

Design & Analysis of Algorithms

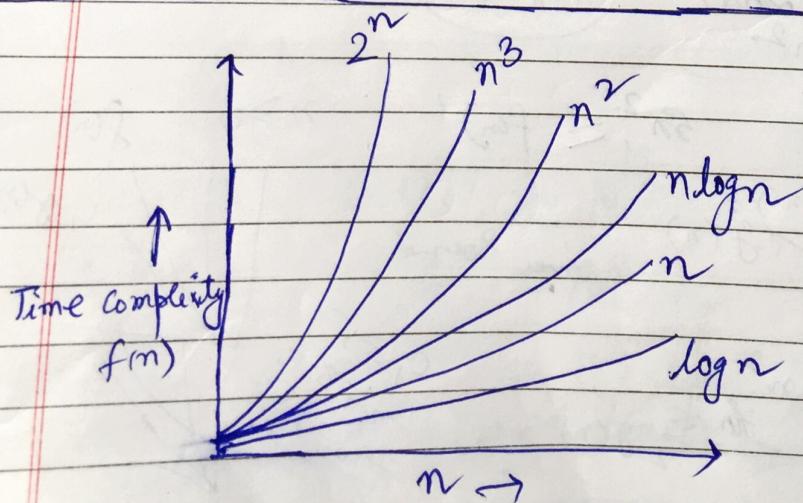
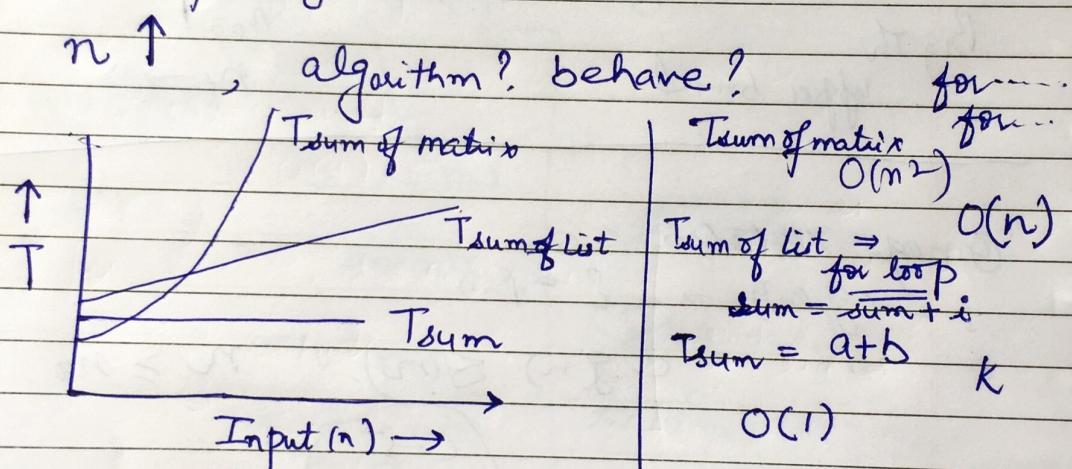
(Prof. Parminder Kaur Wadhwa)

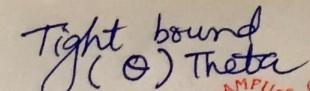
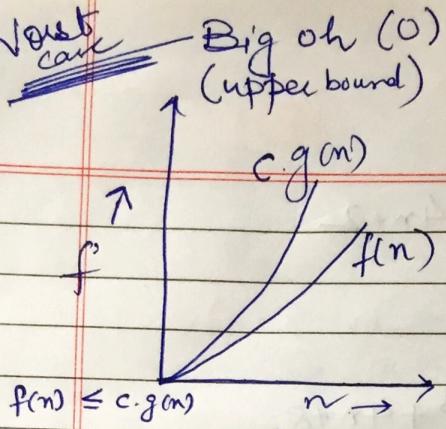
Introduction :-

Algorithm and its importance, mathematical foundations: — Growth functions, Complexity Analysis of Algorithms

Algorithm

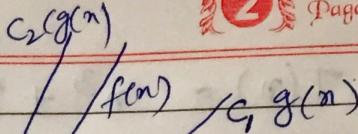
Problem P $\xrightarrow{\text{A1}}$? choose $\xrightarrow{\text{A2}}$
 $n \rightarrow$ input size :



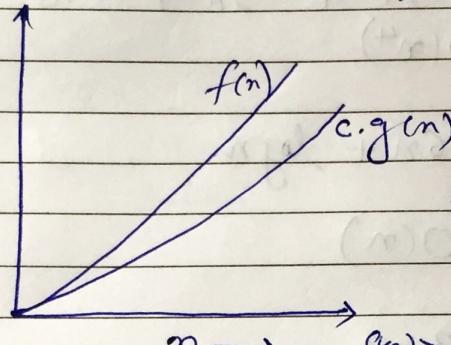


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$$c_1 g(n) \leq f(n) \leq c_2 g(n)$$



Lower Bound
(-2)

