

Tutorial Sheet - 2

Q1 Compute the Big Oh of the algorithm involving one for loop.

Ans)

Time complexity of a loop is considered as $O(n)$ if the loop variables are incremented / decremented by a constant amount.

Example

	Cost	Times
int n=100	C_1	1
for (int i=0; i<n; i++)	C_2	$n+1$
{ cout << "Hello" ; }	C_3	n

Its Time Complexity will be

$$\begin{aligned}
 T_n &= C_1 * 1 + C_2 * (n+1) + C_3 * (n) \\
 &= C_1 * 1 + C_2 n + C_2 + C_3 n \\
 &= (C_2 + C_3) n + (C_1 + C_2)
 \end{aligned}$$

It is in the form of

$$T_n = an + b$$

So its Time complexity will be

$$\Rightarrow O(n)$$

Q2 Compute the time complexity of the following fragments of code.

(a) for ($i=n$; $i>1$; $i=i/2$)
 {
 statement;
 }

Suppose $i=n$ (given)
 when $i=8$

$$i = \frac{8}{2} = 4 = \frac{n}{2^1}$$

$$4 = \frac{n}{2}$$

$$= \frac{n}{2^2}$$

$$\frac{n}{2^3} \Rightarrow 1$$

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

Taking log on both sides

$$\log_2 n = \log_2 2^k$$

$$\log_2 n = k \log_2 2$$

$$\log_2 n = k$$

$$\log_2 2 = 1$$

So its time complexity is $O(\log_2 n)$

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(b) for (i=n; i>1; i=i*2)
    statements;
}
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When $i=1$ $= 2^0$

$i = 1 * 2 = 2$ $= 2^1$

$i = 2 * 2 = 4$ $= 2^2$

$i = 4 * 2 = 8$ $= 2^3$

which implies

$$\Rightarrow 2^k = n$$

Taking \log_2 on both sides

$$\log_2 2^k = \log_2 n$$

$$k \log_2 2 = \log_2 n$$

$$k \log_2 2 = \log_2 n$$

$$\boxed{\log_2 2 = 1}$$

$$\Rightarrow k = \log_2 n$$

So its time complexity is

$$O(\log_2 n)$$