes the block size for H.

For an input message M, he inner hash (left part) first computes $H(K \oplus ipad) \parallel M$ and its result h is passed to the outer hash in order to perform $H((K \oplus opad) \parallel h$, in which ipad and opad are both constants.

In such an algorithm, the MAC key is required in both 2 hash functions. The MAC of a full message calculated by inner hash is required taking by the outer hash function, so if we don't have the internal result, which is invisible to us, we cannot compute the correct final MAC. Due to the sequential design, if the server applies HMAC instead of ordinary MAC methods we discussed above, the attacker cannot directly construct the MAC of an extended message from the final MAC of a legal request only. Therefore, hash length extension attack will fail.