

Day-1

Q-1 Problem Name: Set Matrix zeros

1	1	1
1	0	1
1	1	1

 \Rightarrow

1	0	1
0	0	0
1	0	1

// if there is 0 row in a row or col then make entire row and col to zero.

// first approach (Brute force)

\Rightarrow TC: $O(N \times M) \cdot O(N+M)$
SC: $O(1)$

\Rightarrow // Second approach

- Make two dummy array and traverse in matrix if $mat[i][j] = 0$ then set dummy1[i] = 0 and dummy2[j] = 0
 \downarrow
 row \downarrow for col

// Now traverse again $dummy1[i] = 0$ ||



dummy[j] == 0 then
 dummy[i][j] = 0;

// TC: $O(N * m + N * m)$
 SC: $O(N)$

⇒ // Third approach (optimal)

// Instead of taking dummy array
 we use first row and
 first col and checking particular
 row or col has 0 or not.

// now traverse the remaining
 matrix and take variable col0 = 0
 value in reverse dir.

```
for (int i = row - 1; i >= 0; i--)
{
    for (int j = cols - 1; j >= 1; j--)
    {
        if (matrix[i][0] == 0 ||
            matrix[0][j] == 0)
            matrix[i][j] = 0;
    }
}
```

Remarks



```

if (col == 0)
{
    mat[i][0] = 0;
}

```

Problem Name: Pascal's triangle.

// logic: corner value of row start and row end is one and $mat[i-1][j-1] + mat[i-1][j]$ is sum up for next so remaining value.

```

vector<vector<int>> r(numRows)
int n = numRows;

```

```

for (int i = 0; i < numRows; i++) {

```

```

    r[i].resize(i+1);

```

```

    r[i][0] = r[i][i] = 1;

```

```

    for (int j = 1; j < i; j++)

```

Remarks

```

    {
        r[i][j] = r[i-1][j-1] + r[i-1][j];
    }
    return r;

```



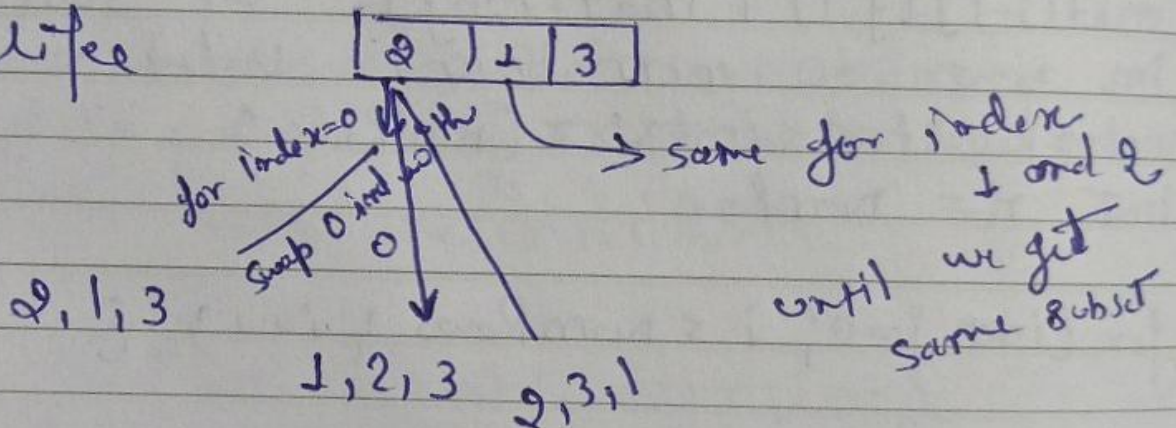
// TC: $O(\text{numRows}^2)$
 SC: $O(\text{numRows}^2)$

// creating 2D-Array.

Problem Name: Next-Permutation

// First approach using recursion
 find all subset

Life



// TC: $O(N! \times N)$
 SC: $O(1)$

Remarks

// Second approach:
 // using inbuilt STL fn



next-permutation (nums.begin(), nums.end())

// Tc: $O(N \log N)$

// Sc: $O(1)$

// Third approach (optional)

// Find break point in a given array.

int k = n - 2;

// if break point (num[k] < num[k + 1])
break;

// if (k < 0)
reverse (nums.begin(), nums.end());

// else {

if (num[i] > num[k])
break;

Remarks else

swap (num[k], num[i])

reverse (nums.begin() + k + 1, nums.end())

Tc: $O(N)$
Sc: $O(1)$



Problem Name: Kadane's Algorithm

// Find contiguous subarray.

// arr[] = { -2, 1, -3, 4, -1, 2, 1, -5, 4 }

// output: 6

// logic: create a variable sum = 0;
traverse whole array.

like

sum += arr[i];

// if (sum < 0) { base }
sum = 0;

TC: $O(N)$

SC: $O(1)$

Remarks

// Brute force using loop (3 pointer)

i = 0 - n

$O(N^3)$

j = 0 - n - 1

k = i to j