

Hamiltonian Cycles

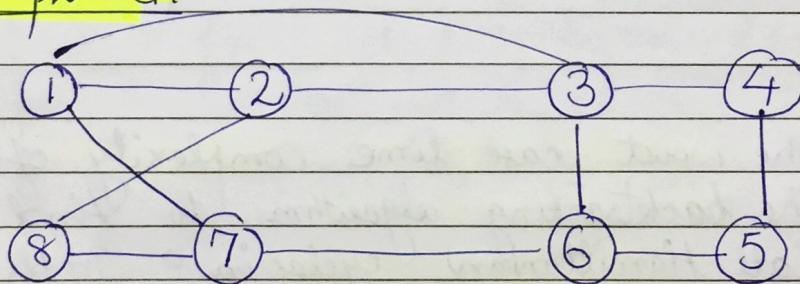
(using backtracking)

Let $G = (V, E)$ be a connected graph with n vertices. A Hamiltonian cycle is a round-trip path along n edges of G that visits every vertex once and returns to its starting position.

(It is same as was the tour in a Travelling Salesperson Problem, but it just finds the cycle and have no concern with costs of edges.)

For example:-

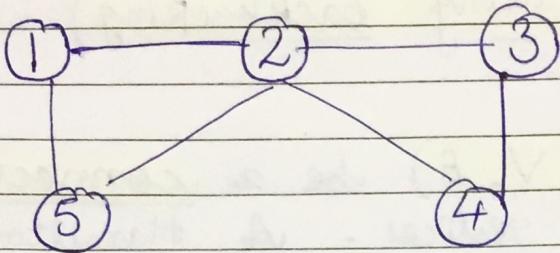
Graph G1:



This graph contains the following Hamiltonian Cycle:-

1, 2, 8, 7, 6, 5, 4, 3, 1

Graph G2 :-



The above graph contains no Hamiltonian cycle.

The backtracking algorithm finds all the Hamiltonian cycles in a graph. The graph may be directed or undirected. Only distinct cycles are output.

The worst case time complexity of the backtracking algorithm to find all Hamiltonian cycles in a graph is
 $\Rightarrow O(n!)$

A backtracking algorithm to find all
the Hamiltonian Cycles in a given graph

G,

Void Hamiltonian (int k)

// The graph is stored as an adjacency
// matrix G [1:n] [1:n]. All cycles begin
// at node 1.

{

do

{

NextValue (k); // Assign a legal
next value (k) to
x[k]

if (! x[k]) return; // no vertex
is possible

if (k == n) // All vertices are
visited

True {
for (int i=1; i<=n; i++)
cout << x[i] << " "; cout << "1";
cout << endl;

False }

else Hamiltonian (k+1);

end the
cycle
at
vertex 1

} while (1); recursive
call

}

The detail of NextValue function for Hamiltonian Cycles problem

Void NextValue(int k)

{
do

{
 $x[R] = (x[R] + 1) \% (n+1); \quad // \text{Next vertex}$

if ($\neg x[k]$) return; $// \text{all vertices are selected}$
 no vertex is left

if ($G[x[k-1]][x[k]]$)

True {

$\rightarrow (k-1)$ edge?
Is there any edge?

for (int j=1; $j \leq k-1$; $j++$)

{ if $x[j] == x[k]$ } break;

\rightarrow
It has already been visited in
the past

reject
the vertex

if ($j == R$)

$// j$ will be R only
if the for loop is
successfully completed
means the vertex has
not appeared in the
past.

{ }

all vertices are
not visited,

All vertices are
visited

CAMPUS No. 2

Date _____

Page _____

if $((k < n) \text{ || } ((k = n) \& \&$

$\in [x[n]] [x[1]]])$

return;

} while (1);

check any
another vertex
to be included
in the cycle.

and there is an
edge from the last
chosen vertex to vertex 1
to complete the cycle

In the above mentioned function,
 $x[k]$ is assigned to the next highest
numbered vertex which:-

(i) does not already appear in
 $x[1], x[2], \dots, x[k-1]$

And

(ii) and is connected by an edge to $x[k-1]$.

$x[k-1] \xrightarrow{\text{edge}} x[k]$
is there
and is not
visited in the past.

Summarizing:- Backtracking Algorithm to find all Hamiltonian Cycles in a graph:-

void Hamiltonian (int k)

{
do {

NextValue (k);

if ($\neg x[k]$) return;

if ($k = n$) // all vertices are visited

{ for (int $i = 1$; $i \leq n$; $i++$)

cout $\ll x[i] \ll " "$;

cout $\ll " 1 " ; cout \ll " \n " ;$

}

else Hamiltonian (k+1);

} while(1);

void NextValue (int k)

do {

$x[k] = (x[k] + 1) \% (n + 1);$

if ($\neg x[k]$) return;

if ($G[x[k-1]][x[k]]$) // if edge is there

true {

for (int $j = 1$; $j \leq k - 1$; $j++$)

if ($x[j] == x[k]$) break;

if ($j == k$)

{ if (($k < n$) || (($k == n$) && $G[x[n]][x[1]]$))

} return;

false

} while(1); // find any other vertex