Omega
Bait case

$$\begin{aligned} f(n) &= 5n^2 + 2n + 1 \\ g(n) &= n^2 \\ c &= 8 \quad = (5+2+1) \end{aligned}$$

$$2n \neq 2n^2$$
$$1 \neq n^2$$

$O(n^2)$

$$f(n) \leq 8n^2$$

for ($n \geq 1$)

$$f(n) = 5n^2 + 2n + 1 = \mathcal{O}(n^2)$$

$n \geq 0$

$$g(n) = n^2$$

$$2n+1 \geq 1$$

$$\frac{C=5}{n_0=0}$$

$$5n^2 \leq f(n), \quad n \geq 0$$

$$5n^2 \leq f(n) \leq 8n^2$$

$$f(n) = \Theta(n^2)$$

$$\begin{aligned}C_1 &= 5 \\C_2 &= 8 \\n_0 &= 1\end{aligned}$$

Examples:-



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1. $T(n) = n^3 + 3n^2 + 4n + 2$
 $= O(n^3)$

2. $T(n) = 17n^4 + 3n^3 + 4n + 8$
 $= O(n^4)$

3. $T(n) = 16n + \log n$
 $= O(n)$

4. $T(n) = 3 + 4$
 $\Rightarrow O(1)$

5. $T(n) = a+b$
 $= O(1)$

	<u>times</u>	<u>cost</u>
6. $\text{for } (i=0; i \leq n; i++)$ { // stmts }	(n+1) times --- n times	c ₁ c ₂

$$T(n) = c_1(n+1) + c_2 n$$

$$= c_1 n + c_1 + c_2 n$$

$$= (c_1 + c_2)n + c_1$$

\Rightarrow It is in the form:-

$$T(n) = an+b \\ \equiv O(n)$$

e. linear relationship

$$c_1(n+1) + c_2 * n(n+1) + c_3((n)(n+1) - 1)$$

$$= c_1 n + c_1 + c_2 n^2 + c_2 n + c_3 (n^2 + n - 1)$$

$$= c_1 n + c_1 + \underline{c_2 n^2} + c_2 n + \underline{c_3 n^2} + c_3 n - c_3$$

$$= (c_2 + c_3) n^2 + (c_2 + c_3)n + (c_1 - c_3)$$

$$= an^2 + bn + c$$

$$\Rightarrow O(n^2)$$

9. fragments: -

for ($0 \rightarrow n$)
{ }

$\rightarrow (n+1)$ times

$O(n)$

for ($0 \rightarrow n$)

{

for ($0 \rightarrow n$)
{ }

$O(n^2)$

$$T(n) = O(n) + O(n^2)$$

$$T(n) \Rightarrow O(n^2)$$

Analysis of :-

p. Algorithm to compute sum

Cost	Times
c_1	1 time
c_2	$(n+1)$ times
c_3	n times
c_4	1 time

$$T(n) = c_1 * 1 + c_2 * (n+1) + c_3 * n + c_4 * 1$$

$$= c_1 + \underline{c_2 n} + c_2 + \underline{c_3 n} + c_4$$

$$= (c_2 + c_3)n + (c_1 + c_2 + c_4)$$

\Rightarrow In the form

$$an+b$$

$$= O(n)$$

11. Analysis of Insertion Sort

{ for $j \leftarrow 2$ to length [A] $c_1 \ n$
 do key $\leftarrow A[j]$ $c_2 \ n-1$
 $i \leftarrow j-1$ $c_3 \ n-1$

while ($i > 0$ and $A[i] > \text{key}$) $c_4 \sum_{j=2}^n K$

{ $c_5 \ n$
 do $A[i+1] \leftarrow A[i]$ $c_6 \sum_{j=2}^{K-1} K-1$
 $i \leftarrow i-1$ $c_7 \sum_{j=2}^n K-1$

} $A[i+1] \leftarrow \text{key}$

$c_7 \ n-1$

}

$$\Rightarrow T(n) = c_1 n + c_2(n-1) + c_3(n-1) + c_4 \sum_{j=2}^n K$$

$$+ c_5 \left(\sum_{j=2}^n K-1 \right) + c_6 \left(\sum_{j=2}^n K-1 \right) + c_7(n-1)$$

$$\text{As } \Rightarrow \sum_{j=1}^n = \frac{n(n+1)}{2} \quad 1+2+3+\dots$$

$$\Rightarrow \sum_{j=2}^n \Rightarrow \frac{n(n+1)}{2} - 1 \quad 2+3+\dots$$

$$\begin{aligned} T(n) &= c_1 n + c_2(n-1) + c_3(n-1) + c_4 \left(\frac{n(n+1)}{2} - 1 \right) \\ &+ c_5 \left(\left(\frac{n(n+1)}{2} - 1 \right) - 1 \right) + c_6 \left(\left(\frac{n(n+1)}{2} - 1 \right) - 1 \right) + c_7(n-1) \\ \Rightarrow & an^2 + bn + c \Rightarrow O(n^2) \end{aligned}$$

Practice Questions :-

1. Do the analysis of Selection Sort algorithm.

(Hint : - $O(n^2)$)

2. Do the analysis of bubble sort algorithm.

(Hint: $\rightarrow O(n^2)$)